(as)

# The Complete <br> <br> SPECTRUM <br> <br> SPECTRUM ROM <br> <br> DISASSEMBLY 

 <br> <br> DISASSEMBLY}

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## Preface

The Sinclair ZX Spectrum is a worthy successor to the $Z X 81$ which in turn replaced the $Z X 80$.
The Spectrum has a 16 K monitor program. This program has been developed directly from the 4 K program of the ZX 80 although there are now so many new features that the differences outweigh the similarities.
We have both enjoyed producing this book. We have learnt a great deal about the techniques of Z80 machine code programming and now feel that between us we have unravelled the 'secrets of the Spectrum'. We would like to thank:
-- Our families.
-- Alfred Milgrom, our publisher who has been extremely helpful.
-- Philip Mitchell whose notes on the cassette format were most informative.
-- Clive Sinclair and his team at Sinclair Research Ltd. who have produced such a 'challenging' and useful machine.

January 1983

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## Introduction

The 16 K monitor program of the Spectrum is a complex Z 80 machine code program. Its overall structure is very clear in that it is divided into three major parts:
a. Input/Output routines
b. BASIC interpreter
c. Expression handling

However these blocks are too large to be managed easily and in this book the monitor program is discussed in ten parts.
Each of these parts will now be 'outlined'.

## The restart routines and tables.

At the start of the monitor program are the various 'restart' routines that are called with the single byte 'RST' instructions. All of the restarts are used. For example 'restart 0008' is used for the reporting of syntax or run-time errors.
The tables in this part of the monitor program hold the expanded forms of the tokens and the 'key-codes'.

## The keyboard routine.

The keyboard is scanned every $1 / 50$ th. of a second (U.K. model) and the keyboard routine returns the required character code. All of the keys of the keyboard 'repeat' if they are held down and the keyboard routine takes this into consideration.

## The loudspeaker routines.

The spectrum has a single on-board loudspeaker and a note is produced by repeatedly using the appropriate 'OUT' instruction. In the controller routine great care has been taken to ensure that the note is held at a given 'pitch' throughout its 'duration'.

## The cassette handling routines.

It was a very unfortunate feature of the ZX 81 that so little of the monitor program for that machine was devoted to the cassette handling.
However in the Spectrum there is an extensive block of code and now the high standard of cassette handling is one of the most successful features of the machine.
BASIC programs or blocks of data are both dealt with in the same manner of having a 'header' block (seventeen bytes) that is SAVEd first. This 'header' describes the 'data block' that is SAVEd after it.
One disadvantage of this system is that it is not possible to produce programs with any 'security' whatsoever.

## The screen and printer handling routines.

All of the remaining input/output routines of the Spectrum are 'vectored' through the 'channel \& stream information areas'.
In the standard Spectrum 'input' is only possible from the keyboard but 'output' can be directed to the printer, the upper part of the T.V. display or the lower part of the T.V. display.
The major 'input' routine in this part of the monitor program is the EDITOR that allows the user to enter characters into the lower part of the T.V. display.
The PRINT-OUT routine is a rather slow routine as the same routine is used for 'all possibilities'. For example, the adding of a single byte to the 'display area' involves considering the present status of OVER and INVERSE on every occasion.

## The executive routines

In this part of the monitor program are to be found the INITIALISATION procedure and the 'main execution loop' of the BASIC interpreter.
In the Spectrum the BASIC line returned by the EDITOR is checked for the correctness of its syntax and then saved in the program area, if it was a line starting with a line number, or 'executed' otherwise.
This execution can in turn lead to further statements being considered. (Most clearly seen as in the case of - RUN.)

## BASIC line and command interpretation.

This part of the monitor program considers a BASIC line as a set of statements and in its turn each statement as starting with a particular command. For each command there is a 'command routine' and it is the execution of the machine code in the appropriate 'command routine' that effects the 'interpretation'.

## Expression evaluation

The Spectrum has a most comprehensive expression evaluator allowing for a wide range of variable types, functions and operations.
Once again this part of the monitor is fairly slow as all the possible alternatives have to be considered.
The handling of strings is particularly well managed. All simple strings are managed 'dynamically' and old copies are 'reclaimed' once they are redundant. This means that there is no 'garbage collecting' to be done.

## The arithmetic routines

The Spectrum has two forms for numbers. Integer values in the range -65535 to +65535 are in an 'integral' or 'short' form whilst all other numbers are in a five byte floating point form.
The present version of the monitor is unfortunately marred by two mistakes in this part.
i. There is a mistake in 'division' whereby the 34th bit of a division is lost.
ii. The value of -65536 is sometimes put in 'short' form and at other times in 'floating-point' and this leads to troubles.

## The floating-point calculator

The CALCULATOR of the Spectrum handles numbers and strings and its operations are specified by 'literals'. It can therefore be considered that there is an internal 'stack operating' language in the CALCULATOR.
This part of the monitor program contains routines for all the mathematical functions. The approximations to SIN X, EXP X, LN X \& ATN $X$ are obtained by developing Chebyshev polynomials and full details are given in the appendix.

Overall the 16 K monitor program offers an extremely wide range of different BASIC commands and functions. The programmers have always however been short of 'room' and hence the program is written for 'compactness' rather than 'speed'.

## THE DISASSEMBLY

## THE RESTART ROUTINES and THE TABLES

## THE 'START'

The maskable interrupt is disabled and the DE register pair set to hold the 'top of possible RAM'.

0000 START \begin{tabular}{llll}

DI \& A \& | Disable the 'keyboard interrupt'. |
| :--- |
| +00 for start (but +FF for | <br>

\& XOR \& 'NEW').
\end{tabular}

## THE 'ERROR' RESTART

The error pointer is made to point to the position of the error.

0008 ERROR-1 | LD | HL,(CH-ADD) |  |
| :--- | :--- | :--- |
|  |  | LD |

The address reached by the interpreter is copied to the error pointer before proceeding.

## THE 'PRINT A CHARACTER' RESTART

The A register holds the code of the character that is to be printed.

0010 | PRINT-A-1 JP | JF2,PRINT-A-2 | Jump forward immediately. |  |
| :--- | :--- | :--- | :--- |
|  |  | DEFB | $+F F,+F F,+F F,+F F,+F F$ | Unused locations.

THE 'COLLECT CHARACTER' RESTART
The contents of the location currently addressed by $\mathrm{CH}-\mathrm{ADD}$ are fetched. A return is made if the value represents a printable character, otherwise CH-ADD is incremented and the tests repeated.

| 0018 | GET-CHAR | LD | HL,(CH-ADD) | Fetch the value that is addressed |
| :--- | :--- | :--- | :--- | :--- |
|  |  | LD | A,(HL) | by CH-ADD. |
| 001 C | TEST-CHAR | CALL | O07D,SKIP-OVER | Find out if the character is |
|  |  | RET | NC | printable. Return if it is so. |

THE 'COLLECT NEXT CHARACTER' RESTART
As a BASIC line is interpreted, this routine is called repeatedly to step along the line.

| 0020 NEXT-CHAR CALL | $0074, \mathrm{CH}-A D D+1$ | CH-ADD needs to be incre- |  |
| :--- | :--- | :--- | :--- |
| JR | 001C,TEST-CHAR | mented. <br> Jump back to test the new |  |
|  | DEFB | + FF,+FF,+FF | value. |

## THE 'CALCULATOR' RESTART

The floating point calculator is entered at 335B.
0028 FP-CALC JP 335B,CALCULATE Jump forward immediately.
DEFB $+F F,+F F,+F F,+F F,+F F \quad$ Unused locations.
THE 'MAKE BC SPACES' RESTART
This routine creates free locations in the work space. The number of locations is determined by the current contents of the BC register pair.

| 0030 BC-SPACES | PUSH | BC | Save the 'number'. |
| :--- | :--- | :--- | :--- |
|  | LD | HL,(WORKSP) | Fetch the present address of the |
|  | PUSH | HL | start of the work space and save |
|  | JP | 169E,RESERVE | that also before proceeding. |

THE 'MASKABLE INTERRUPT' ROUTINE
The real time clock is incremented and the keyboard scanned whenever a maskable interrupt occurs.

0038 MASK-INT | PUSH | AF | Save the current values held in |  |
| :--- | :--- | :--- | :--- |
|  | PUSH | HL | these registers. |

| PUSH | $H L$ | these registers. |
| :--- | :--- | :--- |
| LD | $H L,(F R A M E S)$ | The lower two bytes of the |


|  | INC | HL | frame counter are incremented |
| :--- | :--- | :--- | :--- |
| LD | (FRAMES),HL | every 20 ms. (U.K.) The highest |  |
|  | LD | A,H | byte of the frame counter is |
| 0048 | OR | L | only incremented when the |

## THE 'ERROR-2' ROUTINE

The return address to the interpreter points to the 'DEFB' that signifies which error has occurred. This 'DEFB' is fetched and transferred to ERR-NR. The machine stack is cleared before jumping forward to clear the calculator stack.

| 0053 | ERROR-2 | POP | HL | The address on the stack points |
| :--- | :--- | :--- | :--- | :--- |
| 0055 |  | ERROR-3 | LD | L,(HL) |

## THE 'NON-MASKABLE INTERRUPT' ROUTINE

This routine is not used in the standard Spectrum but the code allows for a system reset to occur following activation of the NMI line. The system variable at 5CB0, named here NMIADD, has to have the value zero for the reset to occur.

| 0066 | RESET | PUSH | AF | Save the current values held |
| :---: | :---: | :---: | :---: | :---: |
|  |  | PUSH | HL | in these registers. |
|  |  | LD | HL,(NMIADD) | The two bytes of NMIADD |
|  |  | LD | A, H | must both be zero for the reset |
|  |  | OR | L | to occur. |
|  |  | JR | NZ,0070,NO-RESET | Note: This should have been 'JR Z'! |
|  |  | JP | (HL) | Jump to START. |
| 0070 | NO-RESET | POP | HL | Restore the current values to |
|  |  | POP RETN | AF | these registers and return. |

## THE 'CH-ADD+1' SUBROUTINE

The address held in CH-ADD is fetched, incremented and restored. The contents of the location now addressed by CH-ADD is fetched. The entry points of TEMP-PTR1 and TEMP-PTR2 are used to set CH-ADD for a temporary period.

| 0074 | CH-ADD+1 | LD | HL,(CH-ADD) | Fetch the address. |
| :--- | :--- | :--- | :--- | :--- |
| 0077 | TEMP-PTR1 | INC | HL | Increment the pointer. |
| 0078 | TEMP-PTR2 | LD | (CH-ADD),HL | Set CH-ADD. |
|  |  | LD | A,(HL) | Fetch he addressed value and |
|  |  | RET |  | then return. |

## THE 'SKIP-OVER' SUBROUTINE

The value brought to the subroutine in the A register is tested to see if it is printable. Various special codes lead to HL being incremented once, or twice, and $\mathrm{CH}-\mathrm{ADD}$ amended accordingly.

| 007D SKIP-OVER | CP | +21 | Return with the carry flag reset |
| :--- | :--- | :--- | :--- |
|  | RET | NC | if ordinary character code. |
|  | CP | $+0 D$ | Return if the end of the line |
|  | RET | Z | has been reached. |
|  | CP | +10 | Return with codes +00 to +0 OF |


|  | RET | C | but with carry set. <br> Return with codes +18 to +20 <br> again with carry set. |
| :--- | :--- | :--- | :--- |
|  | CP | +18 |  |
|  | CCF | C |  |
|  | RET | Ckip-over once. |  |

THE TOKEN TABLE
All the tokens used by the Spectrum are expanded by reference to this table. The last code of each token is 'inverted' by having its bit 7 set.

| 0095 | BF | 52 | 4E | C4 | 49 | 4E | 4B | 45 | '?' | R | N | 'D' | I | N | K | E |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 009D | 59 | A4 | 50 | C9 | 46 | CE | 50 | 4 F | Y | '\$' | P | 'I' | F | 'N' | P | 0 |
| 00A5 | 49 | 4E | D4 | 53 | 43 | 52 | 45 | 45 | I | N | 'T' | S | C | R | E | E |
| OOAD | 4E | A4 | 41 | 54 | 54 | D2 | 41 | D4 | N | '\$' | A | T | T | 'R' | A | T' |
| 00B5 | 54 | 41 | C2 | 56 | 41 | 4C | A4 | 43 | T | A | 'B' | V | A | L | '\$' | C |
| 00 BD | 4 F | 44 | C5 | 56 | 41 | CC | 4C | 45 | 0 | D | 'E' | V | A | 'L' | L | E |
| 00C5 | CE | 53 | 49 | CE | 43 | 4F | D3 | 54 | 'N' | S | I | 'N' | C | $\bigcirc$ | 'S' | T |
| 00CD | 41 | CE | 41 | 53 | CE | 41 | 43 | D3 | A | 'N' | A | S | ' ${ }^{\prime}$ ' | A | C | 'S' |
| 00D5 | 41 | 54 | CE | 4C | CE | 45 | 58 | D0 | A | T | 'N' | L | ' ${ }^{\prime}$ ' | E | X | 'P' |
| 00DD | 49 | 4E | D4 | 53 | 51 | D2 | 53 | 47 | I | N | 'T' | S | Q | 'R' | S | G |
| OOE5 | CE | 41 | 42 | D3 | 50 | 45 | 45 | CB | 'N' | A | B | 'S' | P | E | E | 'K' |
| OOED | 49 | CE | 55 | 53 | D2 | 53 | 54 | 52 | I | 'N' | U | S | 'R' | S | T | R |
| O0F5 | A4 | 43 | 48 | 52 | A4 | 4E | 4F | D4 | '\$' | C | H | R | '\$' | N | 0 | T' |
| 00FD | 42 | 49 | CE | 4F | D2 | 41 | 4E | C4 | B | I | 'N' | 0 | 'R' | A | N | D' |
| 0105 | 3C | BD | 3E | BD | 3C | BE | 4C | 49 | $<$ | ' = ' | , | ' = ' | $<$ | '>' | L | 1 |
| 010D | 4E | C5 | 54 | 48 | 45 | CE | 54 | CF | N | 'E' | T | H | E | 'N' | T | $\mathrm{O}^{\prime}$ |
| 0115 | 53 | 54 | 45 | D0 | 44 | 45 | 46 | 20 | S | T | E | 'P' | D | E | F |  |
| 011D | 46 | CE | 43 | 41 | D4 | 46 | 4 F | 52 | F | 'N' | C | A | 'T' | F | 0 | R |
| 0125 | 4D | 41 | D4 | 4D | 4F | 56 | C5 | 45 | M | A | 'T' | M | $\bigcirc$ | V | 'E' | E |
| 012D | 52 | 41 | 53 | C5 | 4 F | 50 | 45 | 4E | R | A | S | 'E' | 0 | P | E | N |
| 0135 | 20 | A3 | 43 | 4C | 4 F | 53 | 45 | 20 |  | '\#' | C | L | 0 | S | E |  |
| 013 D | A 3 | 4D | 45 | 52 | 47 | C5 | 56 | 45 | '\#' | M | E | R | G | 'E' | V | E |
| 0145 | 52 | 49 | 46 | D9 | 42 | 45 | 45 | D0 | R | I | F | 'Y' | B | E | E | P' |
| 014 D | 43 | 49 | 52 | 43 | 4C | C5 | 49 | 4E | C | I | R | C | L | 'E' | I | N |
| 0155 | CB | 50 | 41 | 50 | 45 | D2 | 46 | 4C | 'K' | P | A | P | E | 'R' | F | L |
| 015 D | 41 | 53 | C8 | 42 | 52 | 49 | 47 | 48 | A | S | 'H' | B | R | I | G | H |
| 0165 | D4 | 49 | 4E | 56 | 45 | 52 | 53 | C5 | 'T' | I | N | V | E | R | S | 'E' |
| 016 D | 4F | 56 | 45 | D2 | 4F | 55 | D4 | 4C | 0 | V | E | 'R' | 0 | U | 'T' | L |
| 0175 | 50 | 52 | 49 | 4E | D4 | 4C | 4C | 49 | P | R | I | N | 'T' | L | L | I |
| 017 D | 53 | D4 | 53 | 54 | 4 F | D0 | 52 | 45 | S | 'T' | S | T | $\bigcirc$ | 'P' | R | E |
| 0185 | 41 | C4 | 44 | 41 | 54 | C1 | 52 | 45 | A | 'D' | D | A | T | 'A' | R | E |
| 018 D | 53 | 54 | 4 F | 52 | C5 | 4E | 45 | D7 | S | T | 0 | R | 'E' | N | E | 'W' |
| 0195 | 42 | 4F | 52 | 44 | 45 | D2 | 43 | 4 F | B | 0 | R | D | E | 'R' | C | 0 |
| 019 D | 4E | 54 | 49 | 4E | 55 | C5 | 44 | 49 | N | T | I | N | U | 'E' | D | I |
| 01A5 | CD | 52 | 45 | CD | 46 | 4 F | D2 | 47 | 'M' | R | E | 'M' | F | $\bigcirc$ | 'R' | G |
| 01AD | 4 F | 20 | 54 | CF | 47 | 4F | 20 | 53 | 0 |  | T | 'O' | G | 0 |  | S |
| 01B5 | 55 | C2 | 49 | 4E | 50 | 55 | D4 | 4C | U | 'B' | I | N | P | U | 'T' | L |
| 01 BD | 4 F | 41 | C4 | 4C | 49 | 53 | D4 | 4C | 0 | A | 'D' | L | I | S | 'T' | L |
| 01C5 | 45 | D4 | 50 | 41 | 55 | 53 | C5 | 4E | E | 'T' | P | A | U | S | 'E' | N |
| 01CD | 45 | 58 | D4 | 50 | 4F | 4B | C5 | 50 | E | X | 'T' | P | 0 | K | 'E' | P |
| 01D5 | 52 | 49 | 4E | D4 | 50 | 4C | 4 F | D4 | R | I | N | 'T' | P | L | $\bigcirc$ | 'T' |
| 01 DD | 52 | 55 | CE | 53 | 41 | 56 | C5 | 52 | R | U | 'N' | S | A | V | 'E' | R |
| 01E5 | 41 | 4E | 44 | 4 F | 4D | 49 | 5A | C5 | A | N | D | 0 | M | I | Z | 'E' |
| 01ED | 49 | C6 | 43 | 4C | D3 | 44 | 52 | 41 | I | 'F' | C | L | 'S' | D | R | A |
| 01F5 | D7 | 43 | 4C | 45 | 41 | D2 |  | 45 | 'W' | C | L | E | A | 'R' | R | E |
| 01FD | 54 | 55 | 52 | CE |  |  |  | D9 | T | U | R | 'N' | C | 0 |  |  |

## THE KEY TABLES

There are six separate key tables. The final character code obtained depends on the particular key pressed and the 'mode' being used.
(a) The main key table - L mode and CAPS SHIFT.


| (b) Extended mode. Letter keys and unshifted |  |  |  |  |  |  |  |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 022 C3 | C4 | E0 | E4 | READ | BIN | LPRINT | DATA |
| 0230 | B4 | BC | BD | BB | TAN | SGN | ABS | SQR

(c) Extended mode. Letter keys and either shift.

| 0246 | 7E | DC | DA | 5C | $\sim$ | BRIGHT | PAPER |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 024A | B7 | 7B | DD | D8 | ATN | $\{$ | $\}$ |
| 024E | BF | AE | AA | AB | IN | VAL\$ | SCREEN\$ | CIRCLE

(d) Control codes. Digit keys and CAPS SHIFT.

| 0260 | $0 C$ | 07 | 06 | 04 | DELETE | EDIT | CAPS LOCK | TRUE VIDEO |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 0264 | 05 | 08 | $0 A$ | OB | INV VIDEO | Cursor left | Cursor down | Cursor up |
| 0268 | 09 | $0 F$ |  | Cursor right | GRAPHICS |  |  |  |

(e) Symbol code. Letter keys and symbol shift.

(f) Extended mode. Digit keys and symbol shift.

| 0284 | D0 | CE A8 CA | FORMAT | DEF FN | FN | LINE |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 0288 | D3 | D4 | D1 | D2 | OPEN | CLOSE |
| 028C | A9 | CF |  |  | POINT | CAT |

## THE KEYBOARD ROUTINES

## THE 'KEYBOARD SCANNING' SUBROUTINE

This very important subroutine is called by both the main keyboard subroutine and the INKEY\$ routine (in SCANNING).
In all instances the E register is returned with a value in the range of +00 to +27 , the value being different for each of the forty keys of the keyboard, or the value +FF , the no-key.
The D register is returned with a value that indicates which single shift key is being pressed. If both shift keys are being pressed then the D and E registers are returned with the values for the CAPS SHIFT and SYMBOL SHIFT keys respectively.
If no keys is being pressed then the DE register pair is returned holding +FFFF.
The zero flag is returned reset if more than two keys are being pressed, or neither key of a pair of keys is a shift key.

028E KEY-SCAN LD \begin{tabular}{ll}

L,+2F \& | The initial key value for each |
| :--- |
| line will be $+2 F,+2 E, \ldots,+28$. |
| (Eight lines.) | <br>

\& LD <br>
\& LD
\end{tabular}

Now enter a loop. Eight passes are made with each pass having a different initial key value and scanning a different line of five keys. (The first line is CAPS SHIFT, $\mathrm{Z}, \mathrm{X}, \mathrm{C}, \mathrm{V}$.)


| CP | +19 | Accept the key value for a pair |
| :---: | :---: | :---: |
| RET | Z | of keys if the 'D' key is SYMBOL SHIFT. |
| LD | A, E | It is however possible for the ' E ' |
| LD | E, D | key of a pair to be SYMBOL |
| LD | D,A | SHIFT - so this has to be |
| CP | +18 | considered. |
| RET |  | Return with the zero flag set if it was SYMBOL SHIFT and 'another key'; otherwise reset. |

## THE 'KEYBOARD' SUBROUTINE

This subroutine is called on every occasion that a maskable interrupt occurs. In normal operation this will happen once every 20 ms . The purpose of this subroutine is to scan the keyboard and decode the key value. The code produced will, if the 'repeat' status allows it, be passed to the system variable LAST-K. When a code is put into this system variable bit 5 of FLAGS is set to show that a 'new' key has been pressed.

| 02BF KEYBOARD | CALL | O28E,KEY-SCAN |
| :--- | :--- | :--- |
| RET | NZ |  |$\quad$| Fetch a key value in the DE |
| :--- |
| register pair but return immedi- |
| ately if the zero pair flag is reset. |

A double system of 'KSTATE system variables' (KSTATE0 - KSTATE 3 and KSTATE4 - KSTATE7) is used from now on. The two sets allow for the detection of a new key being pressed (using one set) whilst still within the 'repeat period' of the previous key to have been pressed (details in the other set).
A set will only become free to handle a new key if the key is held down for about $1 / 10$ th. of a second. i.e. Five calls to KEYBOARD.

| $02 \mathrm{C6}$ | K-ST-LOOP | LD | HL,KSTATE0 | Start with KSTATE0. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | BIT | 7,(HL) | Jump forward if a 'set is free'; |
|  |  | JR | NZ,02D1,K-CH-SET | i.e. KSTATE0/4 holds +FF. |
|  |  | INC | HL | However if the set is not free |
|  |  | DEC | (HL) | decrease its '5 call counter' |
|  |  | DEC | HL | and when it reaches zero signal |
|  |  | JR | NZ,02D1,K-CH-SET | the set as free. |
|  |  | LD | (HL),+FF |  |

After considering the first set change the pointer and consider the second set.

| 02D1 K-CH-SET | LD | A,L | Fetch the low byte of the |
| :--- | :--- | :--- | :--- |
|  | LD | HL,+KSTATE4 | address and jump back if the |
|  | CP | L | second set has still to be |
|  | JR | NZ,02C6,K-ST-LOOP | considered. |

Return now if the key value indicates 'no-key' or a shift key only.

| CALL | 031E,K-TEST | Make the necessary tests and <br> RET |
| :--- | :--- | :--- |
|  |  | return if needed. Also change <br> the key value to a 'main code'. |

A key stroke that is being repeated (held down) is now separated from a new key stroke.

| LD | HL,+KSTATE0 | Look first at KSTATE0. |
| :--- | :--- | :--- |
| CP | (HL) | Jump forward if the codes |
| JR | Z,0310,K-REPEAT | match - indicating a repeat. |
| EX | DE,HL | Save the address of KSTATE0. |
| LD | HL,+KSTATE4 | Now look at KSTATE4. |
| CP | (HL) | Jump forward if the codes |
| JR | Z,0310,K-REPEAT | match - indicating a repeat. |

But a new key will not be accepted unless one of the sets of KSTATE system variables is 'free'.

| BIT | 7,(HL) | Consider the second set. |
| :--- | :--- | :--- |
| JR | NZ,02F1,K-NEW | Jump forward if 'free'. |
| EX | DE,HL | Now consider the first set. |


| BIT | 7,(HL) | Continue if the set is 'free' but <br> exit from the KEYBOARD |
| :--- | :--- | :--- |
| RET | $Z$ | exbroutine if not |

The new key is to be accepted. But before the system variable LAST-K can be filled, the KSTATE system variables, of the set being used, have to be initialised to handle any repeats and the key's code has to be decoded

02F1 K-NEW

| LD | E,A |
| :--- | :--- |
| LD | $(\mathrm{HL}), \mathrm{A}$ |
| INC | HL |
| LD | $(\mathrm{HL}),+05$ |
| INC | HL |
| LD | $\mathrm{A},(\mathrm{REPDEL})$ |
| LD | $(\mathrm{HL}), \mathrm{A}$ |
| INC | HL |

The code is passed to the E register and to KSTATEO/4.
The ' 5 call counter' for this set is reset to ' 5 '.
The third system variable of the set holds the REPDEL value
(normally 0.7 secs.).
Point to KSTATE3/7.
The decoding of a 'main code' depends upon the present state of MODE, bit 3 of FLAGS and the 'shift byte'.

| LD | C,(MODE) | Fetch MODE. |
| :--- | :--- | :--- |
| LD | D,(FLAGS) | Fetch FLAGS. |
| PUSH | HL | Save the pointer whilst the |
| CALL | 0333, K-DECODE | 'main code' is decoded. |
| POP | HL |  |
| LD | (HL),A | The final code value is saved in <br>  |
|  |  | KSTATE3/7; from where it is <br> collected in case of a repeat. |

The next three instruction lines are common to the handling of both 'new keys' and 'repeat keys'.
0308 K-END

| LD | (LAST-K),A |
| :--- | :--- |
| SET | 5,(FLAGS) |
| RET |  |

Enter the final code value into LAST-K and signal 'a new key'. Finally return.

THE 'REPEATING KEY' SUBROUTINE
A key will 'repeat' on the first occasion after the delay period - REPDEL (normally 0.7 secs.) and on subsequent occasions after the delay period - REPPER (normally 0.1 secs.).

0310 K-REPEAT \begin{tabular}{lll}
INC <br>
LD

$\quad$

HL <br>
(HL),+05

$\quad$

Point to the '5 call counter' <br>
of the set being used and reset <br>
it to '5'. <br>
Point to the third system vari- <br>
able - the REPDEL/REPPER <br>
value, and decrement it.
\end{tabular}

## THE 'K-TEST' SUBROUTINE

The key value is tested and a return made if 'no-key' or 'shift-only'; otherwise the 'main code' for that key is found.
031E
K-TEST

| LD | B,D |
| :--- | :--- |
| LD | $\mathrm{D},+00$ |
| LD | A, E |
| CP | +27 |
| RET | NC |

Copy the shift byte.
Clear the D register for later.
Move the key number.
Return now if the key was
'CAPS SHIFT' only or 'no-key'.

| CP | +18 | Jump forward unless the 'E' <br> JR |
| :--- | :--- | :--- |
| NZ,032C, K-MAIN | key was SYMBOL SHIFT. |  |
| BIT | 7,B | However accept SYMBOL SHIFT |
| RET | NZ | and another key; return with |
|  |  | SYMBOL SHIFT only. |

The 'main code' is found by indexing into the main key table.

| 032C | K-MAIN | LD | HL,+0205 |
| :--- | :--- | :--- | :--- | | The base address of the table. |
| :--- |
|  |
|  |
|  |
|  |
|  |
|  |
|  |
|  |
|  |
|  |
|  |
|  |
|  |
| LDD |
| REF |

## THE 'KEYBOARD DECODING' SUBROUTINE

This subroutine is entered with the 'main code' in the E register, the value of FLAGS in the D register, the value of MODE in the C register and the 'shift byte' in the B register.
By considering these four values and referring, as necessary, to the six key tables a 'final code' is produced. This is returned in the A register.

0333
K-DECODE

| LD | A,E |
| :--- | :--- |
| CP | $+3 A$ |
| JR | C,0367,K-DIGIT |
|  |  |
| DEC | C |
| JP | M,034F,K-KLC-LET |
| JR | Z,0341,K-E-LET |

Copy the 'main code'. Jump forward if a digit key is being considered; also SPACE, ENTER \& both shifts.
Decrement the MODE value.
Jump forward, as needed, for modes 'K', 'L', 'C' \& 'E'.

Only 'graphics' mode remains and the 'final code' for letter keys in graphics mode is computed from the 'main code'.

| ADD | $\mathrm{A},+4 \mathrm{~F}$ | Add the offset. <br> RET |
| :--- | :--- | :--- |
| Return with the 'final code'. |  |  |

Letter keys in extended mode are considered next.

| 0341 | K-E-LET | LD | HL,+01EB |
| :--- | :--- | :--- | :--- |$\quad$| The base address for table 'b'. |
| :--- |
|  |
|  |
|  |
|  |
|  |
|  |
|  |
|  |
| INC |$\quad$| LD |
| :--- |
| LD |

Key tables 'b-f' are all served by the following look-up routine. In all cases a 'final code' is found and returned.

| 034A | K-LOOK-UP | LD | D,+00 |
| :--- | :--- | :--- | :--- | | Clear the D register. |
| :--- |
|  |
|  |
|  |
|  |
|  |
|  |
|  |
|  |
|  |
|  |
|  |
|  |
| LDD |

Letter keys in 'K', 'L' or 'C' modes are now considered. But first the special SYMBOL SHIFT codes have to be dealt with.

| 034F | K-KLC-LET | LD <br> BIT <br> JR <br> BIT <br> JR <br> BIT <br> RET <br> INC <br> RET <br> ADD <br> RET | $\begin{aligned} & \text { HL,+0229 } \\ & 0, \mathrm{~B} \\ & \text { Z,034A,K-LOOK-UP } \\ & 3, \mathrm{D} \\ & \text { Z,0364,K-TOKENS } \\ & 3,(F L A G S 2) \\ & \text { NZ } \\ & \text { B } \\ & \text { NZ } \\ & \text { A,+20 } \end{aligned}$ | The base address for table 'e' Jump back if using the SYMBOL SHIFT key and a letter key. Jump forward if currently in 'K' mode. <br> If CAPS LOCK is set then return with the 'main code' Also return in the same manner if CAPS SHIFT is being pressed. However if lower case codes are required then +20 has to be added to the 'main code' to give |
| :---: | :---: | :---: | :---: | :---: |

The 'final code' values for tokens are found by adding +A5 to the 'main code'.

Add the required offset and return.

Next the digit keys; and SPACE, ENTER \& both shifts; are considered.

0367 K-DIGIT \begin{tabular}{lll}
CP <br>
RET

$\quad$

+30 <br>
DEC

$\quad$

Proceed only with the digit keys. <br>
i.e. Return with SPACE (+20),
\end{tabular}

The digit keys ' 0 ' to ' 7 ' in extended mode are to give either a 'paper colour code' or an 'ink colour code' depending on the use of the CAPS SHIFT.

| SUB | +20 | Reduce the range +30 to +37 <br> giving +10 to +17. |
| :--- | :--- | :--- |
| INC | B | Return with this 'paper colour <br> code' if the CAPS SHIFT is |
| RET | $Z$ | not being used. |
| ADD | A,+08 | But if it is then the range is to <br> be +18 to $+1 F$ instead - indicat- <br> ing an 'ink colour code'. |
| RET |  | ind |

The digit keys ' 8 ' and ' 9 ' are to give 'BRIGHT' \& 'FLASH' codes.

| 0382 K-8-\&-9 | SUB | +36 | $+38 \&+39$ go to $+02 \&+03$. |
| :--- | :--- | :--- | :--- |
|  | INC | B | Return with these codes if CAPS |
|  | RET | Z | SHIFT is not being used. (These <br> are 'BRIGHT' codes.) |
|  |  | ADD | Subtract '2' is CAPS SHIFT is |
|  | RET |  | being used; giving +00 \& +01 (as |
|  |  |  |  |
|  |  |  |  |

The digit keys in graphics mode are to give the block graphic characters ( +80 to +8 F ), the GRAPHICS code (+0F) and the DELETE code (+0C)

| 0389 K-GRA-DGT | LD | HL,+0230 | The base address of table 'd'. |
| :--- | :--- | :--- | :--- |
|  | CP | +39 | Use this table directly for |
| JR | Z,034A,K-LOOK-UP | both digit key '9' that is to give |  |
| CP | +30 | GRAPHICS, and digit key ' 0 ' |  |
|  | JR | Z,034A,K-LOOK-UP | that is to give DELETE. |
| AND | +07 | For keys '1' to '8' make the |  |
| ADD | A, +80 | range +80 to +87. |  |
| INC | B | Return with a value from this |  |
| RET | Z | range if neither shift key is |  |
|  |  | being pressed. |  |
|  | XOR | $+0 F$ | But if 'shifted' make the range |
| RET |  | +88 to $+8 F$. |  |

Finally consider the digit keys in 'K', 'L' \& 'C' modes.

| 039D K-KLC-DGT | INC <br> RET | B | Return directly if neither shift <br> key is being used. (Final codes <br> +30 to +39.$)$ |
| :--- | :--- | :--- | :--- |
|  |  |  | Use table 'd' if the CAPS |
|  | BIT | $5, B$ | SHIFT key is also being |
|  | LD | HL,+0230 | NZ,034A,K-LOOK-UP |

The codes for the various digit keys and SYMBOL SHIFT can now be found.

|  |  | SUB | +10 | Reduce the range to give +20 to +29. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | CP | +22 | Separate the '@' character |
|  |  | JR | Z,03B2,K-@-CHAR | from the others. |
|  |  | CP | +20 | The '-' character has also to be separated. |
|  |  | RET | NZ | Return now with the 'final codes' +21, +23 to +29. |
|  |  | $\begin{aligned} & \text { LD } \\ & \text { RET } \end{aligned}$ | A, +5 F | Give the '-' character a code of +5 F . |
| 03B2 | K-@-CHAR | $\begin{aligned} & \text { LD } \\ & \text { RET } \end{aligned}$ | A, +40 | Give the '@' character a code of +40 . |

## THE LOUDSPEAKER ROUTINES

The two subroutines in this section are the BEEPER subroutine, that actually controls the loudspeaker, and the BEEP command routine.
The loudspeaker is activated by having D4 low during an OUT instruction that is using port '254'. When D4 is high in a similar situation the loudspeaker is deactivated. A 'beep' can therefore be produced by regularly changing the level of D4.
Consider now the note 'middle C ' which has the frequency 261.63 hz . In order to get this note the loudspeaker will have to be alternately activated and deactivated every $1 / 523.26^{\text {th }}$. of a second. In the SPECTRUM the system clock is set to run at 3.5 mhz . and the note of 'middle C' will require that the requisite OUT instruction be executed as close as possible to every 6,689 T states. This last value, when reduced slightly for unavoidable overheads, represents the 'length of the timing loop' in the BEEPER subroutine.

## THE 'BEEPER' SUBROUTINE

This subroutine is entered with the DE register pair holding the value ' $f$ 'tt', where a note of given frequency ' $f$ ' is to have a duration of ' t ' seconds, and the HL register pair holding a value equal to the number of T states in the 'timing loop' divided by '4'.
i.e. For the note 'middle C' to be produced for one second DE holds +0105 (INT(261.3 * 1)) and HL holds +066A (derived from 6,689/4 30.125).

| 03B5 | BEEPER | DI |  | Disable the interrupt for the duration of a 'beep'. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | LD | A, L | Save L temporarily. |
|  |  | SRL | L | Each ' 1 ' in the $L$ register is |
|  |  | SRL | L | to count '4' T states, but take INT (L/4) and count '16' T states instead. |
|  |  | CPL |  | Go back to the original value |
|  |  | AND | +03 | in $L$ and find how many were |
|  |  | LD | C,A | lost by taking INT (L/4). |
|  |  | LD | B,+00 |  |
|  |  | LD | IX,+03D1 | The base address of the timing loop. |
|  |  | ADD | IX,BC | Alter the length of the timing loop. Use an earlier starting point for each '1' lost by taking INT (L/4). |
|  |  | LD | A,(BORDCR) | Fetch the present border |
|  |  | AND | +38 | colour and move it to bits |
|  |  | RRCA |  | $2,1 \& 0$ of the A register. |
|  |  | RRCA |  |  |
|  |  | RRCA |  |  |
|  |  | OR | +08 | Ensure the MIC output is 'off'. |

Now enter the sound generation loop. 'DE' complete passes are made, i.e. a pass for each cycle of the note.
The HL register holds the 'length of the timing loop' with '16' T states being used for each '1' in the L register and '1,024' T states for each ' 1 ' in the H register.

| 03D1 | BE-IX+3 | NOP |  |
| :--- | :--- | :--- | :--- |
| 03D2 | BE-IX+2 | NOP |  |
| 03D3 | BE-IX+1 | NOP |  |
| 03D4 | BE-IX+0 | INC | B |
|  |  | INC | C |
| 03D6 | BE-H\&L-LP | DEC | C |
|  |  | JR | NZ,03D6,BE-H\&L-LP |
|  |  | LD | C,+3F |
|  |  | DEC | B |
|  |  | JP | NZ,03D6,BE-H\&L-LP |

Add '4' T states for each
earlier entry port
that is used.
The values in the $B$ \& $C$ registers will come from $\mathrm{H} \& \mathrm{~L}$ registers

- see below.

The 'timing loop'.
i.e. 'BC' * '4' T states.
(But note that at the half-cycle point - C will be equal to 'L+1'.)

The loudspeaker is now alternately activated and deactivated.

| OUT | $(+$ FE),A |
| :--- | :--- |
|  |  |
| LD | B,H |
| LD | C,A |
| BIT | $4, A$ |
| JR | NZ,03F2,BE-AGAIN |

Perform the OUT operation; leaving the border unchanged.
Reset the B register.
Save the A register.
Jump if at the half-cycle point.

After a full cycle the DE register pair is tested.

| LD | A,D | Jump forward if the last |
| :--- | :--- | :--- |
| OR | E | complete pass has been |
| JR | Z,03D6,BE-END | made already. |
| LD | A,C | Fetch the saved value. |
| LD | C,L | Reset the C register. |
| DEC | DE | Decrease the pass counter. <br> JP |
|  | (IX) | Jump back to the required <br> starting location of the loop. |

The parameters for the second half-cycle are set up.

| 03F2 BE-AGAIN | LD | C,L | Reset the $C$ register. <br> INC |
| :--- | :--- | :--- | :--- |
|  | C | Add '16' $T$ states as this path <br> is shorter. |  |
|  | JP | (IX) | Jump back. |

Upon completion of the 'beep' the maskable interrupt has to be enabled.

| 03F6 | EIE-END | Enable interrupt. |
| :--- | :--- | :--- |
|  | RET | Finally return. |

## THE 'BEEP' COMMAND ROUTINE

The subroutine is entered with two numbers on the calculator stack. The topmost number represents the 'pitch' of the note and the number underneath it represents the 'duration'.

| 03F8 BEEP | RST | 0028,FP-CALC | The floating-point calculator is |
| :--- | :--- | :--- | :--- |
|  |  | used to manipulate the two values $-\mathrm{t} \& \mathrm{P}$. <br>  <br>  <br>  <br>  <br>  <br> DEFB | +31, duplicate |

Now perform several tests on I, the integer part of the 'pitch'.

| LD | HL,+5C92 |
| :--- | :--- |
| LD | A,(HL) |
| AND | A |
| JR | NZ,046C,REPORT-B |
| INC | HL |
| LD | $\mathrm{C},(\mathrm{HL})$ |
| INC | HL |
| LD | $\mathrm{A}, \mathrm{B}$ |
| RLA | $\mathrm{A}, \mathrm{A}$ |
| SBC | C |
| CP | $\mathrm{NZ}, 046 \mathrm{C}$, REPORT-B |
| JR | HL |
| INC | CP |
| CP | (HL) |

This is 'mem-0-1st (MEMBOT),
Fetch the exponent of $i$.
Give an error if $i$ is not in the integral (short) form
Copy the sign byte to the
C register.
Copy the low-byte to the
register.
Again give report B if i does not
satisfy the test:
$-128<=$ i<=+127

| JR | NZ,046C,REPORT-B |  |
| :--- | :--- | :--- |
| LD | A,B | Fetch the low-byte and test <br> it further. |
| ADD | A,, $3 C$ |  |
| JP | P,0425,BE-I-OK | Accept $-60<=$ i< $<67$. |
| JP | PO,046C,REPORT-B | Reject -128 to -61. |

Note: The range +70 to +127 will be rejected later on.
The correct frequency for the 'pitch' i can now be found.

| 0425 | BE-I-OK | LD | B,+FA | Start '6' octaves below middle C. |
| :---: | :---: | :---: | :---: | :---: |
| 0427 | BE-OCTAVE | INC | B | Repeatedly reduce i in order to |
|  |  | SUB | +0C | find the correct octave. |
|  |  | JR | NC,0427,BE-OCTAVE |  |
|  |  | ADD | A, +0 C | Ass back the last subtraction. |
|  |  | PUSH | BC | Save the octave number. |
|  |  | LD | HL,+046E | The base address of the 'semitone table'. |
|  |  | CALL | 3406,LOC-MEM | Consider the table and pass the |
|  |  | CALL | 33B4,STACK-NUM | 'A th.' value to the calculator stack (Call it C) |

Now the fractional part of the 'pitch' can be taken into consideration.

| RST | 0028,FP-CALC | $\mathrm{t}, \mathrm{pK}+1, \mathrm{C}$ |
| :--- | :--- | :--- |
| DEFB | +04,multiply | $\mathrm{t}, \mathrm{C}(\mathrm{pK}+1)$ |

The final frequency $f$ is found by modifying the 'last value' according to the octave number.


The number of complete cycles in the 'beep' is given by 'ftt' so this value is now found.

| RST | 0028, FP-CALC | $t$ |
| :--- | :--- | :--- |
| DEFB | + E0,get-mem-0 | $t, f$ |
| DEFB | +04, multiply | $f^{* t}$ |

The result is left on the calculator stack whilst the length of the 'timing loop' required for the 'beep' is computed;

| DEFB | + +E0,get-mem-0 | $f^{*} t$ |
| :--- | :--- | :--- |
| DEFB | +34, stk-data | The value '3.5* $10^{\wedge} 6 / 8^{\prime}$ |
| DEFB | +80, four bytes | is formed on the top of |
| DEFB | +43, exponent +93 | the calculator stack. |
| DEFB | $+55,+9 F,+80,(+00)$ | $f^{*} t, f, 437,500$ (dec.) |
| DEFB | +01, exchange | $f^{\star} t, 437,500, f$ |
| DEFB | +05, division | $f^{*} t, 437,500 / f$ |
| DEFB | +34, stk-data |  |
| DEFB | +35, exponent +85 |  |
| DEFB | $+71,(+00,+00,+00)$ | $f^{\star} t, 437,500 / f, 30.125($ dec. $)$ |
| DEFB | +03, subtract | $f^{*} t, 437,500 / f-30.125$ |
| DEFB | +38, end-calc |  |

Note: The value '437,500/f' gives the 'half-cycle' length of the note and reducing it by '30.125' allows for '120.5' T states in which to actually produce the note and adjust the counters etc.
The values can now be transferred to the required registers.

| CALL | 1E99,FIND-INT2 | The 'timing loop' value is <br> compressed into the BC <br> register pair; and saved. |
| :--- | :--- | :--- |
| PUSH | BC |  |

Note: If the timing loop value is too large then an error will occur (returning via ERROR-1); thereby excluding 'pitch' values of ' +70 to +127 '.

| CALL | 1E99,FIND-INT2 | The 'f*t' value is compressed <br> into the BC register pair. |
| :--- | :--- | :--- |
| POP | HL | Move the 'timing loop' value to <br> the HL register pair. |
| LD | D,B | Move the 'f $f^{\star t}$ value to the |
| LD | E,C | DE register pair. |

However before making the 'beep' test the value 'f*t'.

| LD | A,D |
| :--- | :--- |
| OR | E |
| RET | $Z$ |
| DEC | DE |
| JP | 03B5,BEEPER |

Return if 'f*t' has given the result of 'no cycles'
required.
Decrease the cycle number and jump to the BEEPER subroutine (making, at least, one pass).

Report B - integer out of range

| 046C REPORT-B | RST <br> DEFB | 0008,ERROR-1 <br> $+0 A$ | Call the error handling <br> routine. |
| :--- | :--- | :--- | :--- |

## THE 'SEMI-TONE' TABLE

This table holds the frequencies of the twelve semi-tones in an octave.

| 046E |  |  | frequency hz . | note |
| :---: | :---: | :---: | :---: | :---: |
|  | DEFB | +89,+02,+D0,+12,+86 | 261.63 | C |
|  | DEFB | +89,+0A,+97,+60,+75 | 277.18 | C\# |
|  | DEFB | +89,+12,+D5,+17,+1F | 293.66 | D |
|  | DEFB | +89,+1B,+90,+41,+02 | 311.12 | D\# |
|  | DEFB | +89,+24,+D0,+53,+CA | 329.63 | E |
|  | DEFB | +89,+2E, +9D, +36, +B1 | 349.23 | F |
|  | DEFB | +89,+38,+FF,+49,+3E | 369.99 | F\# |
|  | DEFB | +89,+43,+FF,+6A,+73 | 392 | G |
|  | DEFB | +89,+4F,+A7,+00,+54 | 415.30 | G\# |
|  | DEFB | +89,+5C,+00,+00,+00 | 440 | A |
|  | DEFB | +89,+69,+14,+F6,+24 | 466.16 | A\# |
|  | DEFB | +89,+76,+F1,+10,+05 | 493.88 | B |

THE 'PROGRAM NAME' SUBROUTINE (ZX81)
The following subroutine applies to the ZX81 and was not removed when the program was rewritten for the SPECTRUM.

| 04AA | DEFB | +CD, +FB, +24,+3A |
| :---: | :---: | :---: |
|  | DEFB | +3B,+5C,+87,+FA |
|  | DEFB | +8A, +1C, +E1,+D0 |
|  | DEFB | +E5,+CD,+F1,+2B |
|  | DEFB | +62,+6B,+0D,+F8 |
|  | DEFB | +09,+CB,+FE,+C9 |

## THE CASSETTE HANDLING ROUTINES

The 16 K monitor program has an extensive set of routines for handling the cassette interface. In effect these routines form the SAVE. LOAD, VERIFY \& MERGE command routines.
The entry point to the routines is at SAVE-ETC (0605). However before this point are the subroutines concerned with the actua SAVEing and LOADing (or VERIFYing) of bytes.
In all cases the bytes to be handled by these subroutines are described by the DE register pair holding the 'length' of the block, the IX register pair holding the 'base address' and the A register holding +00 for a header block, or +FF for a program/data block.

## THE 'SA-BYTES' SUBROUTINE

This subroutine is called to SAVE the header information (from 09BA) and later the actual program/data block (from 099E).

| 04C2 | SA-BYTES | $\begin{aligned} & \text { LD } \\ & \text { PUSH } \\ & \text { LD } \end{aligned}$ | $\begin{aligned} & \mathrm{HL},+053 \mathrm{~F} \\ & \mathrm{HL} \\ & \mathrm{HL},+1 \mathrm{~F} 80 \end{aligned}$ | Pre-load the machine stack with the address-SA/LD-RET. <br> This constant will give a leader of about 5 secs for a 'header' |
| :---: | :---: | :---: | :---: | :---: |
|  |  | BIT | 7,A | Jump forward if SAVEing a |
|  |  | JR | Z,04D0,SA-FLAG | header. |
|  |  | LD | HL,+0C98 | This constant will give a leader of about 2 secs. for a program/ data block. |
| 04D0 | SA-FLAG | EX | AF,A'F' | The flag is saved. |
|  |  | INC | DE | The 'length' is incremented |
|  |  | DEC | IX | and the 'base address' reduced to allow for the flag. |
|  |  | DI |  | The maskable interrupt is disabled during the SAVE. |
|  |  | LD | A,+02 | Signal 'MIC on' and border to be RED. |
|  |  | LD | B,A | Give a value to B. |

A loop is now entered to create the pulses of the leader. Both the 'MIC on' and the 'MIC off' pulses are $2,168 \mathrm{~T}$ states in length. The colour of the border changes from RED to CYAN with each 'edge'.

Note: An 'edge' will be a transition either from 'on' to 'off', or from 'off' to 'on'.

04D8 SA-LEADER

| DJNZ | 04D8,SA-LEADER |
| :--- | :--- |
| OUT | $(+$ FE),A |
| XOR | $+0 F$ |
| LD | B,+A4 |
| DEC | L |
| JR | NZ,04D8,SA-LEADER |
| DEC | B |
|  |  |
| DEC | H |
| JP | P,04D8,SA-LEADER |

A sync pulse is now sent.

|  |  | LD | B, +2 F |  |
| :---: | :---: | :---: | :---: | :---: |
| 04EA | SA-SYNC-1 | DJNZ | 04EA,SA-SYNC-1 | MIC off for 667 T states from 'OUT to OUT'. |
|  |  | OUT | (+FE), A | MIC on and RED. |
|  |  | LD | A, +0D | Signal 'MIC off \& CYAN'. |
|  |  | LD | B,+37 | MIC on for 735 T States from |
| 04F2 | SA-SYNC-2 | DJNZ | 04F2,SA-SYNC-2 | 'OUT to OUT'. |
|  |  | OUT | (+FE),A | Now MIC off \& border CYAN. |

The header v. program/data flag will be the first byte to be SAVEd.

The main timing period. MIC on/off, border RED/CYAN, on each pass.
The main timing constant.
Decrease the low counter Jump back for another pulse. Allow for the longer path (-reduce by 13 T states) Decrease the high counter. Jump back for another pulse until completion of the leader.

| LD | BC, +3 BOE | +3B is a timing constant; + +OE |
| :--- | :--- | :--- |
| signals 'MIC off \& YELLOW'. |  |  |
| EX | AF,A'F' | Fetch the flag and pass it to the |
| LD | L,A | L register for 'sending'. |
| JP | 0507, SA-START | Jump forward into the SAVEing <br> loop. |

The byte SAVEing loop is now entered. The first byte to be SAVEd is the flag; this is followed by the actual data byte and the final byte sent is the parity byte that is built up by considering the values of all the earlier bytes.

| 04FE | SA-LOOP | LD | A,D | The 'length' counter is tested <br> and the jump taken when it |
| :--- | :--- | :--- | :--- | :--- |
| has reached zero. |  |  |  |  |

When it is time to send the 'parity' byte then it is transferred to the L register for SAVEing.

| 050E | SA-PARITY | LD | L,H | Get final 'parity' value. |
| :--- | :--- | :--- | :--- | :--- |

The following inner loop produces the actual pulses. The loop is entered at SA-BIT-1 with the type of the bit to be SAVEd indicated by the carry flag. Two passes of the loop are made for each bit thereby making an 'off pulse' and an 'on pulse'. The pulses for a reset bit are shorter by 855 T states.

| 0511 | SA-BIT-2 | LD | A, C | Come here on the second pass and fetch 'MIC off \& YELLOW'. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | BIT | 7,B | Set the zero flag to show 'second pass'. |
| 0514 | SA-BIT-1 | DJNZ | 0514,SA-BIT-1 | The main timing loop; always 801 T states on a 2nd. pass. |
|  |  | JR | NC,051C,SA-OUT | Jump, taking the shorter path, if SAVEing a ' 0 '. |
|  |  | LD DJNZ | $\begin{aligned} & \mathrm{B},+42 \\ & 051 \mathrm{~A}, \mathrm{SA}-\mathrm{SET} \end{aligned}$ | However if SAVEing a ' 1 ' then add 855 T states. |
| 051C | SA-SETSA-OUT | OUT | (+FE),A | On the 1st. pass 'MIC on \& BLUE' and on the 2nd. pass 'MIC off \& YELLOW'. |
|  |  | LD | B, +3 E | Set the timing constant for the second pass. |
|  |  | $\begin{aligned} & \text { JR } \\ & \text { DEC } \end{aligned}$ | NZ,0511,SA-BIT-2 B | Jump back at the end of the first pass; otherwise reclaim 13 T states. |
|  |  | XOR | A | Clear the carry flag and set |
|  |  | INC | A | A to hold +01 (MIC on \& BLUE before continuing into the ' 8 bit loop'. |

The ' 8 bit loop' is entered initially with the whole byte in the $L$ register and the carry flag set. However it is re-entered after each bit has been SAVEd until the point is reached when the 'marker' passes to the carry flag leaving the $L$ register empty.
0525 SA-8-BITS RL L Move bit 7 to the carry and the 'marker' leftwards.

|  | JP | NZ,0514,SA-BIT-1 |
| :--- | :--- | :--- | | SAVE the bit unless finished |
| :--- |
| with the byte. |

Note: A reset bit will give a 'MIC off' pulse of 855 T states followed by a 'MIC on' pulse of 855 T states. Whereas a Set bit will give pulses of exactly twice as long. Note also that there are no gaps either between the sync pulse and the first bit of the flag, or between bytes.

## THE 'SA/LD-RET' SUBROUTINE

This subroutine is common to both SAVEing and LOADing.
The border is set to its original colour and the BREAK key tested for a last time.

| 053F | SA/LD-RET | PUSH | AF | Save the carry flag. (It is reset after a LOADing error.) |
| :---: | :---: | :---: | :---: | :---: |
|  |  | LD | A,(BORDCR) | Fetch the original border colour |
|  |  | AND | +38 | from its system variable. |
|  |  | RRCA |  | Move the border colour |
|  |  | RRCA |  | to bits $2,1 \& 0$. |
|  |  | RRCA |  |  |
|  |  | OUT | (+FE), A | Set the border to its original colour. |
|  |  | LD | A. +7 F | Read the BREAK key for a |
|  |  | IN | A, (+FE) | last time. |
|  |  | RRA |  |  |
|  |  | El |  | Enable the maskable interrupt. |
|  |  | JR | C,0554,SA/LD-END | Jump unless a break is to be |

Report D - BREAK-CONT repeats
0552 REPORT-D RST 0008,ERROR-I Call the error handling
DEFB +0C routine.

Continue here.

0554 SA/LD-END \begin{tabular}{lll}
POP <br>

\& RET \& | Retrieve the carry flag. |
| :--- |
| Return to the calling routine. |

\end{tabular}

THE 'LD-BYTES' SUBROUTINE
This subroutine is called to LOAD the header information (from 07BE) and later LOAD, or VERIFY, an actual block of data (from 0802).

| 0556 | LD-BYTES | INC EX | D AF, $A^{\prime} \mathrm{F}^{\prime}$ | This resets the zero flag. (D cannot hold +FF.) <br> The A register holds +00 for a header and +FF for a block of data. <br> The carry flag is reset for VERIFYing and set for LOADing. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | DEC | D | Restore D to its original value. |

DI The maskable interrupt is now

| LD | A,+0F |
| :--- | :--- |
| OUT | $(+$ FE $)$, A |
| LD | HL,+053F |
| PUSH | HL |
| IN | A,(+FE) |
| RRA |  |
| AND | +20 |
| OR | +02 |
| LD | C,A |
|  |  |
| CP | A |

disabled.
The border is made WHITE.

Preload the machine stack with the address - SA/LD-RET.
Make an initial read of port '254'
Rotate the byte obtained but
keep only the EAR bit,
Signal 'RED' border.
Store the value in the C register. -
(+22 for 'off' and +02 for 'on'

- the present EAR state.)

Set the zero flag.

The first stage of reading a tape involves showing that a pulsing signal actually exist (i.e. 'On/off' or 'off/on' edges.)
\(\left.$$
\begin{array}{llll}\text { 056B } & \text { LD-BREAK } & \text { RET } & \text { NZ }\end{array}
$$ \begin{array}{l}Return if the BREAK key is <br>

being pressed.\end{array}\right]\)| Return with the carry flag reset |
| :--- |
| if there is no 'edge' within |
| approx. 14,000 T states. But if |
| an 'edge' is found the border |
| will go CYAN. |

The next stage involves waiting a while and then showing that the signal is still pulsing.

|  |  | LD | HL,+0415 |
| :--- | :--- | :--- | :--- |$\quad$| The length of this waiting |
| :--- |
| period will be almost one |

Now accept only a 'leader signal'.


After the leader come the 'off' and 'on' part's of the sync pulse.

| 058F LD-SYNC | LD | B,+9C | The timing constant. |
| :--- | :--- | :--- | :--- |
|  | CALL | 05E7,LD-EDGE-1 | Every edge is considered until |
|  | JR | NC,056B,LD-BREAK | two edges are found close |
|  | LD | A,B | together - these will be the |
|  | CP | +D4 | start and finishing edges of |
|  | JR | NC,058F,LD-SYNC | the 'off' sync pulse. |
|  | CALL | 05E7,LD-EDGE-1 | The finishing edge of the |
|  | RET | NC | 'on' pulse must exist. |
|  |  |  | (Return carry flag reset.) |

The bytes of the header or the program/data block can now be LOADed or VERIFied. But the first byte is the type flag.

| LD | A,C | The border colours from now |
| :--- | :--- | :--- |
| XOR | +03 | on will be BLUE \& YELLOW. |


| LD | C,A <br> H,,+00 | Initialise the 'parity matching' <br> byte to zero. |
| :--- | :--- | :--- |
| LD | B,+B0 | Set the timing constant for the <br> flag byte. <br> Jump forward into the byte <br> LOADING loop. |
| JR | 05C8,LD-MARKER | LDA |

The byte LOADing loop is used to fetch the bytes one at a time. The flag byte is first. This is followed by the data bytes and the last byte is the 'parity' byte.

| 05A9 | LD-LOOP | EX | AF, $\mathrm{A}^{\prime} \mathrm{F}^{\prime}$ | Fetch the flags. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | JR | NZ,05B3,LD-FLAG | Jump forward only when handling the first byte. |
|  |  | JR | NC,05BD,LD-VERIFY | Jump forward if VERIFYing a tape. |
|  |  | LD | ( $\mathrm{IX}+00$ ), L | Make the actual LOAD when required. |
|  |  | JR | 05C2,LD-NEXT | Jump forward to LOAD the next byte. |
| 05B3 | LD-FLAG | RL | C | Keep the carry flag in a safe place temporarily. |
|  |  | XOR | L | Return now if the type flag does |
|  |  | RET | NZ | not match the first byte on the tape. (Carry flag reset.) |
|  |  | $\begin{aligned} & \text { LD } \\ & \text { RRA } \end{aligned}$ | A, C | Restore the carry flag now. |
|  |  | LD | C,A |  |
|  |  | INC |  | Increase the counter to |
|  |  | JR | 05CA,LD-DEC | compensate for its 'decrease' after the jump. |

If a data block is being verified then the freshly loaded byte is tested against the original byte.

| 05BD | LD-VERIFY | LD | A,(IX+00) |
| :--- | :--- | :--- | :--- | | Fetch the original byte. |
| :--- |
|  |
|  |
|  |
|  |
|  |
|  |
|  |
| RET |

A new byte can now be collected from the tape.

| 05C2 | LD-NEXT | INC | IX |
| :--- | :--- | :--- | :--- |
| 05C4 | LD-DEC | DEC | DE |
|  |  | EX | AF,A'F' |
|  |  | LD | B,+B2 |
| 05 C 8 | LD-MARKER | LD | L,+01 |

The 'LD-8-BITS' loop is used to build up a byte in the $L$ register.


The 'parity matching' byte has to be updated with each new byte.

| LD | A,H | Fetch the 'parity matching' |
| :--- | :--- | :--- |
| XOR | L | byte and include the new byte. |
| LD | H,A | Save it once again. |

Passes round the loop are made until the 'counter' reaches zero. At that point the 'parity matching' byte should be holding zero.

| LD | A,D | Make a further pass if the DE |
| :--- | :--- | :--- |
| OR | E | register pair does not hold |
| JR | NZ,05A9,LD-LOOP | zero. |
| LD | A,H | Fetch the 'parity matching' <br> byte. |
| CP | +01 | Return with the carry flat set <br> if the value is zero. |
| RET |  | (Carry flag reset if in error.) |

## THE 'LD-EDGE-2' AND 'LD-EDGE-1' SUBROUTINES

These two subroutines form the most important part of the LOAD/VERIFY operation.
The subroutines are entered with a timing constant in the B register, and the previous border colour and 'edge-type' in the C register
The subroutines return with the carry flag set if the required number of 'edges' have been found in the time allowed; and the change to the value in the B register shows just how long it took to find the 'edge(s)'.

The carry flag will be reset if there is an error. The zero flag then signals 'BREAK pressed' by being reset, or 'time-up' by being set.
The entry point LD-EDGE-2 is used when the length of a complete pulse is required and LD-EDGE-1 is used to find the time before the next 'edge'.
\(\left.$$
\begin{array}{lllll}\text { 05E3 } & \text { LD-EDGE-2 } & \text { CALL } & \begin{array}{l}\text { 05E7,LD-EDGE-1 } \\
\text { RET }\end{array} & \text { NC }\end{array}
$$ \begin{array}{l}In effect call LD-EDGE-1 twice; <br>

returning in between if there\end{array}\right]\)| is an error. |
| :--- |

The sampling loop is now entered. The value in the B register is incremented for each pass; 'time-up' is given when B reaches zero.

| 05ED | LD-SAMPLE | INC | B |
| :--- | :--- | :--- | :--- |
|  | RET | Z | Count each pass. <br> Return carry reset \& zero set if |
|  |  |  | 'time-up'. |

A new 'edge' has been found within the time period allowed for the search. So change the border colour and set the carry flag.

| LD | A,C | Change the 'last edge-type' <br> and border colour. |
| :--- | :--- | :--- |
| CPL | C,A | Keep only the border colour. |
| LD | +07 | Signal 'MIC off'. |
| AND | +08 | Change the border colour (RED/ <br> OR$+$ CYAN or BLUE/YELLOW). |
| OUT |  | Signal the successful search <br> before returning. |
| SCF |  |  |

Note: The LD-EDGE-1 subroutine takes 465 T states, plus an additional 58 T states for each unsuccessful pass around the sampling loop.

For example, therefore, when awaiting the sync pulse (see LD-SYNC at 058F) allowance is made for ten additional passes through the sampling loop. The search is thereby for the next edge to be found within, roughly, $1,100 \mathrm{~T}$ states $(465+10 * 58+$ overhead). This will prove successful for the sync 'off' pulse that comes after the long 'leader pulses'.

## THE 'SAVE, LOAD, VERIFY \& MERGE' COMMAND ROUTINES

The entry point SAVE-ETC is used for all four commands. The value held in T-ADDR however distinguishes between the four commands. The first part of the following routine is concerned with the construction of the 'header information' in the work space.

| 0605 | SAVE-ETC | POP | AF | Drop the address - SCAN-LOOP |
| :---: | :---: | :---: | :---: | :---: |
|  |  | LD | A,(T-ADDR-Io) | Reduce T-ADDR-lo by +E0; |
|  |  | SUB | +E0 | giving +00 for SAVE, +01 for |
|  |  | LD | (T-ADDR-Io), A | LOAD, +02 for VERIFY and +03 for MERGE. |
|  |  | CALL | 1C8C,EXPT-EXP | Pass the parameters of the |
|  |  |  |  | 'name' to the calculator stack. |
|  |  | CALL | 2530,SYNTAX-Z | Jump forward if checking |
|  |  | JR | Z,0652,SA-DATA | syntax. |
|  |  | LD | BC, +0011 | Allow seventeen locations |
|  |  | LD | A,(T-ADDR-lo) | for the header of a SAVE but |
|  |  | AND | A | thirty four for the other |
|  |  | JR | Z,0621,SA-SPACE | commands. |
|  |  | LD | C,+22 |  |
| 0621 | SA-SPACE | RST | 0030,BC-SPACES | The required amount of space is made in the work space. |
|  |  | PUSH | DE | Copy the start address to the |
|  |  | POP | IX | IX register pair. |
|  |  | LD | B, +0B | A program name can have |
|  |  | LD | A, +20 | up to ten characters but |
| 0629 | SA-BLANK | LD | (DE), A | first enter eleven space |
|  |  | INC | DE | characters into the prepared |
|  |  | DJNZ | 0629,SA-BLANK | area. |
|  |  | LD | ( $\mathrm{IX}+01$ ), +FF | A null name is +FF only. |
|  |  | CALL | 2BF1,STK-FETCH | The parameters of the name are fetched and its length is tested. |
|  |  | LD | HL,+FFF6 | This is '-10'. |
|  |  | DEC | BC | In effect jump forward if the |
|  |  | ADD | HL, BC | length of the name is not |
|  |  | INC | BC | too long. (i.e. No more than |
|  |  | JR | NC,064B,SA-NAME | ten characters.) |
|  |  | LD | A,(T-ADDR-lo) | But allow for the LOADing, |
|  |  | AND |  | VERIFYing and MERGEing of |
|  |  | JR | NZ,0644,SA-NULL | programs with 'null' names or extra long names. |

Report F - Invalid file name

| 0642 | REPORT-F | RST <br> DEFB | 0008,ERROR-1 <br> +0 E |
| :--- | :--- | :--- | :--- |

Continue to handle the name of the program.

| 0644 | SA-NULL | LD | A,B |
| :--- | :--- | :--- | :--- | | Jump forward if the name |
| :--- |
|  |
|  |
|  |
|  |
|  |
|  |
|  |
|  |
| OR | | LR | Z, | has a 'null' length. |
| :--- | :--- | :--- |
|  | BC,+000A | But truncate longer names. |

The name is now transferred to the work space (second location onwards).

| 064B | SA-NAME | PUSH | IX |
| :--- | :--- | :--- | :--- | | Copy the start address to the |
| :--- |
|  |
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|  |
|  |

The many different parameters, if any, that follow the command are now considered. Start by handling 'xxx "name" DATA'.

| 0652 | SA-DATA | RST | 0018,GET-CHAR | Is the present code the |
| :---: | :---: | :---: | :---: | :---: |
|  |  | CP | +E4 | token 'DATA'? |
|  |  | JR | NZ,06A0,SA-SCR\$ | Jump if not. |
|  |  | LD | A,(T-ADDR-Io) | However it is not possible |
|  |  | CP | +03 | to have 'MERGE name DATA'. |
|  |  | JP | Z,1C8A,REPORT-C |  |
|  |  | RST | 0020,NEXT-CHAR | Advance CH-ADD. |
|  |  | CALL | 28B2,LOOK-VARS | Look in the variables area for the array. |
|  |  | SET | 7,C | Set bit 7 of the array's name. |
|  |  | JR | NC,0672,SA-V-OLD | Jump if handling an existing array. |
|  |  | LD | HL, +0000 | Signal 'using a new array'. |
|  |  | LD | A,(T-ADDR-lo) | Consider the value in T-ADDR |
|  |  | DEC |  | and give an error if trying to |
|  |  | JR | Z,0685,SA-V-NEW | SAVE or VERIFY a new array. |

Report 2 - Variable not found

0670 REPORT-2 RST \begin{tabular}{llll}

RST \& 0008,ERROR-1 \& | Call the error handling |
| :--- |
| routine. |

\end{tabular}

Continue with the handling of an existing array.

| 0672 SA-V-OLD | JP | NZ,1C8A,REPORT-C | Note: This fails to exclude <br> simple strings. |
| :--- | :--- | :--- | :--- |
|  | CALL | 2530,SYNTAX-Z | Jump forward if checking <br> syntax. |
|  | JR | Z,0692,SA-DATA-1 | Point to the 'low length' of the <br> INC |
|  | HL | variable. |  |

The next part is common to both 'old' and 'new' arrays. Note: Syntax path error.

| 0685 | SA-V-NEW | LD | ( $\mathrm{IX}+0 \mathrm{E}$ ), C | Copy the array's name. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | LD | A, +01 | Assume an array of numbers. |
|  |  | BIT | 6,C | Jump if it is so. |
|  |  | JR | Z,068F,SA-V-TYPE |  |
|  |  | INC | A | It is an array of characters. |
| 068F | SA-V-TYPE | LD | (IX+00), A | Save the 'type' in the first |

The last part of the statement is examined before joining the other pathways.
0692

| SA-DATA- | EX | DE,HL |
| :--- | :--- | :--- |
|  | RST | 0020, NEXT-CHAR |
|  | CP | +29 |
|  | JR | NZ,0672,SA-V-OLD |
|  | RST | 0020, NEXT-CHAR |
|  | CALL | 1BEE,CHECK-END |
|  |  |  |
|  | EX | DE,HL |
|  | JP | 075A,SA-ALL |

Save the pointer in DE. Is the next character
a ')' ?
Give report C if it is not. Advance CH-ADD.
Move on to the next statement if checking syntax.
Return the pointer to the HL register pair before jumping forward. (The pointer indicates the start of an existing array's contents.)

Now consider 'SCREEN\$'.
06A0
SA-SCR\$ CP
$+A A$
Is the present code the
token SCREEN ${ }^{\prime}$.

| JR | NZ,06C3,SA-CODE | Jump if not. |
| :---: | :---: | :---: |
| LD | A,(T-ADDR-Io) | However it is not possible to |
| CP | +03 | have 'MERGE name SCREEN\$'. |
| JP | Z,1C8A,REPORT-C |  |
| RST | 0020,NEXT-CHAR | Advance CH-ADD. |
| CALL | 1BEE,CHECK-END | Move on to the next statement if checking syntax. |
| LD | (IX+0B), +00 | The display area and the |
| LD | ( $\mathrm{IX}+0 \mathrm{C}$ ),+1 B | attribute area occupy +1800 locations and these locations |
| LD | HL,+4000 | start at +4000 ; these details |
| LD | (IX+OD), L | are passed to the header area |
| LD | (IX+0E), H | in the work space. |
| JR | 0710,SA-TYPE-3 | Jump forward. |

Now consider 'CODE'.

| 06C3 | SA-CODE | CP |
| :--- | :--- | :--- |
|  |  | +AF |
|  | JR | NZ,0716,SA-LINE |
|  | LD | A,(T-ADDR-lo) |
|  | CP | +03 |
|  | JP | Z,1C8A,REPORT-C |
|  | RST | OO20,NEXT-CHAR |
|  | CALL | 2048,PR-ST-END |
|  | JR | NZ,06E1,SA-CODE-1 |
|  | LD | A,(T-ADDR-lo) |
|  | AND | A |
|  | JP | Z,1C8A,REPORT-C |
|  | CALL | 1CE6,USE-ZERO |
|  |  |  |
|  | JR | 06F0,SA-CODE-2 |

Is the present code the token 'CODE'?

## Jump if not.

However it is not possible to have 'MERGE name CODE'.

Advance CH-ADD.
Jump forward if the statement has not finished.
However it is not possible to have 'SAVE name CODE' by itself.
Put a zero on the calculator stack - for the 'start'.
Jump forward.
Look for a 'starting address'.

| 06E1 | SA-CODE-1 | CALL | 1C82,EXPT-1NUM | Fetch the first number. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | RST | 0018,GET-CHAR | Is the present character a ',' |
|  |  | CP | +2C | or not? |
|  |  | JR | Z,06F5,SA-CODE-3 | Jump if it is - the number was a 'starting address'. |
|  |  | LD | A,(T-ADDR-lo) | However refuse 'SAVE name |
|  |  | AND | A | CODE' that does not have a |
|  |  | JP | Z,1C8A,REPORT-C | 'start' and a 'length'. |
| 06F0 | SA-CODE-2 | CALL | 1CE6,USE-ZERO | Put a zero on the calculator stack - for the 'length'. |
|  |  | JR | 06F9,SA-CODE-4 | Jump forward. |

Fetch the 'length' as it was specified.

06F5 SA-CODE-3 RST
CALL

0020,NEXT-CHAR
1C82,EXPT-1NUM

Advance CH-ADD Fetch the 'length'.

The parameters are now stored in the header area of the work space.

| 06F9 | SA-CODE-4 | CALL | 1BEE,CHECK-END | But move on to the next statement now if checking syntax. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | CALL | 1E99,FIND-INT2 | Compress the 'length' into |
|  |  | LD | (IX+0B), C | the $B C$ register pair and |
|  |  | LD | (IX+0C), B | store it. |
|  |  | CALL | 1E99,FIND-INT2 | Compress the 'starting address' |
|  |  | LD | (IX+0D), C | into the BC register pair |
|  |  | LD | (IX+0E), B | and store it. |
|  |  | LD | H,B | Transfer the 'pointer' to the |
|  |  | LD | L,C | HL register pair as usual. |

'SCREEN\$' and 'CODE' are both of type 3.
SA-TYPE-3 LD
(IX+00),+03
Enter the 'type' number.

JR 075A,SA-ALL Rejoin the other pathways.
Now consider 'LINE'; and 'no further parameters'.

| 0716 | SA-LINE | CP | +CA | Is the present code the token 'LINE'? |
| :---: | :---: | :---: | :---: | :---: |
|  |  | JR | Z,0723,SA-LINE-1 | Jump if it is. |
|  |  | CALL | 1BEE,CHECK-END | Move on to the next statement if checking syntax. |
|  |  | LD | $(\mathrm{IX}+0 \mathrm{E}),+80$ | When there are no further parameters an +80 is entered. |
|  |  | JR | 073A,SA-TYPE-0 | Jump forward. |

Fetch the 'line number' that must follow 'LINE'.
0723

| SA-LINE-1 | LD | A,(T-ADDR-lo) |
| :--- | :--- | :--- |
|  | AND | A |
|  | JP | NZ,1C8A,REPORT-C |
|  | RST | 0020,NEXT-Char |
|  | CALL | 1C82,EXPT-1NUM |
|  |  |  |
|  | CALL | $1 B E E$, CHECK-END |
|  | CALL |  |
|  | LD | 1E99,FIND-INT2 |
|  | LD | (IX+0D),C |
|  |  | (IX+0E),B |

However only allow 'SAVE name LINE number'.

Advance CH-ADD.
Pass the number to the calculator stack.
Move on to the next statement if checking syntax.
Compress the 'line number' into the BC register pair and store it.
'LINE' and 'no further parameters' are both of type 0.
073A
SA-TYPE-0 LD
$(\mathrm{IX}+00),+00$
Enter the 'type' number.
The parameters that describe the program, and its variables, are found and stored in the header area of the work space.

| LD | HL,(E-LINE) | The pointer to the end of the <br> variables area. |
| :--- | :--- | :--- |
| LD | DE,(PROG) | The pointer to the start of the <br> BASIC program. |
| SCF |  | Now perform the subtraction <br> to find the length of the |
| SBC | HL,DE | 'program + variables'; store |
| LD | (IX+0B), | the result. |

In all cases the header information has now been prepared.
The location 'I $X+00$ ' holds the type number.
Locations 'IX+01 to IX+0A' holds the name (+FF in 'IX+01' if null).
Locations ' $I X+0 B \& I X+0 C$ ' hold the number of bytes that are to be found in the 'data block'.
Locations 'IX+0D to IX+10' hold a variety of parameters whose exact interpretation depends on the 'type'.
The routine continues with the first task being to separate SAVE from LOAD, VERIFY and MERGE.
075A SA-ALL

| LD | A,(T-ADDR-lo) |
| :--- | :--- |
| AND | A |
| JP | Z,0970,SA-CONTRL |

Jump forward when handling a SAVE command

In the case of a LOAD, VERIFY or MERGE command the first seventeen bytes of the 'header area' in the work space hold the prepared information, as detailed above; and it is now time to fetch a 'header' from the tape.

| LD | BC, +0011 | Form in the IX register pair |
| :--- | :--- | :--- |
| ADD | IX,BC | the base address of the 'second |
|  |  | header area'. |

Now enter a loop; leaving it only when a 'header' has been LOADed.
0767
LD-LOOK-H

| PUSH | IX |
| :--- | :--- |
| LD | DE,+0011 |
| XOR | A |
| SCF |  |
| CALL | 0556, LD-BYTES |
| POP | IX |
| JR | NC,0767,LD-LOOK-H |

Make a copy of the base address.
LOAD seventeen bytes.
Signal 'header'.
Signal 'LOAD'.
Now look for a header
Retrieve the base address.
Go round the loop until successful

The new 'header' is now displayed on the screen but the routine will only proceed if the 'new' header matches the 'old' header.

|  | LD | A,+FE |
| :--- | :--- | :--- |
|  | CALL | 1601,CHAN-OPEN | | Ensure that channel 'S' |
| :--- |
| is open. |

The appropriate message - 'Program:', 'Number array:', 'Character array:' or 'Bytes:' is printed.

| LD | DE, +09C0 | The base address of the message <br> block. |
| :--- | :--- | :--- |
| PUSH | BC | Save the C register whilst <br> CALL |
| OCOA,PO-MSG | the appropriate message is <br> POP | BC |

The 'new name' is printed and as this is done the 'old' and the 'new' names are compared.

| PUSH | IX | Make the DE register pair |
| :--- | :--- | :--- |
| POP | DE | point to the 'new type' and |
| LD | HL,+FFF0 | the HL register pair to the |
| ADD | HL,DE | 'old name'. |
| LD | B,+0A | Ten characters are to be <br> considered. |
|  |  | Jump forward if the match is |
| LD | A,(HL) | to be against an actual name. |
| INC | A |  |
| JR | NZ,07A6,LD-NAME |  |
| LD | A,C | But if the 'old name' is 'null' |
| ADD | A,B | then signal 'ten characters |
| LD | C,A | already match'. |

A loop is entered to print the characters of the 'new name'. The name will be accepted if the 'counter' reaches zero, at least.

| 07A6 | LD-NAME | INC | DE | Consider each character of the |
| :---: | :---: | :---: | :---: | :---: |
|  |  | LD | A,(DE) | 'new name' in turn. |
|  |  | CP | (HL) | Match it against the appropriate |
|  |  | INC | HL | character of the 'old name'. |
|  |  | JR | NZ,07AD,LD-CH-PR | Do not count it if it does not |
|  |  | INC | C | does not match. |
| 07AD | LD-CH-PR | RST | 0010,PRINT-A-1 | Print the 'new' character. |
|  |  | DJNZ | 07A6,LD-NAME | Loop for ten characters. |
|  |  | BIT | 7,C | Accept the name only if the |
|  |  | JR | NZ,0767,LD-LOOK-H | counter has reached zero. |


| LD | A,+0D | Follow the 'new name' with |
| :--- | :--- | :--- |
| RST | 0010, PRINT-A-1 | a 'carriage return'. |

The correct header has been found and the time has come to consider the three commands LOAD, VERIFY, \& MERGE separately

| POP | HL | Fetch the pointer. |
| :--- | :--- | :--- |
| LD | A,(IX+00) | 'SCREEN\$ and CODE' are |
| CP | +03 | handled with VERIFY. |
| JR | Z,07CB,VR-CONTRL |  |
| LD | A,(T-ADDR-lo) | Jump forward if using a |
| DEC | A | LOAD command. |
| JP | Z,0808,LD-CONTRL |  |
| CP | +02 | Jump forward if using a MERGE |
| JP | Z,08B6,ME-CONTRL | command; continue with a |
|  |  | VERIFY command. |

## THE 'VERIFY' CONTROL ROUTINE

The verification process involves the LOADing of a block of data, a byte at a time, but the bytes are not stored - only checked. This routine is also used to LOAD blocks of data that have been described with 'SCREEN\$ \& CODE'.


The routine continues by considering the 'destination pointer'.


The VERIFY/LOAD flag is now considered and the actual LOAD made.

| 07F4 | VR-CONT-2 | $\begin{aligned} & \text { PUSH } \\ & \text { POP } \\ & \text { LD } \\ & \text { CP } \\ & \text { SCF } \\ & \text { JR } \\ & \text { AND } \end{aligned}$ | $\begin{aligned} & \text { HL } \\ & \text { IX } \\ & \text { A,(T-ADDR-lo) } \\ & +02 \\ & \text { NZ,0800,VR-CONT-3 } \\ & \text { A } \end{aligned}$ | Move the 'pointer' to the IX register pair. <br> Jump forward unless using the VERIFY command; with the carry flag signalling 'LOAD' <br> Signal 'VERIFY'. |
| :---: | :---: | :---: | :---: | :---: |
| 0800 | VR-CONT-3 | LD | A, +FF | Signal 'accept data block only' before LOADing the block. |

THE 'LOAD A DATA BLOCK' SUBROUTINE
This subroutine is common to all the 'LOADing' routines. In the case of LOAD \& VERIFY it acts as a full return from the cassette handling routines but in the case of MERGE the data block has yet to be 'MERGEd'.

| 0802 | LD-BLOCK | CALL | 0556,LD-BYTES |
| :--- | :--- | :--- | :--- |$\quad$| LOAD/VERIFY a data block. |
| :--- |
|  |
|  |
|  |
| RET |

Report R - Tape loading error

| 0806 | REPORT-R | RST 0008,ERROR-1 <br>  DEFB | Call the error handling <br> +1A |
| :--- | :--- | :--- | :--- |

THE 'LOAD' CONTROL ROUTINE
This routine controls the LOADing of a BASIC program, and its variables, or an array.

| 0808 | LD-CONTRL | LD | $\mathrm{E},(\mathrm{IX}+0 \mathrm{O})$ | Fetch the 'number of bytes' |
| :---: | :---: | :---: | :---: | :---: |
|  |  | LD | D,(IX+0C) | as given in the 'new header'. |
|  |  | PUSH | HL | Save the 'destination pointer' |
|  |  | LD | A, H | Jump forward unless trying |
|  |  | OR | L | to LOAD a previously |
|  |  | JR | NZ,0819,LD-CONT-1 | undeclared array. |
|  |  | INC | DE | Add three bytes to the |
|  |  | INC | DE | length - for the name, the |
|  |  | INC | DE | low length \& the high length |
|  |  | EX | DE,HL | of a new variable. |
|  |  | JR | 0825,LD-CONT-2 | Jump forward. |

Consider now if there is enough room in memory for the new data block.

| 0819 | LD-CONT-1 | LD | L,(IX-06) | Fetch the size of the existing |
| :---: | :---: | :---: | :---: | :---: |
|  |  | LD | H,(IX-05) | 'program+variables or array'. |
|  |  | EX | DE,HL |  |
|  |  | SCF |  | Jump forward if no extra |
|  |  | SBC | HL, DE | room will be required; taking |
|  |  | JR | C,082E,LD-DATA | into account the reclaiming o |

Make the actual test for room.
0825
LD-CONT-2 LD
DE,+0005
Allow an overhead of five bytes.
ADD
HL,DE
LD B,H
Move the result to the $B C$ register pair and make the test.

Now deal with the LOADing of arrays.


Space is now made available for the new array - at the end of the present variables area.
084C
LD-DATA-1

| LD | $\mathrm{HL},(\mathrm{E}-\mathrm{LINE})$ |
| :--- | :--- |
| DEC | HL |
|  |  |
| LD | $\mathrm{C},(I \mathrm{X}+0 \mathrm{~B})$ |
| LD | $\mathrm{B},(\mathrm{IX}+0 \mathrm{C})$ |
| PUSH | BC |

Find the pointer to the end-marker of the variables area - the ' 80 -byte'. Fetch the 'length' of the new array.
Save this 'length'.

| INC | BC |
| :--- | :--- |
| INC | BC |
| INC | BC |
| LD | A,(IX-03) |
|  |  |
| PUSH | AF |
| CALL | 1655, MAKE-ROOM |
| INC | HL |
| POP | AF |
| LD | (HL),A |
| POP | DE |
| INC | HL |
| LD | (HL),E |
| INC | HL |
| LD | (HL),D |
| INC | HL |
|  |  |
|  |  |
| PUSH | HL |
| POP | IX |
| SCF |  |
| LD | A,+FF |
| JP | 0802, LD-BLOCK |

Now deal with the LOADing of a BASIC program and its variables 0873

| LD-PROG | $\begin{aligned} & \text { EX } \\ & \text { LD } \\ & \text { DEC } \end{aligned}$ | $\begin{aligned} & \mathrm{DE}, \mathrm{HL} \\ & \mathrm{HL},(\mathrm{E}-\mathrm{LINE}) \\ & \mathrm{HL} \end{aligned}$ |
| :---: | :---: | :---: |
|  | LD | (X-PTR), IX |
|  | LD | C, (IX+OB) |
|  | LD | $\mathrm{B},(\mathrm{IX}+0 \mathrm{C})$ |
|  | PUSH | BC |
|  | CALL | 19E5,RECLAIM-1 |
|  | POP | BC |
|  | PUSH | HL |
|  | PUSH | BC |
|  | CALL | 1655,MAKE-ROOM |
|  | LD | IX, (X-PTR) |
|  | INC | HL |
|  | LD | C, (IX+0F) |
|  | LD | $\mathrm{B},(\mathrm{IX}+10)$ |
|  | ADD | HL, BC |
|  | LD | (VARS), HL |
|  | LD | H,(IX+0E) |
|  | LD | A,H |
|  | AND | +C0 |
|  | JR | NZ,08AD,LD-PROG-1 |
|  | LD | L,(IX+0D) |
|  | LD | (NEWPPC), HL |
|  | LD | (NSPPC),+00 |

The data block can now be LOADed.

| LD-PROG-1 | POP | DE |
| :--- | :--- | :--- |
|  | POP | IX |
|  | SCF |  |
|  | LD | A,+FF |
|  | JP | 0802, LD-BLOCK |

Add three bytes - one for the name and two for the 'length'.
'IX +0 E ' of the old header gives the name of the array. The name is saved whilst the appropriate amount of room is made available. In effect 'BC' spaces before the 'new 80-byte'. The name is entered.
The 'length' is fetched and its two bytes are also entered.

HL now points to the first location that is to be filled with data from the tape. This address is moved to the IX register pair; the carry flag set; 'data block' is signalled; and the block LOADed.

Save the 'destination pointer'
Find the address of the end-marker of the current variables area - the ' 80 -byte'. Save IX temporarily. Fetch the 'length' of the new data block.
Keep a copy of the 'length' whilst the present program and variables areas are reclaimed. Save the pointer to the program area and the length of the new data block.
Make sufficient room available for the new program and its variables.
Restore the IX register pair. The system variable VARS has also to be set for the new program.

If a line number was
specified then it too has to be considered.
Jump if 'no number'; otherwise set NEWPPC \& NSPPC

Fetch the 'length'.
Fetch the 'start'.
Signal 'LOAD'.
Signal 'data block' only.
Now LOAD it.

## THE 'MERGE' CONTROL ROUTINE

There are three main parts to this routine.
I. LOAD the data block into the work space.
II. MERGE the lines of the new program into the old program.
III. MERGE the new variables into the old variables.

Start therefore with the LOADing of the data block.

| 08B6 | ME-CONTRL | LD | C, (IX +0 O ) | Fetch the 'length' of the |
| :---: | :---: | :---: | :---: | :---: |
|  |  | LD | $B,(\mathrm{IX}+0 \mathrm{C})$ | data block. |
|  |  | PUSH | BC | Save a copy of the 'length'. |
|  |  | INC | BC | Now made 'length+1' locations |
|  |  | RST | 0030,BC-SPACES | available in the work space. |
|  |  | LD | (HL), +80 | Place an end-marker in the extra location. |
|  |  | EX | DE,HL | Move the 'start' pointer to the HL register pair. |
|  |  | POP | DE | Fetch the original 'length'. |
|  |  | PUSH | HL | Save a copy of the 'start'. |
|  |  | PUSH | HL | Now set the IX register pair |
|  |  | POP | IX | for the actual LOAD. |
|  |  | SCF |  | Signal 'LOAD'. |
|  |  | LD | A,+FF | Signal 'data block only'. |
|  |  | CALL | 0802,LD-BLOCK | LOAD the data block. |

The lines of the new program are MERGEd with the lines of the old program.

| POP | HL | Fetch the 'start' of the new <br> program. <br> Initialise DE to the 'start' of <br> the old program. |
| :--- | :--- | :--- |
| LD | DE,(PROG) | then |
|  |  |  |

Enter a loop to deal with the lines of the new program.

| 08D2 | ME-NEW-LP | LD | A,(HL) |
| :--- | :--- | :--- | :--- |$\quad$| Fetch a line number and test |
| :--- |
|  |
|  |
|  |
|  |
|  |
|  |
|  |
|  |
|  |

Now enter an inner loop to deal with the lines of the old program.

| 08D7 | ME-OLD-LP | LD | A,(DE) | Fetch the high line number |
| :--- | :--- | :--- | :--- | :--- |
|  |  | INC | DE | byte and compare it. |
|  |  | CP | (HL) | Jump forward if it does not |
|  |  | INC | HL | match but in any case advance |

In a similar manner the variables of the new program are MERGEd with the variables of the old program.
A loop is entered to deal with each of the new variables in turn.

| LD | C,A | turn and test it. |
| :--- | :--- | :--- |
| CP | +80 | Return when all the variables |
| RET | Z | have been considered. |
| PUSH | HL | Save the current new pointer |
| LD | HL,(VARS) | Fetch VARS (for the old <br> program). |

Now enter an inner loop to search the existing variables area.

| 08F9 | ME-OLD-VP | LD | A,(HL) | Fetch each variable name and |
| :---: | :---: | :---: | :---: | :---: |
|  |  | CP | +80 | test it. |
|  |  | JR | Z,0923,ME-VAR-L2 | Jump forward once the end |
|  |  |  |  | marker is found. (Make an 'addition'.) |
|  |  | CP | C | Compare the names 0 st. bytes). |
|  |  | JR | Z,0909,ME-OLD-v2 | Jump forward to consider it further; returning here if it proves not to match fully. |
| 0901 | ME-OLD-V1 | PUSH | BC | Save the new variable's name |
|  |  | CALL | 19B8,NEXT-ONE | whilst the next 'old variable' |
|  |  | POP | BC | is located. |
|  |  | EX | DE,HL | Restore the pointer to the |
|  |  | JR | 08F9,ME-OLD-VP | D E register pair and go round |

The old and new variables match with respect to their first bytes but variables with long names will need to be matched fully.

| 0909 | ME-OLD-V2 | AND | +E0 |
| :--- | :--- | :--- | :--- |
|  | CP | +A0 | Consider bits 7, $6 \& 5$ only. |
|  | JR | NZ,0921,ME-VAR-L1 | Accept all the variable types |
|  | POP | DE | except 'long named variables'. |
|  | PUSH | DE | Make DE point to the first |
|  | PUSH | HL | character of the 'new name'. |
|  |  |  | Save the pointer to the 'old |
|  |  | name'. |  |

Enter a loop to compare the letters of the long names.

| 0912 | ME-OLD-V3 | INC | HL | Update both the 'old' and the |
| :---: | :---: | :---: | :---: | :---: |
|  |  | INC | DE | 'new' pointers. |
|  |  | LD | A,(DE) | Compare the two letters |
|  |  | CP | (HL) |  |
|  |  | JR | NZ,091E,ME-OLD-V4 | Jump forward if the match fails. |
|  |  | RLA |  | Go round the loop until the |
|  |  | JR | NC,0912,ME-OLD-V3 | 'last character' is found. |
|  |  | POP | HL | Fetch the pointer to the start of the 'old' name and |
|  |  | JR | 0921,ME-VAR-L1 | jump forward - successful. |
| 091E | ME-OLD-V4 | POP | $\mathrm{HL}$ | Fetch the pointer and jump |
|  |  | JR | 0901,ME-OLD-V1 | back - unsuccessful. |

Come here if the match was found.

| 0921 | ME-VAR-L1 | LD | A, + FF |
| :--- | :--- | :--- | :--- |
| And here if not. (A holds +80 - variable to be 'added'.) | Signal 'replace' variable. |  |  |
| 0923 | ME-VAR-L2 | POP | DE |

THE 'MERGE A LINE OR A VARIABLE' SUBROUTINE
This subroutine is entered with the following parameters:


The new entry can now be made.

| 093E | ME-ENT-1 | EX <br> PUSH | AF,A'F' DE |
| :---: | :---: | :---: | :---: |
|  |  | CALL | 19B8,NEXT-ONE |
|  |  | LD | (X-PTR), HL |
|  |  | LD | HL,(PROG) |
|  |  | EX | (SP), HL |
|  |  | PUSH | BC |
|  |  | EX | AF, A'F' |
|  |  | JR | C,0955,ME-ENT-2 |
|  |  | DEC | HL |
|  |  | CALL INC | $\begin{aligned} & \text { 1655,MAKE-ROOM } \\ & \text { HL } \end{aligned}$ |
|  |  | JR | 0958,ME-ENT-3 |
| 0955 | ME-ENT-2 | CALL | 1655,MAKE-ROOM |
| 0958 | ME-ENT-3 | INC | HL |
|  |  | POP | BC |
|  |  | POP | DE |
|  |  | LD | (PROG),DE |
|  |  | LD | DE,(X-PTR) |
|  |  | PUSH | BC |
|  |  | PUSH | DE |
|  |  | EX | DE,HL |
|  |  | LDIR |  |

Save the flags.
Make a copy of the
'destination' pointer.
Find the length of the 'new' variable/line.
Save the pointer to the 'new' variable/line.
Fetch PROG - to avoid corruption.
Save PROG on the stack and fetch the 'new' pointer.
Save the length.
Retrieve the flags.
Jump forward if adding a new variable.
A new line is added before the 'destination' location. Make the room for the new line.

Jump forward.
Make the room for the new variable.
Point to the 1st new location.
Retrieve the length.
Retrieve PROG and store it in its correct place.
Also fetch the 'new' pointer.
Again save the length and the new' pointer.
Switch the pointers and copy the 'new' variable/line into the room made for it.

The 'new' variable/line has now to be removed from the work space.

| POP | HL | Fetch the 'new' pointer. <br> POP |
| :--- | :--- | :--- |
| PUSH | DE | Fetch the length. <br> Save the 'old' pointer. (Points <br> to the location after the 'added' |
| CALL | 19E8,RECLAIM-2 | variable/line.) <br> Remove the variable/line from <br> the work space. |
| POP | DE | Return with the 'old' pointer <br> in the DE register pair. |
| RET |  | ind |

## THE 'SAVE' CONTROL ROUTINE

The operation of SAVing a program or a block of data is very straightforward.

## 0970

SA-CONTRL

| PUSH | HL |
| :--- | :--- |
| LD | A,+FD |
| CALL | 1601,CHAN-OPEN |
| XOR | A |
| LD | DE,+09A1 |
| CALL | 0COA,PO-MSG |
| SET | 5,(TV-FLAG) |
| CALL | 15D4,WAIT-KEY |

Upon receipt of a keystroke the 'header' is saved.

| PUSH | IX |
| :--- | :--- |
| LD | DE, +0011 |
| XOR | A |
| CALL | $04 C 2$, SA-BYTES |

Ensure that channel 'K'
Ensure that channel K
is open.
Signal 'first message'.
Print the message - Start tape,
then press any key.'.
Signal 'screen will require to be cleared'.
Wait for a key to be pressed.

Save the base address of the 'header' on the machine stack. Seventeen bytes are to be SAVEd.
Signal 'It is a header'.
Send the 'header'; with a leading 'type' byte and a trailing 'parity' byte.

There follows a short delay before the program/data block is SAVEd.

|  | POP | IX | Retrieve the pointer to the <br> 'header'. |
| :--- | :--- | :--- | :--- |
| 0991 SA-1-SEC | LD | B,+32 | The delay is for fifty <br> interrupts, i.e. one second. |
|  | HALT | 0991, SA-1-SEC |  |
| DJNZ | Fetch the length of the |  |  |
| LD | D,(IX+0B) | data block that is to be SAVEd. | Signal 'data block'. |

THE CASSETTE MESSAGES
Each message is given with the last character inverted (+80 hex.).

| 09A1 | DEFB +80 | - Initial byte is stepped over. |
| :--- | :--- | :--- |
| 09A2 | DEFM | - Start tape, then press any key. |
| 09C1 | DEFM | - 'carriage return' - Program: |
| 09CB | DEFM | - 'carriage return' - Number array: |
| 09DA | DEFM | - 'carriage return' - Character array: |
| 09EC | DEFM | - 'carriage return' - Bytes: |

## THE SCREEN \& PRINTER HANDLING ROUTINES

## THE 'PRINT-OUT' ROUTINES

All of the printing to the main part of the screen, the lower part of the screen and the printer is handled by this set of routines. The PRINT-OUT routine is entered with the A register holding the code for a control character, a printable character or a token.

| 09F4 | PRINT-OUT | CALL <br> CP <br> JP <br> CP <br> JR <br> CP <br> JR <br> LD <br> LD <br> LD <br> ADD <br> LD <br> ADD <br> PUSH <br> JP | $\begin{aligned} & \text { 0B03,PO-FETCH } \\ & \text { +20 } \\ & \text { NC,0AD9,PO-ABLE } \\ & \text { +06 } \\ & \text { C,0A69,PO-QUEST } \\ & \text { +18 } \\ & \text { NC,0A69,PO-QUEST } \\ & \text { HL,+0A0B } \\ & \text { E,A } \\ & \text { D,+00 } \\ & \text { HL,DE } \\ & \text { E,(HL) } \\ & \text { HL,DE } \\ & \text { HL } \\ & \text { 0B03,PO-FETCH } \end{aligned}$ | The current print position. If the code represents a printable character then jump. Print a question mark for codes in the range $+00-+05$. And also for codes $+18-+1$ F. <br> Base of 'control' table. Move the code to the DE register pair. Index into the table and fetch the offset. <br> Add the offset and make an indirect jump to the appropriate subroutine. |
| :---: | :---: | :---: | :---: | :---: |

## THE 'CONTROL CHARACTER' TABLE

| address | offset | character | address | offset | character |
| :--- | :--- | :--- | :--- | :--- | :--- |
|  |  |  |  |  |  |
| OA11 | 4 E | PRINT comma | OA1A | 4 F | not used |
| OA12 | 57 | EDIT | OA1B | 5 F | INK control |
| OA13 | 10 | cursor left | OA1C | 5E | PAPER control |
| OA14 | 29 | cursor right | OA1D | 5D | FLASH control |
| OA15 | 54 | cursor down | OA1E | 5C | BRIGHT control |
| OA16 | 53 | cursor up | OA1F | 5B | INVERSE control |
| OA17 | 52 | DELETE | OA20 | 5A | OVER control |
| OA18 | 37 | ENTER | OA21 | 54 | AT control |
| OA19 | 50 | not used | OA22 | 53 | TAB control |

## THE 'CURSOR LEFT' SUBROUTINE

The subroutine is entered with the $B$ register holding the current line number and the $C$ register with the current column number.

| 0A23 | PO-BACK-1 | INC | C | Move leftwards by one column. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | LD | A, +22 | Accept the change unless |
|  |  | CP | C | up against the lefthand side. |
|  |  | JR | NZ,0A3A,PO-BACK-3 |  |
|  |  | BIT | 1,(FLAGS) | If dealing with the printer |
|  |  | JR | NZ,0A38,PO-BACK-2 | jump forward. |
|  |  | INC | B | Go up one line. |
|  |  | LD | C, +02 | Set column value. |
|  |  | LD | A, +18 | Test against top line. |
|  |  | CP | B | Note: This ought to be +19 . |
|  |  | JR | NZ,0A3A,PO-BACK-3 | Accept the change unless at the top of the screen. |
|  |  | DEC | B | Unacceptable so down a line. |
| 0A38 | PO-BACK-2 | LD | C,+21 | Set to lefthand column. |
| 0A3A | PO-BACK-3 | JP | 0DD9,CL-SET | Make an indirect return via |

THE 'CURSOR RIGHT' SUBROUTINE
This subroutine performs an operation identical to the BASIC statement - PRINT OVER 1;CHR\$ 32; -

| OA3D PO-RIGHT LD | A,(P-FLAG) | Fetch P-FLAG and save it on |  |
| :--- | :--- | :--- | :--- |
|  | PUSH | AF | the machine stack. |


| LD | (P-FLAG),+01 | Set P-FLAG to OVER 1. |
| :--- | :--- | :--- |
| LD | A,+20 | A 'space'. |
| CALL | OB65,PO-CHAR | Print the character. |
| POP | AF | Fetch the old value of |
| LD | (P-FLAG),A | P-FLAG. |
| RET |  | Finished. |
|  |  | Note: The programmer has <br> forgotten to exit via PO-STORE. |

## THE 'CARRIAGE RETURN' SUBROUTINE

If the printing being handled is going to the printer then a carriage return character leads to the printer buffer being emptied. If the printing is to the screen then a test for 'scroll?' is made before decreasing the line number.

| 0A4F | PO-ENTER | BIT | 1,(FLAGS) |
| :--- | :--- | :--- | :--- |
|  | JP | NZ,0ECD,COPY-BUFF | Jump forward if handling |
|  | LD | C,+21 | the printer. |
|  | CALL | 0C55,PO-SCR | Set to lefthand column. |
|  | DEC | B | Scroll if necessary. |
|  | JP | 0DD9,CL-SET | Now down a line. |
|  |  |  | Make an indirect return via |
|  |  |  | CL-SET \& PO-STORE. |

## THE 'PRINT COMMA' SUBROUTINE

The current column value is manipulated and the A register set to hold +00 (for TAB 0 ) or +10 (for TAB 16).
0A5F

| PO-COMMA | CALL | OB03,PO-FETCH |
| :--- | :--- | :--- |
|  | LD | A,C |
|  | DEC | A |
|  | DEC | A |
|  | AND | +10 |
|  |  |  |
|  | JR | $0 A C 3$, PO-FILL |

Why again?
Current column number Move rightwards by two columns and then test.
The A register will be +00 or +10 .
Exit via PO-FILL.

## THE 'PRINT A QUESTION MARK' SUBROUTINE

A question mark is printed whenever an attempt is made to print an unprintable code.

| 0A69 PO-QUEST | LD | A,+3F | The character '?'. |
| :--- | :--- | :--- | :--- |
|  |  | JR | 0AD9,PO-ABLE |

## THE 'CONTROL CHARACTERS WITH OPERANDS' ROUTINE

The control characters from INK to OVER require a single operand whereas the control characters AT \& TAB are required to be followed by two operands.
The present routine leads to the control character code being saved in TVDATA-lo, the first operand in TVDATA-hi or the A register if there is only a single operand required, and the second operand in the A register.

| 0A6D | PO-TV-2 | LD | DE,+0A87 |
| :--- | :--- | :--- | :--- |
|  |  | LD | (TVDATA-hi),A |
|  |  | JR | OA80,PO-CHANGE |

Save the first operand in TVDATA-hi and change the address of the 'output' routine to PO-CONT (+0A87).

Enter here when handling the characters AT \& TAB.
$0 A 75$ PO-2-OPER LD DE,+0A6D

The character code will be saved in TVDATA-lo and the address of the 'output' routine changed to PO-TV-2 (+0A6D).

Enter here when handling the colour items - INK to OVER.

| OA7A | PO-1-OPER | LD | DE,+0A87 | The 'output' routine is to be <br> changed to PO-CONT (+OA87). |
| :--- | :--- | :--- | :--- | :--- |
| 0A7D | PO-TV-1 | LD | (TVDATA-lo),A | Save the control character code. |

The current 'output' routine address is changed temporarily.

HL,(CURCHL)

| LD | $(\mathrm{HL}), \mathrm{E}$ |
| :--- | :--- |
| INC | HL |
| LD | $(\mathrm{HL}), \mathrm{D}$ |

Once the operands have been collected the routine continues.

| 0 0A87 | PO-CONT | LD |
| :--- | :--- | :--- |
|  | CALL | DE,+09F4 |
|  | LD | OA80,PO-CHANGE |
|  |  | HL,(TVDATA) |
|  |  |  |
|  | LD | D,A |
|  | LD | A,L |
|  | CP | +16 |
|  | JP | C,2211,CO-TEMPS |
|  | JR | NZ,0AC2,PO-TAB |

Now deal with the AT control character.

|  |  | LD | B, H |
| :---: | :---: | :---: | :---: |
|  |  | LD | C, D |
|  |  | LD | A,+1F |
|  |  | SUB | C |
|  |  | JR | C,OAAC,PO-AT-ERR |
|  |  | ADD | A,+02 |
|  |  | LD | C, A |
|  |  | BIT | 1,(FLAGS) |
|  |  | JR | NZ,0ABF,PO-AT-SET |
|  |  | LD | A,+16 |
|  |  | SUB | B |
| OAAC | PO-AT-ERR | JP | C,1E9F,REPORT-B |
|  |  | INC | A |
|  |  | LD | B,A |
|  |  | INC | B |
|  |  | BIT | 0,(TV-FLAG) |
|  |  | JP | NZ,0C55,PO-SCR |
|  |  | CP | (DF-SZ) |
|  |  | JP | C,0C86,REPORT-5 |
| OABF | PO-AT-SET | JP | 0D09,CL-SET |

And the TAB control character.

| $0 \mathrm{AC2}$ | PO-TAB | LD | A, H | Fetch the first operand. |
| :---: | :---: | :---: | :---: | :---: |
| 0AC3 | PO-FILL | CALL | OB03,PO-FETCH | The current print position. |
|  |  | ADD | A, C | Add the current column value. |
|  |  | DEC | A | Find how many 'spaces', modulo |
|  |  | AND | +1F | 32 , are required and return |
|  |  | RET | Z | if the result is zero. |
|  |  | LD | D,A | Use 0 as the counter. |
|  |  | SET | 0,(FLAGS) | Suppress 'leading space'. |
| OADO | PO-SPACE | LD | A,+20 | Print 'D number' of |
|  |  | CALL | 0С3B,PO-SAVE | spaces. |
|  |  | DEC | D |  |
|  |  | JR | NZ,0AD0,PO-SPACE |  |
|  |  | RET |  | Now finished. |

PRINTABLE CHARACTER CODES.
The required character (or characters) is printed by calling PO-ANY followed by PO-STORE.

THE 'POSITION STORE' SUBROUTINE
The new position's 'line \& column' values and the 'pixel' address are stored in the appropriate system variables.

| OADC | PO-STORE | BIT | 1,(FLAGS) | Jump forward if handling |
| :---: | :---: | :---: | :---: | :---: |
|  |  | JR | NZ,0AFC,PO-ST-PR | the printer. |
|  |  | BIT | 0,(TV-FLAG) | Jump forward if handling the |
|  |  | JR | NZ,0AF0,PO-ST-E | lower part of the screen. |
|  |  | LD | (S-POSN),BC | Save the values that relate |
|  |  | LD | (DF-CC), HL | to the main part of the |
|  |  | RET |  | screen. Then return. |
| OAFO | PO-ST-E | LD | (S-POSNL),BC | Save the values that relate |
|  |  | LD | (ECHO-E), BC | to the lower part of the |
|  |  | LD | (DF-CCL), HL | screen. |
|  |  | RET |  | Then return. |
| OAFC | PO-ST-PR | LD | (P-POSN), C | Save the values that relate |
|  |  | LD | (PR-CC), HL | to the printer buffer. |
|  |  | RET |  | Then return. |

## THE 'POSITION FETCH' SUBROUTINE

The current position's parameters are fetched from the appropriate system variables.

| 0B03 | PO-FETCH | BIT | 1,(FLAGS) |
| :--- | :--- | :--- | :--- |
|  | JR | NZ,OB1D,PO-F-PR | Jump forward if handling |
|  |  | LD | BC,(S-POSN) |

THE 'PRINT ANY CHARACTER(S)' SUBROUTINE
Ordinary character codes, token codes and user-defined graphic codes, and graphic codes are dealt with separately.

| 0B24 | PO-ANY | CP | +80 | Jump forward with ordinary |
| :---: | :---: | :---: | :---: | :---: |
|  |  | JR | C,0B65,PO-CHAR | character codes. |
|  |  | CP | +90 | Jump forward with token |
|  |  | JR | NC,0B52,PO-T\&UDG | codes and UDG codes. |
|  |  | LD | B,A | Move the graphic code. |
|  |  | CALL | 0B38,PO-GR-1 | Construct the graphic form. |
|  |  | CALL | 0B03,PO-FETCH | HL has been disturbed so 'fetch' again. |
|  |  | LD | DE, +5 C 92 | Make DE point to the start of the graphic form; i.e. MEMBOT. |
|  |  | JR | 0B7F,PO-ALL | Jump forward to print the graphic character. |

Graphic characters are constructed in an Ad Hoc manner in the calculator's memory area; i.e. MEM-0 \& MEM-1.

| 0B38 | PO-GR-1 | LD | HL,+5C92 |  |
| :--- | :--- | :--- | :--- | :--- |
| OB3E | CALL | OB3E,PO-GR-2 | This is MEMBOT. <br> In effect call the following <br> subroutine twice. |  |
|  | PO-GR-2 | RR | B | Determine bit 0 (and later bit 2) <br> of the graphic code. |
|  |  | SBC | AND A A | +0F |


|  |  | LD | C,A | Save the result in C. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | RR | B | Determine bit 1 (and later bit 3) |
|  |  | SBC | A, A | of the graphic code. |
| 0B4C | PO-GR-3 | AND | +F0 | The A register will hold +00 or + FO. |
|  |  | OR | C | The two results are combined. |
|  |  | LD | C, +04 | The A register holds half the |
|  |  | LD | (HL),A | character form and has to be |
|  |  | INC | HL | used four times. |
|  |  | DEC | C | This is done for the upper |
|  |  | JR | NZ,0B4C,PO-GR-3 | half of the character form |
|  |  | RET |  | and then the lower. |

Token codes and user-defined graphic codes are now separated.

| 0B52 | PO-T\&UDG | SUB | +A5 | Jump forward with token codes |
| :---: | :---: | :---: | :---: | :---: |
|  |  | JR | NC,0B5F,PO-T |  |
|  |  | ADD | A,+15 | UDG codes are now +00-+0F. |
|  |  | PUSH | BC | Save the current position values on the machine stack. |
|  |  | LD | BC,(UDG) | Fetch the base address of the |
|  |  | JR | 0B6A,PO-CHAR-2 | UDG area and jump forward. |
| 0B5F | PO-T | CALL | 0C10,PO-TOKENS | Now print the token and return |
|  |  | JP | 0B03,PO-FETCH | via PO-FETCH. |

The required character form is identified.

| 0B65 | PO-CHAR | $\begin{aligned} & \text { PUSH } \\ & \text { LD } \end{aligned}$ | $\begin{aligned} & \mathrm{BC} \\ & \mathrm{BC},(\mathrm{CHARS}) \end{aligned}$ | The current position is saved. The base address of the character area is fetched. |
| :---: | :---: | :---: | :---: | :---: |
| 0B6A | PO-CHAR-2 | EX | DE, HL | The print address is saved. |
|  |  | LD | HL, +5C3B | This is FLAGS. |
|  |  | RES | 0,(HL) | Allow for a leading space |
|  |  | CP | +20 | Jump forward if the character |
|  |  | JR | NZ,0B76,PO-CHAR-3 | is not a 'space'. |
|  |  | SET | 0,(HL) | But 'suppress' if it is. |
| 0B76 | PO-CHAR-3 | LD | H, +00 | Now pass the character code |
|  |  | LD | L,A | to the HL register pair. |
|  |  | ADD | HL, HL | The character code is in |
|  |  | ADD | HL, HL | effect multiplied by 8. |
|  |  | ADD | HL, HL |  |
|  |  | ADD | HL, BC | The base address of the character form is found. |
|  |  | POP | BC | The current position is fetched |
|  |  | EX | DE,HL | and the base address passed to the $D E$ register pair. |

## THE 'PRINT ALL CHARACTERS' SUBROUTINE

This subroutine is used to print all ' $8 * 8$ ' bit characters. On entry the DE register pair holds the base address of the character form, the HL register the destination address and the BC register pair the current 'line \& column' values.

| 0B7F | PR-ALL | LD | A, C | Fetch the column number. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | DEC | A | Move one column rightwards. |
|  |  | LD | A, +21 | Jump forward unless a new |
|  |  | JR | NZ,0893,PR-ALL-1 | line is indicated. |
|  |  | DEC | B | Move down one line. |
|  |  | LD | C,A | Column number is +21. |
|  |  | BIT | 1,(FLAGS) | Jump forward if handling |
|  |  | JR | Z,0B93,PR-ALL-1 | the screen. |
|  |  | PUSH | DE | Save the base address whilst |
|  |  | CALL | OECD,COPY-BUFF | the printer buffer is |
|  |  | POP | DE | emptied. |
|  |  | LD | A,C | Copy the new column number. |
| 0B93 | PR-ALL-1 | CP | C | Test whether a new line is |



The character can now be printed. Eight passes of the loop are made - one for each 'pixel-line'.

| 0BB7 | PR-ALL-4 | EX | AF, $A^{\prime} \mathrm{F}^{\prime}$ | The carry flag is set when using the printer. Save this flag in F'. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | LD | A, (DE) | Fetch the existing 'pixel-line'. |
|  |  | AND | B | Use the 'OVER-mask' and then |
|  |  | XOR | (HL) | XOR the result with the 'pixelline' of the character form. |
|  |  | XOR | C | Finally consider the 'INVERSEmask'. |
|  |  | LD | (DE), A | Enter the result. |
|  |  | EX | AF,A'F' | Fetch the printer flag and |
|  |  | JR | C,0BD3,PR-ALL-6 | jump forward if required. |
|  |  | INC | D | Update the destination address |
| 0BC1 | PR-ALL-5 | INC | HL | Update the 'pixel-line' of the character form. |
|  |  | $\begin{aligned} & \text { DEC } \\ & \text { JR } \end{aligned}$ | A NZ,0BB7,PR-ALL-4 | Decrease the counter and loop back unless it is zero. |

Once the character has been printed the attribute byte is to set as required.

| EX | DE,HL | Make the H register hold a <br> correct high-address for the <br> character area. |
| :--- | :--- | :--- |
| DEC | H | Set the attribute byte only if <br> handling the screen. |
| BIT | 1,(FLAGS) | Restore the original <br> CALL |
| Zestination address and the |  |  |
| POP | HL | position values. <br> POP |
| BC | Decrease the column number <br> and increase the destination <br> address before returning. |  |
| DEC | C | HL |

When the printer is being used the destination address has to be updated in increments of +20 .

| PR-ALL-6 | EX | AF,A'F' |
| :--- | :--- | :--- |
|  | LD | A,+20 |
|  | ADD | A,E |
|  | LD | E,A |
|  | EX | AF,A'F' |
|  | JR | $0 B C 1$, PR-ALL-5 |

Save the printer flag again.
The required increment value
Add the value and pass the result back to the E register

Fetch the flag.
Jump back into the loop.

## THE 'SET ATTRIBUTE BYTE' SUBROUTINE

The appropriate attribute byte is identified and fetched. The new value is formed by manipulating the old value, ATTR-T, MASK-T and P-FLAG. Finally this new value is copied to the attribute area.


## THE 'MESSAGE PRINTING' SUBROUTINE

This subroutine is used to print messages and tokens. The A register holds the 'entry number' of the message or token in a table. The DE register pair holds the base address of the table.

| 0C0A | PO-MSG | PUSH | HL |
| :--- | :--- | :--- | :--- | | The high byte of the last |
| :--- |
|  |
|  |
|  |
|  |
|  |
|  |
|  |
|  |
|  |
|  |
|  |

Enter here when expanding token codes.

| $0 C 10$ | PO-TOKENS LD | DE095 | The base address of the token <br> table. |
| :--- | :--- | :--- | :--- |
|  | PUSH AF | Save the code on the stack. | (Range $+00-+5 A ;$ RND - COPY). |

The table is searched and the correct entry printed.

| 0C14 | PO-TABLE | CALL | OC41,PO-SEARCH |
| :--- | :--- | :--- | :--- | | Locate the required entry. |
| :--- |
|  |
|  |
|  |
|  |
|  |
|  |
|  |
|  |
| LR |


| BIT | 0,(FLAGS) | before the message/token |
| :--- | :--- | :--- |
| CALL | Z,OC3B,PO-SAVE | if required. |

The characters of the message/token are printed in turn.

| 0 C 22 | PO-EACH | LD | A,(DE) |
| :--- | :--- | :--- | :--- |
|  |  | AND | +7 F |
|  |  | CALL | OC3B,PO-SAVE |
|  | LD | A,(DE) |  |
|  | INC | DE |  |
|  |  | ADD | A,A |
|  |  | JR | NC,0C22,PO-EACH |

Collect a code.
Cancel any 'inverted bit'.
Print the character.
Collect the code again.
Advance the pointer.
The 'inverted bit' goes to the carry flag and signals the end of the message/token; otherwise jump back.

Now consider whether a 'trailing space' is required.

|  |  | POP | DE | For messages - D holds +00 ; for tokens - D holds $+00-+5 \mathrm{~A}$ |
| :---: | :---: | :---: | :---: | :---: |
|  |  | CP | +48 | Jump forward if the last |
|  |  | JR | Z,0C35,PO-TRSP | character was a '\$'. |
|  |  | CP | +82 | Return if the last character |
|  |  | RET | C | was any other before ' A '. |
| 0C35 | PO-TR-SP | LD | A, D | Examine the value in D and |
|  |  | CP | +03 | return if it indicates a |
|  |  | RET | C | message, RND, INKEY\$ or PI. |
|  |  | LD | A, +20 | All other cases will require a 'trailing space'. |

## THE 'PO-SAVE' SUBROUTINE

This subroutine allows for characters to be printed 'recursively'. The appropriate registers are saved whilst 'PRINT-OUT' is called.

| 0C3B | PO-SAVE | PUSH | DE |
| :--- | :--- | :--- | :--- |
|  | EXX |  | Save the DE register pair. <br> RST |
|  | EXX | 0010, PRINT-A-1 | Save HL \& BC. |
|  | POP |  | Restore HL \& BC. |
|  | RET | DE | Restore DE. |

## THE 'TABLE SEARCH' SUBROUTINE

The subroutine returns with the DE register pair pointing to the initial character of the required entry and the carry flag reset if a 'leading space' is to be considered.

| 0C41 | PO-SEARCH | $\begin{aligned} & \text { PUSH } \\ & \text { EX } \\ & \text { INC } \end{aligned}$ | $\begin{aligned} & \text { AF } \\ & \text { DE,HL } \\ & \text { A } \end{aligned}$ | Save the 'entry number'. HL now holds the base address. Make the range +01 - ? |
| :---: | :---: | :---: | :---: | :---: |
| 0C44 | PO-STEP | BIT | 7,(HL) | Wait for an 'inverted |
|  |  | INC | HL | character'. |
|  |  | JR | Z,0C44,PO-STEP |  |
|  |  | DEC | A | Count through the entries |
|  |  | JR | NZ,0C44,PO-STEP | until the correct one is found. |
|  |  | EX | DE,HL | DE points to the initial character. |
|  |  | POP | AF | Fetch the 'entry number' and |
|  |  | CP | +20 | return with carry set for the |
|  |  | RET | C | first thirty two entries. |
|  |  | LD | A,(DE) | However if the initial |
|  |  | SUB | +41 | character is a letter then a |
|  |  | RET |  | leading space may be needed. |

## THE 'TEST FOR SCROLL' SUBROUTINE

This subroutine is called whenever there might be the need to scroll the display. This occurs on three occasions; i. when handling a 'carriage return' character; ii. when using AT in an INPUT line; \& iii. when the current line is full and the next line has to be used. On entry the B register holds the line number under test.

| 0 C 55 | PO-SCR | BIT <br> RET <br> LD <br> PUSH <br> LD <br> BIT <br> JP <br> CP <br> JR <br> RET | 1,(FLAGS) <br> NZ <br> DE,+0DD9 <br> DE <br> A, B <br> 0,(TV-FLAG) <br> NZ,0D02,PO-SCR-4 <br> (DF-SZ) <br> C,0C86,REPORT-6 <br> NZ |
| :---: | :---: | :---: | :---: |
|  |  | $\begin{aligned} & \text { BIT } \\ & \text { JR } \\ & \text { LD } \\ & \text { DEC } \\ & \text { JR } \end{aligned}$ | $\begin{aligned} & \text { 4,(TV-FLAG) } \\ & \text { Z,0C88,PO-SCR-2 } \\ & \text { E,(BREG) } \\ & \text { E } \\ & \text { Z,0CD2,PO-SCR-3 } \end{aligned}$ |
|  |  | LD CALL LD RES RET | $\mathrm{A},+00$ <br> 1601,CHAN-OPEN <br> SP,(LIST-SP) <br> 4,(TV-FLAG) |

Report 5 - Out of screen

| $0 C 86$ |  |  |
| :--- | :--- | :--- |
| REPORT-5 | RST | 0008, ERROR-1 |
|  | DEFB | +04 |

Now consider if the prompt 'scroll?' is required.

| $0 C 88$ | PO-SCR-2 | DEC | (SCR-CT) |
| :--- | :--- | :--- | :--- |
|  |  | JR | NZ,0CD2,PO-SCR-3 |

Proceed to give the prompt message.

| LD | A,+18 |
| :--- | :--- |
| SUB | B |
| LD | (SCR-CT),A |
| LD | HL,(ATTR-T) |
| PUSH | HL |
| LD | A,(P-FLAG) |
| PUSH | AF |
| LD | A,+FD |
| CALL | 1601, CHAN-OPEN |
| XOR | A |
| LD | DE,+0CF8 |
| CALL | 0 0C0A,PO-MSG |
| SET | $5,(T V-F L A G)$ |
|  |  |
| LD | HL,+5C3B |
| SET | $3,(H L)$ |
| RES | $5,(H L)$ |
| EXX | $15 D 4$, WAIT-KEY |
| CALL |  |
| EXX | +20 |
| CP | Z,0D00,REPORT-D |
| JR | + E2 |
| CP | Z,0D00,REPORT-D |
| JR | +20 |
| OR | $+6 E$ |
| CP | Z,0D00,REPORT-D |
| JR | A,+FE |
| LD | $1601, C H A N-O P E N ~$ |
| CALL | AF |
| POP |  |

Return immediately if the printer is being used. Pre-load the machine stack with the address of 'CL-SET'.
Transfer the line number. Jump forward if considering 'INPUT ... AT ..'.
Return, via CL-SET, if the line number is greater than the value of DF-SZ; give report 5 if it is less; otherwise continue. Jump forward unless dealing with an 'automatic listing'. Fetch the line counter. Decrease this counter. Jump forward if the listing is to be scrolled. Otherwise open channel ' $K$ ', restore the stack pointer, flag that the automatic listing has finished and return via CL-SET.

Call the error handling routine.

Decrease the scroll counter and proceed to give the prompt only if is becomes zero.

The counter is reset.

The current values of ATTR-T and MASK-T are saved. The current value of P-FLAG is saved.
Channel ' K ' is opened.
The message 'scroll?' is message ' 0 '. This message is now printed.
Signal 'clear the lower screen after a keystroke'.
This is FLAGS.
Signal 'L mode'.
Signal 'no key yet'.
Note: DE should be pushed also.
Fetch a single key code.
Restore the registers.
There is a jump forward to REPORT-D - 'BREAK - CONT repeats' - if the keystroke was 'BREAK', 'STOP', 'N' or ' $n$ '; otherwise accept the keystroke as indicating the need to scroll the display. Open channel 'S'.

Restore the value of

| LD | (P-FLAG),A | P-FLAG. |
| :--- | :--- | :--- |
| POP | HL | Restore the values of ATTR-T |
| LD | (ATTR-T),HL | and MASK-T. |

The display is now scrolled.

| 0CD2 | PO-SCR-3 | CALL LD INC LD PUSH CALL LD RRCA RRCA RRCA AND OR LD | $\begin{aligned} & \text { ODFE,CL-SC-ALL } \\ & \text { B(DF-SZ) } \\ & \text { B } \\ & \text { C,+21 } \\ & \text { BC } \\ & \text { OE9B,CL-ADDR } \\ & \text { A,H } \\ & \\ & \\ & +03 \\ & \text { +58 } \\ & \text { H,A } \end{aligned}$ | The whole display is scrolled. The line and column numbers for the start of the line above the lower part of the display are found and saved. The corresponding attribute byte for this character area is then found. The HL register pair holds the address of the byte. |
| :---: | :---: | :---: | :---: | :---: |

The line in question will have 'lower part' attribute values and the new line at the bottom of the display may have 'ATTR-P' values so the attribute values are exchanged.

|  |  | LD | DE, +5AE0 | DE points to the first attribute byte of the bottom line. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | LD | A, (DE) | The value is fetched. |
|  |  | LD | C, (HL) | The 'lower part' value. |
|  |  | LD | B,+20 | There are thirty two bytes. |
|  |  | EX | DE,HL | Exchange the pointers. |
| OCFO | PO-SCR-3A | LD | (DE), A | Make the first exchange |
|  |  | LD | (HL), C | and then proceed to use the |
|  |  | INC | DE | same values for the thirty |
|  |  | INC | HL | two attribute bytes of the |
|  |  | DJNZ | 0CF0,PO-SCR-3A | two lines being handled. |
|  |  | POP | BC | The line and column numbers of the bottom line of the 'upper part' are fetched before |
|  |  | RET |  | returning. |

The 'scroll?' message
0CF8

| DEFB | +80 |
| :--- | :--- |
| DEFB | $+73,+63,+72,+6 F$ |
| DEFB | $+6 C,+6 C,+B F$ |

Initial marker - stepped over.
s-c-r-o
I-I-? (inverted).
Report 0 - BREAK - CONT repeats
0D00 REPORT-D RST 0008,ERROR-1
DEFB +0C
The lower part of the display is handled as follows:

| 0D02 | PO-SCR-4 | CP | +02 | The 'out of screen' error is |
| :---: | :---: | :---: | :---: | :---: |
|  |  | JR | C,0C86,REPORT-5 | given if the lower part is |
|  |  | ADD | A,(DF-SZ) | going to be 'too large' and a |
|  |  | SUB | +19 | return made if scrolling is |
|  |  | RET | NC | unnecessary. |
|  |  | NEG |  | The A register will now hold 'the number of scrolls to be made'. |
|  |  | PUSH | BC | The line and column numbers are now saved. |
|  |  | LD | B,A | The 'scroll number', ATTR-T |
|  |  | LD | HL,(ATTR-T) | MASK-T \& P-FLAG are all |
|  |  | PUSH | HL | saved. |
|  |  | LD | HL,(P-FLAG) |  |
|  |  | PUSH | HL |  |
|  |  | CALL | 0D40,TEMPS | The 'permanent' colour items are to be used. |

The lower part of the screen is now scrolled ' A ' number of times.

| 0D1C | PO-SCR-4A | PUSH | AF | Save the 'number'. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | LD | HL,+5C6B | This is DF-SZ. |
|  |  | LD | B,(HL) | The value in DF-SZ is |
|  |  | LD | A, B | incremented; the B register |
|  |  | INC | A | set to hold the former value and |
|  |  | LD | (HL), A | the A register the new value. |
|  |  | LD | HL, +5C89 | This is S-POSN-hi. |
|  |  | CP | (HL) | The jump is taken if only the |
|  |  | JR | C,0D2D,PO-SCR-4B | lower part of the display is to be scrolled. ( $\mathrm{B}=$ old DF -SZ). |
|  |  | INC | (HL) | Otherwise S-POSN-hi is |
|  |  | LD | B,+18 | incremented and the whole display scrolled. $(B=+18)$ |
| 0D2D | PO-SCR-4B | CALL | 0E00,CL-SCROLL | Scroll 'B' lines. |
|  |  | POP | AF | Fetch and decrement the |
|  |  | DEC | A | scroll number'. |
|  |  | JR | NZ,0D1C,PO-SCR-4A | Jump back until finished. |
|  |  | POP | HL | Restore the value of |
|  |  | LD | (P-FLAG), L | P-FLAG. |
|  |  | POP | HL | Restore the values of ATTR-T |
|  |  | LD | (ATTR-T), HL | and MASK-T. |
|  |  | LD | BC,(S-POSN) | In case S-POSN has been |
|  |  | RES | 0,(TV-FLAG) | changed CL-SET is called to |
|  |  | CALL | 0DD9,CL-SET | give a matching value to DF-CC. |
|  |  | SET | 0,(TV-FLAG) | Reset the flag to indicate that |
|  |  | POP | BC | the lower screen is being |
|  |  | RET |  | handled, fetch the line and |
|  |  |  |  | column numbers, and then |
|  |  |  |  | return. |

## THE 'TEMPORARY COLOUR ITEMS' SUBROUTINE

This is a most important subroutine. It is used whenever the 'permanent' details are required to be copied to the 'temporary' system variables. First ATTR-T \& MASK-T are considered


## THE 'CLS COMMAND' ROUTINE

In the first instance the whole of the display is 'cleared' - the 'pixels' are all reset and the attribute bytes are set to equal the value in ATTR-P - then the lower part of the display is reformed.

| 0D6B | CLS | CALL | ODAF,CL-ALL | The whole of the display is 'cleared'. |
| :---: | :---: | :---: | :---: | :---: |
| 0D6E | CLS-LOWER | LD | HL,+5C3C | This is TV-FLAG. |
|  |  | RES | 5,(HL) | Signal 'do not clear the lower screen after keystroke'. |
|  |  | SET | 0,(HL) | Signal 'lower part'. |
|  |  | CALL | 0D4D,TEMPS | Use the permanent values. i.e. ATTR-T is copied from BORDCR. |
|  |  | $\begin{aligned} & \text { LD } \\ & \text { CALL } \end{aligned}$ | $\begin{aligned} & \mathrm{B},(\mathrm{DF}-\mathrm{SZ}) \\ & 0 \mathrm{E} 44, \mathrm{CL}-\mathrm{LINE} \end{aligned}$ | The lower part of the screen is now 'cleared' with these values. |

With the exception of the attribute bytes for lines ' 22 ' \& ' 23 ' the attribute bytes for the lines in the lower part of the display will need to be made equal to ATTR-P.

|  |  | LD | HL, +5AC0 | Attribute byte at start of line '22'. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | LD | A,(ATTR-P) | Fetch ATTR-P. |
|  |  | DEC | B | The line counter. |
|  |  | JR | 0D8E,CLS-3 | Jump forward into the loop. |
| 0D87 | CLS-1 | LD | C, +20 | +20 characters per line. |
| 0D89 | CLS-2 | DEC | HL | Go back along the line setting |
|  |  | LD | (HL), A | the attribute bytes. |
|  |  | DEC | C |  |
|  |  | JR | NZ,0D89,CLS-2 |  |
| 0D8E | CLS-3 | DJNZ | 0D87,CLS-1 | Loop back until finished. |

The size of the lower part of the display can now be fixed.
LD (DF-SZ),+02 It will be two lines in size.
It now remains for the following 'house keeping' tasks to be performed.

| 0D94 | CL-CHAN | LD | A,+FD |
| :--- | :--- | :--- | :--- |
|  |  | CALL | 1601,CHAN-OPEN |$\quad$ Open channel 'K'.

THE 'CLEARING THE WHOLE DISPLAY AREA' SUBROUTINE
This subroutine is called from; i. the CLS command routine. ii. the main execution routine, and iii. the automatic listing routine.

| 0DAF | CL-ALL | LD | HL,+0000 <br> (COORDS),HL |
| :--- | :--- | :--- | :--- |
|  | LD | The system variable C00RDS <br> is reset to zero. |  |
|  | RES | O,(FLAGS2) | Signal 'the screen is clear'. |


| LD | $(\mathrm{HL}), \mathrm{E}$ | (PRINT-OUT). |
| :--- | :--- | :--- |
| INC | HL |  |
| LD | $(\mathrm{HL}), \mathrm{D}$ |  |
| LD | $(\mathrm{SCR}-\mathrm{CT}),+01$ | Reset the scroll counter. <br> As the upper part of the display <br> is being handled the 'upper print |
|  |  | line' will be Line ' 0 '. <br> CD |
|  |  | Continue into CL-SET. |

## THE 'CL-SET' SUBROUTINE

This subroutine is entered with the BC register pair holding the line and column numbers of a character areas, or the C register holding the column number within the printer buffer. The appropriate address of the first character bit is then found. The subroutine returns via PO-STORE so as to store all the values in the required system variables.

| 0DD9 | CL-SET | LD | HL, +5B00 | The start of the printer buffer. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | BIT | 1,(FLAGS) | Jump forward if handling the |
|  |  | JR | NZ,0DF4,CL-SET-2 | printer buffer. |
|  |  | LD | A, B | Transfer the line number. |
|  |  | BIT | 0,(TV-FLAG) | Jump forward if handling the |
|  |  | JR | Z,0DEE,CL-SET-1 | main part of the display. |
|  |  | ADD | A,(DF-SZ) | The top line of the lower |
|  |  | SUB | +18 | part of the display is called 'line +18 ' and this has to be converted. |
| ODEE | CL-SET-1 | PUSH | BC | The line \& column numbers are saved. |
|  |  | LD | B,A | The line number is moved. |
|  |  | CALL | 0E9B,CL-ADDR | The address for the start of the line is formed in HL. |
|  |  | POP | BC | The line \& column numbers are fetched back. |
| 0DF4 | CL-SET-2 | LD | A, +21 | The column number is now |
|  |  | SUB | C | reversed and transferred to |
|  |  | LD | E,A | the DE register pair. |
|  |  | LD | D, +00 |  |
|  |  | ADD | HL, DE | The required address is now |
|  |  | JP | OADC,PO-STORE | formed; and the address and |
|  |  |  |  | the line and column numbers |
|  |  |  |  | are stored by jumping to |
|  |  |  |  | PO-STORE. |

## THE 'SCROLLING' SUBROUTINE

The number of lines of the display that are to be scrolled has to be held on entry to the main subroutine in the $B$ register.
ODFE
CL-SC-ALL
LD
$B,+17$
The entry point after 'scroll?'

The main entry point - from above and when scrolling for INPUT..AT.

| 0E00 CL-SCROLL CALL | 0E9B,CL-ADDR | Find the starting address of <br> the line. |
| :---: | :--- | :--- |
| LD | C, +08 | There are eight pixel lines to <br> a complete line. |

Now enter the main scrolling loop. The B register holds the number of the top line to be scrolled, the HL register pair the starting address in the display area of this line and the C register the pixel line counter.

| 0E05 | CL-SCR-1 | $\begin{aligned} & \text { PUSH } \\ & \text { PUSH } \\ & \text { LD } \\ & \text { AND } \\ & \text { LD } \\ & \text { JR } \end{aligned}$ | $\begin{aligned} & \text { BC } \\ & \mathrm{HL} \\ & \mathrm{~A}, \mathrm{~B} \\ & +07 \\ & \text { A,B } \\ & \text { NZ,0E19,CL-SCR-3 } \end{aligned}$ | Save both counters. Save the starting address. Jump forward unless dealing at the present moment with a 'third' of the display. |
| :---: | :---: | :---: | :---: | :---: |

The pixel lines of the top lines of the 'thirds' of the display have to be moved across the 2 K boundaries. (Each 'third' = 2K.)

| OEOD | CL-SCR-2 | EX LD <br> ADD <br> EX <br> LD <br> DEC <br> LDIR | DE,HL <br> HL,+F8E0 <br> HL,DE <br> DE,HL <br> BC,+0020 <br> A | The result of this manipulation is to leave HL unchanged and DE pointing to the required destination. There are +20 characters. Decrease the counter as one line is being dealt with. Now move the thirty two bytes. |
| :---: | :---: | :---: | :---: | :---: |
| The pixel lines within the 'thirds' can now be scrolled. The A register holds, on the first pass, +01-+07, + |  |  |  |  |
| 0E19 | CL-SCR-3 | EX | DE,HL | Again DE is made to point |
|  |  | LD | HL,+FFE0 | to the required destination. |
|  |  | ADD | HL, DE | This time only thirty two |
|  |  | EX | DE,HL | locations away. |
|  |  | LD | B,A | Save the line number in B. |
|  |  | AND | +07 | Now find how many characters |
|  |  | RRCA |  | there are remaining in the |
|  |  | RRCA |  | 'third'. |
|  |  | RRCA |  |  |
|  |  | LD | C,A | Pass the 'character total' to the C register. |
|  |  | LD | A,B | Fetch the line number. |
|  |  | LD | B, +00 | BC holds the 'character total' |
|  |  | LDIR |  | and a pixel line from each of the characters is 'scrolled'. |
|  |  | LD | B,+07 | Now prepare to increment the address to jump across a 'third' boundary. |
|  |  | ADD | HL,BC | Increase HL by +0700 . |
|  |  | AND | +F8 | Jump back if there are any |
|  |  | JR | NZ,0EOD,CL-SCR-2 | 'thirds' left to consider. |

Now find if the loop has been used eight times - once for each pixel line.

| POP | HL | Fetch the original address. |
| :--- | :--- | :--- |
| INC | H | Address the next pixel line. |
| POP | BC | Fetch the counters. |
| DEC | C | Decrease the pixel line counter |
| JR | NZ,0E05,CL-SR-1 | and jump back unless eight lines |

Next the attribute bytes are scrolled. Note that the B register still holds the number of lines to be scrolled and the C register holds zero.

| CALL | 0E88,CL-ATTR | The required address in the <br> attribute area and the number <br> of characters in 'B' lines are <br> found. |
| :--- | :--- | :--- |
| LD | HL,+FFE0 | The displacement for all <br> ADD |
| EX | HL,DE attribute bytes is |  |
| LDIR | DE,HL | thirty two locations away. <br> The attribute bytes are |
| 'scrolled'. |  |  |

It remains now to clear the bottom line of the display.
LD $\quad B,+01$

THE 'CLEAR LINES' SUBROUTINE
This subroutine will clear the bottom ' $B$ ' lines of the display.

| OE44 CL-LINE | BUSH | The line number is saved for the <br> duration of the subroutine. |
| :--- | :--- | :--- |
| CALL | 0E9B,CL-ADDR | The starting address for the line <br> is formed in HL. |

LD C,+08

Now enter a loop to clear all the pixel lines.

| 0E4A | CL-LINE-1 | PUSH | BC | Save the line number and the pixel line counter. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | PUSH | HL | Save the address. |
|  |  | LD | A, B | Save the line number in A. |
| 0E4D | CL-LINE-2 | AND | +07 | Find how many characters are |
|  |  | RRCA |  | involved in 'B mod 8' lines. |
|  |  | RRCA |  | Pass the result to the |
|  |  | RRCA |  | C register. ( C will hold +00 |
|  |  | LD | C,A | i.e. 256 dec. for a 'third'.) |
|  |  | LD | A,B | Fetch the line number. |
|  |  | LD | B, +00 | Make the BC register pair |
|  |  | DEC | C | hold 'one less' than the number of characters. |
|  |  | LD | D, H | Make DE point to the first |
|  |  | LD | E,L | character. |
|  |  | LD | (HL) , +00 | Clear the pixel-byte of the first character. |
|  |  | INC | DE | Make DE point to the second |
|  |  | LDIR |  | character and then clear the pixel-bytes of all the other |
|  |  |  |  | characters. |
|  |  | LD | DE, +0701 | For each 'third' of the |
|  |  | ADD | HL, DE | display HL has to be increased by +0701 . |
|  |  | DEC | A | Now decrease the line number. |
|  |  | AND | +F8 | Discard any extra lines and |
|  |  | LD | B,A | pass the 'third' count to B. |
|  |  | JR | NZ,0E4D,CL-LINE-2 | Jump back if there are still |
|  |  |  |  | 'thirds' to be dealt with. |

Now find if the loop has been used eight times.

| POP | HL | Update the address for each |
| :--- | :--- | :--- |
| INC | H | pixel line. |
| POP | BC | Fetch the counters. |
| DEC | C | Decrease the pixel line |
| JR | NZ,0E4A,CL-LINE-1 | counter and jump back unless <br> finished. |

Next the attribute bytes are set as required. The value in ATTR-P will be used when handling the main part of the display and the value in BORDCR when handling the lower part.

| CALL | OE88,CL-ATTR | The address of the first <br> attribute byte and the number <br> of bytes are found. |
| :--- | :--- | :--- |
| LD | H,D | HL will point to the first <br> attribute byte and DE the |
| LD | L, | second. |

## THE 'CL-ATTR' SUBROUTINE

This subroutine has two separate functions.
i. For a given display area address the appropriate attribute address is returned in the DE register pair. Note that the value on entry points to the 'ninth' line of a character.
ii. For a given line number, in the $B$ register, the number of character areas in the display from the start of that line onwards is returned in the $B C$ register pair.

| 0E88 | CL-ATTR | A,H | Fetch the high byte. <br> Multiply this value by <br> RRCA |
| :--- | :--- | :--- | :--- |
|  | RRCA |  | thirty two. |
|  | RRCA |  | Go back to the 'eight' line. |
|  | DEC | A | Address the attribute area. |
|  | OR | +50 | Restore to the high byte and |
|  | LD | $\mathrm{H}, \mathrm{A}$ | transfer the address to DE. |

## THE 'CL-ADDR' SUBROUTINE

For a given line number, in the B register, the appropriate display file address is formed in the HL register pair.

| 0E9B | CL-ADDR | LD <br> SUB <br> LD <br> RRCA <br> RRCA <br> RRCA <br> AND <br> LD <br> LD <br> AND <br> OR | $\begin{aligned} & \mathrm{A},+18 \\ & \mathrm{~B} \\ & \mathrm{D}, \mathrm{~A} \\ & \\ & + \text { +E0 } \\ & \text { L,A } \\ & \mathrm{A}, \mathrm{D} \\ & +18 \\ & +40 \end{aligned}$ | The line number has to be reversed. <br> The result is saved in $D$. In effect '(A mod 8) * 32 '. In a 'third' of the display the low byte for the: <br> 1 st . line $=+00$, 2nd. line $=+20$, etc. <br> The low byte goes into $L$. The true line number is fetched. In effect '64 +8 * INT (A/8)' For the upper 'third' of the display the high byte $=+40$, middle 'third' $=+48$, and the lower 'third' = +50. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | $\begin{aligned} & \text { LD } \\ & \text { RET } \end{aligned}$ | H,A | The high byte goes to H . Finished. |

THE 'COPY' COMMAND ROUTINE
The one hundred and seventy six pixel lines of the display are dealt with one by one.

| OEAC COPY | DI | The maskable interrupt is <br> disabled during COPY. |  |
| :--- | :--- | :--- | :--- |
|  | LD | $\mathrm{B},+\mathrm{BO}$ | The '176' lines. |

The following loop is now entered.

| 0EB2 | COPY-1 | PUSH |
| :--- | :--- | :--- |
|  |  | HL |
|  |  | CUSH |
|  | CALL | BC |
|  | POP | BC |
|  |  | POP |
|  |  | HL |
|  |  | HNC |
|  |  | H |

[^0]| LD | A,H |
| :--- | :--- |
| AND | +07 |
| JR | NZ,0EC9,COPY-2 |

JR NZ,0EC9,COPY-2

Jump forward and hence round the loop again directly for the eight pixel lines of a character line.

For each new line of characters the base address has to be updated.


THE 'COPY-BUFF' SUBROUTINE
This subroutine is called whenever the printer buffer is to have its contents passed to the printer.

| OECD | COPY-BUFF | DI | HL,+5800 | Disable the maskable interrupt. <br> The base address of the printer |
| :--- | :--- | :--- | :--- | :--- |
| OU |  |  |  |  |

Continue into the COPY-END routine.

| OEDA | COPY-END | LD | A,+04 <br> OUT |
| :--- | :--- | :--- | :--- |
|  | EI |  | Stop the printer motor. |
|  |  |  | Enable the maskable interrupt |
|  |  |  | and continue into CLEAR-PRB. |

THE 'CLEAR PRINTER BUFFER' SUBROUTINE
The printer buffer is cleared by calling this subroutine.

| OEDF | CLEAR-PRB | LD | HL, +5B00 | The base address of the printer buffer. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | LD | (PR-CC-lo), L | Reset the printer 'column'. |
|  |  | XOR | A | Clear the A register. |
|  |  | LD | B,A | Also clear the $B$ register in effect B holds dec.256). |
| 0EE7 | PRB-BYTES | LD | (HL), A | The '256' bytes of the |
|  |  | INC | HL | printer buffer are all |
|  |  | DJNZ | 0EE7,PRB-BYTES | cleared in turn. |
|  |  | RES | 1,(FLAGS2) | Signal 'the buffer is empty'. |
|  |  | LD | C,+21 | Set the printer position and |
|  |  | JP | 0DD9,CL-SET | return via CL-SET \& P0-STORE |

THE 'COPY-LINE' SUBROUTINE
The subroutine is entered with the HL register pair holding the base address of the thirty two bytes that form the pixel-line and the B register holding the pixel-line number.

| 0EF4 | COPY-LINE | LD | A, B | Copy the pixel-line number |
| :---: | :---: | :---: | :---: | :---: |
|  |  | CP | +03 | The A register will hold |
|  |  | SBC | A, A | +00 until the last two lines |
|  |  | AND | +02 | are being handled. |


| OUT | $(+$ FB $), \mathrm{A}$ |
| :--- | :--- |
| LD | $\mathrm{D}, \mathrm{A}$ |

D,A

Slow the motor for the last two pixel lines only
The D register will hold either +00 or +02.

There are three tests to be made before doing any 'printing'.

| OEFD | COPY-L-1 | CALL | 1F54,BREAK-KEY | Jump forward unless the |
| :---: | :---: | :---: | :---: | :---: |
|  |  | JR | C,0F0C,COPY-L-2 | BREAK key is being pressed. |
|  |  | LD | A,+04 | But if it is then; |
|  |  | OUT | (+FB), A | stop the motor, |
|  |  | El |  | enable the maskable interrupt, |
|  |  | CALL | OEDF,CLEAR-PRB | clear the printer buffer and exit |
|  |  | RST | 0008,ERROR-1 | via the error handling routine |
|  |  | DEFB | +0C | - 'BREAK-CONT repeats'. |
| OFOC | COPY-L-2 | IN | A,(+FB) | Fetch the status of the |
|  |  | ADD | A, A | printer. |
|  |  | RET | M | Make an immediate return if the printer is not present. |
|  |  | JR | NC,0EFD, COPY-L-1 | Wait for the stylus. |
|  |  | LD | C, +20 | There are thirty two bytes. |

Now enter a loop to handle these bytes.

| 0F14 | COPY-L-3 | LD | E,(HL) | Fetch a byte. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | INC | HL | Update the pointer. |
|  |  | LD | B,+08 | Eight bits per byte. |
| 0F18 | COPY-L-4 | RL | D | Move D left. |
|  |  | RL | E | Move each bit into the carry. |
|  |  | RR | D | Move D back again, picking up the carry from $E$. |
| 0F1E | COPY-L-5 | IN RRA | A,(+FB) | Again fetch the status of the printer and wait for the |
|  |  | JR | NC,0F1E,COPY-L-5 | signal from the encoder. |
|  |  | LD | $\begin{aligned} & \mathrm{A}, \mathrm{D} \\ & (+\mathrm{FB}), \mathrm{A} \end{aligned}$ | Now go ahead and pass the 'bit' to the printer |
|  |  |  |  | Note: bit 2 - low starts the motor, bit 1 - high slows the motor and bit 7 is high for the actual 'printing'. |
|  |  | DJNZ | 0F18,COPY-L-4 | 'Print' each bit. |
|  |  | DEC |  | Decrease the byte counter. |
|  |  | JR | NZ,0F14,COPY-L-3 | Jump back whilst there are |
|  |  | RET |  | still bytes; otherwise return. |

## THE 'EDITOR' ROUTINES

The editor is called on two occasions:
i. From the main execution routine so that the user can enter a BASIC line into the system.
ii. From the INPUT command routine.

First the 'error stack pointer' is saved and an alternative address provided.

| 0F2C | EDITOR | LD | HL,(ERR-SP) | The current value is saved on |
| :--- | :--- | :--- | :--- | :--- |
| 0F30 | ED-AGAIN | PUSH | HL | the machine stack. |
|  |  | PUSH | HL,+107F | This is ED-ERROR. |
|  |  | LD | (ERR-SP),SP | Any event that leads to the <br> error handling routine being |
|  |  |  | used will come back to |  |
|  |  |  | ED-ERROR. |  |

A loop is now entered to handle each keystroke.

| 0F38 ED-LOOP | CALL | 15D4,WAIT-KEY |
| :--- | :--- | :--- | | Return once a key has been |
| :--- |
|  |
|  |
| PUSH | AF $\quad$| pressed. |
| :--- |
| Save the code temporarily. |


| LD | D,+00 |
| :--- | :--- |
| LD | E,(PIP) |
| LD | HL,+00C8 |
| CALL | $03 B 5$, BEEPER |
| POP | AF |
| LD | HL,+0F38 |
| PUSH | HL |

Fetch the duration of the keyboard click.
And the pitch.
Now make the 'pip'.
Restore the code.
Pre-load the machine stack with the address of ED-LOOP.

Now analyse the code obtained.

| CP | +18 |
| :--- | :--- |
| JR | NC,0F81,ADD-CHAR |
| CP | +07 |
| JR | C,0F81,ADD-CHAR |
| CP | +10 |
| JR | C,0F92,ED-KEYS |

Accept all character codes, graphic codes and tokens.
Also accept ','.
Jump forward if the code represents an editing key.

The control keys - INK to TAB -are now considered.

| LD | BC,+0002 |
| :--- | :--- |
| LD | D,A |
| CP | +16 |
| JR | C,0F6C,ED-CONTR |

INK \& PAPER will require two locations.
Copy the code to 0 .
Jump forward with INK \& PAPER'

AT \& TAB would be handled as follows:

| INC | BC |
| :--- | :--- |
| BIT | 7,(FLAGX) |
| JP | Z,101E,ED-IGNORE |
| CALL | 15D4,WAIT-KEY |
| LD | E,A |

Three locations required. Jump forward unless dealing with INPUT LINE... .
Get the second code. and put it in E .

The other bytes for the control characters are now fetched.

| 0F6C | ED-CONTR | CALL |
| :--- | :--- | :--- |
|  | PUSH | 15D4,WAIT-KEY |
|  | LD | HL,(K-CUR) |
|  | RES | $0,(M O D E)$ |
|  | CALL | $1655, M A K E-R O O M$ |
|  | POP | BC |
|  | INC | HL |
|  | LD | $(\mathrm{HL}), \mathrm{B}$ |
|  | INC | HL |
|  | LD | $(\mathrm{HL}), \mathrm{C}$ |
|  |  |  |
|  |  |  |
|  |  |  |
|  |  |  |

Get another code.
Save the previous codes.
Fetch K-CUR.
Signal 'K mode'.
Make two or three spaces.
Restore the previous codes.
Point to the first location.
Enter first code.
Then enter the second code which will be Overwritten if there are only two codes - i.e. with INK \& PAPER. Jump forward.

THE 'ADDCHAR' SUBROUTINE
This subroutine actually adds a code to the current EDIT or INPUT line.

| 0F81 | ADD-CHAR | RES | 0,(MODE) | Signal 'K mode'. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | LD | HL,(K-CUR) | Fetch the cursor position. |
|  |  | CALL | 1652,ONE-SPACE | Make a single space. |
| 0F8B | ADD-CH-1 | LD | (DE),A | Enter the code into the space |
|  |  | INC | DE | and signal that the cursor is to |
|  |  | LD | (K-CUR),DE | occur at the location after. Then |
|  |  | RET |  | return indirectly to ED-LOOP. |

The editing keys are dealt with as follows:

| 0F92 | ED-KEYS | LD | E,A |
| :--- | :--- | :--- | :--- |
|  | LD | C,+00 | The code is transferred to <br> the DE register pair. |
|  | LD | HL,+0F99 | The base address of the editing <br> key table. |
|  |  | TDD | HL,DE |


| PUSH | HL | routine is saved on the machine <br> stack. |
| :--- | :--- | :--- |
| LD | HL,(K-CUR) | The HL register pair is set and <br> an indirect jump made to the <br> required routine. |
| RET |  |  |

## THE 'EDITING KEYS' TABLE

| address | offset | character | address | offset | character |
| :--- | :--- | :--- | :--- | :--- | :--- |
| OFA0 | 09 | EDIT | 0FA5 | 70 | DELETE |
| OFA1 | 66 | cursor left | OFA6 | 7E | ENTER |
| 0FA2 | $6 A$ | cursor right | OFA7 | CF | SYMBOL SHIFT |
| 0FA3 | 50 | cursor down | OFA8 | D4 | GRAPHICS |
| 0FA4 | 85 | cursor up |  |  |  |

## THE 'EDIT KEY' SUBROUTINE

When in 'editing mode' pressing the EDIT key will bring down the 'current BASIC line'. However in 'INPUT mode' the action of the EDIT key is to clear the current reply and allow a fresh one.

| OFA9 | ED-EDIT | LD | HL,(E-PPC) | Fetch the current line number. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | BIT | 5,(FLAGX) | But jump forward if in |
|  |  | JP | NZ,1097,CLEAR-SP | 'INPUT mode'. |
|  |  | CALL | 196E,LINE-ADDR | Find the address of the start of |
|  |  | CALL | 1695,LINE-NO | the current line and hence its number. |
|  |  | LD | A, D | If the line number returned is |
|  |  | OR | E | zero then simply clear the |
|  |  | JP | Z,1097,CLEAR-SP | editing area. |
|  |  | PUSH | HL | Save the address of the line. |
|  |  | INC | HL | Move on to collect the |
|  |  | LD | C,(HL) | length of the line. |
|  |  | INC | HL |  |
|  |  | LD | B,(HL) |  |
|  |  | LD | HL, +000A | Add +0 A to the length and test |
|  |  | ADD | HL,BC | that there is sufficient room |
|  |  | LD | B, H | for a copy of the line. |
|  |  | LD | C,L |  |
|  |  | CALL | 1F05,TEST-ROOM |  |
|  |  | CALL | 1097,CLEAR-SP | Now clear the editing area. |
|  |  | LD | HL,(CURCHL) | Fetch the current channel |
|  |  | EX | (SP), HL | address and exchange it for the address of the line. |
|  |  | PUSH | HL | Save it temporarily. |
|  |  | LD | A,+FF | Open channel 'A' so that the |
|  |  | CALL | 1601,CHAN-OPEN | line will be copied to the editing area. |
|  |  | POP | HL | Fetch the address of the line. |
|  |  | DEC | HL | Goto before the line. |
|  |  | DEC | (E-PPC-1o) | Decrement the current line |
|  |  |  |  | number so as to avoid printing the cursor. |
|  |  | CALL | 1855,OUT-LINE | Print the BASIC line |
|  |  | INC | (E-PPC-Io) | Increment the current line number. |
|  |  |  |  | Note: The decrementing of the |
|  |  |  |  | line number does not always |
|  |  |  |  | stop the cursor from being printed |
|  |  | LD | HL,(E-LINE) | Fetch the start of the line in |
|  |  | INC | HL | the editing area and step past |
|  |  | INC | HL | the line number and the |
|  |  | INC | HL | length to find the address |
|  |  | INC | HL | for K-CUR. |


| LD | (K-CUR),HL |
| :--- | :--- |
| POP | HL |
| CALL | 1615, CHAN-FLAG |

CALL

1615,CHAN-FLAG
1615,CHAN

Fetch the former channel address and set the appropriate flags before returning to ED-LOOP.

| THE 'CURSOR DOWN EDITING' SUBROUTINE |  |  |  |
| :---: | :---: | :---: | :---: |
| 0FF3 | ED-DOWN | BIT | 5,(FLAGX) |
|  |  | JR | NZ, 1001,ED-STOP |
|  |  | LD | HL, +5C49 |
|  |  | CALL | 190F,LN-FETCH |
|  |  | JR | 106E,ED-LIST |
| 1001 | ED-STOP | LD | (ERR-NR), +10 |
|  |  | JR | 1024,ED-ENTER |

Jump forward if in
'INPUT' mode'.
This is E-PPC.
The next line number is found and a new automatic listing produced.
'STOP in INPUT' report.
Jump forward.
THE 'CURSOR LEFT EDITING' SUBROUTINE
1007 ED-LEFT CALL 1031,ED-EDGE

JR 1011,ED-CUR
The cursor is moved. Jump forward.

THE 'CURSOR RIGHT EDITING' SUBROUTINE

| 100 C | ED-RIGHT | LD | A,(HL) |
| :--- | :--- | :--- | :--- |
|  |  | CP | +0 D |
|  |  | RET | Z |
|  |  | INC | HL |
| 1011 | ED-CUR | LD | (K-CUR),HL |

THE 'DELETE EDITING' SUBROUTINE
1015 ED-DELETE CALL

| CALL | 1031,ED-EDGE |
| :--- | :--- |
| $D$ | BC,+0001 |
| $P$ | $19 E 8, R E C L A I M-2$ |

Move the cursor leftwards. Reclaim the current character.

THE 'ED-IGNORE' SUBROUTINE

| 101E | ED-IGNORE | CALL | 15D4,WAIT-KEY |
| :--- | :--- | :--- | :--- |
|  |  | CALL | 15D4,WAIT-KEY |
| THE 'ENTER EDITING' SUBROUTINE |  |  |  |
| 1024 | ED-ENTER | POP | HL |
|  |  | POP | HL |
| 1026 | ED-END | POP | HL |
|  |  | LD | (ERR-SP),HL |
|  |  | BIT | 7,(ERR-NR) |
|  |  | RET | NZ |
|  |  |  | LD |
|  |  | RET | SP,HL |

The next two codes from the key-input routine are ignored.

The address of ED-LOOP and ED-ERROR are discarded. The old value of ERR-SP is restored.
Now return if there were no errors.
Otherwise make an indirect jump to the error routine.

THE 'ED-EDGE' SUBROUTINE
The address of the cursor is in the HL register pair and will be decremented unless the cursor is already at the start of the line. Care is taken not to put the cursor between control characters and their parameters.

1031 ED-EDGE \begin{tabular}{lll}
SCF <br>
CALL

$\quad$ 1195,SET-DE 

DE will hold either E-LINE
\end{tabular}

| POP | BC | Drop the return address. <br> Return via ED-LOOP if the |
| :--- | :--- | :--- |
| RET | C | carry flag is set. |
| PUSH | BC | Restore the return address. |
| Move the current address of |  |  |
| LD | B,H | the cursor to BC. |

Now enter a loop to check that control characters are not split from their parameters.

| 103E | ED-EDGE-1 | LD | H,D | HL will point to the |
| :---: | :---: | :---: | :---: | :---: |
|  |  | LD | L, E | character in the line after |
|  |  | INC | HL | that addressed by DE. |
|  |  | LD | A,(DE) | Fetch a character code. |
|  |  | AND | +F0 | Jump forward if the code |
|  |  | CP | +10 | does not represent |
|  |  | JR | NZ,1051,ED-EDGE-2 | INK to TAB. |
|  |  | INC | HL | Allow for one parameter. |
|  |  | LD | A,(DE) | Fetch the code anew. |
|  |  | SUB | +17 | Carry is reset for TAB. |
|  |  | ADC | A, +00 | Note: This splits off AT \& TAB but AT \& TAB in this form are not implemented anyway so it makes no difference. |
|  |  | JR | NZ,1051,ED-EDGE-2 | Jump forward unless dealing |
|  |  | INC | HL | with AT \& TAB which would have two parameters, if used. |
| 1051 | ED-EDGE-2 | AND | A | Prepare for true subtraction. |
|  |  | SBC | HL, BC | The carry flag will be reset |
|  |  | ADD | HL, BC | when the 'updated pointer' reaches K-CUR. |
|  |  | EX | DE,HL | For the next loop use the |
|  |  | JR | C,103E,ED-EDGE-1 | 'updated pointer', but if |
|  |  | RET |  | exiting use the 'present pointer' for K-CUR. |
|  |  |  |  | Note: It is the control character that is deleted when using DELETE. |
| THE 'CURSOR UP EDITING' SUBROUTINE |  |  |  |  |
| 1059 | ED-UP | BIT | 5,(FLAGX) | Return if in 'INPUT mode'. |
|  |  | RET | NZ |  |
|  |  | LD | HL,(E-PPC) | Fetch the current line |
|  |  | CALL | 196E,LINE-ADDR | number and its start address. |
|  |  | EX | DE,HL | HL now points to the previous line. |
|  |  | CALL | 1695,LINE-NO | This line's number is fetched. |
|  |  | LD | HL,+5C4A | This is E-PPC-hi. |
|  |  | CALL | 1910,LN-STORE | The line number is stored. |
| 106E | ED-LIST | CALL | 1795,AUTO-LIST | A new automatic listing is |
|  |  | LD | A, +00 | now produced and channel 'K' |
|  |  | JP | 1601,CHAN-OPEN | re-opened before returning to ED-LOOP. |

## THE 'ED-SYMBOL' SUBROUTINE

If SYMBOL \& GRAPHICS codes were used they would be handled as follows:

| 1076 | ED-SYMBOL BIT | 7,(FLAGX) | Jump back unless dealing with |  |
| :--- | :--- | :--- | :--- | :--- |
|  |  | JR | Z,1024,ED-ENTER | INPUT. LINE. |
| 107C | ED-GRAPH | JP | OF81,ADD-CHAR | Jump back. |

THE 'ED-ERROR' SUBROUTINE
Come here when there has been some kind of error.

| 107F ED-ERROR | BIT | 4,(FLAGS2) | Jump back if using other than |
| :--- | :--- | :--- | :--- |
|  | JR | Z,1026,ED-END | channel 'K'. |


| LD | (ERR-NR),+FF |
| :--- | :--- |
| LD | D,+00 |
| LD | E,(RASP) |
| LD | HL,+1A90 |
| CALL | 0385, BEEPER |
| JP | $0 F 30$, ED-AGAIN |

Cancel the error number and give a 'rasp' before going around the editor again.

## THE 'CLEAR-SP' SUBROUTINE

The editing area or the work space is cleared as directed. 1097

| CLEAR-SP | PUSH | HL |
| :--- | :--- | :--- |
|  | CALL | 1190, SET-HL |
|  | DEC |  |
|  | CALL | HL |
|  | LD | (K-CUR),RECLAIM-1 |
|  | LD | (MODE),+00 |
|  | POP | HL |
|  | RET |  |

Save the pointer to the space.
DE will point to the first character and HL the last. The correct amount is now reclaimed.
The system variables K-CUR and MODE ('K mode') are initialised before fetching the pointer and returning.

THE 'KEYBOARD INPUT' SUBROUTINE
This important subroutine returns the code of the last key to have bean pressed but note that CAPS LOCK, the changing of the mode and the colour control parameters are handled within the subroutine.

10A8 KEY-INPUT | BIT | $3,($ TV-FLAG) |  |
| :--- | :--- | :--- |
|  | CALL | NZ,111D,ED-COPY |
|  |  |  |
|  | AND | A |
|  | BIT | $5,($ FLAGS $)$ |
|  | RET | Z |
|  | LD | A,(LAST-K) |
|  | RES | $5,($ FLAGS) |
|  | PUSH | AF |
|  | BIT | $5,(T V-F L A G)$ |
|  | CALL | NZ,0D6E,CLS-LOWER |
|  |  |  |
|  | POP | AF |
|  | CP | +20 |
|  | JR | NC,111B,KEY-DONE |
|  | CP | +10 |
|  | JR | NC,10FA,KEY-CONTR |
|  | CP | +06 |
|  | JR | NC,10DB,KEY=M\&CL |

Copy the edit-linear the INPUT-line to the screen if the mode has changed.
Return with both carry
and zero flags reset if no
new key has been pressed
Otherwise fetch the code and
signal that it has been taken
Save the code temporarily.
Clear the lower part of the
display if necessary;
e.g. after 'scroll?';

Fetch the code.
Accept all characters and token codes.
Jump forward with most of the control character codes Jump forward with the 'mode'. codes and the CAPS LOCK code.

Now deal with the FLASH, BRIGHT\& INVERSE codes.

| LD | B,A |
| :--- | :--- |
| AND | +01 |
| LD | C,A |
|  |  |
| LD | A,B |
| RRA | A,+12 |
| ADD | 1105, KEY-DATA |

Save the code.
Keep only bit 0 .
C holds +00 (= OFF) or
C holds +01 (= ON).
Fetch the code.
Rotate it once (losing bit 0).
Increase it by +12 giving for
FLASH - +12, BRIGHT - +13 and INVERSE - +14.

The CAPS LOCK code and the mode codes are dealt with 'locally'.

| 10DB | KEY-M\&CL | JR | NZ,10E6,KEY-MODE | Jump forward with 'mode' codes. |
| :--- | :--- | :--- | :--- | :--- |
|  |  | LD | HL,+5C6A | This is FLAGS2. |
|  |  | LD | A,+08 | Flip bit 3 of FLAGS2. This is |
|  |  | XOR | (HL) | the CAPS LOCK flag. |
|  |  | LD | (HL),A |  |
|  |  | JR | 10F4,KEY-FLAG | Jump forward. |
| 10E6 | KEY-MODE | CP | $+0 E$ | Check the lower limit. |


|  |  | RET | C |  |
| :---: | :---: | :---: | :---: | :---: |
|  |  | SUB | +0D | Reduce the range. |
|  |  | LD | HL, +5C41 | This is MODE. |
|  |  | CP | (HL) | Has it been changed? |
|  |  | LD | (HL), A | Enter the new 'mode' code. |
|  |  | JR | NZ,10F4,KEY-FLAG | Jump if it has changed; |
|  |  | LD | (HL), +00 | otherwise make it 'L mode'. |
| 10F4 | KEY-FLAG | SET |  | Signal 'the mode might have changed. |
|  |  | CP | A | Reset the carry flag and |
|  |  | RET |  | return. |

The control key codes (apart from FLASH, BRIGHT \& INVERSE) are manipulated.

| 10FA | KEY-CONTR | BD | Save the code. |
| :--- | :--- | :--- | :--- |
|  | AND | +07 | Make the C register hold the |
| LD | C,A | parameter. (+00 to +07) |  |
|  | LD | A,+10 | A now holds the INK code. |
|  | BIT | $3, B$ | But if the code was an |
| JR | NZ,1105,KEY-DATA | 'unshifted' code then make A |  |
|  | INC | A | hold the PAPER code. |

The parameter is saved in K-DATA and the channel address changed from KEY-INPUT to KEY-NEXT.

1105 KEY-DATA | LD | (K-DATA),C | Save the parameter. |  |
| :--- | :--- | :--- | :--- |
|  | LD | DE,+110D | This is KEY-NEXT. |
|  | JR | 1113, KEY-CHAN | Jump forward. |

Note: On the first pass entering at KEY-INPUT the A register is returned holding a control code' and then on the next pass, entering at KEY-NEXT, it is the parameter that is returned.

| 110D KEY-NEXT | LD | A,(K-DATA) | Fetch the parameter. |
| :--- | :--- | :--- | :--- |
|  | LD | DE,+10A8 | This is KEY-INPUT. |

Now set the input address in the first channel area.

| 1113 KEY-CHAN | LD | $\mathrm{HL},(\mathrm{CHANS})$ | Fetch the channel address. |
| :--- | :--- | :--- | :--- |
|  | INC | HL |  |
|  | INC | HL |  |
|  | LD | $(\mathrm{HL}), E$ | Now set the input address. |
|  | INC | HL |  |
|  | LD | $(\mathrm{HL}), \mathrm{D}$ |  |

Finally exit with the required code in the A register.
111B KEY-DONE SCF Show a code has been found
RET and return.

THE 'LOWER SCREEN COPYING' SUBROUTINE
This subroutine is called whenever the line in the editing area or the INPUT area is to be printed in the lower part of the screen.


| LD | HL,(ECHO-E) | Push the value of ECHO-E |
| :---: | :---: | :---: |
| PUSH | HL | on to the stack. |
| SCF |  | Make HL point to the start |
| CALL | 1195,SET-HL | of the space and DE the end. |
| EX | DE,HL |  |
| CALL | 187D,OUT-LINE2 | Now print the line. |
| EX | DE,HL | Exchange the pointers and |
| CALL | 18E1,OUT-CURS | print the cursor. |
| LD | HL,(S-POSNL) | Next fetch the Current value |
| EX | (SP),HL | of S-POSNL and exchange it with ECHO-E. |
| EX | DE,HL | Pass ECHO-E to DE. |
| CALL | 0D4D,TEMPS | Again fetch the permanent |

The remainder of any line that has been started is now completed with spaces printed with the 'permanent' PAPER colour.

| 1150 | ED-BLANK | LD | A,(S-POSNL-hi) |
| :--- | :--- | :--- | :--- | | Fetch the current line number |
| :--- |
| and subtract the old line number. |

New deal with any errors.
1167 ED-FULL $\quad$ LD $\quad$ D,+00 $\quad$ Give out a 'rasp'.

The normal exit upon completion of the copying over of the editor the INPUT line.

| $117 C$ | ED-C-DONE | POP | DE |
| :--- | :--- | :--- | :--- |
|  | POP | HL | The new position value. |
|  | The 'error address'. |  |  |

But come here after an error.
117E ED-C-END POP HL The old value of ERR-SP

| LD | (ERR-SP),HL | restored. <br> POP |
| :--- | :--- | :--- |
|  | BC | Fetch the old value of |
| PUSH | DE | S-POSNL. |
| CALL | ODD9,CL-SET | Save the new position values. |
| POP | HL | Set the system variables. |
| LD | (ECHO-E),HL | The old value of S-POSNL |
| LD | (X-PTR-hi),+00 | goes into ECHO-E. |
| RET |  | XiPTR is cleared in a <br> suitable manner and the return <br> Made. |

## THE 'SET-HL' AND 'SET-DE' SUBROUTINES

These subroutines return with HL pointing to the first location and DE the 'last' location of either the editing area or the work space.
1190
SET-HL
LD
HL,(WORKSP)
Point to the last location

| 1195 | SET-DE | DEC | HL | of the editing area. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | AND | A | Clear the carry flag. |
|  |  | LD | DE,(E-LINE) | Point to the start of the |
|  |  | BIT | 5,(FLAGX) | editing area and return if |
|  |  | RET |  | in 'editing mode'. |
|  |  | LD | DE,(WORKSP) | Otherwise change DE. |
|  |  | RET | C | Return if now intended. |
|  |  | $\stackrel{\text { LD }}{\text { RET }}$ | HL,(STKBOT) | Fetch STKBOT and then |

THE 'REMOVE-FP' SUBROUTINE
This subroutine removes the hidden floating-point forms in a BASIC line.


## THE EXECUTIVE ROUTINES

## THE 'INITIALISATION' ROUTINE

The main entry point to this routine is at START/NEW (11CB). When entered from START (0000), as when power is first applied to the system, the A register holds zero and the DE register the value +FFFF. However the main entry point can also be reached following the execution of the NEW command routine.

| THE 'NEW COMMAND' |  |  |
| :--- | :--- | :--- |
| 11B7 | NOUTINE |  |
|  | DI |  |
|  | LD | A,+FF |
|  | LD | DE,(RAMTOP) |
|  |  |  |
|  | LXX |  |
|  | LD | BC,(P-RAMT) |
|  | LD | DE,(RASP/PIP) |
|  | EXX |  |
|  |  |  |

Disable the maskable interrupt.
The NEW flag.
The existing value of RAMTOP is preserved.
Load the alternate registers
with the following system
variables. All of which will also be preserved.

The main entry point.

| 11 CB | START/NEW | LD | B,A | Save the flag for later. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | LD | A, +07 | Make the border white in |
|  |  | OUT | (+FE), A | colour. |
|  |  | LD | A, +3F | Set the I register to hold |
|  |  | LD | I,A | the value of +3 F . |
|  |  | DEFB | +00,+00,+00 | Wait 24 T states. |
|  |  | DEFB | +00,+00,+00 |  |

Now the memory is checked.

| 11DA | RAM-CHECK | $\begin{aligned} & \text { LD } \\ & \text { LD } \end{aligned}$ | $\begin{aligned} & \mathrm{H}, \mathrm{D} \\ & \mathrm{~L}, \mathrm{E} \end{aligned}$ | Transfer the value in DE (START = +FFFF, NEW = RAMTOP). |
| :---: | :---: | :---: | :---: | :---: |
| 11DC | RAM-FILL | LD | (HL), +02 | Enter the value of +02 into |
|  |  | DEC | HL | every location above +3FFF. |
|  |  | CP | H |  |
|  |  | JR | NZ,11DC,RAM-FILL |  |
| 11E2 | RAM-READ | AND | A | Prepare for true subtraction. |
|  |  | SBC | HL, DE | The carry flag will become |
|  |  | ADD | HL, DE | reset when the top is reached. |
|  |  | INC | HL | Update the pointer. |
|  |  | JR | NC,11EF,RAM-DONE | Jump when at top. |
|  |  | DEC | (HL) | +02 goes to +01. |
|  |  | JR | Z,11EF,RAM-DONE | But if zero then RAM is faulty. Use current HL as top. |
|  |  | DEC | (HL) | +01 goes to +00. |
|  |  | JR | Z,11E2,RAM-READ | Step to the next test unless it fails. |
| 11EF | RAM-DONE | DEC | HL | HL points to the last actual location in working order. |

Next restore the 'preserved' system variables. (Meaningless when coming from START.)

| EXX |  | Switch registers. |
| :--- | :--- | :--- |
| LD | (P-RAMT),BC | Restore P-RAMT,RASP/PIP |
| LD | (RASP/PIP),DE | \&UDG |
| LD | (UDG),HL |  |
| EXX |  |  |
| INC | B | Test the START/NEW flag. |
| JR | Z,1219,RAM-SET | Jump forward if coming from |
|  |  | the NEW command routine. |

Overwrite the system variables when coming from START and initialise the user-defined graphics area.

| LD | (P-RAMT),HL <br> LE, +3EAF | Top of physical RAM. <br> Last byte of 'U' in character <br> set. |
| :--- | :--- | :--- |
| LD | BC,+00A8 | There are this number of bytes <br> in twenty one letters. |
| EX | DE,HL | Switch the pointers. <br> Now copy the character forms <br> LDDR |
| EX | DE,HL the letter 'A' to 'U'. |  |
| INC | HL | Switch the pointers back. |
| LD | (UDG),HL | Point to the first byte. |
| DEC | HL | Now set UDG. |
| LD | BC,+0040 | Down one location. |
| LD | (RASP/PIP),BC | Set the system variables |
| RASP \& PIP. |  |  |

The remainder of the routine is common to both the START and the NEW operations.

1219 RAM-SET | LD | (RAMTOP),HL | Set RAMTOP. |  |
| :--- | :--- | :--- | :--- |
|  | LD | HL,+3C00 | Initialise the system variable |
|  | LD | (CHARS),HL | CHARS. |

Next the machine stack is set up.

| LD | $\mathrm{HL},(\mathrm{RAMTOP})$ |
| :--- | :--- |
| LD | $(\mathrm{HL}),+3 \mathrm{E}$ |
| DEC | HL |
|  |  |
| LD | $\mathrm{SP}, \mathrm{HL}$ |
| DEC | HL |
| DEC | HL |
| LD | (ERR-SP),HL |

The top location is made to hold $+3 E$.
The next location is left holding zero.
These two locations represent the 'last entry'.
Step down two locations to find the correct value for ERR-SP.

The initialisation routine continues with:

| IM | 1 | Interrupt mode 1 is used. |
| :---: | :---: | :---: |
| LD | IY, +5C3A | IY holds +ERR-NR always. |
| El |  | The maskable interrupt can now |
|  |  | be enabled. The real-time clock |
|  |  | will be updated and the keyboard |
|  |  | scanned every 1/50th of a |
|  |  | second. |
| LD | HL,+5CB6 | The base address of the |
| LD | (CHANS), HL | channel information area. |
| LD | DE,15AF | The initial channel data |
| LD | BC, +0015 | is moved from the table |
| EX | DE,HL | (15AF) to the channel |
| LDIR |  | information area. |
| EX | DE,HL | The system variable DATADD |
| DEC | HL | is made to point to the last |
| LD | (DATADD), HL | location of the channel data. |
| INC | HL | And PROG \& VARS to the |
| LD | (PROG), HL | the location after that. |
| LD | (VARS), HL |  |
| LD | (HL), +80 | The end-marker of the variables area |
| INC | HL | Move on one location to find |
| LD | (E-LINE), HL | the value for E-LINE. |
| LF | (HL), +0D | Make the edit-line be a single |
| INC | HL | 'carriage return' character. |
| LD | (HL), +80 | Now enter an end-marker. |
| INC | HL | Move on one location to find |
| LD | (WORKSP), HL | the value for WORKSP, STKBOT |
| LD | (STKBOT), HL | \& STKEND. |


| LD | (STKEND),HL |  |
| :---: | :---: | :---: |
| LD | A, +38 | Initialise the colour system |
| LD | (ATTR-P),A | variables to : FLASH 0, |
| LD | (ATTR-T),A | BRIGHT 0, PAPER 7, \& INK 0. |
| LD | (BORDCR),A |  |
| LD | HL,+0523 | Initialise the system |
| LD | (REPDEL), HL | variables REPDEL \& REPPER. |
| DEC | (KSTATE-0) | Make KSTATE-0 hold +FF |
| DEC | (KSTATE-4) | Make KSTATE-4 hold +FF |
| LD | HL,+15C6 | Next move the initial stream |
| LD | DE, +5 C 10 | data from its table to the |
| LD | BC,+000E | streams area. |
| LDIR |  |  |
| SET | 1,(FLAGS) | Signal 'printer in use' |
| CALL | OEDF,CLEAR-PRB | and clear the printer buffer. |
| LD | (DF-SZ),+02 | Set the size of the lower |
| CALL | 0D6B,CLS | part of the display and clear the whole display. |
| XOR | A | Now print the message |
| LD | DE,+1538 | '© 1982 Sinclair Research Ltd' |
| CALL | 0COA,PO-MSG | on the bottom line. |
| SET | 5,(TV-FLAG) | Signal 'the lower part will required to be cleared. |
| JR | 12A9,MAIN-1 | Jump forward into the main |

THE 'MAIN EXECUTION' LOOP
The main loop extends from location 12A2 to location 15AE and it controls the 'editing mode', the execution of direct commands and the production of reports.

| 12A2 | MAIN-EXEC | LD | (DF-SZ),+02 | The lower part of the screen is to be two lines in size. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | CALL | 1795,AUTO-LIST | Produce an automatic listing. |
| 12 A 9 | MAIN-1 | CALL | 16B0,SET-MIN | All the areas from E-LINE onwards are given their minimum configurations. |
| 12AC | MAIN-2 | LD | A, +00 | Channel ' K ' is opened before |
|  |  | CALL | 1601,CHAN-OPEN | calling the EDITOR. |
|  |  | CALL | OF2C,EDITOR | The EDITOR is called to allow the user to build up a BASIC line |
|  |  | CALL | 1B17,LINE-SCAN | The current line is scanned for correct syntax. |
|  |  | BIT | 7,(ERR-NR) | Jump forward if the syntax is |
|  |  | JR | NZ,12CF,MAIN-3 | correct. |
|  |  | BIT | 4,(FLAGS2) | Jump forward if other than |
|  |  | JR | Z,1303,MAIN-4 | channel ' K ' is being used. |
|  |  | LD | HL,(E-LINE) | Point to the start of the line with the error. |
|  |  | CALL | 11A7,REMOVE-FP | Remove the floating-point forms from this line. |
|  |  | LD | (ERR-NR), +FF | Reset ERR-NR and jump back |
|  |  | JR | 12AC,MAIN-2 | to MAIN-2 leaving the listing unchanged. |

The 'edit-line' has passed syntax and the three types of line that are possible have to be distinguished from each other.

| 12CF | MAIN-3 | HL,(E-LINE) | Point to the start of the line. |
| :--- | :--- | :--- | :--- |
|  | LD | (CH-ADD),HL | Set CH-ADD to the start also. |
|  | CALL | 19FB,E-LINE-NO | Fetch any line number into BC. |
|  | LD | A,B | Is the line number a valid |
|  | OR | C | one? |
|  | JR | NZ,155D,MAIN-ADD | Jump if it is so, and add the new |
|  |  |  | line to the existing program. |


| RST | 0018 | Fetch the first character of <br> CP |
| :--- | :--- | :--- |
| the line and see if the line is |  |  |

The 'edit-line' must start with a direct BASIC command so this line becomes the first line to be interpreted.

| BIT | 0,(FLAGS2) | Clear the whole display unless |
| :--- | :--- | :--- |
| CALL | NZ,0DAF,CL-ALL | the flag says it is unnecessary. |
| CALL | OD6E,CLS-LOWER | Clear the lower part anyway. |
| LD | A,+19 | Set the appropriate value |
| SUB | (S-POSN-hi) | for the scroll counter. |
| LD | (SCR-CT),A |  |
| SET | 7,(FLAGS) | Signal 'line execution'. |
| LD | (ERR-NR),+FF | Ensure ERR-NR is correct. |
| LD | (NSPPC),+01 | Deal with the first statement in <br> the line. |
| CALL | 1B8A,PROG-RUN | Now the line is interpreted. <br> Note: The address 1303 goes on <br>  |
|  |  | to the machine stack and is <br> addressed by ERR-SP. |
|  |  |  |

After the line has been interpreted and all the actions consequential to it have been completed a return is made to MAIN-4, so that a report can be made.

| 1303 | MAIN-4 | HALT |  | The maskable interrupt must be enabled. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | RES | 5,(FLAGS) | Signal 'ready for a new key'. |
|  |  | BIT | 1,(FLAGS2) | Empty the printer buffer if |
|  |  | CALL | NZ,0ECD,COPY-BUFF | it has been used. |
|  |  | LD | A,(ERR-NR) | Fetch the error number and |
|  |  | INC | A | increment it. |
| 1313 | MAIN-G | PUSH | AF | Save the new value. |
|  |  | LD | HL,+0000 | The system variables |
|  |  | LD | (FLAGX), H |  |
|  |  | LD | (X-PTR-hi), H | DEFADD are all set to zero. |
|  |  | LD | (DEFADD), HL |  |
|  |  | LD | HL,+0001 | Ensure that stream +00 |
|  |  | LD | (STRMS-6), HL | points to channel ' K ' |
|  |  | CALL | 16B0,SET-MIN | Clear all the work areas and the calculator stack. |
|  |  | RES | 5,(FLAGX) | Signal 'editing mode'. |
|  |  | CALL | 0D6E,CLS-LOWER | Clear the lower screen. |
|  |  | SET | 5,(TV-FLAG) | Signal 'the lower screen will require clearing'. |
|  |  | POP | AF | Fetch the report value. |
|  |  | LD | B,A | Make a copy in B. |
|  |  | CP | +0A | Jump forward with report |
|  |  | JR | C,133C,MAIN-5 | numbers '0 to 9'. |
|  |  | ADD | A,+07 | Add the ASCII letter offset value. |
| 133C | MAIN-5 | CALL | 15EF,OUT-CODE | Print the report code and |
|  |  | LD | A,+20 | follow it with a 'space'. |
|  |  | RST | 0010,PRINT-A-1 |  |
|  |  | LD | A, B | Fetch the report value and |
|  |  | LD | DE, +1391 | use it to identify the |
|  |  | CALL | 0C0A,PO-MSG | required report message. |
|  |  | XOR | A | Print the message and follow |
|  |  | LD | DE,+1536 | it by a 'comma' and a 'space'. |
|  |  | CALL | 0C0A,PO-MSG |  |
|  |  | LD | BC,(PPC) | Now fetch the current line |
|  |  | CALL | 1A1B,OUT-NUM1 | number and print it as well. |
|  |  | LD | A, +3 A | Follow it by a ':' |
|  |  | RST | 0010,PRINT-A-1 |  |

$\left.\begin{array}{llll} & \text { LD } & \text { C,(SUBPPC) } & \text { Fetch the current statement } \\ & \text { LD } & \text { B,+00 } & \text { number into the BC register }\end{array}\right]$ pair and print it.

## THE REPORT MESSAGES

Each message is given with the last character inverted (+80 hex.)

| 1391 | DEFB +80 | - initial byte is stepped over. |
| :--- | :--- | :--- |
| 1392 | Report 0 | - 'OK' |
| 1394 | Report 1 | - 'NEXT without FOR' |
| 13A4 | Report 2 | - 'Variable not found' |
| 13B6 | Report 3 | - 'Subscript wrong' |
| 13C6 | Report 4 | - 'Out of memory' |
| 13D2 | Report 5 | - 'Out of screen'' |
| 13DF | Report 6 | - 'Number too big' |
| 13ED | Report 7 | - 'RETURN without GOSUB' |
| 1401 | Report 8 | - 'End of file' |
| 140C | Report 9 | - 'STOP statement' |
| 141A | Report A | - 'Invalid argument' |
| 142A | Report B | - 'Integer out of range' |
| 143E | Report C | - 'Nonsense in BASIC' |
| 144F | Report D | - 'BREAK - CONT repeats' |
| 1463 | Report E | - 'Out of DATA' |
| 146E | Report F | - 'Invalid file name' |
| 147F | Report G | - 'No room for line' |
| 148F | Report H | - 'STOP in INPUT' |
| 149C | Report I | - 'FOR without NEXT' |
| 14AC | Report J | - 'Invalid I/O device' |
| 14BE | Report K | - 'Invalid colour' |
| 14CC | Report L | - 'BREAK into program' |
| 14DE | Report M | - 'RAMTOP no good' |
| 14EC | Report N | - 'Statement lost' |
| 14FA | Report O | - 'Invalid stream' |
| 1508 | Report P | - 'FN without DEF' |
| 1516 | Report Q | - 'Parameter error' |
| 1525 | Report R | - 'Tape loading error' |

There are also the following two messages.
1537 1539 $\quad$ '© 1982 Sinclair Research Ltd' $\quad$ - a 'comma' and a 'space'

Report G - No room for line

| 1555 REPORT-G | LD | A, +10 | 'G' has the code '10+07+30' |
| :--- | :--- | :--- | :--- |
|  | LD | BC,+0000 | Clear BC. |
|  | JP | $1313, M A I N-G$ | Jump back to give the report. |

THE 'MAIN-ADD' SUBROUTINE
This subroutine allows for a new BASIC line to be added to the existing BASIC program in the program area. If a line has both an old and a new version then the old one is 'reclaimed'. A new line that consists of only a line number does not go into the program area.

| 155D | MAIN-ADD | LD | (E-PPC), BC | Make the new line number the 'current line'. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | LD | HL,(CH-ADD) | Fetch $\mathrm{CH}-\mathrm{ADD}$ and save the |
|  |  | EX | DE,HL | address in DE. |
|  |  | LD | HL, +1555 | Push the address of REPORT-G |
|  |  | PUSH | HL | on to the machine stack. ERR-SP will now point to REPORT-G. |
|  |  | LD | HL,(WORKSP) | Fetch WORKSP. |
|  |  | SCF |  | Find the length of the line |
|  |  | SBC, | HL, DE | from after the line number to the 'carriage return' character inclusively. |
|  |  | PUSH | HL | Save the length. |
|  |  | LD | H, B | Move the line number to the |
|  |  | LD | L, C | HL register pair. |
|  |  | CALL | 196E,LINE-ADDR | Is there an existing line with this number? |
|  |  | JR | NZ,157D,MAIN-ADD1 | Jump if there was not. |
|  |  | CALL | 19B8,NEXT-ONE | Find the length of the 'old' |
|  |  | CALL | 19E8,RECLAIM-2 | line and reclaim it. |
| 157D | MAIN-ADD1 | POP | BC | Fetch the length of the |
|  |  | LD | A, C | 'new' line and jump forward |
|  |  | DEC | A | if it is only a 'line number |
|  |  | OR | B | and a carriage return'. |
|  |  | JR | 15AB,MAIN-ADD2 |  |
|  |  | PUSH | BC | Save the length. |
|  |  | INC | BC | Four extra locations will be |
|  |  | INC | BC | needed. |
|  |  | INC | BC |  |
|  |  | INC | BC | two for the length. |
|  |  | DEC | HL | Make HL point to the location before the 'destination'. |
|  |  | LD | DE,(PROG) | Save the current value of |
|  |  | PUSH | DE | PROG to avoid corruption when adding a first line. |
|  |  | CALL | 1655,MAKE-ROOM | Space for the new line is created. |
|  |  | POP | HL | The old value of PROG is |
|  |  | LD | (PROG), HL | fetched and restored. |
|  |  | POP | BC | A copy of the line length |
|  |  | PUSH | BC | (without parameters) is taken. |
|  |  | INC | DE | Make DE point to the end location of the new area |
|  |  | LD | HL,(WORKSP) | and HL to the 'carriage |
|  |  | DEC | HL | return' character of the new |
|  |  | DEC | HL | line in the editing area. |
|  |  | LDDR |  | Now copy over the line. |
|  |  | LD | HL,(E-PPC) | Fetch the line's number. |


|  |  | EX | DE,HL | Destination into HL \& number into DE. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | POP | BC | Fetch the new line's length. |
|  |  | LD | (HL), B | The high length byte. |
|  |  | DEC | HL |  |
|  |  | LD | (HL), C | The low length byte. |
|  |  | DEC | HL |  |
|  |  | LD | (HL), E | The low line number byte. |
|  |  | DEC | HL |  |
|  |  | LD | (HL), D | The high line number byte. |
| 15AB | MAIN-ADD2 | POP | AF | Drop the address of REPORT-G. |
|  |  | JP | 12A2,MAIN-EXEC | Jump back and this time do |
|  |  |  |  | produce and automatic listing. |

## THE 'INITIAL CHANNEL INFORMATION'

Initially there are four channels - 'K', 'S', 'R', \& 'P' - for communicating with the 'keyboard', 'screen', 'work space' and 'printer'. For each channel the output routine address comes before the input routine address and the channel's code.

| 15AF | DEFB | F4 09 |  | - PRINT-OUT |
| :---: | :---: | :---: | :---: | :---: |
|  | DEFB | A8 10 |  | - KEY-INPUT |
|  | DEFB | 4B |  | - 'K' |
| 15B4 | DEFB | F4 09 |  | - PRINT-OUT |
|  | DEFB | C4 15 |  | - REPORT-J |
|  | DEFB | 53 |  | - 'S' |
| 15B9 | DEFB | 81 0F |  | - ADD-CHAR |
|  | DEFB | C4 15 |  | - REPORT-J |
|  | DEFB | 52 |  | - 'R' |
| 15BE | DEFB | F4 09 |  | - PRINT-OUT |
|  | DEFB | C4 15 |  | - REPORT-J |
|  | DEFB | 50 |  | - 'P' |
| 15C3 | DEFB | 80 |  | - End marker. |
| Report J - Invalid I/O device |  |  |  |  |
| 15C4 | REPORT-J | $\begin{aligned} & \text { RST } \\ & \text { DEFB } \end{aligned}$ | $\begin{aligned} & \text { 0008,ERROR-1 } \\ & +12 \end{aligned}$ | Call the error handling routine |

THE 'INITIAL STREAM DATA'
Initially there are seven streams - +FD to +03 .

| 15C6 | DEFB | 0100 | - stream +FD leads to channel |  |  | 'K' |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 15C8 | DEFB | 0600 | - stream +FE |  |  | 'S' |
| 15CA | DEFB | OB 00 | - stream +FF |  |  | 'R' |
| 15CC | DEFB | 0100 | - stream +00 | " |  | 'K' |
| 15CE | DEFB | 0100 | - stream +01 |  | , | 'K' |
| 15D0 | DEFB | 0600 | - stream +02 |  | " | 'S' |
| 15D2 | DEFB | 1000 | - stream +03 | " | " | 'P' |

THE 'WAIT-KEY' SUBROUTINE
This subroutine is the controlling subroutine for calling the current input subroutine.

| 15D4 | WAIT-KEY | BIT | 5,(TV-FLAG) | Jump forward if the flag |
| :---: | :---: | :---: | :---: | :---: |
|  |  | JR | NZ,15DE,WAIT-KEY1 | indicates the lower screen does not require clearing. |
|  |  | SET | 3,(TV-FLAG) | Otherwise signal 'consider |
|  |  |  |  | the mode as having changed'. |
| 15DE | WAIT-KEY1 | CALL | 15E6,INPUT-AD | Call the input subroutine indirectly via INPUT-AD. |
|  |  | RET | C | Return with acceptable codes. |
|  |  | JR | Z,15DE,WAIT-KEY1 | Both the carry flag and the zero flag are reset if 'no key is being pressed'; otherwise signal an error. |

Report 8 - End of file
15E4 REPORT-8 RST 0008,ERROR-1 Call the error handling

## THE 'INPUT-AD' SUBROUTINE

The registers are saved and HL made to point to the input address.


THE 'MAIN PRINTING' SUBROUTINE
The subroutine is called with either an absolute value or a proper character code in the A register.

| 15EF | OUT-CODE | LD | E,+30 | Increase the value in the |
| :--- | :--- | :--- | :--- | :--- |
| 15F2 | PRINT-A-2 | ADD | A,E | A register by +30. |
|  |  | PUSH | HL | Again save the registers. |

Now call the actual subroutine. HL points to the output or the input address as directed.

| 15F7 | CALL-SUB | LD | E,(HL) | Fetch the low byte. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | INC | HL |  |
|  |  | LD | D,(HL) | Fetch the high byte. |
|  |  | EX | DE,HL | Move the address to the HL register pair. |
|  |  | CALL | 162C,CALL-JUMP | Call the actual subroutine. |
|  |  | POP | HL | Restore the registers. |
|  |  | EXX |  |  |
|  |  | RET |  | Return will be from here |
|  |  |  |  | unless an error occurred. |

THE 'CHAN-OPEN' SUBROUTINE
This subroutine is called with the A register holding a valid stream number - normally +FD to +03 . Then depending on the stream data a particular channel will be made the current channel.


Using the stream data now find the base address of the channel information associated with that stream.
1610
CHAN-OP-1 DEC
DE
HL,(CHANS)

Reduce the stream data.
The base address of the whole channel information area.

Form the required address in this area.

THE 'CHAN-FLAG' SUBROUTINE
The appropriate flags for the different channels are set by this subroutine.

| 1615 | CHAN-FLAG | LD | (CURCHL), HL | The HL register pair holds the base address for a particular channel. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | RES | 4,(FLAGS2) | Signal 'using other than channel 'K". |
|  |  | INC | HL | Step past the output |
|  |  | INC | HL | and the input addresses and |
|  |  | INC | HL | make HL point to the |
|  |  | INC | HL | channel code. |
|  |  | LD | C, (HL) | Fetch the code. |
|  |  | LD | HL,+162D | The base address of the 'channe code look-up table'. |
|  |  | CALL | 16DC,INDEXER | Index into this table and locate the required offset; but return if |
|  |  | RET | NC | there is not a matching channel code. |
|  |  | LD | D, +00 | Pass the offset to the |
|  |  | LD | E, (HL) | DE register pair. |
|  |  | ADD | HL, DE | Jump forward to the appropriate |
| 162C | CALL-JUMP | JP | (HL) | flag setting routine. |

THE 'CHANNEL CODE LOOK-UP' TABLE

| $162 D$ | DEFB | 4 B 06 | - channel 'K', | offset +06, | address 1634 |
| :--- | :--- | :--- | :--- | :--- | :--- |
| $162 F$ | DEFB | 5312 | - channel 'S', | offset +12, | address 1642 |
| 1631 | DEFB | $501 B$ | - channel 'P', | offset +1B, | address 164D |
| 1633 | DEFB | 00 | - end marker. |  |  |

the 'Channel ' K ' FLAG' SUBROUTINE

1634 CHAN-K | SET | 0,(TV-FLAG) |  |
| :--- | :--- | :--- |
|  |  | RES |
|  | SET | ,(FLAGS) |
|  |  | JR (FLAGS2) |
|  |  | 1646, CHAN-S-1 |

Signal 'using lower screen'.
Signal 'ready for a key'.
$\begin{array}{ll}\text { SET } & \text { 4,(FLAGS2) } \\ \text { JR } & 1646, \text { CHAN-S- }\end{array}$
Signal 'using channel 'K'. Jump forward.

## THE 'CHANNEL 'S' FLAG' SUBROUTINE

| 1642 | CHAN-S | RES | 0,(TV-FLAG) |
| :--- | :--- | :--- | :--- |
| 1646 | CHAN-S-1 | RES | 1,(FLAGS) |
|  |  | JP | 0D4D,TEMPS |

Signal 'using main screen'. Signal 'printer not being used'. Exit via TEMPS so as to set the colour system variables.

## THE 'CHANNEL 'P' FLAG' SUBROUTINE

164D CHAN-P | SET | 1,(FLAGS) $\quad$ Signal 'printer in use'. |
| ---: | :--- |

## THE 'MAKE-ROOM' SUBROUTINE

This is a very important subroutine. It is called on many occasions to 'open up' an area. In all cases the HL register pair points to the location after the place where 'room' is required and the BC register pair holds the length of the 'room' needed. When a single space only is required then the subroutine is entered at ONE-SPACE.

| 1652 | ONE-SPACE LD | BC, +0001 | Just the single extra location <br> is required. |
| :--- | :--- | :--- | :--- |
| 1655 | MAKE-ROOM PUSH | HL | Save the pointer. <br> CALL |
|  | 1F05,TEST-ROOM | Make sure that there is sufficient <br> memory available for the task <br> being undertaken. |  |
|  | POP | HL | Restore the pointer. |


| CALL | 1664,POINTERS | Alter all the pointers before <br> making the 'room'. |
| :--- | :--- | :--- |
| LD | HL,(STKEND) | Make HL hold the new STKEND. |
| EX | DE,HL | Switch 'old' and 'new'. |
| LDDR |  | Now make the 'room' |
| RET |  | and return. |

Note: This subroutine returns with the HL register pair pointing to the location before the new 'room' and the DE register pair pointing to the last of the new locations. The new 'room' therefore has the description: '(HL)+1' to '(DE)' inclusive.
However as the 'new locations' still retain their 'old values' it is also possible to consider the new 'room' as having been made after the original location '(HL)' and it thereby has the description '(HL)+2' to (DE)+1'.
In fact the programmer appears to have a preference for the 'second description' and this can be confusing.

## THE 'POINTERS' SUBROUTINE

Whenever an area has to be 'made' or 'reclaimed' the system variables that address locations beyond the 'position' of the change have to be amended as required. On entry the BC register pair holds the number of bytes involved and the HL register pair addresses the location before the 'position'.

| 1664 | POINTERS | PUSH | AF | These registers are saved. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | PUSH | HL | Copy the address of the 'position'. |
|  |  | LD | HL, +5C4B | This is VARS, the first of the |
|  |  | LD | A, +0 E | fourteen system pointers. |

A loop is now entered to consider each pointer in turn. Only those pointers that point beyond the 'position' are changed.

| 166B | PTR-NEXT | LD | E,(HL) | Fetch the two bytes of the |
| :---: | :---: | :---: | :---: | :---: |
|  |  | INC | HL | current pointer. |
|  |  | LD | D,(HL) |  |
|  |  | EX | (SP), HL | Exchange the system variable with the address of the 'position' |
|  |  | AND | A | The carry flag will become |
|  |  | SBC | HL, DE | set if the system variable's |
|  |  | ADD | HL, DE | address is to be updated. |
|  |  | EX | (SP), HL | Restore the 'position'. |
|  |  | JR | NC,167F,PTR-DONE | Jump forward if the pointer is to be left; otherwise change it. |
|  |  | PUSH | DE | Save the old value. |
|  |  | EX | DE,HL | Now add the value in BC |
|  |  | ADD | HL, BC | to the old value. |
|  |  | EX | DE,HL |  |
|  |  | LD | (HL), D | Enter the new value into the |
|  |  | DEC | HL | system variable - high byte |
|  |  | LD | (HL), E | before low byte. |
|  |  | INC | HL | Point again to the high byte. |
|  |  | POP | DE | Fetch the old value. |
| 167F | PTR-DONE | INC | HL | Point to the next system |
|  |  | DEC | A | variable and jump back until all |
|  |  | JR | NZ,166B,PTR-NEXT | fourteen have been considered. |

Now find the size of the block to be moved.

| EX | DE,HL |
| :--- | :--- |
| POP | DE |
| POP | AF |
| AND | A |
| SBC | HL,DE |
| LD | B,H |
| LD | C,L |
| INC | BC |

Put the old value of STKEND in
HL and restore the other registers.
Now find the difference between the old value of STKEND and the 'position'. Transfer the result to BC and add ' 1 ' for the inclusive byte.

| ADD | HL,DE | Reform the old value of |
| :--- | :--- | :--- |
| EX | DE,HL | STKEND and pass it to DE |
| RET |  | before returning. |

## THE 'COLLECT A LINE NUMBER' SUBROUTINE

On entry the HL register pair points to the location under consideration. If the location holds a value that constitutes a suitable high byte for a line number then the line number is returned in DE. However if this is not so then the location addressed by DE is tried instead; and should this also be unsuccessful line number zero is returned.

| $168 F$ | LINE-ZERO | DEFB | +00 | Line number zero. |
| :--- | :--- | :--- | :--- | :--- |
|  |  | DEFB | +00 |  |
| 1691 | LINE-NO-A | EX | DE,HL | Consider the other pointer. |
|  |  | LD | DE,+168F | Use line number zero. |

The usual entry point is at LINE-NO.

| 1695 | LINE-NO | LD | A,(HL) | Fetch the high byte and |
| :---: | :---: | :---: | :---: | :---: |
|  |  | AND | +C0 | test it. |
|  |  | JR | NZ,1691,LINE-NO-A | Jump back if not suitable. |
|  |  | LD | D,(HL) | Fetch the high byte. |
|  |  | INC | HL |  |
|  |  | LD | E,(HL) | Fetch the low byte and |

## THE 'RESERVE' SUBROUTINE

This subroutine is normally called by using RST 0030,BC-SPACES.
On entry here the last value on the machine stack is WORKSP and the value above it is the number of spaces that is to be 'reserved'. This subroutine always makes 'room' between the existing work space and the calculator stack.

| 169E | RESERVE |  | HL,(STKBOT) <br> HL <br> 1655,MAKE-ROOM <br> HL <br> HL <br> BC <br> (WORKSP),BC <br> BC <br> DE,HL <br> HL | Fetch the current value of STKBOT and decrement it to get the last location of the work space. <br> Now make 'BC spaces'. Point to the first new space and then the second. Fetch the old value of WORKSP and restore it. Restore BC - number of spaces. Switch the pointers, Make HL point to the first of the displaced bytes. Now return. |
| :---: | :---: | :---: | :---: | :---: |

Note: It can also be considered that the subroutine returns with the DE register pair pointing to a 'first extra byte' and the HL register pair pointing to a 'last extra byte', these extra bytes having been added after the original '(HL)+1' location.

## THE 'SET-MIN' SUBROUTINE

This subroutine resets the editing area and the areas after it to their minimum sizes. In effect it 'clears' the areas.

| 16B0 | SET-MIN | LD | HL,(E-LINE) | Fetch E-LINE. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | LD | (HL), +0D | Make the editing area hold |
|  |  | LD | (K-CUR), HL | only the 'carriage return' |
|  |  | INC | HL | character and the end marker. |
|  |  | LD | (HL), +80 |  |
|  |  | INC |  | Move on to clear the work |
|  |  | LD | (WORKSP),HL | space. |

Entering here will 'clear' the work space and the calculator stack.

| 16BF | SET-WORK | LD | HL,(WORKSP) |
| :--- | :--- | :--- | :--- |
|  | LD | (STKBOT),HL | Fetch the WORKSP. |
|  |  | This clears the work space. |  |

Entering here will 'clear' only the calculator stack.

| 16C5 | SET-STK | LD | HL,(STKBOT) |
| :--- | :--- | :--- | :--- |$\quad$ Fetch STKBOT.

In all cases make MEM address the calculator's memory area.

| PUSH | HL | Save STKEND. |
| :--- | :--- | :--- |
| LD | $\mathrm{HL},+5 \mathrm{C} 92$ | The base of the memory area. |
| LD | $(\mathrm{MEM}), \mathrm{HL}$ | Set MEM to this address. |
| POP | HL | Restore STKEND to the HL |
| RET |  | register pair before returning. |

## THE 'RECLAIM THE EDIT-LINE' SUBROUTINE'

16 R4 REC-EDIT LD

DE,(E-LINE) 19E5,RECLAIM-1

Fetch E-LINE. Reclaim the memory.

THE 'INDEXER' SUBROUTINE
This subroutine is used on several occasions to look through tables. The entry point is at INDEXER.

| 16DB | INDEXER-1 | INC | HL | Move on to consider the next <br> pair of entries. |
| :--- | :--- | :--- | :--- | :--- |
| 16DC | INDEXER | LD | A,(HL) | Fetch the first of a pair of <br> entries but return if it is |
|  |  | AND | A | zero - the end marker. |
|  | RET | Z | Compare it to the supplied |  |
|  |  |  | Code. |  |

THE 'CLOSE \#' COMMAND ROUTINE
This command allows the user to CLOSE streams. However for streams +00 to +03 the 'initial' stream data is restored and these streams cannot therefore be CLOSEd.

| 16E5 | CLOSE | CALL | 171E,STR-DATA | The existing data for the stream is fetched. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | CALL | 1701,CLOSE-2 | Check the code in that stream's channel. |
|  |  | LD | BC, +0000 | Prepare to make the stream's data zero. |
|  |  | LD | DE,+A3E2 | Prepare to identify the use of |
|  |  | EX | DE,HL | streams +00 to +03. |
|  |  | ADD | HL, DE | The carry flag will be set with streams +04 to +0 F . |
|  |  | JR | C,16FC,CLOSE-1 | Jump forward with these |
|  |  | LD | BC, +15D4 | streams; otherwise find the |
|  |  | ADD | HL, BC | correct entry in the 'initial stream data' table. |
|  |  | LD | C,(HL) | Fetch the initial data |
|  |  | INC | HL | for streams +00 to +03 . |
|  |  | LD | B,(HL) |  |
| 16FC | CLOSE-1 | EX | DE,HL | Now enter the data; either |
|  |  | LD | (HL), C | zero \& zero, or the initial |
|  |  | INC | HL | values. |
|  |  | LD | (HL), B |  |
|  |  | RET |  |  |

the 'CLOSE-2' SUBROUTINE
The code of the channel associated with the stream being closed has to be ' $K$ ', ' $S$ ', or ' $P$ '.
1701
CLOSE-2
PUSH
HL
Save the address of the stream's data.

| LD | HL,(CHANS) | Fetch the base address of the |
| :---: | :---: | :---: |
| ADD | HL, BC | channel information area and find the channel data for the stream being CLOSEd. |
| INC | HL | Step past the subroutine |
| INC | HL | addresses and pick up |
| INC | HL | the code for that channel. |
| LD | C,(HL) |  |
| EX | DE,HL | Save the pointer. |
| LD | HL,+1716 | The base address of the 'CLOSE stream look-up' table. |
| CALL | 16DC,INDEXER | Index into this table and locate the required offset. |
| LD | C,(HL) | Pass the offset to the BC |
| LD | B,+00 | register pair. |
| ADD | HL,BC | Jump forward to the |
| JP | (HL) | appropriate routine. |

## THE 'CLOSE STREAM LOOK-UP' TABLE

| 1716 | DEFB | $4 B 05$ |
| :--- | :--- | :--- |
| 1718 | DEFB | 5303 |
| $171 A$ | DEFB | 5001 |

- channel ' $K$ ', offset +05 , address 171 C
- channel 'S', offset +03 , address 171C
- channel 'P', offset +01 , address 171C

Note: There is no end marker at the end of this table.

## THE 'CLOSE STREAM' SUBROUTINE.

171C CLOSE-STR POP HL Fetch the channel information

THE 'STREAM DATA' SUBROUTINE
This subroutine returns in the BC register pair the stream data for a given stream.

| 171E | STR-DATA | CALL | 1E94,STK-TO-A | The given stream number is taken off the calculator stack |
| :---: | :---: | :---: | :---: | :---: |
|  |  | $\begin{aligned} & \text { CP } \\ & \text { JR } \end{aligned}$ | $\begin{aligned} & +10 \\ & \text { C,1727,STR-DATA1 } \end{aligned}$ | Give an error if the stream number is greater than +0 F . |

Report O - Invalid stream
1725 REPORT-O RST 0008,ERROR-1 Call the error handling

Continue with valid stream numbers.

## 1727

| STR-DATA1 | ADD | $\mathrm{A},+03$ |
| :--- | :--- | :--- |
|  | RLCA | $\mathrm{HL},+5 \mathrm{C} 10$ |
|  | LD |  |
|  |  | LD |
|  | LD | $\mathrm{C}, \mathrm{A}$ |
|  | ADD | $\mathrm{B},+00$ |
|  | LD | $\mathrm{HL}, \mathrm{BC}$ |
|  | INC | $\mathrm{C},(\mathrm{HL})$ |
|  | LD | HL |
|  | DEC | $\mathrm{B},(\mathrm{HL})$ |
|  | RET | HL |
|  |  |  |

Range now +03 to +12 ;
and now +06 to +24 .
The base address of the stream data area.
Move the stream code to the BC register pair. Index into the data area and fetch the two data bytes into the BC register pair.

Make the pointer address the first of the data bytes before returning.

THE 'OPEN \#' COMMAND ROUTINE
This command allows the user to OPEN streams. A channel code must be supplied and it must be ' K ', ' $k$ ', ' S ', ' s ', ' P ', or ' p '. Note that no attempt is made to give streams +00 to +03 their initial data.

1736 OPEN | RST | 0028,FP-CALC | Use the CALCULATOR. |
| :--- | :--- | :--- |
|  | DEFB | +01,exchange |

|  |  | DEFB | +38, end-calc | and the channel code. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | CALL | 171E,STR-DATA | Fetch the data for the stream. |
|  |  | LD | A, B | Jump forward if both bytes of |
|  |  | OR | C | the data are zero, i.e. the |
|  |  | JR | Z,1756,OPEN-1 | stream was in a closed state. |
|  |  | EX | DE,HL | Save DE. |
|  |  | LD | HL,(CHANS) | Fetch CHANS - the base |
|  |  | ADD | HL,BC | address of the channel |
|  |  | INC | HL | information and find the |
|  |  | INC | HL | code of the channel |
|  |  | INC | HL | associated with the stream |
|  |  | LD | A,(HL) | being OPENed. |
|  |  | EX | DE,HL | Return DE. |
|  |  | CP | +4B | The code fetched from the |
|  |  | JR | Z,1756,OPEN-1 | channel information area |
|  |  | CP | +53 | must be 'K', 'S' or 'P'; |
|  |  | JR | Z,1756,OPEN-1 | give an error if it is not. |
|  |  | CP | +50 |  |
|  |  | JR | NZ,1725,REPORT-O |  |
| 1756 | OPEN-1 | CALL | 175D,OPEN-2 | Collect the appropriate data in DE. |
|  |  | LD | (HL), E | Enter the data into the |
|  |  | INC | HL | two bytes in the stream |
|  |  | LD | (HL), D | information area. |
|  |  | RET |  | Finally return. |

THE 'OPEN-2' SUBROUTINE
The appropriate stream data bytes for the channel that is associated with the stream being OPENed are found.

| 175D | OPEN-2 | PUSH CALL <br> LD <br> OR <br> JR | HL <br> 2BF1,STK-FETCH <br> A,B <br> C <br> NZ,1767,OPEN-3 | Save HL <br> Fetch the parameters of the channel code. <br> Give an error if the expression supplied is a null expression; i.e. OPEN \#5,"". |
| :---: | :---: | :---: | :---: | :---: |
| Report F - Invalid file name |  |  |  |  |
| 1765 | REPORT-F | RST DEFB | $\begin{aligned} & \text { 0008,ERROR-1 } \\ & +0 \mathrm{E} \end{aligned}$ | Call the error handling routine. |
| Continue if no error occurred. |  |  |  |  |
| 1767 | OPEN-3 | PUSH | BC | The length of the expression is saved. |
|  |  | LD | A, (DE) | Fetch the first character. |
|  |  | AND | +DF | Convert lower case codes to upper case ones. |
|  |  | LD | C,A | Move code to the C register. |
|  |  | LD | HL, +177A | The base address of the 'OPEN stream look-up' table. |
|  |  | CALL | 16DC,INDEXER | Index into this table and locate the required offset. |
|  |  | JR | NC,1765,REPORT-F | Jump back if not found. |
|  |  | LD | C, (HL) | Pass the offset to the BC |
|  |  | LD | B, +00 | register pair. |
|  |  | ADD | HL, BC | Make HL point to the start of the appropriate subroutine. |
|  |  | POP | BC |  |
|  |  | JP | (HL) | expression before jumping to the subroutine. |

THE 'OPEN STREAM LOOK-UP' TABLE
177A DEFB 4B 06

- channel 'K', offset +06, address 1781

| $177 C$ | DEFB | 5308 | - channel 'S', offset +08, address 1785 |
| :--- | :--- | :--- | :--- |
| 177 E | DEFB | 500 A | - channel 'P', offset +0 A, address 1789 |
| 1780 | DEFB | 00 | - end marker; |

THE 'OPEN-K' SUBROUTINE
1781 OPEN-K $\quad$ E, $+01 \quad$ The data bytes will be +01

THE 'OPEN-S' SUBROUTINE

1785 OPEN-S LD | E,+06 |  |  |
| :--- | :--- | :--- |
|  |  | JR |

The data bytes will be +06 $\&+00$.

THE 'OPEN-P' SUBROUTINE

| 1789 | OPEN-P | LD | E, +10 | The data bytes will be +10 $\&+00$. |
| :---: | :---: | :---: | :---: | :---: |
| 178B | OPEN-END | DEC | BC | Decrease the length of the |
|  |  | LD | A,B | expression and give an error |
|  |  | OR | C | if it was not a single |
|  |  | JR | NZ,1765,REPORT-F | character; otherwise clear the |
|  |  | LD | D,A | D register, fetch HL and |
|  |  | POP | HL | return. |

## THE 'CAT, ERASE, FORMAT \& MOVE' COMMAND ROUTINES

In the standard SPECTRUM system the use of these commands leads to the production of report O-Invalid stream.

THE 'LIST \& LLIST' COMMAND ROUTINES
The routines in this part of the 16K program are used to produce listings of the current BASIC program. Each line has to have its line number evaluated, its tokens expanded and the appropriate cursors positioned.
The entry point AUTO-LIST is used by both the MAIN EXECUTION routine and the EDITOR to produce a single page of the listing.


The 'automatic' number has now to be altered to give a listing with the 'current' line appearing near the bottom of the screen.

| PUSH | DE | Save the 'automatic' number. |
| :--- | :--- | :--- |
| CALL | 196E,LINE-ADDR | Find the address of the |
| LD | DE,+02C0 | start of the 'current' line |
| EX | DE,HL | and produce an address roughly |


| SBC | HL,DE | a 'screen before it' (negated). <br> EX |
| :--- | :--- | :--- |
| (SP),HL | Save the 'result' on the machine <br> stack whilst the 'automatic' line |  |
| CALL | 196E,LINE-ADDR | address is also found (in HL). <br> POP |
| BC | The 'result' goes to the BC <br> register pair. |  |

A loop is now entered. The 'automatic' line number is increased on each pass until it is likely that the 'current' line will show on a listing.

| 17CE | AUTO-L-1 | PUSH CALL | $\mathrm{BC}$ 19B8,NEXT-ONE |
| :---: | :---: | :---: | :---: |
|  |  | POP | BC |
|  |  | ADD | HL,BC |
|  |  | JR | C,17E4,AUTO-L-3 |
|  |  | EX | DE,HL |
|  |  | LD | D,(HL) |
|  |  | INC | HL |
|  |  | LD | E,(HL) |
|  |  | DEC | HL |
|  |  | LD | (S-TOP),DE |
|  |  | JR | 17CE,AUTO-L-1 |

Save the 'result'.
Find the address of the start of the line after the present 'automatic' line (in DE).
Restore the 'result'.
Perform the computation and jump forward if finished.
Move the next line's address to the HL register pair and collect its line number.

Now S-TOP can be updated and the test repeated with the new line.

When E-PPC is less than S-TOP.
Fetch the top line's number and hence its address. If the line cannot be found use DE instead.
The listing is produced. The return will be to here unless scrolling was needed to show the current line.

Use stream +03 . Jump forward.

THE 'LIST' ENTRY POINT
The 'main screen' channel will need to be opened.

| $\begin{aligned} & \text { 17F9 } \\ & \text { 17FB } \end{aligned}$ | LIST <br> LIST-1 | LD | A, +02 | Use stream +02 . |
| :---: | :---: | :---: | :---: | :---: |
|  |  | LD | (TV-FLAG),+00 | Signal 'an ordinary listing in the main part of the screen'. |
|  |  | CALL | 2530,SYNTAX-Z | Open the channel unless |
|  |  | CALL | NZ,1601,CHAN-OPEN | checking syntax. |
|  |  | RST | 0018,GET-CHAR | With the present character in |
|  |  | CALL | 2070,STR-ALTER | the A register see if the stream is to be changed. |
|  |  | JR | C,181F,LIST-4 | Jump forward if unchanged. |
|  |  | RST | 0018,GET-CHAR | Is the present character |
|  |  | CP | +3B | a ';'? |
|  |  | JR | Z,1814,LIST-2 | Jump if it is. |
|  |  | CP | +2C | Is it a ','? |
|  |  | JR | NZ,181A,LIST-3 | Jump if it is not. |
| 1814 | LIST-2 | RST | 0020,NEXT-CHAR | A numeric expression must |
|  |  | CALL | 1C82,EXPT-1NUM | follow, e.g. LIST \#5,20 |
|  |  | JR | 1822,LIST-5 | Jump forward with it. |
| 181A | LIST-3 | CALL | 1CE6,USE-ZERO | Otherwise use zero and |
|  |  | JR | 1822,LIST-5 | also jump forward. |

Come here if the stream was unaltered.
\(\left.$$
\begin{array}{llll}\text { 181F } & \text { LIST-4 } & \text { CALL } & \text { 1CDE,FETCH-NUM }\end{array}
$$ \begin{array}{l}Fetch any line or use zero if <br>
none supplied. <br>
If checking the syntax of the <br>
edit-line move on to the next <br>

statement.\end{array}\right]\)| Line number to BC. |
| :--- |

Now the controlling loop for printing a series of lines is entered.


## THE 'PRINT A WHOLE BASIC LINE' SUBROUTINE

The HL register pair points to the start of the line - the location holding the high byte of the line number. Before the line number is printed it is tested to determine whether it comes before the 'current' line, is the 'current' line or comes after.



Now enter a loop to print all the codes in the rest of the BASIC line - jumping over floating-point forms as necessary.

| 1894 | OUT-LINE4 | LD | HL,(X-PTR) | Fetch the syntax error |
| :---: | :---: | :---: | :---: | :---: |
|  |  | AND | A | pointer and jump forward |
|  |  | SBC | HL, DE | unless it is time to print |
|  |  | JR | NZ,18A1,OUT-LINE5 | the error marker. |
|  |  | LD | A, +3F | Print the error marker now. |
|  |  | CALL | 18C1,OUT-FLASH | It is a flashing '?'. |
| 18A1 | OUT-LINE5 | CALL | 18E1,OUT-CURS | Consider whether to print the cursor. |
|  |  | EX | DE,HL | Move the pointer to HL now. |
|  |  | LD | A, (HL) | Fetch each character in turn. |
|  |  | CALL | 18B6,NUMBER | If the character is a 'number marker' then the hidden floatingpoint form is not to be printed. |
|  |  | INC | HL | Update the pointer for the next pass. |
|  |  | CP | +0D | Is the character a 'carriage return'. |
|  |  | JR | Z,18B4,OUT-LINE6 | Jump if it is. |
|  |  | EX | DE,HL | Switch the pointer to DE. |
|  |  | CALL | 1937,OUT-CHAR | Print the character. |
|  |  | JR | 1894,OUT-LINE4 | Go around the loop for at least one further pass. |

The line has now been printed.

| 18B4 OUT-LINE6 | POP | DE | Restore the DE register pair <br> and return. |
| :--- | :--- | :--- | :--- |

THE 'NUMBER' SUBROUTINE
If the A register holds the 'number marker' then the HL register pair is advanced past the floating-point form.

| 18B6 | NUMBER | CP | $+0 \mathrm{E}$ | Is the character a 'number |
| :---: | :---: | :---: | :---: | :---: |
|  |  | RET | NZ | marker'. Return if not. |
|  |  | INC | HL | Advance the pointer six |
|  |  | INC | HL | times so as to step past the |
|  |  | INC | HL | 'number marker' and the five |
|  |  | INC | HL | locations holding the |
|  |  | INC | HL | floating-point form. |
|  |  | INC | HL |  |
|  |  | LD | A,(HL) | Fetch the current code before |

THE 'PRINT A FLASHING CHARACTER' SUBROUTINE
The 'error cursor' and the 'mode cursors' are printed using this subroutine.


## THE 'PRINT THE CURSOR' SUBROUTINE

A return is made if it is not the correct place to print the cursor but if it is then either 'C', ' $E$ ', ' $G$ ', ' $K$ ' or ' $L$ ' will be printed.

| 18E1 | OUT-CURS | LD | HL,(K-CUR) | Fetch the address of the |
| :---: | :---: | :---: | :---: | :---: |
|  |  | AND | A | cursor but return if the |
|  |  | SBC | HL, DE | correct place is not being |
|  |  | RET | NZ | considered. |
|  |  | LD | A,(MODE) | The current value of MODE is |
|  |  | RLC | A | fetched and doubled. |
|  |  | JR | Z,18F3,OUT-C-1 | Jump forward unless dealing with Extended mode or Graphics. |
|  |  | ADD | A, +43 | Add the appropriate offset to give 'E' or 'G'. |
|  |  | JR | 1909,OUT-C-2 | Jump forward to print it. |
| 18F3 | OUT-C-1 | LD | HL, +5C3B | This is FLAGS. |
|  |  | RES | 3,(HL) | Signal 'K-mode'. |
|  |  | LD | A, +4B | The character 'K'. |
|  |  | BIT | 2,(HL) | Jump forward to print 'K'. |
|  |  | JR | Z,1909,OUT-C-2 | If 'the printing is to be in K-mode'. |
|  |  | SET | 3,(HL) | The 'printing is to be in L-mode' so signal 'L-MODE'. |
|  |  | INC | A | Form the character 'L'. |
|  |  | BIT | 3,(FLAGS2) | Jump forward if not in |
|  |  | JR | Z,1909,OUT-C-2 | 'C-mode'. |
|  |  | LD | A, +43 | The character ' C '. |
| 1909 | OUT-C-2 | PUSH | DE | Save the DE register pair |
|  |  | CALL | 18C1,OUT-FLASH | whilst the cursor is printed |
|  |  | POP | DE | - FLASHing. |
|  |  | RET |  | Return once it has been done. |

Note: It is the action of considering which cursor-letter is to be printed that determines the mode - 'K' vs. 'L/C'.

## THE 'LN-FETCH' SUBROUTINE

This subroutine is entered with the HL register pair addressing a system variable - S-TOP or E-PPC. The subroutine returns with the system variable holding the line number of the following line.

| 190F LN-FETCH | LD | E,(HL) | The line number held by the <br> INC |
| :--- | :--- | :--- | :--- |
|  | HL | system variable is collected. |  |


| LD | D,(HL) | The pointer is saved. |
| :--- | :--- | :--- |
| PUSH | HL | The line number is moved to the |
| EX | DE,HL | HL register pair and incremented. <br> INC |
| CALL | HL | The address of the start of this <br> line is found, or the next line <br> if the actual line number is not |
|  |  | being used. <br> The number of that line is |
| CALL | 1695,LINE-NO | fetched. <br> The pointer to the system <br> variable is restored. |
| POP | HL |  |

The entry point LN-STORE is used by the EDITOR.
191C LN-STORE BIT

| RET | NZ |
| :--- | :--- |
| LD | $(H L), D$ |
| DEC | $H L$ |
| LD | $(H L), E$ |

Return if in 'INPUT mode'; otherwise proceed to enter the line number into the two locations of the system variable. Return when it has been done.

## THE 'PRINTING CHARACTERS IN A BASIC LINE' SUBROUTINE

All of the character/token codes in a BASIC line are printed by repeatedly calling this subroutine. The entry point OUT-SP-NO is used when printing line numbers which may require leading spaces.

| 1925 | OUT-SP-2 | LD | A, E | The A register will hold +20 fo a space or +FF for no-space. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | AND | A | Test the value and return if |
|  |  | RET | M | there is not to be a space. |
|  |  | JR | 1937,OUT-CHAR | Jump forward to print a space |
| 192A | OUT-SP-NO | XOR | A | Clear the A register. |

The HL register pair holds the line number and the BC register the value for 'repeated subtraction'. (BC holds '-1000, -100 or -10 '.)

| 192B | OUT-SP-1 | ADD <br> INC <br> JR <br> SBC <br> DEC <br> JR <br> JP | HL,BC <br> A <br> C,192B,OUT-SP-1 <br> HL,BC <br> A <br> Z,1925,OUT-SP-2 <br> 15EF,OUT-CODE | The 'trial subtraction'. Count each 'trial'. Jump back until exhausted. Restore last 'subtraction' and discount it. If no 'subtractions' were possible jump back to see if a space is to be printed. Otherwise print the digit. |
| :---: | :---: | :---: | :---: | :---: |

The entry point OUT-CHAR is used for all characters, tokens and control characters.

| 1937 | OUT-CHAR | CALL | 2D1B,NUMERIC |
| :--- | :--- | :--- | :--- | | Return carry reset if handling a |
| :--- |
| digit code. |


| 195A | OUT-CH-1 | CP | +22 | Accept for printing all characters except "'". Save the character code whilst changing the 'quote mode'. Fetch FLAGS2 and flip bit 2. <br> Enter the amended value and restore the character code. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | JR | NZ,1968,OUT-CH-2 |  |
|  |  | PUSH | AF |  |
|  |  | LD | A,(FLAGS2) |  |
|  |  | XOR | +04 |  |
|  |  | LD | (FLAGS2),A |  |
|  |  | POP | AF |  |
| 1968 | OUT-CH-2 | SET | 2,(FLAGS) | Signal 'the next character is to be printed in L-mode'. |
| 196C | OUT-CH-3 | $\begin{aligned} & \text { RST } \\ & \text { RET } \end{aligned}$ | 0010,PRINT-A-1 | The present character is printed before returning. |

Note: It is the consequence of the tests on the present character that determines whether the next character is to be "printed in 'K' or 'L' mode".
Also note how the program does not cater for ':' in REM statements.

## THE 'LINE-ADDR' SUBROUTINE

For a given line number, in the HL register pair, this subroutine returns the starting address of that line or the 'first line after', in the HL register pair, and the start of the previous line in the DE register pair.
If the line number is being used the zero flag will be set. However if the 'first line after' is substituted then the zero flag is returned reset.
196E

LINE-ADDR PUSH
LD
LD
LD E,L

Save the given line number
Fetch the system variable
PROG and transfer the address
to the DE register pair.

Now enter a loop to test the line number of each line of the program against the given line number until the line number is matched or exceeded.

1974 LINE-AD-1 \begin{tabular}{llll}
POP \& BC \& The given line number. <br>

\& CALL \& 1980,CP-LINES \& | Compare the given line number |
| :--- |
| Cogainst the addressed line | <br>

\& RET \& NC \& | aguber. Return if carry reset; |
| :--- |
| numberwise address the next | <br>

\& PUSH \& BC \& | others |
| :--- |
| line's number. | <br>

\& CALL \& 19B8,NEXT-ONE \& Switch the pointers and <br>
\& EX \& DE,HL \& jump back to consider the next
\end{tabular}

## THE 'COMPARE LINE NUMBERS' SUBROUTINE

The given line number in the $B C$ register pair is matched against the addressed line number.

| 1980 | CP-LINES | LD | A,(HL) | Fetch the high byte of the |
| :---: | :---: | :---: | :---: | :---: |
|  |  | CP | B | addressed line number and |
|  |  | RET | NZ | compare it. Return if they do not match. |
|  |  | INC | HL | Next compare the low bytes. |
|  |  | LD | A,(HL) | Return with the carry flag |
|  |  | DEC | HL | set if the addressed line |
|  |  | CP | C | number has yet to reach the |
|  |  | RET |  | given line number. |

## THE 'FIND EACH STATEMENT' SUBROUTINE

This subroutine has two distinct functions.
I. It can be used to find the 'D'th. statement in a BASIC line - returning with the HL register pair addressing the location before the start of the statement and the zero flag set.
II. Also the subroutine can be used to find a statement, if any, that starts with a given token code (in the E register).

| 1988 |  | INC | HL | Not used. |
| :--- | :--- | :--- | :--- | :--- |
|  |  | INC | HL |  |
| 198B | EACH-STMT | LNC | HL | (CH-ADD),HL |
|  |  | LD | C,+00 | Set CH-ADD to the current byte. |

Enter a loop to handle each statement in the BASIC line.


Now enter another loop to consider the individual characters in the line to find where the statement ends

| 1998 | EACH-S-2 | $\begin{aligned} & \text { INC } \\ & \text { LD } \end{aligned}$ | $\begin{aligned} & \mathrm{HL} \\ & \mathrm{~A},(\mathrm{HL}) \end{aligned}$ | Update the pointer and fetch the new code. |
| :---: | :---: | :---: | :---: | :---: |
| 199A | EACH-S-3 | CALL | 18B6,NUMBER | Step over any number. |
|  |  | LD | (CH-ADD), HL | Update CH-ADD. |
|  |  | CP | +22 | Jump forward if the character |
|  |  | JR | NZ,19A5,EACH-S-4 | is not a "'". |
|  |  | DEC | C | Otherwise set the 'quotes flag'. |
| 19A5 | EACH-S-4 | CP | +3A | Jump forward if the character |
|  |  | JR | Z,19AD,EACH-S-5 | is a ' $:$ '. |
|  |  | CP | +CB | Jump forward unless the code |
|  |  | JR | NZ,19B1,EACH-S-6 | is the token 'THEN'. |
| 19AD | EACH-S-5 | BIT | 0, C | Read the 'quotes flag' and |
|  |  | JR | Z,1990,EACH-S-1 | jump back at the end of each statement (including after 'THEN'). |
| 19B1 | EACH-S-6 | CP | +0D | Jump back unless at the end |
|  |  | JR | NZ,1998,EACH-S-2 | of a BASIC line. |
|  |  | DEC | D | Decrease the statement |
|  |  | SCF |  | counter and set the carry |
|  |  | RET |  | flag before returning. |

## THE 'NEXT-ONE' SUBROUTINE

This subroutine can be used to find the 'next line' in the program area or the 'next variable' in the variables area. The subroutine caters for the six different types of variable that are used in the SPECTRUM system.

| $19 \mathrm{B8}$ | NEXT-ONE | PUSH | HL | Save the address of the current line or variable. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | LD | A, (HL) | Fetch the first byte. |
|  |  | CP | +40 | Jump forward if searching |
|  |  | JR | C,19D5,NEXT-O-3 | for a 'next line'. |
|  |  | BIT | 5,A | Jump forward if searching for |
|  |  | JR | Z,19D6,NEXT-O-4 | the next string or array variable. |
|  |  | ADD | A,A | Jump forward with simple |
|  |  | JP | M,19C7,NEXT-O-1 | numeric and FOR-NEXT variables. |
|  |  | CCF |  | Long name numeric variables only. |
| 19 C 7 | NEXT-O-1 | LD | BC, +0005 | A numeric variable will |
|  |  | JR | NC,19CE,NEXT-O-2 | occupy five locations but a |
|  |  | LD | C, +12 | FOR-NEXT control variable |
|  |  |  |  | will need eighteen locations. |
| 19CE | NEXT-O-2 | RLA |  | The carry flag becomes reset for long named variables only; until the final character of the |



In all cases the address of the 'next' line or variable is found.

| 19DB | HEXT-O-5 | ADD | Point to the first byte of the <br> 'next' line or variable. |
| :--- | :--- | :--- | :--- |
|  | POP | DE | Fetch the address of the <br> previous one and continue into <br> the 'difference' subroutine. |

THE 'DIFFERENCE' SUBROUTINE
The 'length' between two 'starts' is formed in the BC register pair. The pointers are reformed but returned exchanged.

| 19DD | DIFFER | AND | A |
| :--- | :--- | :--- | :--- | | Prepare for a true subtraction. |
| :--- |
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## THE 'RECLAIMING' SUBROUTINE

The entry point RECLAIM-1 is used when the address of the first location to be reclaimed is in the DE register pair and the address of the first location to be left alone is in the HL register pair. The entry point RECLAIM-2 is used when the HL register pair points to the first location to be reclaimed and the BC register pair holds the number of the bytes that are to be reclaimed.

19E5 RECLAIM-1 \begin{tabular}{llll}

RE8 \& CALL \& | 19DD,DIFFER |
| :--- |
| RECLAIM-2 |
| to develop the appropriate | <br>

values.
\end{tabular}

## THE 'E-LINE-NO' SUBROUTINE

This subroutine is used to read the line number of the line in the editing area. If there is no line number, i.e. a direct BASIC line, then the line number is considered to be zero.
In all cases the line number is returned in the BC register pair.

| 19FB | E-LINE-NO | LD | HL,(E-LINE) | Pick up the pointer to the edit-line. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | DEC |  | Set the CH-ADD to point to the |
|  |  | LD | (CH-ADD), HL | location before any number. |
|  |  | RST | 0020,NEXT-CHAR | Pass the first code to the A register. |
|  |  | LD | HL, +5C92 | However before considering |
|  |  | LD | (STKEND),HL | the code make the calculator's memory area a temporary calculator stack area. |
|  |  | CALL | 2D3B,INT-TO-FP | Now read the digits of the line number. Return zero if no number exists. |
|  |  | CALL | 2DA2,FP-TO-BC | Compress the line number into the BC register pair. |
|  |  | JR | C,1A15,E-L-1 | Jump forward if the number exceeds '65,536'. |
|  |  | LD | HL,+D8F0 | Otherwise test it against |
|  |  | ADD | HL, BC | '10,000'. |
| $1 \mathrm{A15}$ | E-L-1 | JP | C,1C8A,REPORT-C | Give report C if over '9,999'. |
|  |  | JP | 16C5,SET-STK | Return via SET-STK that restores the calculator stack to its rightful place. |

## THE 'REPORT AND LINE NUMBER PRINTING' SUBROUTINE

The entry point OUT-NUM-1 will lead to the number in the BC register pair being printed. Any value over ' 9,999 ' will not however be printed correctly.
The entry point OUT-NUM-2 will lead to the number indirectly addressed by the HL register pair being printed. This time any necessary leading spaces will appear. Again the limit of correctly printed numbers is '9,999'.

| 1A1B | OUT-NUM-1 | PUSH | DE | Save the other registers |
| :---: | :---: | :---: | :---: | :---: |
|  |  | PUSH | HL | throughout the subroutine. |
|  |  | XOR | A | Clear the A register. |
|  |  | BIT | 7,B | Jump forward to print a zero rather than '-2' when |
|  |  | JR | NZ,1A42,OUT-NUM-4 | reporting on the edit-line. |
|  |  | LD | H,B | Move the number to the |
|  |  | LD | L, C | HL register pair. |
|  |  | LD | E,+FF | Flag 'no leading spaces'. |
|  |  | JR | 1A30,OUT-NUM-3 | Jump forward to print the number. |
| 1 A28 | OUT-NUM-2 | PUSH | DE | Save the DE register pair. |
|  |  | LD | D,(HL) | Fetch the number into the |
|  |  | INC | HL | DE register pair and save |
|  |  | LD | E,(HL) | the pointer (updated). |
|  |  | PUSH | HL |  |
|  |  | EX | DE,HL | Move the number to the HL |
|  |  | LD | E, +20 | register pair and flag 'leading |
|  |  |  |  | space are to be printed'. |

Now the integer form of the number in the HL register pair is printed.

| OUT-NUM-3 | LD | BC,+FC18 |
| :--- | :--- | :--- |
|  | CALL | 192A,OUT-SP-NO |
|  | LD | BC,+FF9C |
|  | CALL | 192A,OUT-SP-NO |
|  | LD | C,+F6 |
|  | CALL | $192 A$, OUT-SP-NO |

This is ' $-1,000$ '
Print a first digit.
This is '-100'.
Print the second digit.
This is '-10'.
Print the third digit.

| 1A42 | OUT-NUM-4 | LD | A, L | Move any remaining part of the number to the A register. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | CALL | 15EF,OUT-CODE | Print the digit. |
|  |  | POP |  | Restore the registers |
|  |  | POP | DE | before returning. |

## BASIC LINE AND COMMAND INTERPRETATION

## THE SYNTAX TABLES

i. The offset table

There is an offset value for each of the fifty BASIC commands

|  |  |  | command | address |  |  |  | command | address |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1A48 | DEFB | +B1 | DEF FN | 1AF9 | 1 A61 | DEFB | +94 | BORDER | 1AF5 |
| 1A49 | DEFB | +CB | CAT | 1B14 | 1 A62 | DEFB | +56 | CONTINUE | 1 AB8 |
| 1A4A | DEFB | +BC | FORMAT | 1 B06 | 1 A63 | DEFB | +3F | DIM | 1AA2 |
| 1A4B | DEFB | +BF | MOVE | 1B0A | 1 A64 | DEFB | +41 | REM | 1AA5 |
| 1A4C | DEFB | +C4 | ERASE | $1 \mathrm{B10}$ | 1 A65 | DEFB | $+2 \mathrm{~B}$ | FOR | 1 A90 |
| 1A4D | DEFB | +AF | OPEN \# | 1AFC | 1 A66 | DEFB | +17 | GO TO | 1A7D |
| 1A4E | DEFB | +B4 | CLOSE \# | 1 B 02 | 1 A67 | DEFB | +1F | GO SUB | 1 A86 |
| 1A4F | DEFB | +93 | MERGE | 1AE2 | 1 A68 | DEFB | +37 | INPUT | 1A9F |
| 1 A50 | DEFB | +91 | VERIFY | 1AE1 | 1 A69 | DEFB | +77 | LOAD | 1 AE0 |
| 1A51 | DEFB | +92 | BEEP | 1AE3 | 1A6A | DEFB | +44 | LIST | 1AAE |
| 1A52 | DEFB | +95 | CIRCLE | 1AE7 | 1A6B | DEFB | +0F | LET | 1A7A |
| 1A53 | DEFB | +98 | INK | 1AEB | 1A6C | DEFB | +59 | PAUSE | 1AC5 |
| 1A54 | DEFB | +98 | PAPER | 1AEC | 1A6D | DEFB | +2B | NEXT | 1 A98 |
| 1A55 | DEFB | +98 | FLASH | 1AED | 1A6E | DEFB | +43 | POKE | $1 \mathrm{AB1}$ |
| 1A56 | DEFB | +98 | BRIGHT | 1AEE | 1A6F | DEFB | +2D | PRINT | 1A9C |
| 1A57 | DEFB | +98 | INVERSE | 1AEF | 1 A70 | DEFB | +51 | PLOT | 1AC1 |
| 1A58 | DEFB | +98 | OVER | 1AF0 | 1A71 | DEFB | $+3 \mathrm{~A}$ | RUN | 1AAB |
| 1A59 | DEFB | +98 | OUT | 1AF1 | 1A72 | DEFB | +6D | SAVE | 1ADF |
| 1A5A | DEFB | +7F | LPRINT | 1AD9 | 1 A73 | DEFB | +42 | RANDOMIZE | 1 AB5 |
| 1A5B | DEFB | +81 | LLIST | 1ADC | 1 A74 | DEFB | +0D | IF | 1A81 |
| 1A5C | DEFB | +2E | STOP | 1A8A | 1A75 | DEFB | +49 | CLS | 1 ABE |
| 1A5D | DEFB | +6C | READ | $1 \mathrm{AC9}$ | 1A76 | DEFB | +5C | DRAW | 1AD2 |
| 1A5E | DEFB | +6E | DATA | 1ACC | 1 A77 | DEFB | +44 | CLEAR | 1ABB |
| 1A5F | DEFB | +70 | RESTORE | 1ACF | 1 A78 | DEFB | +15 | RETURN | 1A8D |
| 1A60 | DEFB | +48 | NEW | 1AA8 | 1A79 | DEFB | +5D | COPY | 1AD6 |

ii. The parameter table

For each of the fifty BASIC commands there are up to eight entries in the parameter table. These entries comprise command class details, required separators and, where appropriate, command routine addresses.

| 1A7A | P-LET | DEFB | +01 | CLASS-01 |
| :---: | :---: | :---: | :---: | :---: |
|  |  | DEFB | +3D | '=' |
|  |  | DEFB | +02 | CLASS-02 |
| 1A7D | P-GO-TO | DEFB | +06 | CLASS-06 |
|  |  | DEFB | +00 | CLASS-00 |
|  |  | DEFB | +67,+1E | GO-TO,1E67 |
| 1 A81 | P-IF | DEFB | +06 | CLASS-06 |
|  |  | DEFB | +CB | 'THEN' |
|  |  | DEFB | +05 | CLASS-05 |
|  |  | DEFB | +F0,+1C | IF,1CF0 |
| 1 A86 | P-GO-SUB | DEFB | +06 | CLASS-06 |
|  |  | DEFB | +00 | CLASS-00 |
|  |  | DEFB | +ED, +1E | GO-SUB,1EED |
| 1A8A | P-STOP | DEFB | +00 | CLASS-00 |
|  |  | DEFB | +EE, +1C | STOP,1CEE |
| 1A8D | P-RETURN | DEFB | +00 | CLASS-00 |
|  |  | DEFB | +23,+1F | RETURN,1F23 |
| 1 A 90 | P-FOR | DEFB | +04 | CLASS-04 |
|  |  | DEFB | +3D | '=' |
|  |  | DEFB | +06 | CLASS-06 |
|  |  | DEFB | +CC | 'TO' |
|  |  | DEFB | +06 | CLASS-06 |
|  |  | DEFB | +05 | CLASS-05 |
|  |  | DEFB | +03,+1D | FOR,1D03 |


| 1A98 | P-NEXT | DEFB | +04 | CLASS-04 |
| :---: | :---: | :---: | :---: | :---: |
|  |  | DEFB | +00 | CLASS-00 |
|  |  | DEFB | +AB,+1D | NEXT,1DAB |
| 1A9C | P-PRINT | DEFB | +05 | CLASS-05 |
|  |  | DEFB | +CD,+1F | PRINT,1FCD |
| 1A9F | P-INPUT | DEFB | +05 | CLASS-05 |
|  |  | DEFB | +89,+20 | INPUT,2089 |
| 1AA2 | P-DIM | DEFB | +05 | CLASS-05 |
|  |  | DEFB | +02,+2C | DIM,2C02 |
| 1AA5 | P-REM | DEFB | +05 | CLASS-05 |
|  |  | DEFB | +B2,+1B | REM,1BB2 |
| 1AA8 | P-NEW | DEFB | +00 | CLASS-00 |
|  |  | DEFB | +B7,+11 | NEW,11B7 |
| 1AAB | P-RUN | DEFB | +03 | CLASS-03 |
|  |  | DEFB | +A1,+1E | RUN,1EA1 |
| 1AAE | P-LIST | DEFB | +05 | CLASS-05 |
|  |  | DEFB | +F9,+17 | LIST,17F9 |
| 1AB1 | P-POKE | DEFB | +08 | CLASS-08 |
|  |  | DEFB | +00 | CLASS-00 |
|  |  | DEFB | +80,+1E | POKE,1E80 |
| 1AB5 | P-RANDOM | DEFB | +03 | CLASS-03 |
|  |  | DEFB | +4F,+1E | RANDOMIZE,1E4F |
| 1AB8 | P-CONT | DEFB | +00 | CLASS-00 |
|  |  | DEFB | +5F,+1E | CONTINUE,1E5F |
| 1ABB | P-CLEAR | DEFB | +03 | CLASS-03 |
|  |  | DEFB | +AC,+1E | CLEAR,1EAC |
| 1ABE | P-CLS | DEFB | +00 | CLASS-00 |
|  |  | DEFB | +6B,+0D | CLS,0D6B |
| 1AC1 | P-PLOT | DEFB | +09 | CLASS-09 |
|  |  | DEFB | +00 | CLASS-00 |
|  |  | DEFB | +DC,+22 | PLOT,22DC |
| 1AC5 | P-PAUSE | DEFB | +06 | CLASS-06 |
|  |  | DEFB | +00 | CLASS-00 |
|  |  | DEFB | +3A,+1F | PAUSE, 1F3A |
| 1AC9 | P-READ | DEFB | +05 | CLASS-05 |
|  |  | DEFB | +ED,+1D | READ,1DED |
| 1ACC | P-DATA | DEFB | +05 | CLASS-05 |
|  |  | DEFB | +27,+1E | DATA,1E27 |
| 1ACF | P-RESTORE | DEFB | +03 | CLASS-03 |
|  |  | DEFB | +42,+1E | RESTORE,1E42 |
| 1AD2 | P-DRAW | DEFB | +09 | CLASS-09 |
|  |  | DEFB | +05 | CLASS-05 |
|  |  | DEFB | +82,+23 | DRAW,2382 |
| 1AD6 | P-COPY | DEFB | +00 | CLASS-00 |
|  |  | DEFB | +AC+0E | COPY,0EAC |
| 1AD9 | P-LPRINT | DEFB | +05 | CLASS-05 |
|  |  | DEFB | +C9,+1F | LPRINT,1FC9 |
| 1ADC | P-LLIST | DEFB | +05 | CLASS-05 |
|  |  | DEFB | +F5,+17 | LLIST,17F5 |
| 1ADF | P-SAVE | DEFB | +0B | CLASS-0B |
| 1AE0 | P-LOAD | DEFB | +0B | CLASS-0B |
| 1AE1 | P-VERIFY | DEFB | +0B | CLASS-0B |
| 1AE2 | P-MERGE | DEFB | +0B | CLASS-0B |
| 1AE3 | P-BEEP | DEFB | +08 | CLASS-08 |
|  |  | DEFB | +00 | CLASS-00 |
|  |  | DEFB | +F8,+03 | BEEP,03F8 |
| 1AE7 | P-CIRCLE | DEFB | +09 | CLASS-09 |
|  |  | DEFB | +05 | CLASS-05 |
|  |  | DEFB | +20,+23 | CIRCLE,2320 |
| 1AEB | P-INK | DEFB | +07 | CLASS-07 |
| 1AEC | P-PAPER | DEFB | +07 | CLASS-07 |
| 1AED | P-FLASH | DEFB | +07 | CLASS-07 |
| 1AEE | P-BRIGHT | DEFB | +07 | CLASS-07 |


| 1AEF | P-INVERSE | DEFB | +07 | CLASS-07 |
| :---: | :---: | :---: | :---: | :---: |
| 1AF0 | P-OVER | DEFB | +07 | CLASS-07 |
| 1AF1 | P-OUT | DEFB | +08 | CLASS-08 |
|  |  | DEFB | +00 | CLASS-00 |
|  |  | DEFB | +7A,+1E | OUT,1E7A |
| 1AF5 | P-BORDER | DEFB | +06 | CLASS-06 |
|  |  | DEFB | +00 | CLASS-00 |
|  |  | DEFB | +94,+22 | BORDER,2294 |
| 1AF9 | P-DEF-FN | DEFB | +05 | CLASS-05 |
|  |  | DEFB | +60,+1F | DEF-FN,1F60 |
| 1AFC | P-OPEN | DEFB | +06 | CLASS-06 |
|  |  | DEFB | $+2 \mathrm{C}$ | ',' |
|  |  | DEFB | $+0 \mathrm{~A}$ | CLASS-0A |
|  |  | DEFB | +00 | CLASS-00 |
|  |  | DEFB | +36,+17 | OPEN,1736 |
| 1B02 | P-CLOSE | DEFB | +06 | CLASS-06 |
|  |  | DEFB | +00 | CLASS-00 |
|  |  | DEFB | +E5,+16 | CLOSE,16E5 |
| $1 \mathrm{B06}$ | P-FORMAT | DEFB | $+0 \mathrm{~A}$ | CLASS-0A |
|  |  | DEFB | +00 | CLASS-00 |
|  |  | DEFB | +93,+17 | CAT-ETC,1793 |
| 1B0A | P-MOVE | DEFB | $+0 \mathrm{~A}$ | CLASS-0A |
|  |  | DEFB | $+2 \mathrm{C}$ | ',' |
|  |  | DEFB | $+0 \mathrm{~A}$ | CLASS-0A |
|  |  | DEFB | +00 | CLASS-00 |
|  |  | DEFB | +93,+17 | CAT-ETC,1793 |
| 1 B 10 | P-ERASE | DEFB | +0A | CLASS-0A |
|  |  | DEFB | +00 | CLASS-00 |
|  |  | DEFB | +93,+17 | CAT-ETC,1793 |
| 1B14 | P-CAT | DEFB | +00 | CLASS-00 |
|  |  | DEFB | +93,+17 | CAT-ETC,1793 |

Note: The requirements for the different command classes are as follows:
CLASS-00 - No further operands.
CLASS-01 - Used in LET. A variable is required.
CLASS-02 - Used in LET. An expression, numeric or string, must follow.
CLASS-03 - A numeric expression may follow. Zero to be used in case of default.
CLASS-04 - A single character variable must follow.
CLASS-05 - A set of items may be given.
CLASS-06 - A numeric expression must follow.
CLASS-07 - Handles colour items.
CLASS-08 - Two numeric expressions, separated by a comma, must follow.
CLASS-09 - As for CLASS-08 but colour items may precede the expressions.
CLASS-0A - A string expression must follow.
CLASS-0B - Handles cassette routines.

## THE 'MAIN PARSER' OF THE BASIC INTERPRETER

The parsing routine of the BASIC interpreter is entered at LINE-SCAN when syntax is being checked, and at LINE-RUN when a BASIC program of one or more statements is to be executed.

Each statement is considered in turn and the system variable $\mathrm{CH}-\mathrm{ADD}$ is used to point to each code of the statement as it occurs in the program area or the editing area.

| $1 B 17$ | LINE-SCAN | RES | 7,(FLAGS) <br> CALL |
| :--- | :--- | :--- | :--- |
|  | 19FB,E-LINE-NO | Signal 'syntax checking'. <br> CH-ADD is made to point to the <br> first code after any line number. |  |
|  | LD | A | (SUBPPC),A |

## THE STATEMENT LOOP.

Each statement is considered in turn until the end of the line is reached.

| $1 \mathrm{B28}$ | STMT-LOOP | RST | 0020,NEXT-CHAR | Advance CH-ADD along the line. |
| :---: | :---: | :---: | :---: | :---: |
| 1B29 | STMT-L-1 | CALL | 16BF,SET-WORK | The work space is cleared. |
|  |  | INC | (SUBPPC) | Increase SUBPPC on each |
|  |  | JP | M,1C8A,REPORT-C | But only '127' statements are allowed in a single line. |
|  |  | RST | 0018,GET-CHAR | Fetch a character. |
|  |  | LD | B,+00 | Clear the register for later. |
|  |  | CP | +0D | Is the character a 'carriage |
|  |  | JR | Z,1BB3,LINE-END | return'; jump if it is. |
|  |  | CP | +3A | Go around the loop again if |
|  |  | JR | Z,1B28,STMT-LOOP | it is a ' $:$ '. |

A statement has been identified so, first, its initial command is considered.

| LD | HL,+1B76 | Pre-load the machine stack |
| :---: | :---: | :---: |
| PUSH | HL | with the return address - STMT-RET. |
| LD | C,A | Save the command temporarily |
| RST | 0020,NEXT-CHAR | in the C register whilst |
| LD | A, C | CH-ADD is advanced again. |
| SUB | +CE | Reduce the command's code by $+C E$; giving the range +00 to <br> +31 for the fifty commands. |
| JP | C,1C8A,REPORT-C | Give the appropriate error if not a command code. |
| LD | C, A | Move the command code to the $B C$ register pair ( $B$ holds +00 ). |
| LD | HL, +1A48 | The base address of the syntax offset table. |
| ADD | HL, BC | The required offset is passed to |
| LD | C,(HL) | the C register and used to |
| ADD | HL, BC | compute the base address for the command's entries in the parameter table. |
| JR | 1B55,GET-PARAM | Jump forward into the scanning loop with this address. |

Each of the command class routines applicable to the present command are executed in turn. Any required separators are also considered.

| 1B52 | SCAN-LOOP | LD | HL,(T-ADDR) | The temporary pointer to the entries in the parameter table. |
| :---: | :---: | :---: | :---: | :---: |
| 1B55 | GET-PARAM | LD | A, (HL) | Fetch each entry in turn. |
|  |  | INC | HL | Update the pointer to the |
|  |  | LD | (T-ADDR), HL | entries for the next pass. |
|  |  | LD | $B C,+1 \mathrm{~B} 52$ | Pre-load the machine stack |
|  |  | PUSH | BC | with the return address -SCAN-LOOP. |
|  |  | LD | C, A | Copy the entry to the C register for later. |
|  |  | CP | +20 | Jump forward if the entry is |
|  |  | JR | NC,1B6F,SEPARATOR | a 'separator'. |
|  |  | LD | HL, +1C01 | The base address of the 'command class' table. |
|  |  | LD | B, +00 | Clear the B register and |
|  |  | ADD | HL, BC | index into the table. |
|  |  | LD | C, (HL) | Fetch the offset and compute |
|  |  | ADD | HL, BC | the starting address of the |
|  |  | PUSH | HL | Push the address on to the machine stack. |


| RST | 0018,GET-CHAR | Before making an indirect <br> DEC |
| :--- | :--- | :--- |
| jump to the command class |  |  |
| RET |  | routine pass the command code <br> to the A register and set the B <br> register to + FF. |

## THE 'SEPARATOR' SUBROUTINE

The report - 'Nonsense in BASIC is given if the required separator is not present. But note that when syntax is being checked the actual report does not appear on the screen - only the 'error marker'.

| 1B6F SEPARATOR RST | RP <br> CP | The current character is <br> fetched and compared to the |
| :--- | :--- | :--- |
| entry in the parameter table. |  |  |

## THE 'STMT-RET' SUBROUTINE

After the correct interpretation of a statement a return is made to this entry point.

| 1B76 STMT-RET | CALL | 1F54,BREAK-KEY |
| :--- | :--- | :--- | | The BREAK key is tested after |
| :--- |
| every statement. |

Report L - 'BREAK into program'

| 1B7B REPORT-L | RST <br> DEFB | 0008,ERROR-1 <br> +14 | Call the error handling <br> routine. |
| :--- | :--- | :--- | :--- |

Continue here as the BREAK key was not pressed.

| 1B7D | STMT-R-1 | BIT | 7,(NSPPC) |
| :--- | :--- | :--- | :--- |
|  | JR | NZ,1BF4,STMT-NEXT | Jump forward if there is not |
|  | LD 'jump' to be made. |  |  |
|  | BIT | HL,(NEWPPC) | 7,H |
|  | JR | Z,1B9E,LINE-NEW | and jump forward unless dealing |
|  |  |  | with a further statement in the |
|  |  |  | editing area. |

## THE 'LINE-RUN' ENTRY POINT

This entry point is used wherever a line in the editing area is to be 'run'. In such a case the syntax/run flag (bit 7 of FLAGS) will be set.
The entry point is also used in the syntax checking of a line in the editing area that has more than one statement (bit 7 of FLAGS will be reset).

1B8A LINE-RUN

| LD | HL,+FFFE |
| :--- | :--- |
| LD | (PPC),HL |
| LD | HL,(WORKSP) |
| DEC | HL |
| LD | DE,(E-LINE) |
| DEC | DE |
| LD | A,(NSPPC) |
|  |  |
| JR | 1BD1,NEXT-LINE |

> A line in the editing area is considered as line '-2'. Make HL point to the end marker of the editing area and DE to the location before the start of that area. Fetch the number of the next statement to be handled before jumping forward.

## THE 'LINE-NEW' SUBROUTINE

There has been a jump in the program and the starting address of the new line has to be found.


|  |  | JR | NZ,1BEC,REPORT-N | check the validity of the statement number - must be zero. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | LD | B,A | Also check that the 'first |
|  |  | LD | A,(HL) | line after' is not after the |
|  |  | AND | +C0 | actual 'end of program'. |
|  |  | LD | A,B |  |
|  |  | JR | Z,1BBF,LINE-USE | Jump forward with valid addresses; otherwise signal the error 'OK'. |
| Report 0 - 'OK' |  |  |  |  |
| $1 \mathrm{BB0}$ | REPORT-0 | RST | 0008,ERROR-1 | Use the error handling |
|  |  | DEFB | +FF | routine. |

Note: Obviously not an error in the normal sense - but rather a jump past the program.

## THE 'REM' COMMAND ROUTINE

The return address to STMT-RET is dropped which has the effect of forcing the rest of the line to be ignored.
1BB2
REM
POP
BC
Drop the address - STMT-RET.

## THE 'LINE-END' ROUTINE

If checking syntax a simple return is made but when 'running' the address held by NXTLIN has to be checked before it can be used.

| 1BB3 LINE-END | CALL | 2530,SYNTAX-Z | Return if syntax is being <br> checked; otherwise fetch |
| :--- | :--- | :--- | :--- |
|  | RET | Z | the address in NXTLIN. |
|  | LD | HL,(NXTLIN) | Return also if the address is |
|  | LD | A,+C0 | after the end of the program |

## THE 'LINE-USE' ROUTINE

This short routine has three functions; i. Change statement zero to statement '1'; ii. Find the number of the new line and enter it into PPC; \& iii. Form the address of the start of the line after.


## THE 'NEXT-LINE' ROUTINE

On entry the HL register pair points to the location after the end of the 'next' line to be handled and the DE register pair to the location before the first character of the line. This applies to lines in the program area and also to a line in the editing area - where the next line will be the same line again whilst there are still statements to be interpreted.

| 1BD1 | (NXTLIN),HL | SINE | Set NXTLIN for use once the <br> current line has been completed. |
| :--- | :--- | :--- | :--- |
|  | EX | DE,HL | As usual CH-ADD points to the |


| LD | (CH-ADD),HL | location before the first <br> character to be considered. |
| :--- | :--- | :--- |
| LD | D,A | The statement number is <br> fetched. <br> The E register is cleared in case |
| LD | E,+00 | EACH-STMT is used. |
| LD | (NSPPC),+FF | Signal 'To jump'. <br> The statement number minus <br> one goes into SUBPPC. |
| DEC | D | (SUBPPC),D |
| LD | Z,1B28,STMT-LOOP | A first statement can now be <br> considered. |
| INC | D | However for later statements <br> the 'starting address' has to be <br> found. |
| CALL | 198B,EACH-STMT | Jump forward unless the state- <br> ment does not exist. |
| JR | Z,1BF4,STMT-NEXT |  |

Report N - 'Statement lost'

| 1BEC REPORT-N | RST <br> DEFB | 0008,ERROR-1 <br> +16 | Call the error handling <br> routine. |
| :--- | :--- | :--- | :--- |

## THE 'CHECK-END' SUBROUTINE

This is an important routine and is called from many places in the monitor program when the syntax of the edit-line is being checked. The purpose of the routine is to give an error report if the end of a statement has not been reached and to move on to the next statement if the syntax is correct.

1BEE CHECK-END CALL 2530,SYNTAX-Z $\begin{array}{lll} & \text { NZ } & \text { Do not proceed unless } \\ & & \text { checking syntax }\end{array}$

| RET | NZ |
| :--- | :--- |
| POP | BC |
| POP | BC | checking syntax.

Drop the addresses of SCAN-LOOP \& STMT-RET before continuing into STMT-NEXT.

## THE 'STMT-NEXT' ROUTINE

If the present character is a 'carriage return' then the 'next statement' is on the 'next line'; if ' : ' it is on the same line; but if any other character is found then there is an error in syntax.

| 1BF4 | STMT-NEXT | RST | 0018,GET-CHAR |
| :--- | :--- | :--- | :--- |
| CP | +0D | Fetch the present character. |  |
|  | JR | Z,1BB3,LINE-END | Consider the 'next line' if |
| CP | $+3 A$ | it is a 'carriage return'. |  |
|  | JP | Z,1B28,STMT-LOOP | Consider the 'next statement' |
|  | JP it is a' $:$ '. |  |  |

THE 'COMMAND CLASS' TABLE

| address | offset | class number | address | offset | class number |
| :--- | :--- | :--- | :--- | :--- | :--- |
| 1C01 | 0F | CLASS-00-1C10 | 1C07 | 7B | CLASS-06,1C82 |
| 1C02 | 1D | CLASS-01,1C1F | 1 C08 | 8 E | CLASS-07,1C96 |
| 1C03 | 4B | CLASS-02,1C4E | 1 C09 | 71 | CLASS-08,1C7A |
| 1C04 | 09 | CLASS-03,1C0D | 1C0A | B4 | CLASS-09,1CBE |
| 1C05 | 67 | CLASS-04,1C6C | 1C0B | 81 | CLASS-0A,1C8C |
| 1C06 | $0 B$ | CLASS-05,1C11 | 1C0C | CF | CLASS-0B,1CDB |

THE 'COMMAND CLASSES - 00, 03 \& 05'
The commands of class-03 may, or may not, be followed by a number. e.g. RUN \& RUN 200.

1C0D CLASS-03 CALL 1CDE,FETCH-NUM | A number is fetched but zero |
| :--- |
| is used in cases of default. |

The commands of class-00 must not have any operands. e.g. COPY \& CONTINUE.
1 C10 CLASS-00 A Set the zero flag for later.

The commands of class-05 may be followed by a set of items. e.g. PRINT \& PRINT "222".


THE 'JUMP-C-R' ROUTINE
After the command class entries and the separator entries in the parameter table have been considered the jump to the appropriate command routine is made.

| 1 C 16 JUMP-C-R | LD | HL,(T-ADDR) | Fetch the pointer to the |
| :--- | :--- | :--- | :--- |
|  | LD | $\mathrm{C},(\mathrm{HL})$ | entries in the parameter table |
|  | INC | HL | and fetch the address of the |

THE 'COMMAND CLASSES -01, 02 \& 04'
These three command classes are used by the variable handling commands - LET, FOR \& NEXT and indirectly by READ \& INPUT. Command class 01 is concerned with the identification of the variable in a LET, READ or INPUT statement.

1C1F CLASS-01
CALL 28B2,LOOK-VARS
Look in the variables area to determine whether or not the variable has been used already.

## THE 'VARIABLE IN ASSIGNMENT' SUBROUTINE

This subroutine develops the appropriate values for the system variables DEST \& STRLEN.

| 1 C 22 | VAR-A-1 | LD | (FLAGX),+00 | Initialise FLAGX to +00 . |
| :---: | :---: | :---: | :---: | :---: |
|  |  | JR | NC,1C30,VAR-A-2 | Jump forward if the variable has been used before. |
|  |  | SET | 1,(FLAGX) | Signal 'a new variable'. |
|  |  | JR | NZ,1C46,VAR-A-3 | Give an error if trying to use an 'undimensioned array'. |

Report 2 - Variable not found
1C2E REPORT-2 RST 0008,ERROR-1 Call the error handling

DEFB
$+01$ routine.

Continue with the handling of existing variables.

1C30 VAR-A-2 \begin{tabular}{lll}

CALL \& Z,2996,STK-VARS \& | The parameters of simple string |
| :--- |
| variables and all array variables |
| are passed to the calculator | <br>

stack. (STK-VARS will 'slice' a <br>
string if required.)
\end{tabular}

| OR | (HL) | Bit 0 is set only when handling <br> complete simple strings' thereby <br> signalling 'old copy to be |
| :--- | :--- | :--- |
| (HL),A | deleted'. |  |

The pathways now come together to set STRLEN \& DEST as required. For all numeric variables and 'new' string \& string array variables STRLEN-lo holds the 'letter' of the variable's name. But for 'old' string \& string array variables whether 'sliced' or complete it holds the 'length' in 'assignment'.
1 C 46
VAR-A-3
LD
(STRLEN),BC
Set STRLEN as required.

DEST holds the address for the 'destination of an 'old' variable but in effect the 'source' for a 'new' variable.

| LD | (DEST),HL | Set DEST as required and |
| :--- | :--- | :--- |
| RET |  |  |

Command class 02 is concerned with the actual calculation of the value to be assigned in a LET statement.

| 1C4E | PLASS-02 | POP |
| :--- | :--- | :--- | | The address - SCAN-LOOP is |
| :--- |
| dropped. |

## THE 'FETCH A VALUE' SUBROUTINE

This subroutine is used by LET, READ \& INPUT statements to first evaluate and then assign values to the previously designated variable.

The entry point VAL-FET-1 is used by LET \& READ and considers FLAGS whereas the entry point VAL-FET-2 is used by INPUT and considers FLAGX

| 1 C 56 | VAL-FET-1 | LD | A,(FLAGS) | Use FLAGS. |
| :---: | :---: | :---: | :---: | :---: |
| 1 C 59 | VAL-FET-2 | PUSH | AF | Save FLAGS or FLAGX. |
|  |  | CALL | 24FB,SCANNING | Evaluate the next expression. |
|  |  | POP | AF | Fetch the old FLAGS or FLAGX. |
|  |  | LD | D,(FLAGS) | Fetch the new FLAGS. |
|  |  | XOR | D | The nature - numeric or string |
|  |  | AND | +40 | of the variable and the expression must match. |
|  |  | JR | NZ,1C8A,REPORT-C | Give report C if they do not. |
|  |  | BIT | 7,D | Jump forward to make the |
|  |  | JP | NZ,2AFF,LET | actual assignment unless |
|  |  |  |  | checking syntax when simply |
|  |  | RET |  | return. |

## THE 'COMMAND CLASS 04' ROUTINE

The command class 04 entry point is used by FOR \& NEXT statements.

| 1C6C | CLASS-04 | CALL | 28B2,LOOK-VARS |
| :--- | :--- | :--- | :--- | | Look in the variables area for |
| :--- |
| the variable being used. |

## THE 'EXPECT NUMERIC/STRING EXPRESSIONS' SUBROUTINE

There is a series of short subroutines that are used to fetch the result of evaluating the next expression. The result from a single expression is returned as a 'last value' on the calculator stack

The entry point NEXT-2NUM is used when CH-ADD needs updating to point to the start of the first expression.
1 C79 NEXT-2NUM RST 0020,NEXT-CHAR Advance CH-ADD.

The entry point EXPT-2NUM (EQU. CLASS-08) allows for two numeric expressions, separated by a comma, to be evaluated.

| 1C7A | EXPT-2NUM <br> (CLASS-08) | CALL | 1C82,EXPT-1NUM |
| :--- | :--- | :--- | :--- | | Evaluate each expression in |
| :--- |
|  |
|  |
|  |
|  |
|  |
|  |
|  |
|  |
|  |
|  |
|  |
|  |
|  |
|  |
|  |
| RST |

The entry point EXPT-1NUM (EQU. CLASS-06) allows for a single numeric expression to be evaluated.

| 1C82 | EXPT-1NUM <br> (CLASS-06) | CALL | 24FB,SCANNING |
| :--- | :--- | :--- | :--- |$\quad$| Evaluate the next expression. |
| :--- |
|  |
|  |
|  |
|  |
|  |
|  |
|  |
|  |
|  |
|  |

Report C - Nonsense in BASIC

| $1 C 8 A$ |  |  |  |
| :--- | :--- | :--- | :--- |
| REPORT-C | RST <br> DEFB | 0008,ERROR-1 <br> $+0 B$ | Call the error handling <br> routine. |

The entry point EXPT-EXP (EQU. CLASS-OA) allows for a single string expression to be evaluated.

| EXPT-EXP | CALL | 24FB,SCANNING | Evaluate the next expression. |  |
| :--- | :--- | :--- | :--- | :--- |
|  |  |  |  |  |
|  |  | BIT | 6,(FLAGS) | This time return if the result |
|  | RET | Z | indicates a string; otherwise |  |
|  |  | 1C8A,REPORT-C | give an error report. |  |

THE 'SET PERMANENT COLOURS' SUBROUTINE (EQU. CLASS-07)
This subroutine allows for the current temporary colours to be made permanent. As command class 07 it is in effect the command routine for the six colour item commands.

| 1 C 96 | PERMS <br> (CLASS-07) | BIT | 7,(FLAGS) | The syntax/run flag is read. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | RES | 0,(TV-FLAG) | Signal 'main screen'. |
|  |  | CALL | NZ,0D4D,TEMPS | Only during a 'run' call TEMPS to ensure the temporary colours are the main screen colours. |
|  |  | POP | AF | Drop the return address -SCAN-LOOP. |
|  |  | LD | A,(T-ADDR) | Fetch the low byte of T-ADDR and subtract +13 to give the |
|  |  | SUB | +13 | range +D 9 to +DE which are the token codes for INK to OVER. |
|  |  | CALL | 21FC,CO-TEMP-4 | Jump forward to change the temporary colours as directed by the BASIC statement. |
|  |  | CALL | 1BEE,CHECK-END | Move on to the next statement if checking syntax. |
|  |  | LD | HL,(ATTR-T) | Now the temporary colour |
|  |  | LD | (ATTR-P), HL | values are made permanent (both ATTR-P \& MASK-P). |
|  |  | LD | HL, +5C91 | This is P-FLAG; and that too |
|  |  | LD | A,(HL) | has to be considered. |

The following instructions cleverly copy the even bits of the supplied byte to the odd bits. In effect making the permanent bits the same as the temporary ones.

| RLCA |  | Move the mask leftwards. <br> XOR |
| :--- | :--- | :--- |
| AND | (HL) | Impress onto the mask |
| XOR | (HL) | only the even bits of the |
| LD | $(H L), A$ | other byte. |
| RET |  | Restore the result. |

THE 'COMMAND CLASS 09' ROUTINE
This routine is used by PLOT, DRAW \& CIRCLE statements in order to specify the default conditions of 'FLASH 8; BRIGHT 8; PAPER 8;' that are set up before any embedded colour items are considered.

| 1CBE | CLASS-09 | CALL | 2530,SYNTAX-Z | Jump forward if |
| :---: | :---: | :---: | :---: | :---: |
|  |  | JR | Z,1CD6,CL-09-1 | checking syntax. |
|  |  | RES | 0,(TV-FLAG) | Signal 'main screen'. |
|  |  | CALL | 0D4D,TEMPS | Set the temporary colours for the main screen. |
|  |  | LD | HL, +5C90 | This is MASK-T. |
|  |  | LD | A,(HL) | Fetch its present value but |
|  |  | OR | +F8 | keep only its INK part 'unmasked'. |
|  |  | LD | (HL), A | Restore the value which now indicates 'FLASH 8; BRIGHT 8; PAPER 8;'. |
|  |  | RES | 6,(P-FLAG) | Also ensure NOT 'PAPER 9'. |
|  |  | RST | 0018,GET-CHAR | Fetch the present character before continuing to deal with embedded colour items. |
| $1 \mathrm{CD6}$ | CL-09-1 | CALL | 21E2,CO-TEMP | Deal with the locally dominant colour items. |
|  |  | JR | 1C7A,EXPT-2NUM | Now get the first two operands for PLOT, DRAW or CIRCLE. |

THE 'COMMAND CLASS OB' ROUTINE
This routine is used by SAVE, LOAD, VERIFY \& MERGE statements.
1CDB CLASS-0B JP 0605,SAVE-ETC Jump to the cassette handling routine.

THE 'FETCH A NUMBER' SUBROUTINE
This subroutine leads to a following numeric expression being evaluated but zero being used instead if there is no expression.

| 1CDE | FETCH-NUM | CP | +OD |
| :--- | :--- | :--- | :--- |
|  | JR | Z,1CE6,USE-ZERO | Jump forward if at the end |
|  | CP | of a line. | But jump to EXPT-1NUM unless |
|  |  | JR | NZ,1C82,EXPT-1NUM |

The calculator is now used to add the value zero to the calculator stack.

| 1CE6 | USE-ZERO | CALL | 2530,SYNTAX-Z |
| :--- | :--- | :--- | :--- |
|  | RET | Z | Do not perform the operation |
| RST | 0028, FP-CALC | if syntax is being checked. |  |
|  | DEFB | +A0,stk-zero | Use the calculator. |
|  | DEFB | +38,end-calc | The 'last value' is now zero. |
|  | RET |  | Return with zero added to the <br> stack. |

## THE COMMAND ROUTINES

The section of the 16K monitor program from 1CEE to 23FA contains most of the command routines of the BASIC interpreter.

## THE 'STOP' COMMAND ROUTINE

The command routine for STOP contains only a call to the error handling routine.

| 1CEE | STOP | RST | $0008, E R R O R-1$ | Call the error handling |
| :--- | :--- | :--- | :--- | :--- |
|  | (REPORT-9) | DEFB | +08 | routine. |

THE 'IF' COMMAND ROUTINE
On entry the value of the expression between the IF and the THEN is the 'last value' on the calculator stack. If this is logically true then the next statement is considered; otherwise the line is considered to have been finished.
1 CF0
IF

| POP | BC |
| :--- | :--- |
| CALL | 2530, SYNTAX-Z |
| JR | Z,1D00,IF-1 |

Drop the return address

- STMT-RET.
Jump forward if checking
syntax.

Now use the calculator to 'delete' the last value on the calculator stack but leave the DE register pair addressing the first byte of the value.

|  | RST | 0028,FP-CALC |
| :--- | :--- | :--- |
| DEFB | +02,delete | Use the calculator. |
| DEFB | +38,end-calc | The present 'last value' is |
| EX | DE,HL | deleted. |

## THE 'FOR' COMMAND ROUTINE

This command routine is entered with the VALUE and the LIMIT of the FOR statement already on the top of the calculator stack.

| 1 D03 FOR | CP | +CD | Jump forward unless a 'STEP' |
| :--- | :--- | :--- | :--- |
|  | JR | NZ,1D10,F-USE-1 | is given. |
|  | RST | 0020,NEXT-CHAR | Advance CH-ADD and fetch the |
|  | CALL | 1C82,EXPT-1NUM | value of the STEP. |
|  | CALL | 1BEE,CHECK-END | Move on to the next statement |
|  | JR | 1D16,F-REORDER | if checking syntax; otherwise <br>  <br>  |

There has not been a STEP supplied so the value ' 1 ' is to be used.
$\left.\begin{array}{llll}1 \text { F10 } & \text { F-USE-1 } & \text { CALL } & \text { 1BEE,CHECK-END }\end{array} \begin{array}{l}\text { Move on to the next statement } \\ \text { if checking syntax; otherwise }\end{array}\right]$ use the calculator to place a '1'

The three values on the calculator stack are the VALUE (v), the LIMIT (I) and the STEP ( s ). These values now have to be manipulated.

| $1 D 16$ | F-REORDER | RST | 0028,FP-CALC |
| :--- | :--- | :--- | :--- |$\quad \mathrm{v}, \mathrm{l}, \mathrm{s}$, (mem-0 = s)

A FOR control variable is now established and treated as a temporary calculator memory area.

| CALL | 2AFF,LET | The variable is found, or created <br> if needed (v is used). |
| :--- | :--- | :--- |
| LD | $(M E M), H L$ | Make it a 'memory area'. |

The variable that has been found may be a simple numeric variable using only six locations in which case it will need extending.

| DEC | HL |
| :--- | :--- |
| LD | A,(HL) |
| SET | $7,(H L)$ |
| LD | BC,+0006 |
| ADD | HL,BC |
| RLCA | C,1D34,F-L\&S |
| JR | C,+0D |
| LD | $1655, M A K E-R O O M$ |
| CALL | HL |
| INC | HL |

Fetch the variable's single
character name.
Ensure bit 7 of the name is set
It will have six locations at least.
Make HL point after them.
Rotate the name and jump if
it was already a FOR variable.
Otherwise create thirteen
more locations.
Again make HL point to the LIMIT position.

The initial values for the LIMIT and the STEP are now added.
1D34 F-L\&S

| PUSH | HL |
| :--- | :--- |
| RST | 0028, FP-CALC |
| DEFB | +02, delete |
| DEFB | +02, delete |
| DEFB | +38, end-calc |
| POP | HL |
| EX | $\mathrm{DE}, \mathrm{HL}$ |
| LD | $\mathrm{C},+0 \mathrm{~A}$ |
| LDIR |  |

The pointer is saved.
I, s
-
DE still points to 'I'.
The pointer is restored and both pointers exchanged.
The ten bytes of the LIMIT and the STEP are moved.

The looping line number and statement number are now entered.

| LD | HL,(PPC) | The current line number. |
| :--- | :--- | :--- |
| EX | DE,HL | Exchange the registers before |
| LD | (HL),E | adding the line number to the |
| INC | HL | FOR control variable. |
| LD | (HL),D |  |
| LD | D,(SUBPPC) | The looping statement is |
| INC | D | always the next statement - |
| INC | $H L$ | whether it exists or not. |
| LD | (HL),D |  |

The NEXT-LOOP subroutine is called to test the possibility of a 'pass' and a return is made if one is possible; otherwise the statement after for FOR - NEXT loop has to be identified.

| CALL | 1DDA,NEXT-LOOP | Is a 'pass' possible? |
| :--- | :--- | :--- |
| RET | NC | Return now if it is. |
| LD | B,(STRLEN-lo) | Fetch the variable's name. |
| LD | HL,(PPC) | Copy the present line number |
| LD | (NEWPPC),HL | to NEWPPC. |
| LD | A,(SUBPPC) | Fetch the current statement <br> number and two's complement it. <br> NEG |
| LD | D,A | Transfer the result to the D <br> register. |
| LD | HL,(CH-ADD) | Fetch the current value of <br> CH-ADD. |
| LD | E,+F3 | The search will be for 'NEXT'. |

Now a search is made in the program area, from the present point onwards, for the first occurrence of NEXT followed by the correct variable.

| F-LOOP | PUSH | BC |
| :--- | :--- | :--- |
| LD | BC,(NXTLIN) | Save the variable's name. <br> Fetch the current value of |
|  | CALL | 1D86,LOOK-PROG | | NXTLIN. |
| :--- |
|  |
|  |
|  |
|  |
|  |
| LD |
| POP |
| JR |


| RST | 0020,NEXT-CHAR | Advance past the NEXT that <br> was found. |
| :--- | :--- | :--- |
| OR | Allow for upper and lower |  |
| CP | B | case letters before the new <br> variable name is tested. |
| JR | Z,1D7C,F-FOUND | Jump forward it it matches. <br> Advance CH-ADD again and <br> RST |
| JR | 0020,NEXT-CHAR | jump back if not the correct <br> variable. |

NEWPPC holds the line number of the line in which the correct NEXT was found. Now the statement number has to be found and stored in NSPPC.

| 1D7C | F-FOUND | RST | 0020,NEXT-CHAR | Advance CH-ADD. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | LD | A,+01 | The statement counter in the |
|  |  | SUB | D | D register counted statements |
|  |  |  |  | back from zero so it has to be subtracted from '1'. |
|  |  | LD | (NSPPC),A | The result is stored. |
|  |  | RET |  | Now return - to STMT-RET. |

REPORT I - FOR without NEXT

| 1 R84 | REPORT-I | RST <br> DEFB | $0008, E R R O R-1$ <br> +11 |
| :--- | :--- | :--- | :--- | | Call the error handling |
| :--- |
| routine. |

## THE 'LOOK-PROG' SUBROUTINE

This subroutine is used to find occurrences of either DATA, DEF FN or NEXT. On entry the appropriate token code is in the E register and the HL register pair points to the start of the search area.

| 1 LOO |  |  |
| :--- | :--- | :--- |
|  | A,(HL) | Fetch the present character. |
|  | CP | +3A | | Jump forward if it is a ' $: '$ |
| :--- |
| JR |

Now a loop is entered to examine each further line in the program.


## THE 'NEXT' COMMAND ROUTINE

The 'variable in assignment' has already been determined (see CLASS-04,1C6C); and it remains to change the VALUE as required.

| 1DAB | NEXT | BIT | (FLAGX) |
| :--- | :--- | :--- | :--- |
|  | JP | NZ,1C2E,REPORT-2 | Jump to give the error report |
|  | LD the variable was not found. |  |  |

Next the variable's VALUE and STEP are manipulated by the calculator.

| INC | HL <br> LD | Step past the name. <br> Make the variable a |
| :--- | :--- | :--- |
| temporary 'memory area'. |  |  |

The result of adding the VALUE and the STEP is now tested against the LIMIT by calling NEXT-LOOP.

| CALL | 1DDA,NEXT-LOOP | Test the new VALUE against <br> the LIMIT |
| :--- | :--- | :--- |
| RET | C | Return now if the FOR-NEXT <br> loop has been completed. |

Otherwise collect the 'looping' line number and statement.

| LD | HL,(MEM) | Find the address of the |
| :--- | :--- | :--- |
| LD | $\mathrm{DE},+000 \mathrm{~F}$ | low byte of the looping |
| ADD | $\mathrm{HL}, \mathrm{DE}$ | line number. |
| LD | $\mathrm{E},(\mathrm{HL})$ | Now fetch this line number. |
| INC | HL |  |
| LD | $\mathrm{D},(\mathrm{HL})$ |  |
| INC | HL | Followed by the statement <br> LD |
|  | $\mathrm{H},(\mathrm{HL})$ | number. |
| EX | $\mathrm{DE}, \mathrm{HL}$ | Exchange the numbers before <br> jumping forward to treat them <br> JP |
|  | $1 \mathrm{E} 73, \mathrm{GO}-\mathrm{TO}-2$ | as the destination line of a <br> GO TO command. |

Report 1 - NEXT without FOR
1 RD8 REPORT-1 RST 0008,ERROR-1 Call the error handling
DEFB +00 routine.

THE 'NEXT-LOOP SUBROUTINE
This subroutine is used to determine whether the LIMIT has been exceeded by the present VALUE. Note has to be taken of the sign of the STEP.

The subroutine returns the carry flag set if the LIMIT is exceeded.

| 1DDA | NEXT-LOOP | RST | 0028,FP-CALC | - |
| :---: | :---: | :---: | :---: | :---: |
|  |  | DEFB | +E1,get-mem-1 | I |
|  |  | DEFB | +E0,get-mem-0 | I, v |
|  |  | DEFB | +E2,get-mem-2 | $\mathrm{I}, \mathrm{v}, \mathrm{s}$ |
|  |  | DEFB | +36, less-0 | I, v,(1/0) |
|  |  | DEFB | +00,jump-true | I, v,(1/0) |
|  |  | DEFB | +02,to NEXT-1 | I, v,(1/0) |
|  |  | DEFB | +01,exchange | v , I |
| 1 DE2 | NEXT-1 | DEFB | +03, subtract | v-I or I-v |
|  |  | DEFB | +37, greater-0 | (1/0) |
|  |  | DEFB | +00,jump-true | (1/0) |


| DEFB | +04, to NEXT-2 | - |
| :--- | :--- | :--- |
| DEFB | +38, end-calc | - |
| AND | A | Clear the carry flag and |
| RET |  | return - loop is possible. |

However if the loop is impossible the carry flag has to be set.

| 1 NE9 | DEXT-2 | DEFB |
| :--- | :--- | :--- |
|  | SCF | +38,end-calc |
|  | RET |  |
|  |  | Set the carry flag and |
|  | return. |  |

## THE 'READ' COMMAND ROUTINE

The READ command allows for the reading of a DATA list and has an effect similar to a series of LET statements.
Each assignment within a single READ statement is dealt with in turn. The system variable X-PTR is used as a storage location for the pointer to the READ statement whilst $\mathrm{CH}-\mathrm{ADD}$ is used to step along the DATA list.

| 1DEC | READ-3 | RST | 0020,NEXT-CHAR | Come here on each pass, after the first, to move along the READ statement. |
| :---: | :---: | :---: | :---: | :---: |
| 1DED | READ | CALL | 1C1F,CLASS-01 | Consider whether the variable has been used before; find the existing entry if it has. |
|  |  | CALL | 2530,SYNTAX-Z | Jump forward if checking |
|  |  | JR | Z,1E1E,READ-2 | syntax. |
|  |  | RST | 0018,GET-CHAR | Save the current pointer |
|  |  | LD | (X-PTR), HL | CH-ADD in X-PTR. |
|  |  | LD | HL,(DATADD) | Fetch the current DATA list |
|  |  | LD | A,(HL) | pointer and jump forward |
|  |  | CP | +2C | unless a new DATA statement |
|  |  | JR | Z,1E0A,READ-1 | has to be found. |
|  |  | LD | E,+E4 | The search is for 'DATA'. |
|  |  | CALL | 1D86,LOOK-PROG | Jump forward if the search is |
|  |  | JR | NC,1E0A,READ-1 | successful. |

Report E - Out of DATA
$1 E 08$ REPORT-E RST 0008,ERROR-
DEFB +0D

Continue - picking up a value from the DATA list.

| 1E0A | READ-1 | CALL | 0077,TEMP-PTR1 | Advance the pointer along the DATA list and set CH-ADD. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | CALL | 1C56,VAL-FET-1 | Fetch the value and assign it to the variable. |
|  |  | $\begin{aligned} & \text { RST } \\ & \text { LD } \end{aligned}$ | 0018,GET-CHAR (DATADD),HL | Fetch the current value of $\mathrm{CH}-\mathrm{ADD}$ and store it in DATADD. |
|  |  | LD | HL,(X-PTR) | Fetch the pointer to the |
|  |  | LD | (X-PTR-hi),+00 | READ statement and clear X-PTR. |
|  |  | CALL | 0078,TEMP-PTR2 | Make CH-ADD once again point to the READ statement. |
| 1E1E | READ-2 | RST | 0018,GET-CHAR | GET the present character and |
|  |  | CP | +2C | see if it is a ','. |
|  |  | JR | Z,1DEC,READ-3 | If it is then jump back as there are further items; |
|  |  | CALL | 1BEE,CHECK-END | otherwise return either via |
|  |  | RET |  | CHECK-END (if checking |
|  |  |  |  | syntax) or the RET instruction (to STMT-RET). |

## THE 'DATA' COMMAND ROUTINE

During syntax checking a DATA statement is checked to ensure that it contains a series of valid expressions, separated by commas.
But in 'run-time' the statement is passed by.

| $1 E 27$ | DATA | CALL | 2530,SYNTAX-Z |
| :--- | :--- | :--- | :--- |$\quad$ Jump forward unless checking

A loop is now entered to deal with each expression in the DATA statement.


THE 'PASS-BY' SUBROUTINE
On entry the A register will hold either the token 'DATA' or the token 'DEF FN' depending on the type of statement that is being 'passedby'.


## THE 'RESTORE' COMMAND ROUTINE

The operand for a RESTORE command is taken as a line number, zero being used if no operand is given.
The REST-RUN entry point is used by the RUN command routine.

| 1E42 | RESTORE | CALL | 1E99,FIND-INT2 | Compress the operand into the <br> BC register pair. |
| :--- | :--- | :--- | :--- | :--- |
| 1E45 | REST-RUN | LD | H,B | Transfer the result to the |
|  |  | LD | L,C | HL register pair. |

## THE 'RANDOMIZE' COMMAND ROUTINE

Once again the operand is compressed into the BC register pair and transferred to the required system variable. However if the operand is zero the value in FRAMES1 and FRAMES2 is used instead.

| 1E4F | RANDOMIZE CALL | 1E99,FIND-INT2 | Fetch the operand. |  |
| :--- | :--- | :--- | :--- | :--- |
|  |  | LD | A,B | Jump forward unless the |
|  |  | OR | C | value of the operand is |

## THE 'CONTINUE' COMMAND ROUTINE

The required line number and statement number within that line are made the object of a jump.

| 1E5F CONTINUE | LD | HL,(OLDPPC) | The line number. |
| :--- | :--- | :--- | :--- |
|  | LD | D,(OSPPC) | The statement number. |
|  | JR | 1E73,GO-TO-2 | Jump forward. |

THE 'GO TO' COMMAND ROUTINE
The operand of a GO TO ought to be a line number in the range ' 1 ' to ' 9999 ' but the actual test is against an upper value of ' 61439 '
1 E67 GO-TO CALL 1E99,FIND-INT2 Fetch the operand and transfer

| LD | H,B | it to the HL register pair. |
| :--- | :--- | :--- |
| LD | L,C |  |
| LD | D,+00 | Set the statement number to |
|  |  | zero. |
| LD | A,H | Give the error message |
| CP | + FO | - Integer out of range - |
| JR | NC,1E9F,REPORT-B | with lines over '614139' |

The entry point GO-TO-2 is used to determine the line number of the next line to be handled in several instances.

| $1 E 73$ | GO-TO-2 | LD | (NEWPPC), HL |
| :--- | :--- | :--- | :--- |$\quad$| Enter the line number and |
| :--- |
|  |
|  |
|  |
|  |
|  |
|  |
|  |
|  |
|  |
| RET |

the 'OUT' COMMAND ROUTINE
The two parameters for the OUT instruction are fetched from the calculator stack and used as directed.

| 1E7A OUT | CALL | 1E85,TWO-PARAM | The operands are fetched. |
| :--- | :--- | :--- | :--- |
|  | OUT | (C),A | The actual OUT instruction. |

THE 'POKE' COMMAND ROUTINE
In a similar manner the POKE operation is performed.

| 1E80 | POKE | CALL | 1E85,TWO-PARAM | The operands are fetched. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | LD | (BC),A | The actual POKE operation |
|  |  | RET | Return; - to STMT-R |  |

## THE 'TWO-PARAM' SUBROUTINE

The topmost parameter on the calculator stack must be compressible into a single register. It is two's complemented if it is negative. The second parameter must be compressible into a register pair.

| 1E85 | TWO-PARAM | CALL | 2DD5,FP-TO-A | The parameter is fetched. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | JR | C,1E9F,REPORT-B | Give an error if it is too high a number. |
|  |  | JR | Z,1E8E,TWO-P-1 | Jump forward with positive |
|  |  | NEG |  | numbers but two's complement negative numbers. |
| 1E8E | TWO-P-1 | PUSH | AF | Save the first parameter |
|  |  | CALL | 1E99,FIND-INT2 | whilst the second is fetched. |
|  |  | POP | AF | The first parameter is |
|  |  | RET |  | restored before returning. |

## THE 'FIND INTEGERS' SUBROUTINE

The 'last value' on the calculator stack is fetched and compressed into a single register or a register pair by entering at FIND-INT1 AND FIND-INT2 respectively.

| 1E94 | FIND-INT1 | CALL | 2DD5,FP-TO-A | Fetch the 'last value'. |
| :--- | :--- | :--- | :--- | :--- |
|  |  | JR | 1E9C,FIND-I-1 | Jump forward. |
| 1E99 | FIND-INT2 | CALL | 2DA2,FP-TO-BC | Fetch the 'last value'. |
| 1E9C | FIND-I-1 | JR | C,1E9F,REPORT-B | In both cases overflow |


|  |  | RET | Z | indicated by a set carry flag Return with all positive numbers that are in range. |
| :---: | :---: | :---: | :---: | :---: |
| Report B - Integer out of range |  |  |  |  |
| 1E9F | REPORT-B | $\begin{aligned} & \text { RST } \\ & \text { DEFB } \end{aligned}$ | $\begin{aligned} & \text { 0008,ERROR-1 } \\ & +0 \mathrm{~A} \end{aligned}$ | Call the error handling routine. |

## THE 'RUN' COMMAND ROUTINE

The parameter of the RUN command is passed to NEWPPC by calling the GO TO command routine. The operations of 'RESTORE 0 ' and 'CLEAR 0 ' are then performed before a return is made.

| 1EA1 RUN | CALL | 1E67,GO-TO | Set NEWPPC as required. |
| :--- | :--- | :--- | :--- |
|  | LD | BC,+0000 | Now perform a 'RESTORE 0'. |
|  | CALL | 1E45,REST-RUN |  |
|  | JR | 1EAF,CLEAR-1 | Exit via the CLEAR command routine. |

THE 'CLEAR' COMMAND ROUTINE
This routine allows for the variables area to be cleared, the display area cleared and RAMTOP moved. In consequence of the last operation the machine stack is rebuilt thereby having the effect of also clearing the GO SUB stack.

| 1EAC | CLEAR | CALL | 1E99,FIND-INT2 | Fetch the operand - using zero by default. |
| :---: | :---: | :---: | :---: | :---: |
| 1EAF | CLEAR-RUN | LD | A, B | Jump forward if the operand is |
|  |  | OR | C | other than zero. When called |
|  |  | JR | NZ,1EB7,CLEAR-1 | from RUN there is no jump. |
|  |  | LD | BC,(RAMTOP) | If zero use the existing value in RAMTOP. |
| 1EB7 | CLEAR-1 | PUSH | BC | Save the value. |
|  |  | LD | DE,(VARS) | Next reclaim all the bytes |
|  |  | LD | HL,(E-LINE) | of the present variables |
|  |  | DEC | HL | area. |
|  |  | CALL | 19E5,RECLAIM-1 |  |
|  |  | CALL | 0D6B,CLS | Clear the display area. |

The value in the BC register pair which will be used as RAMTOP is tested to ensure it is neither too low nor too high.

| LD | HL,(STKEND) | The current value of STKEND. |
| :--- | :--- | :--- |
| LD | DE,+0032 | is increased by '50' before |
| ADD | HL,DE | being tested. This forms the |
| POP | DE | lower limit. |
| SBC | HL,DE |  |
| JR | NC,1EDA,REPORT-M | RAMTOP will be too low. |
| LD | HL,(P-RAMT) | For the upper test the value |
| AND | A | for RAMTOP is tested against |
| SBC | HL,DE | P-RAMT. |
| JR | NC,1EDC,CLEAR-2 | Jump forward if acceptable. |

Report M - RAMTOP no good

| 1EDA | REPORT-M | $\begin{aligned} & \text { RST } \\ & \text { DEFB } \end{aligned}$ | $\begin{aligned} & \text { 0008,ERROR-1 } \\ & +15 \end{aligned}$ | Call the error handling routine. |
| :---: | :---: | :---: | :---: | :---: |
| Continue with the CLEAR operation. |  |  |  |  |
| 1EDC | CLEAR-2 | EX | DE,HL | Now the value can actually be |
|  |  | LD | (RAMTOP), HL | passed to RAMTOP. |
|  |  | POP | DE | Fetch the address - STMT-RET. |
|  |  | POP | BC | Fetch the 'error address'. |
|  |  | LD | (HL), +3E | Enter a GO SUB stack end marker. |
|  |  | DEC | HL | Leave one location. |
|  |  | LD | SP,HL | Make the stack pointer point |


| PUSH | BC | to an empty GO SUB stack. |
| :--- | :--- | :--- |
| LD | (ERR-SP),SP | Next pass the 'error address' <br> to the stack and save its |
| EX | DE,HL | address in ERR-SP. |
| JP | (HL) | An indirect return is now <br> made to STMT-RET. |

Note: When the routine is called from RUN the values of NEWPPC \& NSPPC will have been affected and no statements coming after RUN can ever be found before the jump is taken.

## THE 'GO SUB' COMMAND ROUTINE

| The present value of PPC and the incremented value of SUBPPC are stored on the GO SUB stack. |  |  |
| :--- | :--- | :--- |
| 1EED | GO-SUB | POP |

THE 'TEST-ROOM' SUBROUTINE
A series of tests is performed to ensure that there is sufficient free memory available for the task being undertaken.


THE 'FREE MEMORY' SUBROUTINE
There is no BASIC command 'FRE' in the SPECTRUM but there is a subroutine for performing such a task. An estimate of the amount of free space can be found at any time by using:
'PRINT 65536-USR 7962'

| 1F1A FREE-MEM | LD | BC,+0000 | Do not allow any overhead. |
| :--- | :--- | :--- | :--- |
|  | CALL | 1F05,TEST-ROOM | Make the test and pass the |


| LD | $B, H$ | result to the $B C$ register |
| :--- | :--- | :--- |
| LD | $C, L$ | before returning. |

## THE 'RETURN' COMMAND ROUTINE

The line number and the statement number that are to be made the object of a 'return' are fetched from the GO SUB stack.

| 1F23 | RETURN | POP | BC | Fetch the address - STMT-RET. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | POP | HL | Fetch the 'error address'. |
|  |  | POP | DE | Fetch the last entry on the GO SUB stack |
|  |  | LD | A, D | The entry is tested to see if |
|  |  | CP | +3E | it is the GO SUB stack end |
|  |  | JR | Z,1F36,REPORT-7 | marker; jump if it is. |
|  |  | DEC | SP | The full entry uses three locations only. |
|  |  | EX | (SP), HL | Exchange the statement number with the 'error address'. |
|  |  | EX | DE,HL | Move the statement number. |
|  |  | LD | (ERR-SP),SP | Reset the error pointer. |
|  |  | PUSH | BC | Replace the address -STMT-RET. |
|  |  | JP | 1E73,GO-TO-2 | Jump back to change NEWPPC \& NSPPC. |
| Report 7 - RETURN without GOSUB |  |  |  |  |
| 1F36 | REPORT-7 | PUSH | DE | Replace the end marker and |
|  |  | PUSH | HL | the 'error address'. |
|  |  | RST | 0008,ERROR-1 | Call the error handling |
|  |  | DEFB | +06 | routine. |

## THE 'PAUSE' COMMAND ROUTINE

The period of the PAUSE is determined by counting the number of maskable interrupts as they occur every $1 / 50$ th. of a second. A PAUSE is finished either after the appropriate number of interrupts or by the system Variable FLAGS indicating that a key has been pressed.

| 1F3A | PAUSE | CALL | 1E99,FIND-INT2 | Fetch the operand. |
| :---: | :---: | :---: | :---: | :---: |
| 1F3D | PAUSE-1 | HALT |  | Wait for a maskable interrupt. |
|  |  | DEC | BC | Decrease the counter. |
|  |  | LD | A,B | If the counter is thereby |
|  |  | OR | C | reduced to zero the PAUSE |
|  |  | JR | Z,1F4F,PAUSE-END | has come to an end. |
|  |  | LD | A, B | If the operand was zero BC |
|  |  | AND | C | will now hold +FFFF and this |
|  |  | INC | A | value will be returned to |
|  |  | JR | NZ,1F49,PAUSE-2 | zero. Jump will all other |
|  |  | INC | BC | operand values. |
| 1F49 | PAUSE-2 | BIT | 5,(FLAGS) | Jump back unless a key has |
|  |  | JR | Z,1F3D,PAUSE-1 | been pressed. |
| The period of the PAUSE has now finished. |  |  |  |  |
| 1F4F | PAUSE-EN | RES | 5,(FLAGS) | Signal 'no key pressed'. |
|  |  | RET |  | Now return; - to STMT-RET. |

## THE 'BREAK-KEY' SUBROUTINE

This subroutine is called in several instances to read the BREAK key. The carry flag is returned reset only if the SHIFT and the BREAK keys are both being pressed.

| 1F54 | BREAK-KEY | LD |
| :--- | :--- | :--- |
|  |  | IN |
|  |  | RRA |

Form the port address
+7FFE and read in a byte.
Examine only bit 0 by shifting it into the carry position.

| RET | C | Return if the BREAK key is <br> not being pressed. |
| :--- | :--- | :--- |
| LD | A,+FE | Form the port address |
| IN | A, + FE $)$ | +FEFE and read in a byte. |
| RRA |  | Again examine bit 0. |
| RET |  | Return with carry reset if |
|  |  | both keys are being pressed. |

## THE 'DEF FN' COMMAND ROUTINE

During syntax checking a DEF FN statement is checked to ensure that it has the correct form. Space is also made available for the result of evaluating the function.
But in 'run-time' a DEF FN statement is passed-by.

| 1F60 | DEF-FN | CALL JR LD JP | $\begin{aligned} & \text { 2530,SYNTAX-Z } \\ & \text { Z,1F6A,DEF-FN-1 } \\ & \text { A,+CE, } \\ & \text { 1E39,PASS-BY } \end{aligned}$ | Jump forward if checking syntax. <br> Otherwise bass-by the 'DEF FN' statement. |
| :---: | :---: | :---: | :---: | :---: |
| First consider the variable of the function. |  |  |  |  |
| 1F6A | DEF-FN-1 | SET | 6,(FLAGS) | Signal 'a numeric variable'. |
|  |  | CALL | 2C8D,ALPHA | Check that the present code is a letter. |
|  |  | JR | NC,1F89,DEF-FN-4 | Jump forward if not. |
|  |  | RST | 0020,NEXT-CHAR | Fetch the next character. |
|  |  | CP | +24 | Jump forward unless it is |
|  |  | JR | NZ,1F7D,DEF-FN-2 | a '\$'. |
|  |  | RES | 6,(FLAGS) | Change bit 6 as it is a string variable. |
|  |  | RST | 0020,NEXT-CHAR | Fetch the next character. |
| 1F7D | DEF-FN-2 | CP | +28 | A '(' must follow the |
|  |  | JR | NZ,1FBD,DEF-FN-7 | variable's name. |
|  |  | RST | 0020,NEXT-CHAR | Fetch the next character. |
|  |  | CP | +29 | Jump forward if it is a |
|  |  | JR | Z,1FA6,DEF-FN-6 | ')' as there are no |
|  |  |  |  | parameters of the function. |
| A loop is now entered to deal with each parameter in turn. |  |  |  |  |
| 1F86 | DEF-FN-3 | CALL | 2C8D,ALPHA | The present code must be |
| 1F89 | DEF-FN-4 | JP | NC,1C8A,REPORT-C | a letter. |
|  |  | EX | DE,HL | Save the pointer in DE. |
|  |  | RST | 0020,NEXT-CHAR | Fetch the next character. |
|  |  | CP | +24 | Jump forward unless it is |
|  |  | JR | NZ,1F94,DEF-FN-5 | a '\$'. |
|  |  | EX | DE,HL | Otherwise save the new pointer in DE instead. |
|  |  | RST | 0020,NEXT-CHAR | Fetch the next character. |
| 1F94 | DEF-FN-5 | EX | DE,HL | Move the pointer to the last character of the name to the HL register pair. |
|  |  | LD | BC,+0006 | Now make six locations after |
|  |  | CALL | 1655,MAKE-ROOM | that last character and |
|  |  | INC | HL | enter a 'number marker' into |
|  |  | INC | HL | the first of the new |
|  |  | LD | (HL), +0E | locations. |
|  |  | CP | +2C | If the present character is |
|  |  | JR | NZ,1FA6,DEF-FN-6 | a ',' then jump back as |
|  |  | RST | 0020,NEXT-CHAR | there should be a further |
|  |  | JR | 1F86,DEF-FN-3 | parameter; otherwise jump |
| Next the definition of the function is considered. |  |  |  |  |
| 1FA6 | DEF-FN-6 | CP | +29 | Check that the ')' does |
|  |  | JR | NZ,1FBD,DEF-FN-7 | exist. |


|  |  | RST | 0020,NEXT-CHAR | The next character is fetched. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | CP | +3D | It must be an ' $=$ '. |
|  |  | JR | NZ,1FBD,DEF-FN-7 |  |
|  |  | RST | 0020,NEXT-CHAR | Fetch the next character. |
|  |  | LD | A,(FLAGS) | Save the nature - numeric or |
|  |  | PUSH | AF | string - of the variable. |
|  |  | CALL | 2F4B,SCANNING | Now consider the definition as an expression. |
|  |  | POP | AF | Fetch the nature of the |
|  |  | XOR | (FLAGS) | variable and check that it |
|  |  | AND | +40 | is of the same type as found for the definition. |
| 1FBD | DEF-FN-7 | JP | NZ,1C8A,REPORT-C | Give an error report if it is required. |
|  |  | CALL | 1BEE,CHECK-END | Exit via the CHECK-END subroutine. (Thereby moving on to consider the next statement in the line.) |

## THE 'UNSTACK-Z' SUBROUTINE

This subroutine is called in several instances in order to 'return early' from a subroutine when checking syntax. The reason for this is to avoid actually printing characters or passing values to/from the calculator stack.

| 1FC3 | UNSTACK-Z | CALL | 2530,SYNTAX-Z |
| :--- | :--- | :--- | :--- |
|  |  | POP | HL | | Is syntax being checked? |
| :--- |
|  |
|  |
|  |
|  |
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|  |
|  |

THE 'LPRINT \& PRINT' COMMAND ROUTINES
The appropriate channel is opened as necessary and the items to be printed are considered in turn.

| 1FC9 | LPRINT | LD | A, +03 | Prepare to open channel ' $P$ '. Jump forward. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | JR | 1FCF,PRINT-1 |  |
| $\begin{aligned} & 1 \mathrm{FCD} \\ & 1 \mathrm{FCF} \end{aligned}$ | PRINT PRINT-1 | LD | A,+02 | Prepare to open channel 'S'. |
|  |  | CALL | 2530,SYNTAX-Z | Unless syntax is being |
|  |  | CALL | NZ,1601,CHAN-OPEN | checked open a channel. |
|  |  | CALL | 0D4D,TEMPS | Set the temporary colour system variables. |
|  |  | CALL | 1FDF,PRINT-2 | Call the print controlling subroutine. |
|  |  | CALL | 1BEE,CHECK-END | Move on to consider the next |
|  |  | RET |  | statement; via CHECK-END IF checking syntax. |
| The print controlling subroutine is called by the PRINT, LPRINT and INPUT command routines. |  |  |  |  |
| 1FDF | PRINT-2 | RST | 0018,GET-CHAR | Get the first character. |
|  |  | CALL | 2045,PR-END-Z | Jump forward if already at the |
|  |  | JR | Z,1FF2,PRINT-4 | end of the item list. |
| Now enter a loop to deal with the 'position controllers' and the print items. |  |  |  |  |
| 1FE5 | PRINT-3 | CALL | 204E,PR-POSN-1 | Deal with any consecutive |
|  |  | JR | Z,1FE5,PRINT-3 | position controllers. |
|  |  | CALL | 1FFC, PR-ITEM-1 | Deal with a single print item. |
|  |  | CALL | 204E,PR-POSN-1 | Check for further position |
|  |  | JR | Z,1FE5,PRINT-3 | controllers and print items |
|  |  |  |  | until there are none left. |
| 1FF2 | PRINT-4 | CP | +29 | Return now if the present |
|  |  | RET | Z | character is a ')'; otherwise |
|  |  |  |  | consider performing a 'carriage |
|  |  |  |  | return'. |

## THE 'PRINT A CARRIAGE RETURN' SUBROUTINE

CALL
LD
RST
RET

1FC3,UNSTACK-Z
A,+0D
0010,PRINT-A-1

Return if changing syntax
Print a carriage return character and then return

## THE 'PRINT ITEMS' SUBROUTINE

This subroutine is called from the PRINT, LPRINT and INPUT command routines.
The various types of print item are identified and printed.

| 1FFC | PR-ITEM-1 | RST | 0018,GET-CHAR | The first character is fetched. Jump forward unless it is an 'AT'. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | CP | +AC |  |
|  |  | JR | NZ,200E,PR-ITEM-2 |  |
| Now deal with an 'AT'. |  |  |  |  |
|  |  | CALL | 1C79,NEXT-2NUM | The two parameters are transferred to the calculator stack. |
|  |  | CALL | 1FC3,UNSTACK-Z | Return now if checking syntax. |
|  |  | CALL | 2307,STK-TO-BC | The parameters are compressed into the $B C$ register pair. |
|  |  | LD | A, +16 | The A register is loaded with the AT control character before the jump is taken. |
|  |  | JR | 201E,PR-AT-TAB |  |
| Next look for a 'TAB'. |  |  |  |  |
| 200E | PR-ITEM-2 | CP | +AD | Jump forward unless it is |
|  |  | JR | NZ,2024,PR-ITEM-3 | a 'TAB'. |
| Now deal with a 'TAB'. |  |  |  |  |
|  |  | RST | 0020,NEXT-CHAR | Get the next character. Transfer one parameter to the calculator stack. |
|  |  | CALL | 1C82,EXPT-1NUM |  |
|  |  | CALL | 1FC3,UNSTACK-Z | Return now if checking syntax. The value is compressed into the |
|  |  | CALL | 1E99,FIND-INT2 |  |
|  |  |  |  | BC register pair. |
|  |  | LD | A, +17 | The A register is loaded with the TAB control character. |
|  |  |  |  |  |
| The 'AT' and the 'TAB' print items are printed by making three calls to PRINT-OUT. |  |  |  |  |
| 201E | PR-AT-TAB | RST | 0010,PRINT-A-1 | Print the control character. |
|  |  | LD | A, C | Follow it with the first |
|  |  | RST | 0010,PRINT-A-1 |  |
|  |  | LD | A,B | Finally print the second |
|  |  | RST | 0010,PRINT-A-1 | value; then return. |
|  |  | RET |  |  |
| Next consider embedded colour items. |  |  |  |  |
| 2024 | PR-ITEM-3 | CALL | 21F2,CO-TEMP-3 | Return with carry reset if a colour items was found. |
|  |  | RET | NC | Continue if none were found. Next consider if the stream is to be changed. |
|  |  | CALL | 2070,STR-ALTER |  |
|  |  | RET | NC | Continue unless it was altered. |
| The print item must now be an expression, either numeric or string. |  |  |  |  |
|  |  | CALL | 24FB,SCANNING | Evaluate the expression but return now if checking syntax. Test for the nature of the expression. |
|  |  | CALL | 1FC3,UNSTACK-Z |  |
|  |  | BIT | 6,(FLAGS) |  |
|  |  | CALL | Z,2BF1,STK-FETCH | If it is string then fetch the necessary parameters; but if it is |
|  |  |  |  |  |
|  |  | JP | NZ,2DE3,PRINT-FP | numeric then exit via PRINT-FP. |

A loop is now set up to deal with each character in turn of the string

| 203C | PR-STRING | LD |
| :--- | :--- | :--- |
|  | OR | A,B |
|  | DEC | C |
|  | RET | B |
|  | LD | A,(DE) |
|  | INC | DE |
|  | RST | 0010 ,PRINT-A-1 |
|  | JR | 203C,PR-STRING |

THE 'END OF PRINTING' SUBROUTINE
The zero flag will be set if no further printing is to be done.

| 2045 | PR-END-Z | CP | +29 |
| :--- | :--- | :--- | :--- |
| 2048 |  | PR-ST-END | RET |
|  |  | CP | Z |
|  |  | RET | Z |
|  |  | CP | Z |
|  |  |  | RET |

THE 'PRINT POSITION' SUBROUTINE
The various position controlling characters are considered by this subroutine.

| 204E | PR-POSN-1 | RST | 0018,GET-CHAR | Get the present character. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | CP | +3B | Jump forward if it is |
|  |  | JR | Z,2067,PR-POSN-3 |  |
|  |  | CP | +2C | Also jump forward with a |
|  |  | JR | NZ,2061,PR-POSN-2 | character other than a ','; |
|  |  | CALL | 2530,SYNTAX-Z | but do not actually print the |
|  |  | JR | Z,2067,PR-POSN-3 | character if checking syntax. |
|  |  | LD | A, +06 | Load the A register with |
|  |  | RST | 0010,PRINT-A-1 | the 'comma' control code and |
|  |  | JR | 2067,PR-POSN-3 | print it; then jump forward. |
| 2061 | PR-POSN-2 | CP | +27 | Is it a "'? |
|  |  | RET | NZ | Return now if not any of the position controllers. |
|  |  | CALL | 1FF5,PR-CR | Print 'carriage return' unless checking syntax. |
| 2067 | PR-POSN-3 | RST | 0020,NEXT-CHAR | Fetch the next character. |
|  |  | CALL | 2045,PR-END-Z | If not at the end of a print |
|  |  | JR | NZ,206E,PR-POSN-4 | statement then jump forward; |
|  |  | POP | BC | otherwise return to the |
| 206E | PR-POSN-4 | CP | A | calling routine. |
|  |  | RET |  | The zero flag will be reset if the end of the print statement has not been reached. |

THE 'ALTER STREAM' SUBROUTINE
This subroutine is called whenever there is the need to consider whether the user wishes to use a different stream

2070 STR-ALTER \begin{tabular}{lll}

CP \& +23 \& | Unless the present character |
| :--- |
| is a '\#' return with the | <br>

\& SCF \& nZ <br>
RETry flag set.
\end{tabular}

| JP | NC, 160E,REPORT-O | over +FF. <br> CALL |
| :--- | :--- | :--- |
| 1601,CHAN-OPEN | Use the channel for the |  |
| stream in question. |  |  |

THE 'INPUT' COMMAND ROUTINE
This routine allows for values entered from the keyboard to be assigned to variables. It is also possible to have print items embedded in the INPUT statement and these items are printed in the lower part of the display.

| 2089 | INPUT | CALL | 2530,SYNTAX-Z | Jump forward if syntax is |
| :---: | :---: | :---: | :---: | :---: |
|  |  | JR | Z,2096,INPUT-1 | being checked. |
|  |  | LD | A, +01 | Open channel 'K'. |
|  |  | CALL | 1601,CHAN-OPEN |  |
|  |  | CALL | 0D6E,CLS-LOWER | The lower part of the display is cleared. |
| 2096 | INPUT-1 | LD | (TV-FLAG),+01 | Signal that the lower screen is being handled. Reset all other bits. |
|  |  | CALL | 20C1,IN-ITEM-1 | Call the subroutine to deal with the INPUT items. |
|  |  | CALL | 1BEE,CHECK-END | Move on to the next statement if checking syntax. |
|  |  | LD | BC,(S-POSN) | Fetch the current print position. |
|  |  | LD | A,(DF-SZ) | Jump forward if the current |
|  |  | CP | B | position is above the lower |
|  |  | JR | C,20AD,INPUT-2 | screen. |
|  |  | LD | C,+21 | Otherwise set the print position |
|  |  | LD | B,A | to the top of the lower screen. |
| 20AD | INPUT-2 | LD | (S-POSN), BC | Reset S-POSN. |
|  |  | LD | A, +19 | Now set the scroll counter. |
|  |  | SUB | B |  |
|  |  | LD | (SCR-CT),A |  |
|  |  | RES | 0,(TV-FLAG) | Signal 'main screen'. |
|  |  | CALL | 0DD9,CL-SET | Set the system variables |
|  |  | JP | 0D6E,CLS-LOWER | and exit via CLS-LOWER. |
| The INPUT items and embedded PRINT items are dealt with in turn by the following loop. |  |  |  |  |
| 20C1 | IN-ITEM-1 | CALL | 204E,PR-POSN-1 | Consider first any position |
|  |  | JR | Z,20C1,IN-ITEM-1 | control characters. |
|  |  | CP | +28 | Jump forward if the present |
|  |  | JR | NZ,20D8,IN-ITEM-2 | character is not a '('. |
|  |  | RST | 0020,NEXT-CHAR | Fetch the next character. |
|  |  | CALL | 1FDF,PRINT-2 | Now call the PRINT command routine to handle the items inside the brackets. |
|  |  | RST | 0018,GET-CHAR | Fetch the present character. |
|  |  | CP | +29 | Give report C unless the |
|  |  | JP | NZ,1C8A,REPORT-C | character is a ')'. |
|  |  | RST | 0020,NEXT-CHAR | Fetch the next character and |
|  |  | JP | 21B2,IN-NEXT-2 | jump forward to see if there are any further INPUT items. |
| Now consider whether INPUT LINE is being used. |  |  |  |  |
| 20D8 | IN-ITEM-2 | CP | +CA | Jump forward if it is not |
|  |  | JR | NZ,20ED,IN-ITEM-3 | 'LINE'. |
|  |  | RST | 0020,NEXT-CHAR | Advance CH-ADD. |
|  |  | CALL | 1C1F,CLASS-01 | Determine the destination address for the variable. |
|  |  | SET | 7,(FLAGX) | Signal 'using INPUT LINE'. |
|  |  | BIT | 6,(FLAGS) | Give report C unless using |
|  |  | JP | NZ,1C8A,REPORT-C | a string variable. |


| JR |  |  | 20FA,IN-PROMPT | Jump forward to issue the prompt message. |
| :---: | :---: | :---: | :---: | :---: |
| Proceed to handle simple INPUT variables. |  |  |  |  |
| 20ED | IN-ITEM-3 | CALL | 2C8D,ALPHA | Jump to consider going round |
|  |  | JP | NC,21AF-IN-NEXT-1 | the loop again if the present character is not a letter. |
|  |  | CALL | 1C1F,CLASS-01 | Determine the destination |
|  |  |  |  | address for the variable. |
|  |  | RES | 7,(FLAGX) | Signal 'not INPUT LINE'. |
| The prompt message is now built up in the work space. |  |  |  |  |
| 20FA | IN-PROMPT | CALL | 230,SYNTAX-Z | Jump forward if only checking |
|  |  | JP | Z,21B2,IN-NEXT-2 | syntax. |
|  |  | CALL | 16BF,SET-WORK | The work space is set to null. |
|  |  | LD | HL, +5C71 | This is FLAGX. |
|  |  | RES | 6,(HL) | Signal 'string result'. |
|  |  | SET | 5,(HL) | Signal 'INPUT mode'. |
|  |  | LD | BC, +0001 | Allow the prompt message only a single location. |
|  |  | BIT | 7,(HL) | Jump forward if using 'LINE'. |
|  |  | JR | NZ,211C,IN-PR-2 |  |
|  |  | LD | A,(FLAGS) | Jump forward if awaiting |
|  |  | AND | +40 | a numeric entry. |
|  |  | JR | NZ,211A,IN-PR-1 |  |
|  |  | LD | C, +03 | A string entry will need three |
| 211A | IN-PR-1 | OR | (HL) | Bit 6 of FLAGX will become |
|  |  | LD | (HL),A | set for a numeric entry. |
| 211C | IN-PR-2 | RST | 0030,BC-SPACES | The required number of |
|  |  |  |  | locations is made available. |
|  |  | LD | (HL), +0D | A 'carriage return' goes into |
|  |  |  |  | the last location. |
|  |  | LD | A, C | Test bit 6 of the C register |
|  |  | RRCA |  | and jump forward if only |
|  |  | RRCA |  | one location was required. |
|  |  | JR | NC,2129,IN-PR-3 |  |
|  |  | LD | A,+22 | A 'double quotes' character |
|  |  | LD | (DE), A | goes into the first and |
|  |  | DEC | HL | second locations. |
|  |  | LD | (HL), A |  |
| 2129 | IN-PR-3 | LD | (K-CUR), HL | The position of the cursor |
|  |  | In the case of INPUT LINE the EDITOR can be called without further preparation but for other types of INPUT the error stack has to be |  |  |  |  |
|  |  |  |  |  |  |  |
|  |  | BIT | 7,(FLAGX) | Jump forward with INPUT |
|  |  | JR | NZ,215E,IN-VAR-3 | LINE' |
|  |  | LD | HL,(CH-ADD) | Save the current values of |
|  |  | PUSH | HL | CH-ADD \& ERR-SP on the |
|  |  | LD | HL, (ERR-SP) | machine stack. |
|  |  | PUSH | HL |  |
| 213A | IN-VAR-1 | LD | HL, +213A | This will be the 'return |
|  |  | PUSH | HL | point' in case of errors. |
|  |  | BIT | 4,(FLAGS2) | Only change the error |
|  |  | JR | Z,2148,IN-VAR-2 | stack pointer if using channel |
|  |  | LD | (ERR-SP),SP | 'K'. |
| 2148 | IN-VAR-2 | LD | HL,(WORKSP) | Set HL to the start of the |
|  |  | CALL | 11A7,REMOVE-FP | INPUT line and remove any |
|  |  |  |  | floating-point forms. (There will |
|  |  |  |  | not be any except perhaps after |
|  |  |  |  | an error.) |
|  |  | LD | (ERR-NR),+FF | Signal 'no error yet'. |


|  |  | CALL RES CALL JR | $\begin{aligned} & \text { 0F2C,EDITOR } \\ & \text { 7,(FLAGS) } \\ & \text { 21B9,IN-ASSIGN } \\ & \text { 2161,IN-VAR-4 } \end{aligned}$ | Now get the INPUT and with the syntax/run flag indicating syntax, check the INPUT for errors; jump if in order; return to IN-VAR-1 if not. |
| :---: | :---: | :---: | :---: | :---: |
| 215E | IN-VAR-3 | CALL | OF2C,EDITOR | Get a 'LINE'. |
| All the system variables have to be reset before the actual assignment of a value can be made. |  |  |  |  |
| 2161 | IN-VAR-4 | LD | (K-CUR-hi),+00 | The cursor address is reset. |
|  |  | CALL | 21D6,IN-CHAN-K | The jump is taken if using |
|  |  | JR | NZ,2174,IN-VAR-5 | other than channel ' K '. |
|  |  | CALL | 111D,ED-COPY | The input-line is copied to |
|  |  | LD | BC,(ECHO-E) | the display and the position |
|  |  | CALL | 0DD9,CL-SET | in ECHO-E made the current position in the lower screen. |
| 2174 | IN-VAR-5 | LD | HL, +5C71 | This is FLAGX. |
|  |  | RES | 5,(HL) | Signal 'edit mode'. |
|  |  | BIT | 7,(HL) | Jump forward if handling an |
|  |  | RES | 7,(HL) | INPUT LINE. |
|  |  | JR | NZ,219B,IN-VAR-6 |  |
|  |  | POP | HL | Drop the address IN-VAR-1. |
|  |  | POP | HL | Reset the ERR-SP to its |
|  |  | LD | (ERR-SP), HL | original address. |
|  |  | POP | HL | Save the original CH-ADD |
|  |  | LD | (X-PTR), HL | address in X-PTR. |
|  |  | SET | 7,(FLAGS) | Now with the syntax/run flag |
|  |  | CALL | 21B9,IN-ASSIGN | indicating 'run' make the assignment. |
|  |  | LD | HL,(X-PTR) | Restore the original address |
|  |  | LD | (X-PTR-hi),+00 | to CH-ADD and clear X-PTR. |
|  |  | LD | (CH-ADD), HL |  |
|  |  | JR | 21B2,IN-NEXT-2 | Jump forward to see if there are further INPUT items. |
| 219B | IN-VAR-6 | LD | HL,(STKBOT) | The length of the 'LINE' in |
|  |  | LD | DE,(WORKSP) | the work space is found. |
|  |  | SCF |  |  |
|  |  | SBC, | HL, DE |  |
|  |  | LD | B, H | DE points to the start and |
|  |  | LD | C,L | $B C$ holds the length. |
|  |  | CALL | 2AB2,STK-ST-\$ | These parameters are stacked |
|  |  | CALL | 2AFF,LET | and the actual assignment made. |
|  |  | JR | 21B2,IN-NEXT-2 | Also jump forward to consider |
| Further items in the INPUT statement are considered. |  |  |  |  |
| 21AF | IN-NEXT-1 | CALL | 1FFC,PR-ITEM-1 | Handle any print items. |
| 21B2 | IN-NEXT-2 | CALL | 204E,PR-POSN-1 | Handle any position controllers. |
|  |  | JP | Z,20C1,IN-ITEM-1 | Go around the loop again if there are further items; |
|  |  | RET |  | otherwise return. |

THE 'IN-ASSIGN' SUBROUTINE
This subroutine is called twice for each INPUT value. Once with the syntax/run flag reset (syntax) and once with it set (run).

| LD | HL,(WORKSP) |
| :--- | :--- |
| LD | (CH-ADD),HL |
| RST | 0018, GET-CHAR |
| CP | +E2 |
| JR | Z,21D0,IN-STOP |
| LD | A,(FLAGX) |
| CALL | 1 C59,VAL-FET-2 |

Set CH-ADD to point to the
first location of the work
space and fetch the character.
Is it a 'STOP'?
Jump if it is.
Otherwise make the assignment
of the 'value' to the variable.


## THE 'COLOUR ITEM' ROUTINES

This set of routines can be readily divided into two parts:
i. The embedded colour item' handler.
ii. The 'colour system variable' handler.
i. Embedded colour items are handled by calling the PRINT-OUT subroutine as required.

A loop is entered to handle each item in turn. The entry point is at CO-TEMP-2.

| 21E1 | CO-TEMP-1 | RST | 0020,NEXT-CHAR | Consider the next character in the BASIC statement. |
| :---: | :---: | :---: | :---: | :---: |
| 21E2 | CO-TEMP-2 | CALL | 21F2,CO-TEMP-3 | Jump forward to see if the present code represents an embedded 'temporary' colour |
|  |  | RET | C | item. Return carry set if not a colour item. |
|  |  | RST | 0018,GET-CHAR | Fetch the present character. |
|  |  | CP | +2C | Jump back if it is either a |
|  |  | JR | Z,21E1,CO-TEMP-1 | ',' or a ';'; otherwise |
|  |  | CP | +3B | there has been an error. |
|  |  | JR | Z,21E1,CO-TEMP-1 |  |
|  |  | JP | 1C8A,REPORT-C | Exit via 'report C'. |
| 21F2 | CO-TEMP-3 | CP | +D9 | Return with the carry flag |
|  |  | RET | C | Set if the code is not in the |
|  |  | CP | +DF | range +D9 to +DE (INK to OVER). |
|  |  | CCF |  |  |
|  |  | RET | C |  |
|  |  | PUSH | AF | The colour item code is |
|  |  | RST | 0020,NEXT-CHAR | preserved whilst CH-ADD is |
|  |  | POP | AF | advanced to address the |
|  |  |  |  | parameter that follows it. |
| The colour item code and the parameter are now 'printed' by calling PRINT-OUT on two occasions. |  |  |  |  |
| 21FC | CO-TEMP-4 | SUB | +C9 | The token range (+D9 to +DE) |


|  | is reduced to the control <br> character range $(+10$ to +15$)$. |  |
| :--- | :--- | :--- |
| PUSH | AF | The control character code is |
| CALL | 1 C882,EXPT-1NUM | preserved whilst the parameter |
| POP | AF moved to the calculator stack. |  |

ii. The colour system variables - ATTR-T, MASK-T \& P-FLAG - are altered as required. This subroutine is called by PRINT-OUT. On entry the control character code is in the A register and the parameter is in the D register.
Note that all changes are to the 'temporary' system variables.


|  |  | LD | A, D | Fetch the original value. Only allow PAPER/INK a range of '0' to ' 9 '. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | CP | +0A |  |
|  |  | JR | C,2246,CO-TEMP-9 |  |
| Report K - Invalid colour |  |  |  |  |
| 2244 | REPORT-K | RST | 0008,ERROR-1 | Call the error handling |
|  |  | DEFB | +13 |  |
| Continue to handle PAPER \& INK; |  |  |  |  |
| 2246 | CO-TEMP-9 | LD | HL, +5C8F | Prepare to alter ATTR-T, <br> MASK-T \& P-FLAG. <br> Jump forward with PAPER/INK |
|  |  | CP | +08 |  |
|  |  | JR | C,2258,CO-TEMP-B | '0' to 5'7'. |
|  |  | LD | A,(HL) | Fetch the current value of |
|  |  | JR | Z,2257,CO-TEMP-A | ATTR-T and use it unchanged, by jumping forward, with PAPER/INK ' 8 '. |
|  |  | OR | B | But for PAPER/INK '9' the |
|  |  | CPL |  | PAPER and INK colours |
|  |  | AND | +24 | have to be black and white. |
|  |  | JR | Z,2257,CO-TEMP-A | Jump for black INK/PAPER; |
|  |  | LD | A,B | but continue for white INK/ |
|  |  |  |  | PAPER. |
| 2257 | CO-TEMP-A | LD | C,A | Move the value to the C register. |
| The mask (B) and the value ( C ) are now used to change ATTR-T. |  |  |  |  |
| 2258 | CO-TEMP-B | LD | A,C | Move the value. |
|  |  | CALL | 226C,CO-CHANGE | Now change ATTR-T as needed. |
| Next MASK-T is considered. |  |  |  |  |
|  |  | LD | A,+07 | The bits of MASK-T are set |
|  |  | CP | D | only when using PAPER/INK |
|  |  | SBC | A,A | '8' or '9'. |
|  |  | CALL | 226C,CO-CHANGE | Now change MASK-T as needed. |
| Next P-FLAG is considered. |  |  |  |  |
|  |  | RLCA |  | The appropriate mask is |
|  |  | RLCA |  | built up in the B register |
|  |  | AND | +50 |  |
|  |  | LD | B,A | 6 as necessary. |
|  |  | LD | A,+08 | The bits of P-FLAG are set |
|  |  | CP | D | only when using PAPER/INK |
|  |  | SBC | A,A | '9'. |
|  |  |  |  | Continue into CO-CHANGE to manipulate P-FLAG. |

## THE 'CO-CHANGE' SUBROUTINE

This subroutine is used to 'impress' upon a system variable the 'nature' of the bits in the A register, The B register holds a mask that shows which bits are to be 'copied over' from A to (HL).
226C CO-CHANGE XOR $\quad 1 \mathrm{HL})$

FLASH \& BRIGHT are handled by the following routine.

The bits, specified by the mask in the B register, are changed in the value and the result goes to form the system variable.
Move on to address the next system variable.
Return with the mask in the A register.

The zero flag will be set for BRIGHT
The parameter is fetched and rotated.
Prepare the mask for FLASH
Jump forward with FLASH.
Rotate an extra time and
prepare the mask for BRIGHT.
Save the value in the C register.
Fetch the parameter and test
its range; only '0', '1
JR Z,2287,CO-TEMP-E
CP +02
JR NC,2244,REPORT-K
\& '8' are allowable.

The system variable ATTR-T can now be altered
2287 CO-TEMP-E LD A,C
A,C
HL,+5C8F
226C

The value in MASK-T is now considered

| LD | A,C |
| :--- | :--- |
| RRCA |  |
| RRCA |  |
| RRCA |  |
| JR |  |

Fetch the value.
This is ATTR-T.
Now change the system variable
The value is fetched anew.
The set bit of FLASH/BRIGHT
' 8 ' (bit 3 ) is moved to
bit 7 (for FLASH) or bit 6
(for BRIGHT).
Exit via CO-CHANGE

## THE 'BORDER' COMMAND ROUTINE

The parameter of the BORDER command is used with an OUT command to actually alter the colour of the border. The parameter is then saved in the system variable BORDCR.


## THE 'PIXEL ADDRESS' SUBROUTINE

This subroutine is called by the POINT subroutine and by the PLOT command routine. Is is entered with the co-ordinates of a pixel in the BC register pair and returns with HL holding the address of the display file byte which contains that pixel and A pointing to the position of the pixel within the byte.

| 22AA | PIXEL-ADD | LD | A, +AF | Test that the y co-ordinate (in |
| :---: | :---: | :---: | :---: | :---: |
|  |  | SUB | B | B ) is not greater than 175. |
|  |  | JP | C,24F9,REPORT-B |  |
|  |  | LD | B,A | B now contains 175 minus y. |
|  |  | AND | A | A holds b7b6b5b4b3b2b1b0, |
|  |  | RRA |  | the bite of $B$. And now |
|  |  |  |  | 0b7b6b5b4b3b2b1. |
|  |  | SCF |  |  |
|  |  | RRA |  | Now 10b7b6b5b4b3b2. |


| AND | A |  |
| :---: | :---: | :---: |
| RRA |  | Now 010b7b6b5b4b3. |
| XOR | B |  |
| AND | +F8 | Finally 010b7b6b2b1b0, so that |
| XOR | B | H becomes $64+8^{*} \mathrm{INT}(\mathrm{B} / 64)+$ |
| LD | H,A | $B(\bmod 8)$, the high byte of the |
| LD | A,C | pixel address. C contains X . |
| RLCA |  | A starts as c 7 c 6 c 5 c 4 c 3 c 2 c 1 c 0. |
| RLCA |  |  |
| RLCA |  | And is now c2c 1 c 0 c 7 c 6 c 5 c 4 c 3 . |
| XOR | B |  |
| AND | +C7 |  |
| XOR | B | Now c2c1b5b4b3c5c4c3. |
| RLCA |  |  |
| RLCA |  | Finally b5b4b3c7c6c5c4c3, so |
| LD | L,A | that L becomes $32^{*} \mathrm{INT}$ ( B (mod |
| LD | A, C | $64) / 8)+\mathrm{INT}(\mathrm{x} / 8)$, the low byte. |
| AND | +07 | A holds $x(\bmod 8)$ : so the pixel |
| RET |  | is bit ( $\mathrm{A}-7)$ within the byte. |

## THE 'POINT' SUBROUTINE

This subroutine is called by the POINT function in SCANNING. It is entered with the co-ordinates of a pixel on the calculator stack, and returns a last value of 1 if that pixel is ink colour, and 0 if it is paper colour.

| 22CB | POINT-SUB | CALL | 2307,STK-TO-BC | Y co-ordinate to B, x to C. |
| :--- | :--- | :--- | :--- | :--- |
|  |  | CALL | 22AA,PIXEL-ADD | Pixel address to HL. |
|  |  | LD | B,A | B will count A+1 loops to get |
|  |  | INC | B | the wanted bit of (HL) to |
|  |  | LD | A,(HL) | location 0. |

## THE 'PLOT' COMMAND ROUTINE

This routine consists of a main subroutine plus one line to call it and one line to exit from it. The main routine is used twice by CIRCLE and the subroutine is called by DRAW. The routine is entered with the co-ordinates of a pixel on the calculator stack. It finds the address of that pixel and plots it, taking account of the status of INVERSE and OVER held in the P-FLAG.

| 22DC | PLOT | CALL | 2307,STK-TO-BC | $Y$ co-ordinate to $B, x$ to $C$ |
| :---: | :---: | :---: | :---: | :---: |
|  |  | CALL | 22E5,PLOT-SUB | The subroutine is called. |
|  |  | JP | 0D4D,TEMPS | Exit, setting temporary colours. |
| 22E5 | PLOT-SUB | LD | (COORDS),BC | The system variable is set. |
|  |  | CALL | 22AA,PIXEL-ADD | Pixel address to HL. |
|  |  | LD | B,A | $B$ will count $A+1$ loops to get a |
|  |  | INC | B | zero to the correct place in A. |
|  |  | LD | A, +FE | The zero is entered. |
| 22F0 | PLOT-LOOP | RRCA |  | Then lined up with the pixel |
|  |  | DJNZ | 22F0,PLOT-LOOP | bit position in the byte. |
|  |  | LD | B,A | Then copied to B. |
|  |  | LD | A,(HL) | The pixel-byte is obtained in A. |
|  |  | LD | C,(P-FLAG) | P-FLAG is obtained and first |
|  |  | BIT | 0,C | tested for OVER. |
|  |  | JR | NZ,22FD,PL-TST-IN | Jump if OVER 1. |
|  |  | AND | B | OVER 0 first makes the pixel zero. |
| 22FD | PL-TST-IN | BIT | 2,C | Test for INVERSE. |
|  |  | JR | NZ,2303,PLOT-END | INVERSE 1 just leaves the pixel as it was (OVER 1) or zero (OVER 0). |
|  |  | XOR | B | INVERSE 0 leaves the pixel |


|  | CPL |  |
| :--- | :--- | :--- |
| 2303 |  | complemented (OVER 1) or <br> 1 (OVER 0). |
|  | PLOT-END | LD |

## THE 'STK-TO-BC' SUBROUTINE

This subroutine loads two floating point numbers into the BC register pair. It is thus used to pick up parameters in the range $+00-+\mathrm{FF}$. It also obtains in DE the 'diagonal move' values $(+/-1,+/-1)$ which are used in the line drawing subroutine of DRAW.

2307

| STK-TO-BC | CALL | 2314,STK-TO-A |
| :--- | :--- | :--- |
|  | LD | B,A |
|  | PUSH | BC |
|  | CALL | 2314, STK-TO-A |
|  | LD | E,C |
|  | POP | BC |
|  | LD | D,C |
|  | LD | C,A |
|  | RET |  |

First number to A .
Hence to B. Save it briefly.
Second number to A.
Its sign indicator to E .
Restore first number.
Its signs indicator to D.
Second number to C.
$B C, D E$ are now as required.

## THE 'STK-TO-A' SUBROUTINE

This subroutine loads the A register with the floating point number held at the top of the calculator stack. The number must be in the range 00-FF.
2314 STK-TO-A

| CALL | 2DD5,FP-TO-A |
| :--- | :--- |
| JP | C,24F9,REPORT-B |
| LD | C,+01 |
| RET | Z |
| LD | C,+FF |
| RET |  |

Modulus of rounded last value to
STK-TO-A A if possible; else, report error. One to C for positive last value. Return if value was positive.
Else change C to +FF (i.e. minus one). Finished.

## THE 'CIRCLE' COMMAND ROUTINE

This routine draws an approximation to the circle with centre co-ordinates $X$ and $Y$ and radius $Z$. These numbers are rounded to the nearest integer before use. Thus $Z$ must be less than 87.5 , even when $(X, Y)$ is in the centre of the screen. The method used is to draw a series of arcs approximated by straight lines. It is illustrated in the BASIC program in the appendix. The notation of that program is followed here.

CIRCLE has four parts:
I. Tests the radius. If its modulus is less than 1 , just plot $\mathrm{X}, \mathrm{Y}$;
II. Calls CD-PRMS-1 at 2470-24B6, which is used to set the initial parameters for both CIRCLE and DRAW;
III. Sets up the remaining parameters for CIRCLE, including the initial displacement for the first 'arc' (a straight line in fact);
IV. Jumps into DRAW to use the arc-drawing loop at 2420-24FA.

Parts i. to iii. will now be explained in turn
i. 2320-23AA. The radius, say $Z^{\prime}$, is obtained from the calculator stack. Its modulus $Z$ is formed and used from now on. If $Z$ is less than 1 , it is deleted from the stack and the point $X, Y$ is plotted by a jump to PLOT.

| 2320 CIRCLE | RST | 0017,GET-CHAR | Get the present character. |
| :--- | :--- | :--- | :--- |
|  | CP | +2C | Test for comma. |
|  | JP | NZ,1C8A,REPORT-C | If not so, report the error. |
|  | RST | 0020,NEXT-CHAR | Get next character (the radius). |
|  | CALL | 1C82,EXPT-1NUM | Radius to calculator stack. |
|  | CALL | 1BEE,CHECK-END | Move to consider next statement |
|  |  |  | if checking syntax. |
|  | RST | 0028,FP-CALC | Use calculator: the stack holds: |
|  | DEFB | +2A,abs | X, Y, Z |
|  | DEFB | +3D,re-stack | Z is re-stacked; its exponent |
|  | DEFB | +38,end-calc | is therefore available. |
|  | LD | A,(HL) | Get exponent of radius. |


| CP | +81 | Test whether radius less than 1. |
| :--- | :--- | :--- |
| JR | NC,233B,C-R-GRE-1 | If not, jump. |
| RST | 0028,FP-CALC | If less, delete it from the stack. |
| DEFB | +02,delete | The stack holds X, Y. |
| DEFB | +38,end-calc |  |
| JR | 22DC,PLOT | Just plot the point $X, Y$. |

ii. 233B-2346 and the call to CD-PRMS1. 2*PI is stored in mem-5 and CD-PRMS1 is called. This subroutine stores in the B register the number of arcs required for the circle, viz. $\mathrm{A}=4^{*} \mathrm{INT}\left(\mathrm{P} \mathrm{I}^{*} \mathrm{SQR} \mathrm{Z} / 4\right)+4$, hence $4,8,12 \ldots$, up to a maximum of 32 . It also stores in mem- 0 to mem-4 the quantities $2^{\star} \mathrm{P} / / \mathrm{A}, \operatorname{SIN}(\mathrm{PI} / \mathrm{A}), 0, \operatorname{COS}\left(2^{\star} \mathrm{P} / / \mathrm{A}\right)$ and $\operatorname{SIN}\left(2^{\star} \mathrm{P} / / \mathrm{A}\right)$.

## 233B

| C-R-GRE-1 | RST | 0028,FP-CALC |
| :--- | :--- | :--- |
|  | DEFB | +A3,stk-pi/2 |
|  | DEFB | +38,end-calc |
|  | LD | (HL),+83 |
|  | RST | 002,FP-CALC |
|  | DEFB | +C5,st-mem-5 |
|  | DEFB | +02,delete |
|  | DEFB | +38,end-calc |
|  | CALL | 247D,CD-PRMS |

$\mathrm{X}, \mathrm{Y}, \mathrm{Z}, \mathrm{PI} / 2$.
Now increase exponent to 83
hex, changing PI/2 into $2^{*} \mathrm{PI}$.
X, Y, Z, 2*PI.
DEFB +C5,st-mem-5
(2 ${ }^{*} \mathrm{P}$ I is copied to mem-5).
DEFB +38,end-calc
CALL 247D,CD-PRMS
Set the initial parameters.
iii. 2347-2381: the remaining parameters and the jump to DRAW. A test is made to see whether the initial 'arc' length is less than 1 . If it is, a jump is made simply to plot X , Y . Otherwise, the parameters are set: $\mathrm{X}+\mathrm{Z}$ and $\mathrm{X}-\mathrm{Z}^{*} \mathrm{SIN}(\mathrm{PI} / \mathrm{A})$ are stacked twice as start and end point, and copied to COORDS as well; zero and $2^{*} Z^{*}$ SIN (PI/A) are stored in mem-1 and mem-2 as initial increments, giving as first 'arc' the vertical straight line joining $X+Z, y-Z^{*} S I N(P I / A)$ and $X+Z, Y+Z^{*} \operatorname{SIN}(P I / A)$. The arc-drawing loop of DRAW will ensure that all subsequent points remain on the same circle as these two points, with incremental angle $2^{*} \mathrm{P} / / \mathrm{A}$. But it is clear that these 2 points in fact subtend this angle at the point $\mathrm{X}+\mathrm{Z}^{*}(1-\operatorname{COS}(\mathrm{PI} / \mathrm{A}))$, Y not at $\mathrm{X}, \mathrm{Y}$. Hence the end points of each arc of the circle are displaced right by an amount $2^{*}(1-\operatorname{COS}(\mathrm{PI} / \mathrm{A}))$, which is less than half a pixel, and rounds to one pixel at most.

| 2347 | C-ARC-GE1 | PUSH | BC | Save the arc-count in B. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | RST | 0028,FP-CALC | X,Y,Z |
|  |  | DEFB | +31,duplicate | X,Y,Z,Z |
|  |  | DEFB | +E1,get-mem-1 | X,Y,Z,Z,SIN (PI/A) |
|  |  | DEFB | +04,multiply | X,Y,Z,Z*SIN (PI/A) |
|  |  | DEFB | +38,end-calc | $\mathrm{Z}^{*} \mathrm{SIN}(\mathrm{P} / / \mathrm{A})$ is half the initial |
|  |  | LD | $\mathrm{A}, \mathrm{HL}$ ) | 'arc' length; it is tested to see |
|  |  | CP | +80 | whether it is less than 0.5. |
|  |  | JR | NC,235A,C-ARC-GE1 | If not, the jump is made. |
|  |  | RST | 0028,FP-CALC | Otherwise, Z is deleted from the |
|  |  | DEFB | +02,delete | stack, with the half-arc too; the |
|  |  | DEFB | +02,delete | machine stack is cleared; and a |
|  |  | DEFB | +38,end-calc | jump is made to plot $\mathrm{X}, \mathrm{Y}$. |
|  |  | POP | BC |  |
|  |  | JP | 22DC,PLOT |  |
| 235A |  | RST | 0028,FP-CALC | X,Y,Z,Z* ${ }^{\text {SIN (PI/A) }}$ |
|  |  | DEFB | +C2,st-mem-2 | (Z ${ }^{*}$ SIN (Pl/A) to mem-2 for now). |
|  |  | DEFB | +01, exchange | X,Y,Z*SIN (PI/A), Z |
|  |  | DEFB | +C0,st-mem-0 | X,Y,Z*SIN (PI/A),Z |
|  |  | DEFB | +02,delete X,Y,Z*SIN (PI/A) |  |
|  |  | DEFB | +03, subtract | X, Y-Z*SIN (PI/A) |
|  |  | DEFB | +01, exchange | $Y$ - ${ }^{*} \mathrm{SIN}(P / / A), X$ |
|  |  | DEFB | +E0,get-mem-0 | $Y-Z^{*} \operatorname{SIN}(P / / A), \mathrm{X}, \mathrm{Z}$ |
|  |  | DEFB | +0F,addition | Y - Z*SIN (PI/A), $\mathrm{X}+\mathrm{Z}$ |
|  |  | DEFB | +CO,st-mem-0 | ( $X+Z$ is copied to mem-0) |
|  |  | DEFB | +01, exchange | X +Z, Y-Z*SIN (P/A) |
|  |  | DEFB | +31,duplicate | $X+Z, Y-Z * S I N(P I / A), Y-Z * S I N$ (Pl/A) |
|  |  | DEFB | +E0,get-mem-0 | sa,sb,sb,sa |


| DEFB | +01,exchange | sa,sb,sa,sb |
| :--- | :--- | :--- |
| DEFB | +31,duplicate | sa,sb,sa,sb,sb |
| DEFB | +E0,get-mem-0 | sa,sb,sa,sb,sb,sa |
| DEFB | +A0,stk-zero | sa,sb,sa,sb,sb,sa,0 |
| DEFB | +C1,st-mem-1 | (mem-1 is set to zero) |
| DEFB | +02,delete | sa,sb,sa,sb,sb,sa |
| DEFB | +38,end-calc |  |

(Here sa denotes $X+Z$ and sb denotes $Y-Z * S I N(P I / A)$ ).

| INC | (mem-2-1st) | Incrementing the exponent byte <br> of mem-2 sets mem-2 to |
| :--- | :--- | :--- |
| CALL | 1E94,FIND-INT1 | $2^{*} Z^{*} \mathrm{SIN}(\mathrm{Pl} / \mathrm{A})$. <br> The last value X+Z is moved <br> from the stack to A and copied |
| LD | L,A | to L. |
| PUSH | HL | It is saved in HL. |
| CALL | $1 E 94$, FIND-INT1 | Y $-Z^{*} \mathrm{SIN}(\mathrm{PI} / \mathrm{A})$ goes from the <br> stack to A and is copied to H. |
| POP | HL | HL now holds the initial point. |
| LD | $\mathrm{H}, \mathrm{A}$ | It is copied to COORDS. |
| LD | (COORDS),HL | The arc-count is restored. |
| POP | BC | The jump is made to DRAW. |

(The stack now holds $X+Z, Y-Z^{*}$ SIN (PI/A), $\left.Y-Z^{*} \operatorname{SIN}(P I / A), X+Z\right)$.

## THE DRAW COMMAND ROUTINE

This routine is entered with the co-ordinates of a point $\mathrm{XO}, \mathrm{Y} 0$, say, in COORDS. If only two parameters $\mathrm{X}, \mathrm{Y}$ are given with the DRAW command, it draws an approximation to a straight line from the point $\mathrm{X} 0, \mathrm{Y} 0$ to $\mathrm{X} 0+\mathrm{X}, \mathrm{Y} 0+\mathrm{Y}$. If a third parameter G is given, it draws an approximation to a circular arc from $\mathrm{X} 0, \mathrm{Y} 0$ to $\mathrm{X} 0+\mathrm{X}, \mathrm{Y} 0+\mathrm{Y}$ turning anti-clockwise through an angle G radians.

The routine has four parts:
I. Just draws a line if only 2 parameters are given or if the diameter of the implied circle is less than 1 ;
II. Calls CD-PRMS1 at 247D-24B6 to set the first parameters;
III. Sets up the remaining parameters, including the initial displacements for the first arc;
IV. Enters the arc-drawing loop and draws the arc as a series of smaller arcs approximated by straight lines, calling the linedrawing subroutine at 24B7-24FA as necessary.

Two subroutines, CD-PRMS1 and DRAW-LINE, follow the main routine. The above 4 parts of the main routine will now be treated in turn.
i. If there are only 2 parameters, a jump is made to LINE-DRAW at 2477. A line is also drawn if the quantity $Z=(A B S X+A B S Y) / A B S$ $\operatorname{SIN}(\mathrm{G} / 2)$ is less than 1. Z lies between 1 and 1.5 times the diameter of the implied circle. In this section mem- 0 is set to $\operatorname{SIN}(\mathrm{G} / 2)$, mem-1 to $Y$, and mem-5 to $G$.

| 2382 | DRAW | RST | 0018,GET-CHAR | Get the current character. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | CP | +2C | If it is a comma, |
|  |  | JR | Z,238D,DR-3-PRMS | then jump. |
|  |  | CALL | 1BEE,CHECK-END | Move on to next statement if checking syntax. |
|  |  | JP | 2477,LINE-DRAW | Jump to just draw the line. |
| 238D | DR-3-PRMS | RST | 0020,NEXT-CHAR | Get next character (the angle). |
|  |  | CALL | 1C82,EXPT-1NUM | Angle to calculator stack. |
|  |  | CALL | 1BEE,CHECK-END | Move on to next statement if checking syntax. |
|  |  | RST | 0028,FP-CALC | $\mathrm{X}, \mathrm{Y}, \mathrm{G}$ are on the stack. |
|  |  | DEFB | +C5,st-mem-5 | ( $G$ is copied to mem-5) |
|  |  | DEFB | +A2,stk-half | X, Y, G, 0.5 |
|  |  | DEFB | +04,multiply | X, Y, G/2 |
|  |  | DEFB | +1F,sin | X, Y, SIN (G/2) |
|  |  | DEFB | +31,duplicate | X, Y, SIN (G/2), SIN (G/2) |


|  |  | DEFB | +30, not | X, Y, SIN (G/2), (0/1) |
| :---: | :---: | :---: | :---: | :---: |
|  |  | DEFB | +30, not | X, Y, SIN (G/2), (1/0) |
|  |  | DEFB | +00,jump-true | X, Y, SIN (G/2) |
|  |  | DEFB | +06,to DR-SIN-NZ | (If SIN (G/2)=0 i.e. G = 2*N*PI |
|  |  | DEFB | +02,delete | just draw a straight line). |
|  |  | DEFB | +38,end-calc | X, Y |
|  |  | JP | 2477,LINE-DRAW | Line $\mathrm{X} 0, \mathrm{Y} 0$ to $\mathrm{X} 0+\mathrm{X}, \mathrm{Y} 0+\mathrm{Y}$. |
| 23A3 | DR-SIN-NZ | DEFB | +C0,st-mem-0 | ( $\mathrm{SIN}(\mathrm{G} / 2)$ is copied to mem-0) |
|  |  | DEFB | +02, delete | $\mathrm{X}, \mathrm{Y}$ are now on the stack. |
|  |  | DEFB | +C1,st-mem-1 | ( Y is copied to mem-1). |
|  |  | DEFB | +02,delete | $X$ |
|  |  | DEFB | +31, duplicate | X, X |
|  |  | DEFB | +2A, abs | $X, X^{\prime}\left(X^{\prime}=A B S ~ X\right)$ |
|  |  | DEFB | +E1,get-mem-1 | $X, X^{\prime}, Y$ |
|  |  | DEFB | +01, exchange | X, Y, X' |
|  |  | DEFB | +E1,get-mem-1 | X, Y, X', Y |
|  |  | DEFB | +2A, abs | $X, \mathrm{Y}, \mathrm{X}^{\prime}, \mathrm{Y}^{\prime}\left(\mathrm{Y}^{\prime}=\mathrm{ABS} \mathrm{Y}\right)$ |
|  |  | DEFB | +0F,addition | $X, Y, X^{\prime}+Y^{\prime}$ |
|  |  | DEFB | +E0,get-mem-0 | X, Y, X'+Y', SIN (G/2) |
|  |  | DEFB | +05,division | X, Y, ( $\left.\mathrm{X}^{\prime}+\mathrm{Y}^{\prime}\right) / \mathrm{SIN}(\mathrm{G} / 2)=\mathrm{Z}^{\prime}$, say |
|  |  | DEFB | +2A, abs | $X, Y, Z(Z=A B S ~ Z ') ~$ |
|  |  | DEFB | +E0,get-mem-0 | X, Y, Z, SIN (G/2) |
|  |  | DEFB | +01,exchange | X, Y, SIN (G/2), Z |
|  |  | DEFB | +3D,re-stack | ( $Z$ is re-stacked to make sure |
|  |  | DEFB | +38, end-calc | that its exponent is available). |
|  |  | LD | A, (HL) | Get exponent of $Z$. |
|  |  | CP | +81 | If $Z$ is greater than or equal |
|  |  | JR | NC,23C1,DR-PRMS | to 1, jump. |
|  |  | RST | 0028,FP-CALC | X, Y, SIN (G/2), Z |
|  |  | DEFB | +02,delete | $X, Y, \operatorname{SIN}(\mathrm{G} / 2)$ |
|  |  | DEFB | +02,delete | X, Y |
|  |  | DEFB | +38, end-calc | Just draw the line from $\mathrm{X} 0, \mathrm{Y} 0$ |
|  |  | JP | 2477,LINE-DRAW | to $\mathrm{X} 0+\mathrm{X}, \mathrm{Y} 0+\mathrm{Y}$. |

ii. Just calls CD-PRMS1. This subroutine saves in the $B$ register the number of shorter arcs required for the complete arc, viz. $A=4^{*}$ INT ( $\mathrm{G}^{\prime *}$ SQR $\left.\mathrm{Z} / 8\right)+4$, where $\mathrm{G}^{\prime}=\bmod \mathrm{G}$, or 252 if this expression exceeds 252 (as can happen with a large chord and a small angle). So $A$ is $4,8,12, \ldots$, up to 252 . The subroutine also stores in mem-0 to mem-4 the quantities $G / A, S I N\left(G / 2^{*} A\right), 0, \operatorname{COS}(G / A), S I N(G / A)$

23C1 DR-PRMS CALL 247D,CD-PRMS1 The subroutine is called.
iii. Sets up the rest of the parameters as follow. The stack will hold these 4 items, reading up to the top: $\mathrm{X} 0+\mathrm{X}$ and $\mathrm{Y} 0+\mathrm{Y}$ as end of last arc; then XO and YO as beginning of first arc. Mem-0 will hold XO and mem- 5 YO . Mem-1 and mem-2 will hold the initial displacements for the first arc, $U$ and $V$; and mem-3 and mem-4 will hold COS (G/A) and SIN (G/A) for use in the arc-drawing loop.

The formulae for U and V can be explained as follows. Instead of stepping along the final chord, of length $L$, say, with displacements X and $Y$, we want to step along an initial chord (which may be longer) of length $L^{*} W$, where $W=S I N\left(G / 2^{*} A\right) / S I N(G / 2)$, with displacements $X^{\star} W$ and $Y^{\star} W$, but turned through an angle $-\left(G / 2-G / 2^{*} A\right)$, hence with true displacements:

```
U = Y*W*SIN (G/2 - G/2*A) + X*W* COS (G/2 - G/2*A)
Y= Y*W*}\operatorname{COS}(G/2-G/2*A)-\mp@subsup{X}{}{*}W*NIN (G/2-G/2*A)
```

These formulae cam be checked from a diagram, using the normal expansion of COS (P-Q) and SIN (P-Q), where $Q=G / 2-G / 2^{*} A$

| $23 C 4$ | PUSH | BC | Save the arc-counter in B. |
| :--- | :--- | :--- | :--- |
|  | RST | 0028, FP-CALC | X,Y,SIN(G/2),Z |
|  | DEFB | +02,delete | X,Y,SIN(G/2) |
|  | DEFB | +E1,get-mem-1 | X,Y,SIN(G/2),SIN(G/2*A) |
|  | DEFB | +01,exchange | X,Y,SIN(G/2*A),SIN(G/2) |
|  | DEFB | +05,division | X,Y,SIN(G/2*A)/SIN(G/2)=W |
|  | DEFB | $+C 1$, st-mem-1 | (W is copied to mem-1). |


| DEFB | +02, delete | X, Y |
| :---: | :---: | :---: |
| DEFB | +01,exchange | Y, X |
| DEFB | +31, duplicate | Y, X, X |
| DEFB | +E1,get-mem-1 | Y,X,X,W |
| DEFB | +04,multiply | $\mathrm{Y}, \mathrm{X}, \mathrm{X}{ }^{+} \mathrm{W}$ |
| DEFB | +C2,st-mem-2 | ( $\mathrm{X}^{*} \mathrm{~W}$ is copied to mem-2). |
| DEFB | +02,delete | Y, X |
| DEFB | +01, exchange | X, Y |
| DEFB | +31, duplicate | X,Y,Y |
| DEFB | +E1, get-mem-1 | $X, Y, Y, W$ |
| DEFB | +04,multiply | $X, Y, Y^{*} W$ |
| DEFB | +E2,get-mem-2 | $X, Y, Y^{*} W, X^{*} W$ |
| DEFB | +E5,get-mem-5 | $X, Y, Y^{*} W, X^{*} W, G$ |
| DEFB | +E0,get-mem-0 | X,Y, $\mathrm{Y}^{*} \mathrm{~W}, \mathrm{X}^{*} \mathrm{~W}, \mathrm{G}, \mathrm{G} / \mathrm{A}$ |
| DEFB | +03,subtract | $X, Y, Y^{*} W, X^{*} W, G-G / A$ |
| DEFB | +A2,stk-half | $X, Y, Y^{*} W, X^{*} W, G-G / A, 1 / 2$ |
| DEFB | +04,multiply | $\mathrm{X}, \mathrm{Y}, \mathrm{Y}^{*} \mathrm{~W}, \mathrm{X}^{*} \mathrm{~W}, \mathrm{G} / 2-\mathrm{G} / 2^{*} \mathrm{~A}=\mathrm{F}$ |
| DEFB | +31, duplicate | $X, Y, Y^{*} W, X^{*} W, F, F$ |
| DEFB | +1F,sin | $X, Y, Y^{*} W, X^{*} W, ~ F, ~ S I N ~ F ~$ |
| DEFB | +C5,st-mem-5 | (SIN F is copied to mem-5). |
| DEFB | +02,delete | $X, Y, Y^{*} W, X^{*} W, F$ |
| DEFB | +20, cos | $\mathrm{X}, \mathrm{Y}, \mathrm{Y}^{*} \mathrm{~W}, \mathrm{X}^{*} \mathrm{~W}, \operatorname{COS} \mathrm{~F}$ |
| DEFB | +C0,st-mem-0 | (COS F is copied to mem-0). |
| DEFB | +02,delete | X, Y, $\mathrm{Y}^{*} \mathrm{~W}, \mathrm{X}^{*} \mathrm{~W}$ |
| DEFB | +C2,st-mem-2 | ( $\mathrm{X}^{*} \mathrm{~W}$ is copied to mem-2). |
| DEFB | +02,delete | $\mathrm{X}, \mathrm{Y}, \mathrm{Y}^{*} \mathrm{~W}$ |
| DEFB | +C1,st-mem-1 | ( $Y^{*} \mathrm{~W}$ is copied to mem-1). |
| DEFB | +E5,get-mem-5 | $X, Y, Y^{*} W, S$ IN $F$ |
| DEFB | +04, multiply | $X, Y, Y^{*} W^{*}$ SIN F |
| DEFB | +E0,get-mem-0 | $X, Y, Y^{*} W^{*}$ SIN F, ${ }^{*}$ W |
| DEFB | +E2,get-mem-2 | $X, Y, Y^{*} W^{*}$ SIN F, $X^{*}$ W, COS F |
| DEFB | +04, multiply | $X, Y, Y^{*} W^{*}$ SIN F, $\mathrm{X}^{*} \mathrm{~W}^{*} \mathrm{COS} F$ |
| DEFB | +0F,addition | $X, Y, Y^{*} W^{*}$ SIN F+${ }^{*} W^{*}$ COS F $=U$ |
| DEFB | +E1,get-mem-1 | $X, Y, U, Y^{*} W$ |
| DEFB | +01, exchange | X,Y, $\mathrm{Y}^{*} \mathrm{~W}, \mathrm{U}$ |
| DEFB | +C1,st-mem-1 | ( U is copied to mem-1) |
| DEFB | +02, delete | $X, Y, Y^{*} W$ |
| DEFB | +E0,get-mem-0 | $X, Y, Y^{*} W, \operatorname{COS} F$ |
| DEFB | +04, multiply | $X, Y, Y^{*} W^{*} \operatorname{COS} F$ |
| DEFB | +E2,get-mem-2 | $X, Y, Y^{*} W^{*} \operatorname{COS} F, X^{*}$ W |
| DEFB | +E5,get-mem-5 | $X, Y, Y^{*} W^{*} \operatorname{COS} F, X^{*} W$, SIN F |
| DEFB | +04,multiply | $X, Y, Y^{*} W^{*} \operatorname{COS~F} F, X^{*} W^{*}$ SIN F |
| DEFB | +03,subtract | $\begin{aligned} & X, Y, Y^{*} W^{*} \operatorname{COS} F-X^{*} W^{*} \operatorname{SIN} \\ & F=V \end{aligned}$ |
| DEFB | +C2,st-mem-2 | ( V is copied to mem-2). |
| DEFB | +2A,abs | $\mathrm{X}, \mathrm{Y}, \mathrm{V}^{\prime}\left(\mathrm{V}^{\prime}=\mathrm{ABS} \mathrm{V}\right)$ |
| DEFB | +E1,get-mem-1 | X, Y, V', U |
| DEFB | +2A,abs | $\mathrm{X}, \mathrm{Y}, \mathrm{V}^{\prime}, \mathrm{U}$ ( $\left.\mathrm{U}^{\prime}=\mathrm{ABS} \mathrm{U}\right)$ |
| DEFB | +0F,addition | $\mathrm{X}, \mathrm{Y}, \mathrm{U}{ }^{\prime}+\mathrm{V}^{\prime}$ |
| DEFB | +02,delete | X, Y |
| DEFB | +38,end-calc | (DE now points to U' + V'). |
| LD | A,(DE) | Get exponent of $\mathrm{U}^{\prime}+\mathrm{V}^{\prime}$ |
| CP | +81 | If $\mathrm{U}^{\prime}+\mathrm{V}$ ' is less than 1 , just |
| POP | BC | tidy the stack and draw the line |
| JP | C,2477,LINE-DRAW | from $\mathrm{X} 0, \mathrm{Y} 0$ to $\mathrm{X} 0+\mathrm{X}, \mathrm{Y} 0+\mathrm{Y}$. |
| PUSH | BC | Otherwise, continue with the |
| RST | 0028,FP-CALC | parameters: $\mathrm{X}, \mathrm{Y}$, on the stack. |
| DEFB | +01,exchange | Y, X |
| DEFB | +38,end-calc |  |
| LD | A,(COORDS-lo) | Get XO into A and so |
| CALL | 2D28,STACK-A | on to the stack. |
| RST | 0028,FP-CALC | Y, X, XO |


| DEFB | +C0,st-mem-0 | ( X 0 is copied to mem-0). |
| :---: | :---: | :---: |
| DEFB | +0F,addition | $\mathrm{Y}, \mathrm{XO}+\mathrm{X}$ |
| DEFB | +01, exchange | X $0+\mathrm{X}, \mathrm{Y}$ |
| DEFB | +38,end-calc |  |
| LD | A,(COORDS-hi) | Get YO into $A$ and so |
| CALL | 2D28,STACK-A | on to the stack. |
| RST | 0028,FP-CALC | X0+X, Y, Y0 |
| DEFB | +C5,st-mem-5 | ( YO is copied to mem-5). |
| DEFB | +0F,addition | $\mathrm{X} 0+\mathrm{X}, \mathrm{Y} 0+\mathrm{Y}$ |
| DEFB | +E0,get-mem-0 | X0+X, Y0+Y, X0 |
| DEFB | +E5,get-mem-5 | X0+X, Y0+Y, X0, Y0 |
| DEFB | +38,end calc |  |
| POP | BC | Restore the arc-counter in B. |

iv. The arc-drawing loop. This is entered at 2439 with the co-ordinates of the starting point on top of the stack, and the initial displacements for the first arc in mem-1 and mem-2. It uses simple trigonometry to ensure that all subsequent arcs will be drawn to points that lie on the same circle as the first two, subtending the same angle at the centre. It can be shown that if 2 points $\mathrm{X} 1, \mathrm{Y} 1$ and X 2 , Y 2 lie on a circle and subtend an angle N at the centre, which is also the origin of co-ordinates, then $\mathrm{X} 2=\mathrm{X} 1^{*} \mathrm{COS} N-\mathrm{Y} 1^{*} \mathrm{SIN} \mathrm{N}$, and $Y 2=X 1^{*}$ SIN $N+Y 1^{*}$ COS $N$. But because the origin is here at the increments, say $U n=X n+1-X n$ and $V n=Y n+1-Y n$, thus achieving the desired result. The stack is shown below on the $(\mathrm{n}+1)$ th pass through the loop, as Xn and Yn are incremented by Un and Vn , after these are obtained from Un-1 and $\mathrm{Vn}-1$. The 4 values on the top of the stack at 2425 are, in DRAW, reading upwards, X0 X , $Y 0+Y, \mathrm{Xn}$ and Yn but to save space these are not shown until 2439. For the initial values in CIRCLE, see the end of CIRCLE, above. In CIRCLE too, the angle $G$ must be taken to be 2*PI.

| 2420 | DRW-STEPS | DEC | B | B counts the passes through the loop. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | JR | Z,245F,ARC-END | Jump when B has reached zero. |
|  |  | JR | 2439,ARC-START | Jump into the loop to start. |
| 2425 | ARC-LOOP | RST | 0028,FP-CALC | (See text above for the stack). |
|  |  | DEFB | +E1,get-mem-1 | Un-1 |
|  |  | DEFB | +31, duplicate | Un-1,Un-1 |
|  |  | DEBF | +E3,get-mem-3 | Un-1,Un-1, COS(G/A) |
|  |  | DEFB | +04,multiply | Un-1,Un-1*COS(G/A) |
|  |  | DEFB | +E2,get-mem-2 | Un-1,Un-1*COS(G/A), Vn-1 |
|  |  | DEFB | +E4,get-mem-4 | Un-1,Un-1*COS(G/A),Vn-1, SIN(G/A) |
|  |  | DEFB | +04,multiply | Un-1,Un-1* $\operatorname{COS}(\mathrm{G} / \mathrm{A}), \mathrm{Vn}-1^{*}$ SIN(G/A) |
|  |  | DEFB | +03,subtract | $\begin{aligned} & U n-1, U n-1^{*} \operatorname{COS}(G / A)-V n-1^{*} \\ & \operatorname{SIN}(G / A)=U n \end{aligned}$ |
|  |  | DEFB | +C1,st-mem-1 | ( Un is copied to mem-1). |
|  |  | DEFB | +02,delete | Un-1 |
|  |  | DEFB | +E4,get-mem-4 | Un-1,SIN(G/A) |
|  |  | DEFB | +04,multiply | Un-1*SIN(G/A) |
|  |  | DEFB | +E2,get-mem-2 | Un-1*SIN(G/A), Vn-1 |
|  |  | DEFB | +E3,get-mem-3 | Un-1*SIN(G/A),Vn-1, COS(G/A) |
|  |  | DEFB | +04,multiply | Un-1*SIN(G/A), Vn-1*COS(G/A) |
|  |  | DEFB | +0F, addition | $\begin{aligned} & U n-1 * \operatorname{SIN}(\mathrm{G} / \mathrm{A})+V n-1^{*} \mathrm{COS} \\ & (\mathrm{G} / \mathrm{A})=\mathrm{Vn} \end{aligned}$ |
|  |  | DEFB | +C2,st-mem-2 | ( Vn is copied to mem-2). |
|  |  | DEFB | +02,delete | (As noted in the text, the stack |
|  |  | DEFB | +38,end-calc | in fact holds $\mathrm{X} 0+\mathrm{X}, \mathrm{Y} 0+\mathrm{Y}, \mathrm{Xn}$ and $Y n$ ). |
| 2439 | ARC-START | PUSH | BC | Save the arc-counter. |
|  |  | RST | 0028,FP-CALC | X0+X, Y0+y, Xn, Yn |
|  |  | DEFB | +C0,st-mem-0 | ( Yn is copied to mem-0). |
|  |  | DEFB | +02,delete | X0+X, Y0+Y, Xn |
|  |  | DEFB | +E1,get-mem-1 | X0+X, Y0+Y, Xn, Un |
|  |  | DEFB | +0F, addition | $X 0+X, Y 0+Y, X n+U n=X n+1$ |



THE 'INITIAL PARAMETERS' SUBROUTINE
This subroutine is called by both CIRCLE and DRAW to set their initial parameters. It is called by CIRCLE with $\mathrm{X}, \mathrm{Y}$ and the radius Z on the top of the stack, reading upwards. It is called by DRAW with its own $\mathrm{X}, \mathrm{Y}, \mathrm{SIN}(\mathrm{G} / 2)$ and Z , as defined in DRAW i. above, on the top of the stack. In what follows the stack is only shown from $Z$ upwards.
The subroutine returns in B the arc-count A as explained in both CIRCLE and DRAW above, and in mem-0 to mem- 5 the quantities $\mathrm{G} / \mathrm{A}, \operatorname{SIN}\left(\mathrm{G} / 2^{*} \mathrm{~A}\right), 0, \operatorname{COS}(\mathrm{G} / \mathrm{A}), \operatorname{SIN}(\mathrm{G} / \mathrm{A})$ and G . For a circle, G must be taken to be equal to $2^{*} \mathrm{PI}$.

247D CD-PRMS1 | RST | 0028,FP-CALC | Z |  |
| :--- | :--- | :--- | :--- |
|  | DEFB | +31,duplicate | Z, Z |
|  | DEFB | +28,sqr | Z, SQR Z |
|  | DEFB | +34,stk-data | Z, SQR Z, 2 |

|  |  | DEFB | +32,exponent +82 |  |
| :---: | :---: | :---: | :---: | :---: |
|  |  | DEFB | +00,(+00,+00,+00) |  |
|  |  | DEFB | +01, exchange | Z, 2, SQR Z |
|  |  | DEFB | +05,division | Z, 2/SQR Z |
|  |  | DEFB | +E5,get-mem-5 | Z, 2/SQR Z, G |
|  |  | DEFB | +01, exchange | Z, G, 2/SQR Z |
|  |  | DEFB | +05, division | Z, G*SQR Z/2 |
|  |  | DEFB | +2A,abs | Z, G'*SQR Z/2 (G' = mod G) |
|  |  | DEFB | +38,end-calc | Z, G'*SQR Z/2 = A1, say |
|  |  | CALL | 2DD5,FP-TO-A | A1 to A from the stack, if possible. |
|  |  | JR | C,2495,USE-252 | If A1 rounds to 256 or more, use 252. |
|  |  | AND | +FC | $4^{*} \operatorname{INT}(\mathrm{~A} 1 / 4)$ to A . |
|  |  | ADD | A,+04 | Add 4, giving the arc-count A. |
|  |  | JR | NC,2497,DRAW-SAVE | Jump if still under 256. |
| 2495 | USE-252 | LD | A,+FC | Here, just use 252 decimal. |
| 2497 | DRAW-SAVE | PUSH | AF | Now save the arc-count. |
|  |  | CALL | 2D28,STACK-A | Copy it to calculator stack too. |
|  |  | RST | 0028,FP-CALC | Z, A |
|  |  | DEFB | +E5,get-mem-5 | Z, A, G |
|  |  | DEFB | +01, exchange | Z, G, A |
|  |  | DEFB | +05,division | Z, G/A |
|  |  | DEFB | +31, duplicate | Z,G/A, G/A |
|  |  | DEFB | +1F,sin | Z, G/A, SIN (G/A) |
|  |  | DEFB | +C4,st-mem-4 | ( $\mathrm{SIN}(\mathrm{G} / \mathrm{A}$ ) is copied to mem-4). |
|  |  | DEFB | +02,delete | Z, G/A |
|  |  | DEFB | +31,duplicate | Z, G/A, G/A |
|  |  | DEFB | +A2,stk-half | Z, G/A, G/A, 0.5 |
|  |  | DEFB | +04,multiply | Z, G/A, G/2*A |
|  |  | DEFB | +1F,sin | $\mathrm{Z}, \mathrm{G} / \mathrm{A}, \mathrm{SIN}\left(\mathrm{G} / 2^{*} \mathrm{~A}\right)$ |
|  |  | DEFB | +C1,st-mem-1 | $\left(\operatorname{SIN}\left(G / 2^{*} A\right)\right.$ is copied to mem-1). |
|  |  | DEFB | +01, exchange | Z, SIN (G/2*A), G/A |
|  |  | DEFB | +C0,st-mem-0 | ( $\mathrm{G} / \mathrm{A}$ is copied to mem-0). |
|  |  | DEFB | +02, delete | Z, SIN (G/2*A) = S |
|  |  | DEFB | +31,duplicate | Z, S, S |
|  |  | DEFB | +04,multiply | Z, S*S |
|  |  | DEFB | +31,duplicate | Z, S*S, S*S |
|  |  | DEFB | +0F,addition | Z, 2*S*S |
|  |  | DEFB | +A1,stk-one | Z, 2* ${ }^{*}$ S, 1 |
|  |  | DEFB | +03, subtract | Z, 2*S*S - 1 |
|  |  | DEFB | +1B,negate | $\mathrm{Z}, 1-2^{*} \mathrm{~S}^{*} \mathrm{~S}=\operatorname{COS}(\mathrm{G} / \mathrm{A})$ |
|  |  | DEFB | +C3,st-mem-3 | ( $\mathrm{COS}(\mathrm{G} / \mathrm{A})$ is copied to mem-3) |
|  |  | DEFB | +02,delete | Z |
|  |  | DEFB | +38,end-calc |  |
|  |  | $\begin{aligned} & \text { POP } \\ & \text { RET } \end{aligned}$ | BC | Restore the arc-count to B . Finished. |

## THE LINE-DRAWING SUBROUTINE

This subroutine is called by DRAW to draw an approximation to a straight line from the point XO, Y0 held in COORDS to the point $X 0+\mathrm{X}, \mathrm{Y} 0+\mathrm{Y}$, where the increments X and Y are on the top of the calculator stack. The subroutine was originally intended for the ZX 80 and ZX81 8K ROM, and it is described in a BASIC program on page 121 of the ZX81 manual. It is also illustrated here in the Circle program in the appendix.
The method is to intersperse as many horizontal or vertical steps as are needed among a basic set of diagonal steps, using an algorithm that spaces the horizontal or vertical steps as evenly as possible.

| 24B7 DRAW-LINE | CALL | 2307,STK-TO-BC |
| :--- | :--- | :--- | | ABS Y to B; ABS X to C; |
| :--- |
| SGN Y to D; SGN X to E. |


|  |  | JR | NC,24C4,DL-X-GE-Y | smaller goes to L , and the |
| :---: | :---: | :---: | :---: | :---: |
|  |  | LD | L,C | larger (later) goes to H . |
|  |  | PUSH | DE | Save diag. step ( $\pm 1, \pm 1$ ) in DE. |
|  |  | XOR | A | Insert a vertical step ( $\pm 1,0$ ) |
|  |  | LD | E,A | into DE (D holds SGN Y). |
|  |  | JR | 24CB,DL-LARGER | Now jump to set H . |
| 24C4 | DL-X-GE-Y | OR | C | Return if ABS X and ABS Y |
|  |  | RET | Z | are both zero. |
|  |  | LD | L,B | The smaller (ABS $Y$ here) goes to L . |
|  |  | LD | B, C | ABS $X$ to $B$ here, for H . |
|  |  | PUSH | DE | Save the diagonal step here too. |
|  |  | LD | D, +00 | Hor. step ( $0, \pm 1$ ) to DE here. |
| 24CB | DL-LARGER | LD | H,B | Larger of ABS X, ABS Y to $H$ |

The algorithm starts here. The larger of $A B S X$ and $A B S Y$, say $H$, is put into $A$ and reduced to $I N T(H / 2)$. The $H$ - L horizontal or vertical steps and $L$ diagonal steps are taken (where $L$ is the smaller of ABS $X$ and $A B S Y$ ) in this way: $L$ is added to $A$; if $A$ now equals or exceeds H , it is reduced by H and a diagonal step is taken; otherwise a horizontal or vertical step is taken. This is repeated H times ( B also holds H). Note that meanwhile the exchange registers H' and L' are used to hold COORDS.

|  |  | $\begin{aligned} & \text { LD } \\ & \text { RRA } \end{aligned}$ | A, B | $B$ to $A$ as well as to $H$. A starts at INT (H/2). |
| :---: | :---: | :---: | :---: | :---: |
| 24CE | D-L-LOOP | ADD | A,L | $L$ is added to $A$. |
|  |  | JR | C,24D4,D-L-DIAG | If 256 or more, jump - diag. step. |
|  |  | CP | H | If $A$ is less than $H$, jump for |
|  |  | JR | C,24DB,D-L-HR-VT | horizontal or vertical step. |
| 24D4 | D-L-DIAG | SUB | H | Reduce A by H. |
|  |  | LD | C,A | Restore it to C. |
|  |  | EXX |  | Now use the exchange resisters. |
|  |  | POP | BC | Diag. step to $\mathrm{B}^{\prime} \mathrm{C}^{\prime}$. |
|  |  | PUSH | BC | Save it too. |
|  |  | JR | 24DF,D-L-STEP | Jump to take the step. |
| 24DB | D-L-HR-VT | LD | C,A | Save A (unreduced) in C. |
|  |  | PUSH | DE | Step to stack briefly. |
|  |  | EXX |  | Get exchange registers. |
|  |  | POP | BC | Step to B'C' now. |
| 24DF | D-L-STEP | LD | HL,(COORDS) | Now take the step: first, |
|  |  |  |  | COORDS to H'L' as the start point. |
|  |  | LD | A, B | Y-step from B' to A. |
|  |  | ADD | A, H | Add in $\mathrm{H}^{\prime}$. |
|  |  | LD | B,A | Result to B' |
|  |  | LD | A, C | Now the X-step; it will be tested |
|  |  | INC | A | for range ( Y will be tested in PLOT). |
|  |  | ADD | A,L | Add L' to $\mathrm{C}^{\prime}$ in A , jump on |
|  |  | JR | C,24F7,D-L-RANGE | carry for further test. |
|  |  | JR | Z,24F9,REPORT-B | Zero after no carry denotes |
|  |  |  |  | X-position-1, out of range. |
| 24EC | D-L-PLOT | DEC | A | Restore true value to A . |
|  |  | LD | C, A | Value to C ' for plotting. |
|  |  | CALL | 22E5,PLOT-SUB | Plot the step. |
|  |  | EXX |  | Restore main registers. |
|  |  | LD | A,C | $C$ back to $A$ to continue algorithm. |
|  |  | DJNZ | 24CE,D-L-LOOP | Loop back for 8 steps (i.e. H steps). |
|  |  | POP | DE | Clear machine stack. |
|  |  | RET |  | Finished. |

Report B - Integer out of range
24F9 REPORT-B RST 0008,ERROR-1 Call the error handling

## EXPRESSION EVALUATION

## THE 'SCANNING' SUBROUTINE

This subroutine is used to produce an evaluation result of the 'next expression'.
The result is returned as the 'last value' on the calculator stack. For a numerical result, the last value will be the actual floating point number. However, for a string result the last value will consist of a set of parameters. The first of the five bytes is unspecified, the second and third bytes hold the address of the start of the string and the fourth and fifth bytes hold the length of the string

Bit 6 of FLAGS is set for a numeric result and reset for a string result
When a next expression consists of only a single operand, e.g. ... A ..., ... RND ..., ... A\$ (4, 3 TO 7) ... , then the last value is simply the value that is obtained from evaluating the operand

However when the next expression contains a function and an operand, e.g. ... CHR\$ A..., ... NOT A ... , SIN 1 ..., the operation code of the function is stored on the machine stack until the last value of the operand has been calculated. This last value is then subjected to the appropriate operation to give a new last value.

In the case of there being an arithmetic or logical operation to be performed, e.g. ... A+B ... , $A^{*} B \ldots, \ldots A=B \ldots$, then both the last value of the first argument and the operation code have to be kept until the last value of the second argument has been found. Indeed the calculation of the last value of the second argument may also involve the storing of last values and operation codes whilst the calculation is being performed.

It can therefore be shown that as a complex expression is evaluated, e.g. ... CHR $\$\left(T+A-26^{*} I N T((T+A) / 26)+65\right) \ldots$, a hierarchy of operations yet to be performed is built up until the point is reached from which it must be dismantled to produce the final last value.

Each operation code has associated with it an appropriate priority code and operations of higher priority are always performed before those of lower priority.

The subroutine begins with the A register being set to hold the first character of the expression and a starting priority marker - zero being put on the machine stack.


Four subroutines follow; they are called by routines from the scanning function table. The first one, the 'scanning quotes subroutine', is used by S-QUOTE to check that every string quote is matched by another one.

| S-QUOTE-S | CALL | $0074, \mathrm{CH}-\mathrm{ADD}+1$ |
| :--- | :--- | :--- |
|  | INC | BC |
|  | CP | +0 l |
|  | JP | Z, 1C8A,REPORT-C |
|  | CP | +22 |
|  | JR | NZ,250F,S-QUOTE-S |
|  | CALL | $0074, \mathrm{CH}-A D D+1$ |
|  | CP | +22 |
|  | RET |  |

Point to the next character. Increase the length count by one.
Is it a carriage return?
Report the error if so.
Is it another "'"?
Loop back if it is not.
Point to next character; set zero
flag if it is another "'"
Finished.

The next subroutine, the 'scanning: two co-ordinates' subroutine, is called by S-SCREEN\$, S-ATTR and S-POINT to make sure the required two co-ordinates are given in their proper form.

0020, NEXT-CHAR
+28

Fetch the next character.
Is it a '('?

|  | JR | NZ,252D,S-RPORT-C | Report the error if it is not. |
| :--- | :--- | :--- | :--- |
|  | CALL | 1C79,NEXT-2NUM | Co-ordinates to calculator stack. |
|  | RST | 0018,GET-CHAR | Fetch the current character. |
| CP | +29 | Is it a ')'? |  |
| 252D S-RPORT-C | JP | NZ,1C8A,REPORT-C | Report the error if it is not. |

## THE 'SYNTAX-Z' SUBROUTINE

At this point the 'SYNTAX-Z' subroutine is interpolated. It is called 32 times, with a saving of just one byte each call. A simple test of bit 7 of FLAGS will give the zero flag reset during execution and set during syntax checking.
i.e. SYNTAX gives $Z$ set.
2530 SYNTAX-Z BIT (FLAGS) Test bit 7 of FLAGS.

The next subroutine is the 'scanning SCREEN\$ subroutine', which is used by S-SCREENS\$ to find the character that appears at line $x$, column y of the screen. It only searches the character set 'pointed to' to CHARS.

Note: This is normally the characters +20 (space) to +7 F (©) although the user can alter CHARS to match for other characters, including user-defined graphics.



Note: This exit, via STK-STO-\$, is a mistake as it leads to 'double storing' of the string result (see S-STRING, 25DB). The instruction line should be 'RET'.

The last of these four subroutines is the 'scanning attributes subroutine'. It is called by S-ATTR to return the value of ATTR ( $\mathrm{x}, \mathrm{y}$ ) which codes the attributes of line x , column y on the television screen.

| S-ATTR-S | CALL | 2307,STK-TO-BC |
| :--- | :--- | :--- |
|  | LD | A,C |
|  | RRCA |  |
|  | RRCA |  |
|  | RRCA |  |
|  | LD | C,A |
|  | AND | +E0 |
|  | XOR | B |
|  | LD | L,A |
|  |  |  |
|  | LD | A,C |
|  | AND | +03 |
|  | XOR | +58 |
|  | LD | H,A |
|  | LD | A,(HL) |
|  |  |  |
|  | JP | 2D28,STACK-A |

$x$ to $C$, $y$ to $B$. Again, $0<=x<=23$
decimal; $0<=y<=31$ decimal. x is copied to A and the number $32(\text { decimal })^{*} x(\bmod 8)+y$ is formed in A and copied to L. $32^{*} x(\bmod 8)+$ INT $(x / 8)$ is also copied to C.

L holds low byte of attribute address.
$32^{*} x(\bmod 8)+$ INT $(x / 8)$ is copied to A.
88 (decimal)+INT ( $\mathrm{x} / 8$ ) is
formed in A and copied to H . H holds high byte of attribute address.
The attribute byte is copied to A.

Exit, stacking the required byte.

## THE SCANNING FUNCTION TABLE

This table contains 8 functions and 4 operators. It thus incorporates 5 new Spectrum functions and provides a neat way of accessing some functions and operators which already existed on the ZX81.

| location | code | offset | name | address of handling routine |
| :---: | :---: | :---: | :---: | :---: |
| 2596 | 22 | 1 C | S-QUOTE | 25B3 |
| 2598 | 28 | 4 F | S-BRACKET | 25E8 |
| 259A | 2 E | F2 | S-DECIMAL | 268D |
| 259C | 2B | 12 | S-U-PLUS | 25AF |
| 259E | A8 | 56 | S-FN | 25F5 |
| 25A0 | A5 | 57 | S-AND | 25F8 |
| 25A2 | A7 | 84 | S-PI | 2627 |
| 25A4 | A6 | 8F | S-INKEY\$ | 2634 |
| 25A6 | C4 | E6 | S-BIN (EQU. S-DECIMAL) | 268D |
| 25A8 | AA | BF | S-SCREEN\$ | 2668 |
| 25AA | AB | C7 | S-ATTR | 2672 |


| $25 A C$ | A9 | CE | S-POINT | $267 B$ |
| :--- | :--- | :--- | :--- | ---: |
| $25 A E$ | 00 |  |  | End-marker |

THE SCANNING FUNCTION ROUTINES

| 25AF | S-U-PLUS | RST <br> JP | 0020,NEXTCHAR <br> 24FF,S-LOOP-1 |
| :--- | :--- | :--- | :--- | | For unary plus, simply move on |
| :--- |
| to the next character and jump |
| back to the main re-entry |
| of SCANNING. |

The 'scanning QUOTE routine': This routine deals with string quotes, whether simple like "name" or more complex like "a ""white"" lie" or the seemingly redundant VAL\$ """a""".

| 25B3 | S-QUOTE | RST | O018,GET-CHAR | Fetch the current character. |
| :--- | :--- | :--- | :--- | :--- |
|  |  | INC | HL | Point to the start of the string. |

Note that the first quote was not counted into the length; the final quote was, and is discarded now. Inside the string, the first, third, fifth, etc., quotes were counted in but the second, fourth, etc., were not.

|  |  | POP | DE | Restore start of copied string. |
| :--- | :--- | :--- | :--- | :--- |
|  | S-STRING | LD | HL,+5C3B | This is FLAGS; this entry point |
|  |  | RES | 6,(HL) | is used whenever bit 6 is to be |
|  |  | BIT | 7,(HL) | reset and a string stacked if exe- |
|  |  | CALL | NZ,2AB2,STK-STO-S | cuting a line. This is done now. |
|  |  | JP | 2712,S-CONT-2 | Jump to continue scanning the |

Note that in copying the string to the work space, every two pairs of string quotes inside the string ("") have been reduced to one pair of string quotes(").

| 25E8 | S-BRACKET | RST 0020,NEXT-CHAR |
| :--- | :--- | :--- | :--- |
|  | CALL | 24FB,SCANNING |
|  | CP | +29 |
|  |  |  |
|  | JP | NZ,1C8A,REPORT-C |
|  | RST | 0020, NEXT-CHAR |

The 'scanning BRACKET
routine' simply gets the
character and calls SCANNING recursively.
Report the error if no matching
bracket; then continue scanning.

|  | JP | 2712,S-CONT-2 |
| :--- | :--- | :--- |
| 25F5 S-FN | JP | 27BD,S-FN-SBRN The 'scanning FN routine'. |

This routine, for user-defined functions, just jumps to the 'scanning FN subroutine'.

| 25F8 | S-RND | $\begin{aligned} & \text { CALL } \\ & \text { JR } \end{aligned}$ | 2530,SYNTAX-Z <br> Z,2626,S-RND-END | Unless syntax is being checked, jump to calculate a random number. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | LD | BC,(SEED) | Fetch the current value of SEED. |
|  |  | CALL | 2D2B,STACK-BC | Put it on the calculator stack. |
|  |  | RST | 0028,FP-CALC | Now use the calculator, |
|  |  | DEFB | +A1,stk-one | The 'last value' is now |
|  |  | DEFB | +0F,addition | SEED+1. |
|  |  | DEFB | +34,stk-data | Put the decimal number 75 |
|  |  | DEFB | +37,exponent+87 | on the calculator stack. |
|  |  | DEFB | +16,(+00,+00,+00) |  |
|  |  | DEFB | +04,multiply | 'last value' (SEED+1)*75. |
|  |  | DEFB | +34,stk-data | See STACK LITERALS to see |
|  |  | DEFB | +80,(four bytes) | how bytes are expanded so as to |
|  |  | DEFB | +41, exponent +91 | put the decimal number 65537 |
|  |  | DEFB | +00,+00,+80,(+00) | on the calculator stack. |
|  |  | DEFB | +32,n-mod-m | Divide (SEED+1)*75 by 65537 to give a 'remainder' and an 'answer'. |
|  |  | DEFB | +02, delete | Discard the 'answer'. |
|  |  | DEFB | +A1,stk-one | The 'last value' is now |
|  |  | DEFB | +03, subtract | 'remainder' - 1. |
|  |  | DEFB | +31, duplicate | Make a copy of the 'last value'. |
|  |  | DEFB | +38,end-calc | The calculation is finished. |
|  |  | CALL | 2DA2,FP-TO-BC | Use the 'last value' to give the |
|  |  | LD | (SEED),BC | new value for SEED. |
|  |  | LD | A,(HL) | Fetch the exponent of 'last value'. |
|  |  | AND | A | Jump forward if the exponent is |
|  |  | JR | Z,2625,S-RND-END | zero. |
|  |  | SUB | +10 | Reduce the exponent, i.e. divide |
|  |  | LD | (HL), A | 'last value' by 65536 to give the required 'last value'. |
| 2625 | S-RND-END | JR | 2630,S-PI-END | Jump past the 'PI' routine. |

The 'scanning-PI routine': unless syntax is being checked the value of 'PI' is calculated and forms the 'last value' on the calculator stack.

| 2627 | S-PI | CALL | 2530,SYNTAX-Z | Test for syntax checking. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | JR | Z,2630,S-PI-END | Jump if required. |
|  |  | RST | 0028,FP-CALC | Now use the calculator. |
|  |  | DEFB | +A3,stk-pi/2 | The value of $\mathrm{Pl} / 2$ is put on the |
|  |  | DEFB | +3B,end-calc | calculator stack as the 'last value'. |
|  |  | INC | (HL) | The exponent is incremented thereby doubling the 'last value' giving PI. |
| 2630 | S-PI-END | RST | 0020,NEXT-CHAR | Move on to the next character. |
|  |  | JP | 26C3,S-NUMERIC | Jump forward. |
| 2634 | S-INKEY\$ | LD | BC,+105A | Priority +10 hex, operation |
|  |  | RST | 0020,NEXT-CHAR | code +5 A for the 'read-in' subroutine. |
|  |  | CP | +23 | If next char. is '\#', jump. |
|  |  | JP | Z,270D,S-PUSH-PO | There will be a numerical argument. |
|  |  | LD | HL, +5C3B | This is FLAGS. |
|  |  | RES | 6,(HL) | Reset bit 6 for a string result. |
|  |  | BIT | 7,(HL) | Test for syntax checking. |


|  |  | JR | Z,2665,S-INK\$-EN | Jump if required. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | CALL | 028E,KEY-SCAN | Fetch a key-value in DE. |
|  |  | LD | C,+00 | Prepare empty string; stack it if |
|  |  | JR | NZ,2660,S-IK\$-STK | too many keys pressed. |
|  |  | CALL | 031E,K-TEST | Test the key value; stack empty |
|  |  | JR | NC,2660,S-IK\$-STK | string if unsatisfactory. |
|  |  | DEC | D | +FF to D for L made (bit 3 set). |
|  |  | LD | E,A | Key-value to E for decoding. |
|  |  | CALL | 0333,K-DECODE | Decode the key-value. |
|  |  | PUSH | AF | Save the ASCII value briefly. |
|  |  | LD | BC,+0001 | One space is needed in the work space. |
|  |  | RST | 0030,BC-SPACES | Make it now. |
|  |  | POP | AF | Restore the ASCII value. |
|  |  | LD | (DE), A | Prepare to stack it as a string. |
|  |  | LD | C, +01 | Its length is one. |
| 2660 | S-IK\$-STK | LD | B, +00 | Complete the length parameter. |
|  |  | CALL | 2AB2,STK-STO-\$ | Stack the required string. |
| 2665 | S-INK\$-EN | JP | 2712,S-CONT-2 | Jump forward. |
| 2668 | S-SCREEN\$ | CALL | 2522,S-2-COORD | Check that 2 co-ordinates are given. |
|  |  | CALL | NZ,2535,S-SCRN\$-S | Call the subroutine unless |
|  |  | RST | 0020,NEXT-CHAR | checking syntax; then get next |
|  |  | JP | 25DB,S-STRING | character and jump back. |
| 2672 | S-ATTR | CALL | 2522,5-2-COORD | Check that 2 co-ordinates are given. |
|  |  | CALL | NZ,2580,S-ATTR-S | Call the subroutine unless |
|  |  | RST | 0020,NEXT-CHAR | checking syntax; then get the |
|  |  | JR | 26C3,S-NUMERIC | next character and jump forward. |
| 267B | S-POINT | CALL | 2522,S-2-COORD | Check that 2 co-ordinates are given. |
|  |  | CALL | NZ,22CB,POINT-SUB | Call the subroutine unless |
|  |  | RST | 0020,NEXT-CHAR | checking syntax; then get the |
|  |  | JR | 26C3,S-NUMERIC | next character and jump forward. |
| 2684 | S-ALPHNUM | CALL | 2C88,ALPHANUM | Is the character alphanumeric? |
|  |  | JR | NC,26DF,S-NEGATE | Jump if not a letter or a digit. |
|  |  | CP | +41 | Now jump if it a letter; |
|  |  | JR | NC,26C9,S-LETTER | otherwise continue on into |
|  |  |  |  | S-DECIMAL. |

The 'scanning DECIMAL routine' which follows deals with a decimal point or a number that starts with a digit. It also takes care of the expression 'BIN', which is dealt with in the 'decimal to floating-point' subroutine.

| 268D | S-DECIMAL | CALL | 2530,SYNTAX-Z | Jump forward if a line is |
| :--- | :--- | :--- | :--- | :--- |
|  | (EQU. S-BIN) | JR | NZ,2685,S-STK-DEC | being executed. |

The action taken is now very different for syntax checking and line execution. If syntax is being checked then the floating-point form has to be calculated and copied into the actual BASIC line. However when a line is being executed the floating-point form will always be available so it is copied to the calculator stack to form a 'last value'.

During syntax checking:

| CALL | 2C9B,DEC-TO-FP | The floating-point form is found. |
| :---: | :---: | :---: |
| RST | 0018,GET-CHAR | Set HL to point one past the last digit. |
| LD | BC, +0006 | Six locations are required. |
| CALL | 1655,MAKE-ROOM | Make the room in the BASIC line. |
| INC | HL | Point to the first free space. |
| LD | (HL), +0E | Enter the number marker code |
| INC | HL | Point to the second location. |
| EX | DE,HL | This pointer is wanted in DE. |


|  |  | LD | HL,(STKEND) | Fetch the 'old' STKEND. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | LD | C, +05 | There are 5 bytes to move. |
|  |  | AND | A | Clear the carry flag. |
|  |  | SBC | HL,BC | The 'new' STKEND='old' STKEND -5. |
|  |  | $\begin{aligned} & \text { LD } \\ & \text { LDIR } \end{aligned}$ | (STKEND),HL | Move the floating-point number from the calculator stack to |
|  |  |  |  | the line. |
|  |  | EX | DE,HL | Put the line pointer in HL. |
|  |  | DEC | HL | Point to the last byte added. |
|  |  | CALL | 0077,TEMP-PTR1 | This sets CH-ADD. |
|  |  | JR | 26C3,S-NUMERIC | Jump forward. |
| During line execution: |  |  |  |  |
| 26B5 | S-STK-DEC | RST | 0018,GET-CHAR | Get the current character. |
| 26B6 | S-SD-SKIP | INC |  | Now move on to the next |
|  |  | LD | A, (HL) | character in turn until |
|  |  | CP | +0E | the number marker code |
|  |  | JR | NZ,26B6,S-SD-SKIP | is found. |
|  |  | INC | HL | Point to the first byte of the number. |
|  |  | $\begin{aligned} & \text { CALL } \\ & \text { LD } \end{aligned}$ | 33B4,STACK-NUM (CH-ADD),HL | Move the floating-point number. Set CH-ADD. |

A numeric result has now been identified, coming from RND, PI, ATTR, POINT or a decimal number, therefore bit 6 of FLAGS must be set.

26C3 S-NUMERIC SET 6,(FLAGS) Set the numeric marker flag. JR 26DD,S-CONT-1 Jump forward.

THE SCANNING VARIABLE ROUTINE
When a variable name has been identified a call is made to LOOK-VARS which looks through those variables that already exist in the variables area (or in the program area at DEF FN statements for a user-defined function FN). If an appropriate numeric value is found then it is copied to the calculator stack using STACK-NUM. However a string or string array entry has to have the appropriate parameters passed to the calculator stack by the STK-VAR subroutine (or in the case of a user-defined function, by the STK-F-ARG subroutine as called from LOOK-VARS).

| 26C9 | S-LETTER | CALL | 28B2,LOOK-VARS |
| :--- | :--- | :--- | :--- | | Look in the existing variables |
| :--- |
| for the matching entry. |

The character is tested against the code for '-', thus identifying the 'unary minus' operation.
Before the actual test the B register is set to hold the priority +09 and the C register the operation code +D 8 that are required for this operation.

| 26DF S-NEGATE | LD | BC,+09DB | Priority +09, operation code |
| :--- | :--- | :--- | :--- |
|  | CP | +2D | Is it a '-'? |
|  | JR | Z,270D,S-PUSH-PO | Jump forward if it is 'unary |
|  |  |  | minus'. |

Next the character is tested against the code for 'VAL\$', with priority 16 decimal and operation code 18 hex.

| LD | BC,+1018 | Priority 16 dec, operation code <br>  <br> +18 hex. |
| :--- | :--- | :--- |
| CP | + AE | Is it 'VAL\$'? |
| JR | Z,270D,S-PUSH-PO | Jump forward if it is 'VAL\$'. |

The present character must now represent one of the functions CODE to NOT, with codes +AF to +C 3 .

| SUB | + AF | The range of the functions is <br> changed from +AF to +C 3 to <br> range +00 to +14 hex. |
| :--- | :--- | :--- |
| JP | C,1C8A,REPORT-C | Report an error if out of range. |

The function 'NOT' is identified and dealt with separately from the others.

| LD | BC,$+04 F 0$ | Priority +04, operation code <br> + F0. |
| :--- | :--- | :--- |
| CP | +14 | Is it the function 'NOT'? |
| JR | Z,270D,S-PUSH-PO | Jump if it is so. |
| JP | NC,1C8A,REPORT-C | Check the range again. |

The remaining functions have priority 16 decimal. The operation codes for these functions are now calculated. Functions that operate on strings need bit 6 reset and functions that give string results need bit 7 reset in their operation codes.

|  |  | LD | B, +10 | Priority 16 decimal. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | ADD | A,+DC | The function range is now +DC +EF. |
|  |  | LD | C, A | Transfer the operation code. |
| 2707 | S-NO-TO-S | CP | +DF | Separate CODE, VAL and LEN |
|  |  | JR | NC,2707.S-NO-TO-S | which operate on strings to give |
|  |  | RES | 6,C | numerical results. |
|  |  | CP | +EE | Separate STR\$ and CHR\$ |
|  |  | JR | C,2700,S-PUSH-PO | which operate on numbers to give string results. |
|  |  | RES | 7,C | Mark the operation codes. |
|  |  |  |  | The other operation codes have bits 6 and 7 both set. |

The priority code and the operation code for the function being considered are now pushed on to the machine stack. A hierarchy of operations is thereby built up.

270D S-PUSH-PO PUSH BC Stack the priority and operation
RST 0020,NEXT-CHAR codes before moving on to
JP 24FF,S-LOOP-1 consider the next part of the expression.

The scanning of the line now continues. The present argument may be followed by a '(', a binary operator or, if the end of the expression has been reached, then e.g. a carriage return character or a colon, a separator or a 'THEN'.

| 2712 | S-CONT-2 | RST | 0018,GET-CHAR | Fetch the present character. <br> 2713 |
| :--- | :--- | :--- | :--- | :--- |
| S-CONT-3 | CP | +28 | Jump forward if it is not a '(', |  |
|  |  | JR | NZ,2723,S-OPERTR | which indicates a parenthesised |
|  |  |  | expression. |  |

If the 'last value' is numeric then the parenthesised expression is a true sub-expression and must be evaluated by itself. However if the 'last value' is a string then the parenthesised expression represents an element of an array or a slice of a string. A call to SLICING modifies the parameters of the string as required.

| BIT | 6,(FLAGS) | Jump forward if dealing with a <br> numeric parenthesised |
| :--- | :--- | :--- |
| JR | NZ,2734,S-LOOP | expression. <br> Modify the parameters of the |
| CALL | 2A52,SLICING | 'last value'. |
| RST | 0020,NEXT-CHAR | Move on to consider the next <br> JR |
| 2713,S-CONT-3 | chacter. |  |

If the present character is indeed a binary operator it will be given an operation code in the range $+\mathrm{C} 3-+\mathrm{CF}$ hex, and the appropriate priority code.

| S-OPERTR | LD | $\mathrm{B},+00$ |
| :--- | :--- | :--- |
|  | LD | C,A |
|  | LD | $\mathrm{HL},+2795$ |
|  | JALL | $16 D \mathrm{C}$, INDEXER |
|  | LD | $\mathrm{NC}, 2734$, SLOOP |
|  | LD | $\mathrm{C},(\mathrm{HL})$ |
|  |  | $\mathrm{HL},+26 \mathrm{ED}$ |
|  | ADD |  |
|  | LD | $\mathrm{HL}, \mathrm{BC}$ |
|  |  | $\mathrm{B},(\mathrm{HL})$ |

Original code to BC to index into table of operators.
The pointer to the table. Index into the table.
Jump forward if no operation found.
Get required code from the
table.
The pointer to the priority
table: i.e. 26ED +C3 gives 27B0
as the first address.
Index into the table.
Fetch the appropriate priority.
The main loop of this subroutine is now entered. At this stage there are:
I. A 'last value' on the calculator stack.
II. The starting priority market on the machine stack below a hierarchy, of unknown size, of function and binary operation codes. This hierarchy may be null.
III. The BC register pair holding the 'present' operation and priority, which if the end of an expression has been reached will be priority zero.
Initially the 'last' operation and priority are taken off the machine stack and compared against the 'present' operation and priority.
If the 'present' priority is higher than the 'last' priority then an exit is made from the loop as the 'present' priority is considered to bind tighter than the 'last' priority.

However, if the present priority is less binding, then the operation specified as the 'last' operation is performed. The 'present' operation and priority go back on the machine stack to be carried round the loop again. In this manner the hierarchy of functions and binary operations that have been queued are dealt with in the correct order.

| 2734 S-LOOP | POP | DE |
| :--- | :--- | :--- |
|  | LD | Get the 'last' operation and <br> priority. |
|  | CP | B |
|  | The priority goes to the A |  |
| register. |  |  |

Before the 'last' operation is performed, the 'USR' function is separated into 'USR number' and 'USR string' according as bit 6 of FLAGS was set or reset when the argument of the function was stacked as the 'last value'.


Stack the 'present' values.
This is FLAGS
The 'last' operation is compared with the code for USR, which will give 'USR number' unless modified; jump if not 'USR'. Test bit 6 of FLAGS. Jump if it is set ('USR number'). Modify the 'last' operation code: 'offset' 19, +80 for string input and numerical result ('USR string').
Stack the 'last' values briefly.
Do not perform the actual operation if syntax is being checked.

| LD | A, E | The 'last' operation code. |
| :---: | :---: | :---: |
| AND | +3F | Strip off bits 6 and 7 to convert |
| LD | B,A | the operation code to a calculator offset. |
| RST | 0028,FP-CALC | Now use the calculator. |
| DEFB | +3B,fp-calc-2 | Perform the actual operation |
| DEFB | +38,end-calc | It has been done. |
| JR | 2764,S-RUNTEST | Jump forward. |

An important part of syntax checking involves the testing of the operation to ensure that the nature of the 'last value' is of the correct type for the operation under consideration.

| 275B | S-SYNTEST | LD <br> XOR <br> AND | A,E <br> (FLAGS) <br> +40 |
| :--- | :--- | :--- | :--- |
| 2761 | S-RPORT-C JP | Set the 'last' operation code. <br> This tests the nature of the 'last <br> value' against the requirement <br> of the operation. They are to be <br> the same for correct syntax. |  |
| Jump if syntax fails. |  |  |  |

Before jumping back to go round the loop again the nature of the 'last value' must be recorded in FLAGS.

| 2764 | S-RUNTEST | POP | DE | Get the 'last' operation code. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | LD | HL, +5 C3B | This is FLAGS. |
|  |  | SET | 6,(HL) | Assume result to be numeric. |
|  |  | BIT | 7,E | Jump forward if the nature of |
|  |  | JR | NZ,2770,S-LOOPEND | 'last value' is numeric. |
|  |  | RES | 6,(HL) | It is string. |
| 2770 | S-LOOPEND | POP | $B C$ | Get the 'present' values into BC : |
|  |  | JR | 2734,S-LOOP | Jump back. |

Whenever the 'present' operation binds tighter, the 'last' and the 'present' values go back on the machine stack. However if the 'present' operation requires a string as its operand then the operation code is modified to indicate this requirement.

| 2773 | S-TIGHTER | $\begin{aligned} & \text { PUSH } \\ & \text { LD } \end{aligned}$ | $\begin{aligned} & \mathrm{DE} \\ & \mathrm{~A}, \mathrm{C} \end{aligned}$ | The 'last' values go on the stack. Get the 'present' operation code. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | BIT | 6,(FLAGS) | Do not modify the operation |
|  |  | JR | NZ,2790,S-NEXT | code if dealing with a numeric operand. |
|  |  | AND | +3F | Clear bits 6 and 7. |
|  |  | ADD | A,+08 | Increase the code by +08 hex. |
|  |  | LD | C,A | Return the code to the C register. |
|  |  | CP | +10 | Is the operation 'AND'? |
|  |  | JR | NZ,2788,S-NOT-AND | Jump if it is not so. |
|  |  | SET | 6,C | 'AND' requires a numeric operand. |
|  |  | JR | 2790,S-NEXT | Jump forward. |
| 2788 | S-NOT-AND | JR | C,2761,S-RPORT-C | The operations -, ${ }^{*}, /, \wedge$ and OR are not possible between strings. |
|  |  | CP | +17 | Is the operation a ' + '? |
|  |  | JR | Z,2790,S-NEXT | Jump if it is so. |
|  |  | SET | 7,C | The other operations yield a numeric result. |
| 2790 | S-NEXT | PUSH | BC | The 'present' values go on the machine stack. |
|  |  | $\begin{aligned} & \text { RST } \\ & \text { JP } \end{aligned}$ | 0020,NEXT-CHAR 24FF S-LOOP- 1 | Consider the next character. Go around the loop again. |

THE TABLE OF OPERATORS

| location | code | operator code | operator | location | code | operator code | operator |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 2795 | 2B | CF | + | 27A3 | 3C | CD | < |
| 2797 | 2D | C3 | - | 27A5 | C7 | C9 | <= |
| 2799 | 2A | C4 | * | 27A7 | C8 | CA | >= |
| 279B | 2F | C5 | / | 27A9 | C9 | CB | <> |
| 279D | 5E | C6 | $\wedge$ | 27AB | C5 | C7 | OR |
| 279F | 3D | CE | = | 27AD | C6 | C8 | AND |
| 27A1 | 3E | CC | > | 27AF | 00 | End marker |  |

## THE TABLE OF PRIORITIES (precedence table)

| location | priority | operator | location | priority | operator |
| :--- | :--- | :--- | :--- | :--- | :--- |
| 27B0 | 06 | - | $27 B 7$ | 05 | $>=$ |
| 27B1 | 08 | $*$ | $27 B 8$ | 05 | $<>$ |
| 27B2 | 08 | $/$ | $27 B 9$ | 05 | $>$ |
| 27B3 | $0 A$ | A | $27 B A$ | 05 | $<$ |
| 27B4 | 02 | OR | $27 B B$ | 05 | $=$ |
| $27 B 5$ | 03 | AND | $27 B C$ | 06 | + |
| 27B6 | 05 | $<=$ |  |  |  |

## THE 'SCANNING FUNCTION' SUBROUTINE

This subroutine is called by the 'scanning FN routine' to evaluate a user defined function which occurs in a BASIC line. The subroutine can be considered in four stages:
I. The syntax of the FN statement is checked during syntax checking.
II. During line execution, a search is made of the program area for a DEF FN statement, and the names of the functions are compared, until a match is found - or an error is reported.
III. The arguments of the FN are evaluated by calls to SCANNING.
IV. The function itself is evaluated by calling SCANNING, which in turn calls LOOK-VARS and so the 'STACK FUNCTION ARGUMENT' subroutine.

| 27BD | S-FN-SBRN | CALL JR RST | 2530,SYNTAX-Z NZ,27F7,SF-RUN 0020,NEXT-CHAR | Unless syntax is being checked, a jump is made to SF-RUN. Get the first character of the name. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | CALL | 2C8D,ALPHA | If it is not alphabetic, then |
|  |  | JP | NC,1C8A,REPORT-C | report the error. |
|  |  | RST | 0020,NEXT-CHAR | Get the next character. |
|  |  | CP | +24 | Is it a '\$'? |
|  |  | PUSH | AF | Save the zero flag on the stack. |
|  |  | JR | NZ,27D0,SF-BRKT-1 | Jump if it was not a '\$'. |
|  |  | RST | 0020,NEXT-CHAR | But get the next character if it was. |
| 27D0 | SF-BRKT-1 | CP | +28 | If the character is not a '(', then |
|  |  | JR | NZ,27E6,SF-RPRT-C | report the error. |
|  |  | RST | 0020,NEXT-CHAR | Get the next character. |
|  |  | CP | +29 | Is it a ')'? |
|  |  | JR | Z,27E9,SF-FLAG-6 | Jump if it is; there are no arguments. |
| 27D9 | SF-ARGMTS | CALL | 24FB,SCANNING | Within the loop, call SCANNING to check the syntax of each argument and to insert floatingpoint numbers. |
|  |  | RST | 0018,GET-CHAR | Get the character which follows |
|  |  | CP | $+2 \mathrm{C}$ | the argument; if it is not a ',' |
|  |  | JR | NZ,27E4,SF-BRKT-2 | then jump - no more arguments. |
|  |  | RST | 0020,NEXT-CHAR | Get the first character in the next argument. |


|  |  | JR | 27D9,SF-ARGMTS |
| :--- | :--- | :--- | :--- |
| 27E4 | SF-BRKT-2 | CP | +29 |
| 27E6 | SF-RPRT-C | JP | NZ,1C8A,REPORT-C |
| 27E9 | SF-FLAG-6 | RST | 0020,NEXT-CHAR |
|  |  |  |  |
|  |  | LD | HL.+5C3B |
|  |  | RES | $6,(H L)$ |
|  |  |  | POP |
|  |  | JR | AF |
|  |  | SET | Z,27F4,SF-SYN-EN |
| 27F4 | SF-SYN-EN | JP | $6,(H L)$ |
|  |  |  |  |

Loop back to consider this argument.
Is the current character a ')'?
Report the error if it is not.
Point to the next character in the BASIC line.
This is FLAGS; assume a stringvalued function and reset bit 6 of FLAGS.
Restore the zero flag, jump if the FN is indeed string valued.
Otherwise, set bit 6 of FLAGS
Jump back to continue scanning the line.
ii. During line execution, a search must first be made for a DEF FN statement.

| 27F7 | SF-RUN | RST | 0020,NEXT-CHAR | Get the first character of the name. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | AND | +DF | Reset bit 5 for upper case. |
|  |  | LD | B,A | Copy the name to B. |
|  |  | RST | 0020,NEXT-CHAR | Get the next character. |
|  |  | SUB | +24 | Subtract 24 hex, the code for '\$'. |
|  |  | LD | C,A | Copy the result to C (zero for a string, non-zero for a numerical function). |
|  |  | JR | NZ,2802,SF-ARGMT1 | Jump if non-zero: numerical function. |
|  |  | RST | 0020,NEXT-CHAR | Get the next character, the '('. |
| 2802 | SF-ARGMT1 | RST | 0020,NEXT-CHAR | Get 1st character of 1st argument. |
|  |  | PUSH | HL | Save the pointer to it on the stack. |
|  |  | LD | HL,(PROG) | Point to the start of the program. |
|  |  | DEC | HL | Go back one location. |
| 2808 | SF-FND-DF | LD | DE,+00CE | The search will be for 'DEF FN'. |
|  |  | PUSH | BC | Save the name and 'string status'. |
|  |  | CALL | 1D86,LOOK-PROG | Search the program now. |
|  |  | POP | BC | Restore the name and status. |
|  |  | JR | NC,2814,SF-CP-DEF | Jump if a DEF FN statement found. |

REPORT P - FN without DEF.

2812 REPORT-P RST
DEFB

0008,ERROR+18

Call the error handling routine.

When a DEF FN statement is found, the name and status of the two functions are compared: if they do not match, the search is resumed.

2814 SF-CP-DEF PUSH \begin{tabular}{ll}

HL \& | Save the pointer to the DEF FN |
| :--- |
| character in case the search has |
| to be resumed. | <br>

Get the name of the DEF FN
\end{tabular}

|  | JR | Z,2831,SF-VALUES | Jump if complete match now <br> found. |
| :--- | :--- | :--- | :--- |
| SF-NOT-FD | POP | HL | Restore the pointer to the 'DEF <br> FN'. |
|  | DEC | HL | Step back one location. |

ii. The correct DEF FN statement has now been found. The arguments of the FN statement will be evaluated by repeated calls of SCANNING, and their 5 byte values (or parameters, for strings) will be inserted into the DEF FN statement in the spaces made there at syntax checking. HL will be used to point along the DEF FN statement (calling FN-SKPOVR as needed) while CH-ADD points along the FN statement (calling RST 0020, NEXT-CHAR, as needed).

| 2831 | SF-VALUES | AND | A | If HL is now pointing to a '\$', |
| :---: | :---: | :---: | :---: | :---: |
|  |  | CALL | Z,28AB,FN-SKPOVR | move on to the '('. |
|  |  | POP | DE | Discard the pointer to 'DEF FN'. |
|  |  | POP | DE | Get the pointer to the first |
|  |  | LD | (CH-ADD), DE | argument of FN, and copy it to CH-ADD. |
|  |  | CALL | 28AB,FN-SKPOVR | Move past the '(' now. |
|  |  | PUSH | HL | Save this pointer on the stack. |
|  |  | CP | +29 | Is it pointing to a ')'? |
|  |  | JR | Z,2885,SF-R-BR-2 | If so, jump: FN has no arguments. |
| 2843 | SF-ARG-LP | INC | HL | Point to the next code. |
|  |  | LD | A,(HL) | Put the code into A. |
|  |  | CP | +0E | Is it the 'number marker' code, 0E hex? |
|  |  | LD | D,+40 | Set bit 6 of $D$ for a numerical argument. |
|  |  | JR | Z,2852,SF-ARG-VL | Jump on zero: numerical argument. |
|  |  | DEC | HL | Now ensure that HL is pointing to the '\$' character (not e.g. to a |
|  |  | CALL | 28AB,FN-SKPOVR | control code). |
|  |  | INC | HL | HL now points to the 'number marker'. |
|  |  | LD | D, +00 | Bit 6 of $D$ is reset: string argument. |
| 2852 | SF-ARG-VL | INC | HL | Point to the 1st of the 5 bytes in DEF FN. |
|  |  | PUSH | HL | Save this pointer on the stack. |
|  |  | PUSH | DE | Save the 'string status' of the argument. |
|  |  | CALL | 24FB,SCANNING | Now evaluate the argument. |
|  |  | POP | AF | Get the no./string flag into A. |
|  |  | XOR | (FLAGS) | Test bit 6 of it against the result |
|  |  | AND | +40 | of SCANNING. |
|  |  | JR | NZ,288B,REPORT-Q | Give report Q if they did not match. |
|  |  | POP | HL | Get the pointer to the first of |
|  |  | EX | DE,HL | the 5 spaces in DEF FN into DE. |
|  |  | LD | HL,(STKEND) | Point HL at STKEND. |
|  |  | LD | BC,+0005 | BC will count 5 bytes to be moved. |
|  |  | SBC | HL, BC | First, decrease STKEND by 5 , |


iv. Finally, the function itself is evaluated by calling SCANNING, after first setting DEFADD to hold the address of the arguments as they occur in the DEF FN statement. This ensures that LOOK-VARS, when called by SCANNING, will first search these arguments for the required values, before making a search of the variables area.

| 288D | SF-VALUE | POP | DE | Restore pointer to ')' in DEF FN. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | EX | DE,HL | Get this pointer into HL. |
|  |  | LD | (CH-ADD), HL | Insert it into CH-ADD. |
|  |  | LD | HL,(DEFADD) | Get the old value of DEFADD. |
|  |  | EX | (SP), HL | Stack it, and get the start |
|  |  | LD | (DEFADD), HL | address of the arguments area of DEF FN into DEFADD. |
|  |  | PUSH | DE | Save address of ')' in FN. |
|  |  | RST | 0020,NEXT-CHAR | Move CH-ADD on past ')' and |
|  |  | RST | 0020,NEXT-CHAR | ' $=$ ' to the start of the expression for the function in DEF FN. |
|  |  | CALL | 24FB,SCANNING | Now evaluate the function. |
|  |  | POP | HL | Restore the address of ')' in FN. |
|  |  | LD | (CH-ADD), HL | Store it in $\mathrm{CH}-\mathrm{ADD}$. |
|  |  | POP | HL | Restore original value of DEFADD |
|  |  | LD | (DEFADD), HL | Put it back into DEFADD. |


| RST | 0020,NEXT-CHAR | Get the next character in the <br> JASIC line. |
| :--- | :--- | :--- |
| JP | 2712,S-CONT-2 | Jump back to continue <br> Scanning. |

## THE 'FUNCTION SKIPOVER' SUBROUTINE

This subroutine is used by FN and by STK-F-ARG to move HL along the DEF FN statement while leaving CM-ADD undisturbed, as it points along the FN statement.

| 28AB | FN-SKPOVR | INC | Point to the next code in the <br> statement. |
| :--- | :--- | :--- | :--- |
|  | LD | A,(HL) | Copy the code to A. |
|  | CP | +21 | Cump back to skip over it if it is |
| JR |  | a control code or a space. |  |

## THE 'LOOK-VARS' SUBROUTINE

This subroutine is called whenever a search of the variables area or of the arguments of a DEF FN statement is required. The subroutine is entered with the system variable CH -ADD pointing to the first letter of the name of the variable whose location is being sought. The name will be in the program area or the work space. The subroutine initially builds up a discriminator byte, in the C register, that is based on the first letter of the variable's name. Bits $5 \& 6$ of this byte indicate the type of the variable that is being handled.

The $B$ register is used as a bit register to hold flags.


Now find the end character of a name that has more than one character.

| 28D4 | V-CHAR | CALL | 2C88,ALPHANUM | Is the character alphanumeric? |
| :---: | :---: | :---: | :---: | :---: |
|  |  | JR | NC,28EF,V-RUN/SYN | Jump out of the loop when the end of the name is found. |
|  |  | RES | 6,C | Mark the discriminator byte. |
|  |  | RST | 0020,NEXT-CHAR | Get the next character. |
|  |  | JR | 28D4,V-CHAR | Go back to test it. |

Simple strings and arrays of strings require that bit 6 of FLAGS is reset.

| $28 D E$ | V-STR-VAR | RST | 0020,NEXT-CHAR |
| :--- | :--- | :--- | :--- |$\quad$| Step CH-ADD past the '\$'. |
| :--- |
|  |
|  |
|  |
|  |
|  | string.

If DEFADD-hi is non-zero, indicating that a 'function' (a 'FN') is being evaluated, and if in 'run-time', a search will be made of the arguments in the DEF FN statement.
28E3
V-TEST-FN
LD
A,(DEFADD-hi)
Is DEFADD-hi zero?

| AND | A |  |
| :--- | :--- | :--- |
| JR | Z,28EF,V-RUN/SYN | If so, jump forward. |
| CALL | 2530,SYNTAX-Z | In 'run-time'? |
| JP | NZ,2951,STK-F-ARG | If so, jump forward to search |
|  |  | the DEF FN statement. |

Otherwise (or if the variable was not found in the DEF FN statement) a search of variables area will be made, unless syntax is being checked.

| 28EF V-RUN/SYN | LD | B,C | Copy the discriminator bytes to <br> the B register. |
| :--- | :--- | :--- | :--- |
|  | CALL | 2530,SYNTAX-Z | Jump forward if in |

A BASIC line is being executed so make a search of the variables area.
28FD
V-RUN
LD
HL,(VARS)
Pick up the VARS pointer.
Now enter a loop to consider the names of the existing variables.

| 2900 | V-EACH | LD <br> AND JR <br> CP JR RLA ADD JP JR | $\begin{aligned} & \text { A,(HL) } \\ & +7 \mathrm{~F} \\ & \text { Z,2932,V-80-BYTE } \\ & \text { C } \\ & \text { NZ,292A,V-NEXT } \\ & \\ & \text { A,A } \\ & \text { P,293F,V-FOUND-2 } \\ & \text { C,293F,V-FOUND-2 } \end{aligned}$ | The 1st. letter of each existing variable. Match on bits 0 to 6. Jump when the ' 80 -byte' is reached. <br> The actual comparison. Jump forward if the 1st characters do not match. Rotate A leftwards and then double it to test bits 5 \& 6 . Strings and array variables. Simple numeric and FOR-NEXT variables. |
| :---: | :---: | :---: | :---: | :---: |
| Long names are required to be matched fully. |  |  |  |  |
|  |  | POP | DE | Take a copy of the pointer |
|  |  | PUSH | DE | to the 2nd. character. |
|  |  | PUSH | HL | Save the 1st letter pointer. |
| 2912 | V-MATCHES | INC | HL | Consider the next character. |
| 2913 | V-SPACES | LD | A,(DE) | Fetch each character in turn. |
|  |  | INC | DE | Point to the next character. |
|  |  | CP | +20 | Is the character a 'space'? |
|  |  | JR | Z,2913,V-SPACES | Ignore the spaces. |
|  |  | OR | +20 | Set bit 5 so as to match lower and upper case letters. |
|  |  | CP | (HL) | Make the comparison. |
|  |  | JR | Z,2912,V-MATCHES | Back for another character if it does match. |
|  |  | OR | +80 | Will it match with bit 7 set? |
|  |  | CP | (HL) | Try it. |
|  |  | JR | NZ,2929,V-GET-PTR | Jump forward if the 'last characters' do not match. |
|  |  | LD | A,(DE) | Check that the end of the |
|  |  | CALL JR | 2C88,ALPHANUM NC,293E, V-FOUND-1 | name has been reached before jumping forward. |

In all cases where the names fail to match the HL register pair has to be made to point to the next variable in the variables area.

| 2929 | V-GET-PTR | POP | HL | Fetch the pointer. |
| :--- | :--- | :--- | :--- | :--- |
| $292 A$ | V-NEXT | PUSH | BC | Save B \& C briefly. |
|  |  | CALL | 19B8,NEXT-ONE | DE is made to point to the |
|  |  |  |  | next variable. |


| EX | DE,HL | Switch the two pointers. |
| :--- | :--- | :--- |
| POP | BC | Get B \& C back. |
| JR | $2900, V-E A C H$ | Go around the loop again. |

Come here if no entry was found with the correct name.
2932 V-80-BYTE SET 7,B Signal 'variable not found'.

Come here if checking syntax.

3934 V-SYNTAX \begin{tabular}{llll}

\& POP \& DE \& | Drop the pointer to the |
| :--- |
| 2nd. character. | <br>

\& RST \& $0018, G E T-C H A R$ \& Fetch the present character. <br>
\& CP \& +28 \& Is it a '('?
\end{tabular}

Come here when an entry with the correct name was found.

| 293E | V-FOUND-1 | POP | DE | Drop the saved variable pointer. |
| :--- | :--- | :--- | :--- | :--- |
| $293 F$ | V-FOUND-2 | POP | DE | Drop the 2nd character pointer. |
|  |  | POP | DE | Drop the first letter pointer. |
|  |  | PUSH | HL | Save the 'last' letter pointer. |
|  |  | RST | $0018, G E T-C H A R$ | Fetch the current character. |

If the matching variable name has more than a single letter then the other characters must be passed-over.
Note: This appears to have been done already at V-CHAR.

| V-PASS | CALL | 2C88,ALPHANUM |
| :--- | :--- | :--- |
|  | JR | NC,294B,V-END |
|  |  |  |
|  | RST | 0020, NEXT-CHAR |
|  | JR | 2943, V-PASS |

Is it alphanumeric?
Jump when the end of the name
has been found.
RST 0020,NEXT-CHAR
Fetch the next character.
Go back and test it.
The exit-parameters are now set.
294B V-END \(\left.$$
\begin{array}{lll}\text { POP } & \text { HL } & \begin{array}{l}\text { HL holds the pointer to the } \\
\text { letter of a short name or the }\end{array}
$$ <br>
'last' character of a long name. <br>

Rotate the whole register.\end{array}\right\}\)| Specify the state of bit 6. |
| :--- |
| Rinished. |

The exit-parameters for the subroutine can be summarised as follows: The system variable CH-ADD points to the first location after the name of the variable as it occurs in the BASIC line.

When 'variable not found':
I. The carry flag is set.
II. The zero flag is set only when the search was for an array variable.
III. The HL register pair points to the first letter of the name of the variable as it occurs in the BASIC line.

When 'variable found':
I. The carry flag is reset.
II. The zero flag is set for both simple string variables and all array variables.
III. The HL register pair points to the letter of a 'short' name, or the last character of a 'long' name, of the existing entry that was found in the variables area.

In all cases bits $5 \& 6$ of the $C$ register indicate the type of variable being handled. Bit 7 is the complement of the SYNTAX/RUN flag. But only when the subroutine is used in 'runtime' will bits 0 to 4 hold the code of the variable's letter.

In syntax time the return is always made with the carry flag reset. The zero flag is set for arrays and reset for all other variables, except that a simple string name incorrectly followed by a '\$' sets the zero flag and, in the case of SAVE "name" DATA a\$(), passes syntax as well.

## THE 'STACK FUNCTION ARGUMENT' SUBROUTINE

This subroutine is called by LOOK-VARS when DEFADD-hi in non-zero, to make a search of the arguments area of a DEF FN statement, before searching in the variables area. If the variable is found in the DEF FN statement, then the parameters of a string variable are stacked and a signal is given that there is no need to call STK/VAR. But it is left to SCANNING to stack the value of a numerical variable at 26DA in the usual way.

| 2951 | STK-F-ARG | LD | HL,(DEFADD) | Point to the 1st character in the |
| :---: | :---: | :---: | :---: | :---: |
|  |  | LD | A, (HL) | arguments area and put it into A. |
|  |  | CP | +29 | Is it a ')'? |
|  |  | JP | Z,28EF,V-RUN/SYN | Jump to search the variables |
| 295A | SFA-LOOP | LD | A, (HL) | Get the next argument in the loop. |
|  |  | OR | +60 | Set bits 5 \& 6, assuming a |
|  |  | LD | B,A | simple numeric variable; copy it to B. |
|  |  | INC | HL | Point to the next code. |
|  |  | LD | A,(HL) | Put it into the A register. |
|  |  | CP | +0E | Is it the 'number marker' code OE hex? |
|  |  | JR | Z,296B,SFA-CP-VR | Jump if so: numeric variable. |
|  |  | DEC |  | Ensure that HL points to the |
|  |  | CALL | 28AB,FN-SKPOVR | character, not to a space or control code. |
|  |  | INC | HL | HL now points to the 'number marker'. |
|  |  | RES | 5,B | Reset bit 5 of B: string variable. |
| 296B | SFA-CP-VR | LD | A, B | Get the variable name into $A$. |
|  |  | CP | C | Is it the one we are looking for? |
|  |  | JR | Z,2981,SFA-MATCH | Jump if it matches. |
|  |  | INC | HL | Now pass over the 5 bytes of |
|  |  | INC | HL | the floating-point number or |
|  |  | INC | HL | string parameters to get to the |
|  |  | INC | HL | next argument. |
|  |  | INC | HL |  |
|  |  | CALL | 28AB,FN-SKPOVR | Pass on to the next character. |
|  |  | CP | +29 | Is it a ')'? |
|  |  | JP | Z,28EF,V-RUN/SYN | If so, jump to search the variables area. |
|  |  | CALL | 28AB,FN-SKPOVR | Point to the next argument. |
|  |  | JR | 295A,SFA-LOOP | Jump back to consider it. |

A match has been found. The parameters of a string variable are stacked, avoiding the need to call the STK-VAR subroutine.

| 2981 | SFA-MATCH | BIT | 5,C | Test for a numeric variable. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | JR | NZ,2991,SFA-END | Jump if the variable is numeric; SCANNING will stack it. |
|  |  | INC | HL | Point to the first of the 5 bytes to be stacked. |
|  |  | LD | DE,(STKEND) | Point DE to STKEND. |
|  |  | CALL | 33C0,MOVE-FP | Stack the 5 bytes. |
|  |  | EX | $\begin{aligned} & \text { DE,HL } \\ & \text { (STKEND).HL } \end{aligned}$ | Point HL to the new position of STKEND, and reset the |
|  |  |  |  | system variable. |
| 2991 | SFA-END | POP | DE | Discard the LOOK-VARS |
|  |  | POP | DE | pointers (2nd \& 1st character pointers). |
|  |  | XOR | A | Return from the search with |
|  |  | INC | A | both the carry and zero flags |
|  |  |  |  | reset - signalling that a call |
|  |  |  |  | STK-VAR is not required. |
|  |  | RET |  | Finished. |

## THE 'STK-VAR' SUBROUTINE

This subroutine is usually used either to find the parameters that define an existing string entry in the variables area or to return in the HL register pair the base address of a particular element or an array of numbers. When called from DIM the subroutine only checks the syntax of the BASIC statement.

Note that the parameters that define a string may be altered by calling SLICING if this should be specified.
Initially the A and the B registers are cleared and bit 7 of the C register is tested to determine whether syntax is being checked.

| 2996 | STK-VAR | XOR | A |
| :--- | :--- | :--- | :--- |
|  | LD | B,A | Clear the array flag. |
|  | BIT | 7,C | Clear the B register for later. |
|  | JR | NZ,29E7,SV-COUNT | Jump forward if syntax is |
|  |  | being checked. |  |

Next, simple strings are separated from array variables.

| BIT | 7,(HL) | Jump forward if dealing with |
| :--- | :--- | :--- |
| JR | NZ,29AE,SV-ARRAYS | an array variable. |

The parameters for a simple string are readily found.

|  |  | INC | A | Signal 'a simple string'. |
| :---: | :---: | :---: | :---: | :---: |
| 29A1 | SV-SIMPLE\$ | INC | HL | Move along the entry. |
|  |  | LD | C,(HL) | Pick up the low length counter. |
|  |  | INC | HL | Advance the pointer. |
|  |  | LD | B,(HL) | Pick up the high length pointer. |
|  |  | INC | HL | Advance the pointer. |
|  |  | EX | DE,HL | Transfer the pointer to the actual string. |
|  |  | CALL | 2AB2,STK-STORE | Pass these parameters to the calculator stack. |
|  |  | RST | 0018,GET-CHAR | Fetch the present character and jump forward to see if a |
|  |  | JP | 2A49,SV-SLICE? | 'slice' is required. |

The base address of an element in an array is now found. Initially the 'number of dimensions' is collected.

| 29AE | SV-ARRAYS | INC | HL |
| :--- | :--- | :--- | :--- |
|  | INC | HL | Step past the length bytes. |
|  | INC | HL |  |
|  | LD | B,(HL) | Collect the 'number of <br> dimensions'. |
|  | BIT | $6, C$ | Jump forward if handling an <br> array of numbers. |
|  | JR | Z,29C0,SV-PTR |  |

If an array of strings has its 'number of dimensions' equal to '1' then such an array can be handled as a simple string.

| DEC | B | Decrease the 'number of <br> JR |
| :--- | :--- | :--- |
| Z,29A1,SV-SIMPLE\$ | dimensions' and jump if the <br> number is now zero. |  |

Next a check is made to ensure that in the BASIC line the variable is followed by a subscript.

| EX | DE,HL | Save the pointer in DE. |
| :--- | :--- | :--- |
| RST | 0018,GET-CHAR | Get the present character. |
| CP | +28 | Is it a '('? |
| JR | NZ,2A20, REPORT-3 | Report the error if it is not so. |
| EX | DE,HL | Restore the pointer. |

For both numeric arrays and arrays of strings the variable pointer is transferred to the $D E$ register pair before the subscript is evaluated.

| $29 C 0$ | SV-PTR | EX | DE,HL | Pass the pointer to DE. |
| :--- | :--- | :--- | :--- | :--- |
|  | JR | 29E7,SV-COUNT | Jump forward. |  |

The following loop is used to find the parameters of a specified element within an array. The loop is entered at the mid-point - SVCOUNT -, where the element count is set to zero.
The loop is accessed ' B ' times, this being, for a numeric array, equal to the number of dimensions that are being used, but for an array of strings ' B ' is one less than the number of dimensions in use as the last subscript Is used to specify a 'slice' of the string.

| 29 C 3 | SV-COMMA | $\begin{aligned} & \text { PUSH } \\ & \text { RST } \\ & \text { POP } \\ & \text { CP } \\ & \text { JR } \\ & \text { BIT } \\ & \text { JR } \\ & \text { BIT } \\ & \text { JR } \\ & \text { CP } \\ & \text { JR } \\ & \text { RST } \\ & \text { RET } \end{aligned}$ | HL <br> 0018,GET-CHAR <br> HL <br> $+2 \mathrm{C}$ <br> Z,29EA,SV-LOOP <br> 7,C <br> Z,2A20,REPORT-3 <br> 6,C <br> NZ,29D8,SV-CLOSE <br> +29 <br> NZ,2A12,SV-RPT-C <br> 0020,NEXT-CHAR | Save the counter. <br> Get the present character. Restore the counter. Is the present character a ','? Jump forward to consider another subscript. <br> If a line is being executed then there is an error. Jump forward if dealing with an array of strings. <br> Is the present character a ')'? <br> Report an error if not so. <br> Advance CH-ADD. <br> Return as the syntax is correct. |
| :---: | :---: | :---: | :---: | :---: |

For an array of strings the present subscript may represent a 'slice', or the subscript for a 'slice' may yet be present in the BASIC line.

| 29D8 | SV-CLOSE | CP | +29 | Is the present character a ')'? |
| :---: | :---: | :---: | :---: | :---: |
|  |  | JR | Z,2A48,SV-DIM | Jump forward and check whether there is another subscript. |
|  |  | CP | +CC | Is the present character a 'TO'? |
|  |  | JR | NZ,2A12,SV-RPT-C | It must not be otherwise. |
| 29E0 | SV-CH-ADD | RST | 0018,GET-CHAR | Get the present character. |
|  |  | DEC | HL | Point to the preceding |
|  |  | LD | (CH-ADD), HL | character and set CH-ADD. |
|  |  | JR | 2A45,SV-SLICE | Evaluate the 'slice'. |

Enter the loop here.

| 29E7 | SV-COUNT | LD | HL, +0000 | Set the counter to zero. |
| :---: | :---: | :---: | :---: | :---: |
| 29EA | SV-LOOP | PUSH | HL | Save the counter briefly. |
|  |  | RST | 0020,NEXT-CHAR | Advance CH-ADD. |
|  |  | POP | HL | Restore the counter. |
|  |  | LD | A,C | Fetch the discriminator byte. |
|  |  | CP | +C0 | Jump unless checking the |
|  |  | JR | NZ,29FB,SV-MULT | syntax for an array of strings. |
|  |  | RST | 0018,GET-CHAR | Get the present character. |
|  |  | CP | +29 | Is it a ')'? |
|  |  | JR | Z,2A48,SV-DIM | Jump forward as finished counting elements. |
|  |  | CP | +CC | Is to 'TO'? |
|  |  | JR | Z,29E0,SV-CH-ADD | Jump back if dealing with a 'slice'. |
| 29FB | SV-MULT | PUSH | BC | Save the dimension-number counter and the discriminator byte. |
|  |  | PUSH | HL | Save the element-counter. |
|  |  | CALL | 2AEE,DE, (DE+1) | Get a dimension-size Into DE. |
|  |  | EX | (SP), HL | The counter moves to HL and the variable pointer is stacked. |
|  |  | EX | DE,HL | The counter moves to DE and the dimension-size to HL. |
|  |  | CALL | 2ACC,INT-EXP1 | Evaluate the next subscript. |
|  |  | JR | C,2A20,REPORT-3 | Give an error if out of range. |
|  |  | DEC | BC | The result of the evaluation is decremented as the counter is to |


| CALL | 2AF4,GET-HL*DE | before the specified element. <br> Multiply the counter by the <br> dimension-size. |
| :--- | :--- | :--- |
| ADD | HL,BC | Add the result of 'INT-EXP1' <br> to the present counter. |
| POP | DE | Fetch the variable pointer. |
| POP | BC | Fetch the dimension-number <br> and the discriminator byte. <br> Keep going round the loop <br> until 'B' equals zero. |

The SYNTAX/RUN flag is checked before arrays of strings are separated from arrays of numbers.

|  |  | BIT | 7,C |
| :--- | :--- | :--- | :--- |
| 2A12 | SV-RPT-C | JR | NZ,2A7A,SL-RPT-C |
|  | PUSH | HL | Report an error if checking |
|  | BIT | $6, C$ | syntax at this point. |
|  |  | JR | NZ,2A2C,SV-ELEM\$ |

When dealing with an array of numbers the present character must be a ')'.

| LD | B,D | Transfer the variable pointer |
| :--- | :--- | :--- |
| LD | C,E | to the BC register pair. |
| RST | $0018, G E T-C H A R$ | Fetch the present character. |
| CP | +29 | Is it a ')'? |
| JR | Z,2A22,SV-NUMBER | Jump past the error report |
|  |  | unless it is needed. |

Report 3 - Subscript out of range

| 2A20 REPORT-3 | RST <br> DEFB | $0008, E R R O R-1$ <br> +02 | Call the error handling <br> routine. |
| :--- | :--- | :--- | :--- |

The address of the location before the actual floating-point form can now be calculated.

## 2A22

| SV-NUMBER | RST | 0020, NEXT-CHAR |
| :--- | :--- | :--- |
| POP | HL |  |
| LD | DE,+0005 |  |
|  | CALL | $2 A F 4, G E T-H L * D E$ |
|  |  |  |
|  | ADD | $H L, B C$ |
|  | RET |  |

Advance $\mathrm{CH}-\mathrm{ADD}$
Fetch the counter
There are 5 bytes to each element in an array of numbers.
Compute the total number of bytes before the required element.
Make HL point to the location before the required element. Return with this address.

When dealing with an array of strings the length of an element is given by the last 'dimension-size'. The appropriate parameters are calculated and then passed to the calculator stack.
\(\left.$$
\begin{array}{lll}\text { 2A2C SV-ELEM\$ } & \begin{array}{l}\text { CALL } \\
\text { EX }\end{array} & \begin{array}{l}\text { 2AEE,DE,(DE+1) } \\
\text { (SP),HL }\end{array} \\
& \text { CALL } & \text { 2AF4,GET-HL*DE }\end{array}
$$ \begin{array}{l}Fetch the last dimension-size. <br>
The variable printer goes on the <br>
stack and the counter to HL. <br>

Multiply 'counter' by\end{array}\right]\)| 'dimension-size'. |
| :--- |

and hence the existing entry is not to be reclaimed.

There are three possible forms of the last subscript. The first is illustrated by - $\mathrm{A} \$(2,4$ TO 8$)$-, the second by - $\mathrm{A} \$(2)(4 \mathrm{TO} 8)-$ and the third by - $\mathrm{A} \$(2)$ - which is the default form and indicates that the whole string is required.

|  |  | RST | 0018,GET-CHAR | Get the present character. Is it a ')'? <br> Jump if it is so. <br> Is it a ','? |
| :---: | :---: | :---: | :---: | :---: |
|  |  | CP | +29 |  |
|  |  | JR | Z,2A48,SV-DIM |  |
|  |  | CP | +2C |  |
|  |  | JR | NZ,2A20,REPORT-3 | Report the error if not so. |
| 2A45 | SV-SLICE | CALL | 2A52,SLICING | Use SLICING to modify the set of parameters. |
| 2A48 | SV-DIM | RST | 0020,NEXT-CHAR | Fetch the next character. |
| 2A49 | SV-SLICE? | CP | +28 | Is It a '('? |
|  |  | JR | Z,2A45,SV-SLICE | Jump back if there is a 'slice' to be considered. |

When finished considering the last subscript a return can be made.

| RES | 6,(FLAGS) |
| :--- | :--- |
| RET | Signal - string result. |
|  | Return with the parameters of |
| the required string forming a |  |
|  | 'last value' on the calculator |
|  | stack. |

THE 'SLICING' SUBROUTINE
The present string can be sliced using this subroutine. The subroutine is entered with the parameters of the string being present on the top of the calculator stack and in the registers A, B, C, D \& E. Initially the SYNTAX/RUN flag is tested and the parameters of the string are fetched only if a line is being executed.

| 2452 | SLICING | CALL |
| :--- | :--- | :--- |$\quad$ 2530,SYNTAX-Z

Check the flag.
Take the parameters off the stack in 'run-time'.

The possibility of the 'slice' being '()' has to be considered.

| RST | 0020,NEXT-CHAR |
| :--- | :--- |
| CP | +29 |
| JR | Z,2AAD,SL-STORE |

Before proceeding the registers are manipulated as follows:

| PUSH | DE |
| :--- | :--- |
| XOR | A |
| PUSH | AF |
| PUSH | BC |
| LD | DE,+0001 |
|  |  |
| RST | 0018, GET-CHAR |
| POP | HL |

The 'start' goes on the machine stack.
The A register is cleared and saved.
The 'length' is saved briefly. Presume that the 'slice' is to begin with the first character. Get the first character. Pass the 'length' to HL.

The first parameter of the 'slice' is now evaluated.

| CP | +CC | Is the present character a <br> 'TO'? |
| :--- | :--- | :--- |
| JR | Z,2A81,SL-SECOND | The first parameter, by default, <br> will be '1' if the jump is taken. |
| POP | AF | At this stage A is zero. <br> BC is made to hold the first <br> parameter. A will hold +FF if <br> there has been an 'out of range' |
|  | 2ACD,INT-EXP2 | error. |
| PUSH | AF | Save the value anyway. <br> LD |
| Transfer the first parameter |  |  |
| LD | D,B | to DE. |
| PUSH | E,C | SL |


|  | RST | O018,GET-CHAR |
| :--- | :--- | :--- |
|  | POP | HL | | Get the present character. |
| :--- |
|  |
|  |
|  |
|  |
| CP |

At this point a 'slice' of a single character has been identified. e.g. - $\mathrm{A} \$(4)$.

| LD | H,D |
| :--- | :--- |
| LD | L,E |
| JR | 2A94,SL-DEFINE |

The last character of the 'slice' is also the first character. Jump forward.

The second parameter of a 'slice' is now evaluated.

2A81

| SL-SECOND | PUSH <br> RST <br> POP <br> CP | $\begin{aligned} & \text { HL } \\ & \text { OO20,NEXT-CHAR } \\ & \text { HL } \\ & +29 \end{aligned}$ |
| :---: | :---: | :---: |
|  | JR | Z,2A94,SL-DEFINE |
|  | POP | AF |
|  | CALL | 2ACD,INT-EXP2 |
|  | PUSH | AF |
|  | RST | 0018,GET-CHAR |
|  | LD | H,B |
|  | LD | L,C |
|  | CP | +29 |
|  | JR | NZ,2A7A,SL-RPT-C |

The 'new' parameters are now defined.
$\left.\begin{array}{lll}\text { 2A94 SL-DEFINE } & \begin{array}{l}\text { POP } \\ \text { EX }\end{array} & \begin{array}{l}\text { AF } \\ \text { (SP),HL }\end{array} \\ & & \\ & \text { ADD } & \mathrm{HL}, \mathrm{DE}\end{array}\right]$

Save the 'length' briefly.
Get the next character.
Restore the 'length'.
Is the present character a
')'?
Jump if there is not a second parameter.
If the first parameter was in range A will hold zero; otherwise +FF.
Make BC hold the second parameter.
Save the 'error register'.
Get the present character. Pass the result obtained from INT-EXP2 to the HL register pair.
Check that there is a closing bracket now.

Fetch the 'error register'.
The second parameter goes on the stack and the 'start' goes to HL
The first parameter is added to the 'start'.
Go back a location to get it correct.
The 'new start' goes on the stack and the second parameter goes to HL.
Subtract the first parameters
from the second to find the length of the 'slice'.
Initialise the 'new length'.
A negative 'slice' is a 'null string' rather than an error condition. (See manual.)
Allow for the inclusive byte.
Only now test the 'error register'.
Jump if either parameter was out of range for the string Transfer the 'new length' to BC.
Get the 'new start'.
Ensure that a string is still indicated.

| 2AAD | SL-STORE | CALL | 2530,SYNTAX-Z <br>  <br> RET |
| :--- | :--- | :--- | :--- |
| $Z$ |  |  |  |

Return at this point if checking syntax; otherwise continue into the STK-STORE subroutine.

## THE 'STK-STORE' SUBROUTINE

This subroutine passes the values held in the $A, B, C, D \& E$ registers to the calculator stack. The stack thereby grows in size by 5 bytes with every call to this subroutine.
The subroutine is normally used to transfer the parameters of strings but it is also used by STACK-BC and LOG ( $2^{\wedge}$ A) to transfer 'smal integers' to the stack.
Note that when storing the parameters of a string the first value stored (coming from the A register) will be a zero if the string comes from an array of strings or is a 'slice' of a string. The value will be ' 1 ' for a complete simple string. This 'flag' is used in the 'LET' command routine when the ' 1 ' signals that the old copy of the string is to be 'reclaimed'.

| 2AB1 | STK-ST-0 | XOR | A | Signal - a string from an array of strings or a 'sliced' string. |
| :---: | :---: | :---: | :---: | :---: |
| 2AB2 | STK-STO-\$ | RES | 6,(FLAGS) | Ensure the flag Indicates a string result. |
| 2AB6 | STK-STORE | PUSH |  | Save B \& C briefly. |
|  |  | CALL | 33A9,TEST-5-SP | Is there room for 5 bytes? Do not return here unless there is room available. |
|  |  | POP | BC | Restore B \& C. |
|  |  | LD | HL,(STKEND) | Fetch the address of the first location above the present stack. |
|  |  | LD | (HL), A | Transfer the first byte. |
|  |  | INC | HL | Step on. |
|  |  | LD | (HL), E | Transfer the second and |
|  |  | INC | HL | third bytes; for a string |
|  |  | LD | (HL), D | these will be the 'start'. |
|  |  | INC | HL | Step on. |
|  |  | LD | (HL), C | Transfer the fourth and |
|  |  | INC | HL | fifth bytes; for a string |
|  |  | LD | (HL), B | these will be the 'length'. |
|  |  | INC | HL | Step on so as to point to the location above the stack. |
|  |  | $\begin{aligned} & \text { LD } \\ & \text { RET } \end{aligned}$ | (STKEND), HL | Save this address In STKEND and return. |

## THE 'INT-EXP' SUBROUTINE

This subroutine returns the result of evaluating the 'next expression' as an integer value held in the BC register pair. The subroutine also tests this result against a limit-value supplied in the HL register pair. The carry flag becomes set if there is an 'out of range' error The A register is used as an 'error register' and holds +00 of there is no 'previous error' and +FF if there has been one.

| 2ACC | INT-EXP1 | XOR | A |
| :--- | :--- | :--- | :--- |
| 2ACD | INT-EXP2 | PUSH | DE |$\quad$| Clear the 'error register'. |
| :--- |
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| JR | Z,2AEB,I-RESTORE | syntax. |
| :---: | :---: | :---: |
| PUSH | AF | Save the error register again. |
| CALL | 1E99,FIND-INT2 | The 'last value' is compressed Into BC. |
| POP | DE | Error register to D. |
| LD | A, B | A 'next expression' that |
| OR | C | gives zero is always in |
| SCF |  | error so jump forward if it |
| JR | Z,2AE8,I-CARRY | is so. |
| POP | HL | Take a copy of the |
| PUSH | HL | limit-value. This will be a 'dimension-size' a 'DIM-limit' or a 'string length'. |
| AND | A | Now compare the result of |
| SBC | HL,BC | evaluating the expression against the limit. |

The state of the carry flag and the value held in the D register are now manipulated so as to give the appropriate value for the 'error register'.

| 2AE8 | I-CARRY | LD | A,D |
| :--- | :--- | :--- | :--- |
|  |  | SBC | $A,+00$ |

Fetch the 'old error value'
Form the 'new error value'; +00 if no error at anytime/ + FF or less if an 'out of range' error on this pass or on previous ones.

Restore the registers before returning.

| $2 A E B$ | I-RESTORE | POP | $H L$ | Restore HL \& DE. |
| :--- | :--- | :--- | :--- | :--- |
|  |  | POP | DE |  |
|  | RET |  | Return; 'error register' is the <br> A register. |  |

## THE 'DE,(DE+1)' SUBROUTINE

This subroutine performs the construction - LD DE,(DE+1) - and returns HL pointing to 'DE+2'.

| 2AEE | DE,(DE+1) | EX INC LD INC LD RET | DE,HL <br> HL <br> E,(HL) <br> HL <br> D, (HL) | Use HL for the construction. <br> Point to 'DE+1'. <br> In effect-LDE,(DE+1). <br> Point to 'DE+2'. <br> In effect - LD D,(DE+2). <br> Finished. |
| :---: | :---: | :---: | :---: | :---: |

## THE 'GET-HL*DE' SUBROUTINE

Unless syntax is being checked this subroutine calls 'HL=HL*DE' which performs the implied construction.
Overflow of the 16 bits available in the HL register pair gives the report 'out of memory'. This is not exactly the true situation but it implies that the memory is not large enough for the task envisaged by the programmer.

| CALL | 2530, SYNTAX-Z |
| :--- | :--- |
| RET | $Z$ |
| CALL | $30 A 9, H L=H L * D E$ |
| JP | C,1F15,REPORT-4 |
| RET |  |

Return directly if syntax is being checked
Perform the multiplication
Report 'Out of memory'. Finished.

## THE 'LET' COMMAND ROUTINE

This is the actual assignment routine for the LET, READ and INPUT commands.
When the destination variable is a 'newly declared variable' then DEST will point to the first letter of the variable's name as it occurs in the BASIC line. Bit 1 of FLAGX will be set.
However if the destination variable 'exists already' then bit 1 of FLAGX will be reset and DEST will point for a numeric variable to the location before the five bytes of the
'old number'; and for a string variable to the first location of the 'old string'. The use of DEST in this manner applies to simple variables and to elements of arrays.
Bit 0 of FLAGX is set if the destination variable is a 'complete' simple string variable. (Signalling - delete the old copy.) Initially the current value of DEST is collected and bit 1 of FLAGS tested.

| 2AFF LET | HD | HL,(DEST) | Fetch the present address in |
| :--- | :--- | :--- | :--- |
|  | BIT | 1,(FLAGX) | DEST. |
|  | JR | Z,2B66,L-EXISTS | Jump if handling a variable |
|  | that 'exists already'. |  |  |

A 'newly declared variable' is being used. So first the length of its name is found.

LD BC,+0005 | Presume dealing with a |
| :--- |
| numeric variable -5 bytes |

Enter a loop to deal with the characters of a long name. Any spaces or colour codes in the name are ignored.

| 2B0B | L-EACH-CH | INC | BC | Add '1' to the counter for each character of a name. |
| :---: | :---: | :---: | :---: | :---: |
| 2B0C | L-NO-SP | INC | HL | Move along the variable's name. |
|  |  | LD | A, (HL) | Fetch the 'present code'. |
|  |  | CP | +20 | Jump back if it is a 'space'; |
|  |  | JR | Z,2B0C,L-NO-SP | thereby Ignoring spaces. |
|  |  | JR | NC,2B1F,L-TEST-CH | Jump forward if the code is +21 to +FF . |
|  |  | CP | +10 | Accept, as a final code, those |
|  |  | JR | C,2B29,L-SPACES | in the range +00 to +0 F . |
|  |  | CP | +16 | Also accept the range |
|  |  | JR | NC,2B29,L-SPACES | +16 to +1F. |
|  |  | INC | HL | Step past the control code after any of INK to OVER. |
|  |  | JR | 2B0C,L-NO-SP | Jump back as these control codes are treated as spaces. |

Separate 'numeric' and 'string' names.


The 'newly declared numeric variable' presently being handled will require 'BC' spaces in the variables area for its name and its value. The room is made available and the name of the variable is copied over with the characters being 'marked' as required.

## 2B29

L-SPACES

| LD | A,C |
| :--- | :--- |
| LD | HL,(E-LINE) |
| DEC | $H L$ |
| CALL | 1655, MAKE-ROOM |


| INC | HL |
| :--- | :--- |
| INC | HL |
| EX | $\mathrm{DE}, \mathrm{HL}$ |
| PUSH | DE |
| LD | $\mathrm{HL},(\mathrm{DEST})$ |
|  |  |
| DEC | DE |
|  |  |
| SUB | +06 |
| LD | $\mathrm{B}, \mathrm{A}$ |

Copy the 'length' to A.
Make HL point to the
'80-byte' at the end of the variables area.
Now open up the variables area.
Note: In effect 'BC' spaces are
made before the displaced
'80-byte'.
Point to the first 'new' byte.
Make DE point to the second
'new' byte.
Save this pointer.
Fetch the pointer to the
start of the name.
Make DE point to the first
'new' byte.
Make $B$ hold the 'number of extra letters' that are found in a 'long name'.

The 'extra' codes of a long name are passed to the variables area.


Point to each 'extra' code.
Fetch the code.
Accept codes from +21 to +FF;
ignore codes +00 to +20 .
Set bit 5, as for lower case letters.
Transfer the codes in turn
to the 2nd 'new' byte
onwards.
Go round the loop for all the 'extra' codes.

The last code of a 'long' name has to be ORed with +80 .

| OR | +80 | Mark the code as required |
| :--- | :--- | :--- |
| LD | (DE),A | and overwrite the last code. |

The first letter of the name of the variable being handled is now considered.

|  | LD | A, +C0 | Prepare the mark the letter <br> of a 'long' name. |
| :--- | :--- | :--- | :--- |
| 2B4F L-SINGLE | LD | HL,(DEST) <br> (HL) | Fetch the pointer to the letter. <br> A holds +00 for a 'short' |
| name and + C0 for a 'long' name. |  |  |  |

The subroutine L-FIRST is now called to enter the 'letter' into its appropriate location.

| CALL 2BEA,L-FIRST | Enter the letter and return <br> with HL pointing to 'new <br> 80-byte'. |
| :--- | :--- |

The 'last value' can now be transferred to the variables area. Note that at this point HL always points to the location after the five locations allotted to the number.
A 'RST 0028' instruction is used to call the CALCULATOR and the 'last value' is deleted. However this value is not overwritten.
2B59

| L-NUMERIC | PUSH | HL |
| :--- | :--- | :--- |
|  | RST | 0028,FP-CALC |
|  | DEFB | +02,delete |
|  | DEFB | +38,end-calc |
|  | POP | HL |
|  | LD | BC, +0005 |
|  |  |  |
|  | AND | A |
|  | SBC | HL,BC |
|  | JR | 2BA6,L-ENTER |

Save the 'destination' pointer.
Use the calculator.
This moves STKEND back five bytes.
Restore the pointer.
Give the number a 'length' of five bytes.
Make HL point to the first of the five locations and jump forward to make the actual transfer.

Come here if considering a variable that 'exists already'. First bit 6 of FLAGS is tested so as to separate numeric variables from string or array of string variables.
2 L-EXISTS BIT 6,(FLAGS) Jump forward if handling any Z,2B72,L-DELETES kind of string variable.

For numeric variables the 'new' number overwrites the 'old' number. So first HL has to be made to point to the location after the five bytes of the existing entry. At present HL points to the location before the five bytes.

| LD | DE, +0006 | The five bytes of a number +'1'. |
| :--- | :--- | :--- |
| ADD | HL,DE | HL now points 'after'. |
| JR | 2B59,L-NUMERIC | Jump back to make the actual |
|  |  | transfer. |

The parameters of the string variable are fetched and complete simple strings separated from 'sliced' strings and array strings.
2B72

| L-DELETE\$ | LD | HL,(DEST) |
| :--- | :--- | :--- |
|  |  |  |
|  | LD | BC,(STRLEN) |
|  | BIT | $0,($ FLAGX) |
|  | JR | NZ,2BAF,L-ADD\$ |

Fetch the 'start'. Note: This
line is redundant.
Fetch the 'length'.
Jump if dealing with a complete simple string; the old string will need to be 'deleted' in this case only.

When dealing with a 'slice' of an existing simple string, a 'slice' of a string from an array of strings or a complete string from an array of strings there are two distinct stages involved. The first is to build up the 'new' string in the work space, lengthening or shortening it as required. The second stage is then to copy the 'new' string to its allotted room in the variables area.
However do nothing if the string has no 'length'.

| LD | A,B | Return if the string is |
| :--- | :--- | :--- |
| OR | $C$ | a null string. |
| RET | $Z$ |  |

Then make the required number of spaces available in the work space.

| PUSH | HL | Save the 'start' (DEST). |
| :---: | :---: | :---: |
| RST | 0030,BC-SPACES | Make the necessary amount of room in the work space. |
| PUSH | DE | Save the pointer to the first location. |
| PUSH | BC | Save the 'length' for use later on. |
| LD | D,H | Make DE point to the last |
| LD | E,L | location. |
| INC | HL | Make HL point 'one past' the new locations. |
| $\begin{aligned} & \text { LD } \\ & \text { LDDR } \end{aligned}$ | (HL), +20 | Enter a 'space' character. Copy this character into all the new locations. Finish with HL pointing to the first new location. |

The parameters of the string being handled are now fetched from the calculator stack.

| PUSH | HL | calculator stack. |
| :--- | :--- | :--- |
| CALL | 2BF1,STK-FETCH | Fetch the pointer briefly. |
| POP | HL | Restore the pointer. |

Note: At this point the required amount of room has been made available in the work space for the 'variable in assignment'. e.g. For statement - LET A\$(4 to 8)="abcdefg" - five locations have been made.
The parameters fetched above as a 'last value' represent the string that is to be copied into the new locations with Procrustean lengthening or shortening as required.
The length of the 'new' string is compared to the length of the room made available for it.

|  |  | EX | (SP), HL | 'Length' of new area to HL. <br> 'Pointer' to new area to stack. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | AND | A | Compare the two 'lengths' |
|  |  | SBC | HL,BC | and jump forward if the 'new' |
|  |  | ADD | HL,BC | string will fit into the room. |
|  |  | JR | NC,2B9B,L-LENGTH | i.e. No shortening required. |
|  |  | LD | B,H | However modify the 'new' |
|  |  | LD | C,L | length if it is too long. |
| 2B9B | L-LENGTH | EX | (SP), HL | 'Length' of new area to stack. 'Pointer' to new area to HL. |

As long as the new string is not a 'null string' it is copied into the work space. Procrustean lengthening is achieved automatically if the 'new' string is shorter than the room available for it.

| EX | DE,HL | 'Start' of new string to HL. 'Pointer' to new area to DE |
| :---: | :---: | :---: |
| LD | A, B | Jump forward if the |
| OR | C | 'new' string is a 'null' |
| JR | Z,2BA3,L-IN-W/S | string. |
| LDIR |  | Otherwise move the 'new' string to the work space. |

The values that have been saved on the machine stack are restored.

| 2BA3 | L-IN-W/S | $\begin{aligned} & \text { POP } \\ & \text { POP } \\ & \text { POP } \end{aligned}$ | $\begin{aligned} & \mathrm{BC} \\ & \mathrm{DE} \\ & \mathrm{HL} \end{aligned}$ | 'Length' of new area. 'Pointer' to new area. The start - the pointer to the 'variable in assignment' which was originally in DEST. L-ENTER is now used to pass the 'new' string to the variables area. |
| :---: | :---: | :---: | :---: | :---: |

## THE 'L-ENTER' SUBROUTINE

This short subroutine is used to pass either a numeric value, from the calculator stack, or a string, from the work space, to its appropriate position in the variables area.
The subroutine is therefore used for all except 'newly declared' simple strings and 'complete \& existing' simple strings.

| 2BA6 | L-ENTER | EX | DE,HL | Change the pointers over. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | LD | A, B | Check once again that the |
|  |  | OR | C | length is not zero. |
|  |  | RET | Z |  |
|  |  | PUSH <br> LDIR | DE |  |
|  |  |  |  | Move the numeric value or the string |
|  |  | $\begin{aligned} & \text { POP } \\ & \text { RET } \end{aligned}$ | HL | Return with the HL register pair pointing to the first byte of the numeric value or the string. |

## THE LET SUBROUTINE CONTINUES HERE

When handling a 'complete \& existing' simple string the new string is entered as if it were a 'newly declared' simple string before the existing version is 'reclaimed'.

| 2BAF | DEC | HL | Make HL point to the letter <br> of the variable's name. |
| :--- | :--- | :--- | :--- |
|  | DEC | HL | i.e. DEST -3. |

'Newly declared' simple strings are handled as follows:
2BC0 L-NEW\$ LD A,+DF Prepare for the marking of the variable's letter.

Fetch the pointer to the
letter.
Mark the letter as required.
L-STRING is now used to add the new string to the variables area.

## THE 'L-STRING' SUBROUTINE

The parameters of the 'new' string are fetched, sufficient room is made available for it and the string is then transferred.

| 2BC6 | L-STRING | PUSH | AF | Save the variable's letter |
| :---: | :---: | :---: | :---: | :---: |
|  |  | CALL | 2BF1,STK-FETCH | Fetch the 'start' and the 'length' of the 'new' string. |
|  |  | EX | DE, HL | Move the 'start' to HL. |
|  |  | ADD | HL, BC | Make HL point 'one-past' the string. |
|  |  | PUSH | BC | Save the 'length'. |
|  |  | DEC | HL | Make HL point to the end of the string. |
|  |  | LD | (DEST), HL | Save the pointer briefly. |
|  |  | INC | BC | Allow one byte for the letter |
|  |  | INC | BC | and two bytes for the length. |
|  |  | INC | BC |  |
|  |  | LD | HL,(E-LINE) | Make HL point to the |
|  |  | DEC | HL | '80-byte' at the end of the variables area. |
|  |  | CALL | 1655,MAKE-ROOM | Now open up the variables area. Note: In effect 'BC' spaces are made before the displaced '80-byte'. |
|  |  | LD | HL,(DEST) | Restore the pointer to the end of the 'new' string. |
|  |  | POP | BC | Make a copy of the length |
|  |  | PUSH | BC | of the 'new' string. |
|  |  | INC | BC | Add one to the length in case the 'new' string is a 'null' string. |
|  |  | LDDR |  | Now copy the 'new' string + one byte. |
|  |  | EX | DE,HL | Make HL point to the byte |
|  |  | INC | HL | that is to hold the high-length. |
|  |  | POP | BC | Fetch the 'length'. |
|  |  | LD | (HL), B | Enter the high-length. |
|  |  | DEC | HL | Back one. |
|  |  | LD | (HL), C | Enter the low-length. |
|  |  | POP | AF | Fetch the variable's letter. |

## THE 'L-FIRST' SUBROUTINE

This subroutine is entered with the letter of the variable, suitably marked, in the A register. The letter overwrites the 'old 80-byte' in the variables area. The subroutine returns with the HL register pair pointing to the 'new 80 -byte'.
\(\left.$$
\begin{array}{lll}\text { 2BEA L-FIRST } & \text { DEC } & \begin{array}{l}\text { Make HL point to the 'old } \\
\text { 80-byte'. }\end{array}
$$ <br>
Lt is overwritten with the <br>

letter of the variable.\end{array}\right]\)| Make HL point to the 'new |
| :--- |
| LD |

## THE 'STK-FETCH' SUBROUTINE

This important subroutine collects the 'last value' from the calculator stack. The five bytes can be either a floating-point number, in 'short' or 'long' form, or set of parameters that define a string.

| 2BF1 | STK-FETCH | LD | HL,(STKEND) |
| :--- | :--- | :--- | :--- |
|  | DEC | HL | Get STKEND. |
|  | LD | $\mathrm{B},(\mathrm{HL})$ | Back one; |
|  | DEC | HL | The fifth value. |
|  | LD | $\mathrm{C},(\mathrm{HL})$ | Back one. |
|  | DEC | HL | The fourth one. |
|  | LD | $\mathrm{D},(\mathrm{HL})$ | Back one. |
|  | DEC | HL | The third value. |
|  | LD | $\mathrm{E},(\mathrm{HL})$ | Back one. |
|  | DEC | HL | The second value. |
|  | LD | $\mathrm{A,(HL)}$ | Back one. |
|  | LD | (STKEND),HL | The first value. |
|  |  |  | Reset STKEND to its new |
|  | RET |  | position |
|  |  | Finished. |  |

## THE 'DIM' COMMAND ROUTINE

This routine establishes new arrays in the variables area. The routine starts by searching the existing variables area to determine whether there is an existing array with the same name. If such an array is found then it is 'reclaimed' before the new array is established.
A new array will have all its elements set to zero, if it is a numeric array, or to 'spaces', if it is an array of strings.

| $2 \mathrm{C02}$ | DIM | CALL | 28B2,LOOK-VARS | Search the variables area. |
| :---: | :---: | :---: | :---: | :---: |
| 2C05 | D-RPORT-C | JP | NZ,1C8A,REPORT-C | Give report C as there has been an error. |
|  |  | CALL | 2530,SYNTAX-Z | Jump forward if in |
|  |  | JR | NZ,2C15,D-RUN | 'run time'. |
|  |  | RES | 6,C | Test the syntax for string arrays as if they were numeric. |
|  |  | CALL | 2996,STK-VAR | Check the syntax of the parenthesised expression. |
|  |  | CALL | 1BEE,CHECK-END | Move on to consider the next statement as the syntax was satisfactory. |

An 'existing array' is reclaimed.

| 2C15 D-RUN | JR | C,2C1F,D-LETTER |
| :--- | :--- | :--- |
|  | PUSH | BC |
|  | CALL | 19B8,NEXT-ONE |
|  |  |  |
|  | CALL | 19E8,RECLAIM-2 |
|  |  | BOP |

Jump forward if there is no 'existing array'.
Save the discriminator byte.
Find the start of the next variable
Reclaim the 'existing array'.
Restore the discriminator byte.
The initial parameters of the new array are found.

| 2C1F | D-LETTER | SET | 7,C | Set bit 7 in the discriminator byte. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | LD | $B,+00$ | Make the dimension counter zero. |
|  |  | PUSH | BC | Save the counter and the discriminator byte. |
|  |  | LD | HL, +0001 | The HL register pair is to |
|  |  | BIT | 6,C | hold the size of the elements |
|  |  | JR | NZ,2C2D,D-SIZE | in the array, '1' for a string |
|  |  | LD | L,+05 | array/ '5' for a numeric array. |
| 2C2D | D-SIZE | EX | DE,HL | Element size DE. |

The following loop is accessed for each dimension that is specified in the parenthesised expression of the DIM statement. The total number of bytes required for the elements of the array is built up in the DE register pair.

2C2E

| D-NO-LOOP | RST | 0020,NEXT-CHAR |
| :--- | :--- | :--- |
|  | LD | H,+FF |
|  | CALL | $2 A C C$, INT-EXP1 |
|  | JP | C,2A20,REPORT-3 |
|  | POP | HL |
|  | PUSH | BC |
|  | INC | H |
|  |  |  |
|  | LUSH | HL |
|  | LD | H,B |
|  | CALL | L,C |
|  | EX | DE,HL,GET-HL*DE |
|  | RST |  |
|  | CP | $0018, G E T-C H A R ~$ |
|  | JR | Z,2C,2C2E,D-NO-LOOP |

Advance $\mathrm{CH}-\mathrm{ADD}$ on each pass.
Set a 'limit value'.
Evaluate a parameter.
Give an error if 'out of range'.
Fetch the dimension-counter
and the discriminator byte.
Save the parameter on each
pass through the loop.
Increase the dimension counter on each pass also.
Restack the dimension-counter and the discriminator byte.
The parameter is moved to the HL register pair.
The byte total is built up
in HL and the transferred to DE.
Get the present character and go around the loop again if there is another dimension.

Note: At this point the DE register pair indicates the number of bytes required for the elements of the new array and the size of each dimension is stacked, on the machine stack.
Now check that there is indeed a closing bracket to the parenthesised expression.

| CP | +29 | Is it a ')'? |
| :--- | :--- | :--- |
| JR | NZ,2C05,D-REPORT-C | Jump back if not so. |
| RST | 0020,NEXT-CHAR | Advance CH-ADD past it. |

Allowance is now made for the dimension-sizes.

| POP | BC | Fetch the dimension-counter <br> and the discriminator byte. <br> Pass the discriminator byte <br> to the A register for later. |
| :--- | :--- | :--- |
| LD | A,C | Move the counter to L. |
| LD | Lear the H register. |  |
| LD | H, +00 | Increase the dimension- <br> counter by two and double the <br> result and form the <br> correct overall length for <br> the variable by adding the <br> element byte total. |
| INC | HL | Give the report 'Out of <br> memory' if required. |
| ADD | HL,HL | Save the element byte total. |
| JP | C,1F15,REPORT-4 | Save the dimension counter <br> and the discriminator byte. <br> Save the overall length also. |
| PUSH | DE | Move the overall length to BC. |
| PUSH | BC | HL |

The required amount of room is made available for the new array at the end of the variables area.

| LD | HL,(E-LINE) | Make the HL register pair |
| :--- | :--- | :--- |
| DEC | HL | point to the '80-byte'. |
| CALL | 1655,MAKE-ROOM | The room is made available. |
| INC | HL | HL is made to point to the first |

The parameters are now entered.

| LD | $(\mathrm{HL}), A$ | The letter, suitably marked, <br> is entered first. |
| :--- | :--- | :--- |
| POP | BC | The overall length is fetched |
| DEC | BC | and decreased by '3'. |
| DEC | BC |  |
| DEC | BC |  |
| INC | HL | Advance HL. |
| LD | $(\mathrm{HL}), C$ | Enter the low length. |
| INC | HL | Advance HL. |
| LD | $(\mathrm{HL}), B$ | Enter the high length. |
| POP | BC | Fetch the dimension counter. |
| LD | $\mathrm{A}, \mathrm{B}$ | Move it to the A register. |
| INC | HL | Advance HL. |
| LD | $(\mathrm{HL}), A$ | Enter the dimension count. |

The elements of the new array are now 'cleared'.


The 'dimension-sizes' are now entered.

| 2C7F | DIM-SIZES | POP | BC |
| :--- | :--- | :--- | :--- |
|  | LD | $(\mathrm{HL}), \mathrm{B}$ | Get a dimension-size. |
|  | DEC | HL | Enter the high byte. |
|  | LD | $(\mathrm{HL}), \mathrm{C}$ | Back one. |

## THE 'ALPHANUM' SUBROUTINE

This subroutine returns with the carry flag set if the present value of the A register denotes a valid digit or letter.

| 2C88 ALPHANUM | CALL | 2D1B,NUMERIC |
| :--- | :--- | :--- | | Test for a digit; carry will be |
| :--- |
| reset for a digit. |

## THE 'ALPHA' SUBROUTINE

This subroutine returns with the carry flag set if the present value of the A register denotes a valid letter of the alphabet.

| 2C8D ALPHA | CP | +41 | Test against 41 hex, the code <br> for 'A' |
| :--- | :--- | :--- | :--- |
| CCF | NC | Complement the carry flag. <br> Return if not a valid character <br> code. |  |
| RET | $+5 B$ | Test against $5 B$ hex, 1 more <br> than code for 'Z'. |  |
|  | CP | C | Return if an upper case letter. <br> Test against 61 hex, the code <br> for 'a'. |


| CCF |  | Complement the carry flag. <br> RET |
| :--- | :--- | :--- |
| NC | Return if not a valid character <br> code. |  |
| CP | $+7 B$ | Test against 7B hex, 1 more <br> than the code for 'z'. |
| RET |  | Finished. |

## THE 'DECIMAL TO FLOATING POINT' SUBROUTINE

As part of syntax checking decimal numbers that occur in a BASIC line are converted to their floating-point forms. This subroutine reads the decimal number digit by digit and gives its result as a 'last value' on the calculator stack. But first it deals with the alternative notation BIN , which introduces a sequence of 0 's and 1 's giving the binary representation of the required number.

| 2C9B | DEC-TO-FP | CP | +C4 |
| :---: | :---: | :---: | :---: |
|  |  | JR | NZ,2CB8,NOT-BIN |
|  |  | LD | DE,+0000 |
| 2CA2 | BIN-DIGIT | RST | 0020,NEXT-CHAR |
|  |  | SUB | +31 |
|  |  | ADC | A, +00 |
|  |  | JR | NZ,2CB3,BIN-END |
|  |  | EX | DE,HL |
|  |  | CCF |  |
|  |  | ADC | HL,HL |
|  |  | JP | C,31AD,REPORT-6 |
|  |  | EX | DE,HL |
|  |  | JR | 2CA2,BIN-DIGIT |
| 2CB3 | BIN-END | LD | B, D |
|  |  | LD | C, E |
|  |  | JP | 2D2B,STACK-BC |

For other numbers, first any integer part is converted; considered

| 2CB8 | NOT-BIN | CP | +2E | Is the first character a '.'? |
| :--- | :--- | :--- | :--- | :--- |


| 2CDA | NXT-DGT-1 | $\begin{aligned} & \text { DEFB } \\ & \text { DEFB } \end{aligned}$ | +02 , delete <br> +38 , end-calc | save it in the memory area. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | RST | 0018,GET-CHAR | Get the present character. |
|  |  | CALL | 2D22,STK-DIGIT | If it is a digit then stack it. |
|  |  | JR | C,2CEB,E-FORMAT | If not jump forward. |
|  |  | RST | 0028,FP-CALC | Now use the calculator. |
|  |  | DEFB | +E0,get-mem-0 | For each passage of the loop, |
|  |  | DEFB | +A4,stk-ten | the number saved in the memory |
|  |  | DEFB | +05,division | area is fetched, divided by 10 |
|  |  | DEFB | +C0,st-mem-0 | and restored: i.e. going from . 1 to .01 to .001 etc. |
|  |  | DEFN | +04,multiply | The present digit is multiplied |
|  |  | DEFB | +0F, addition | by the 'saved number' and |
|  |  | DEFB | +38, end-calc | added to the 'last value'. |
|  |  | RST | 0020,NEXT-CHAR | Get the next character. |
|  |  | JR | 2CDA,NXT-DGT-1 | Jump back (one more byte than |

Next consider any 'E notation', i.e. the form xEm or xem where $m$ is a positive or negative integer.

| 2CEB | E-FORMAT | CP | +45 | Is the present character an 'E'? |
| :---: | :---: | :---: | :---: | :---: |
|  |  | JR | Z,2CF2,SIGN-FLAG | Jump forward if it is. |
|  |  | CP | +65 | Is it an 'e'? |
|  |  | RET | NZ | Finished unless it is so. |
| 2CF2 | SIGN-FLAG | LD | B,+FF | Use B as a sign flag, FF for ' + '. |
|  |  | RST | 0020,NEXT-CHAR | Get the next character. |
|  |  | CP | +2B | Is it a '+'? |
|  |  | JR | Z,2CFE,SIGN-DONE | Jump forward. |
|  |  | CP | +2D | Is it a '-'? |
|  |  | JR | NZ,2CFF,ST-E-PART | Jump if neither '+' not '-'. |
|  |  | INC | B | Change the sign of the flag. |
| 2CFE | SIGN-DONE | RST | 0020,NEXT-CHAR | Point to the first digit. |
| 2CFF | ST-E-PART | CALL | 2D1B,NUMERIC | Is it indeed a digit? |
|  |  | JR | C,2CCF,DEC-RPT-C | Report the error if not. |
|  |  | PUSH | BC | Save the flag in B briefly. |
|  |  | CALL | 2D3B,INT-TO-FP | Stack ABS $m$, where $m$ is the exponent. |
|  |  | CALL | 2DD5,FP-TO-A | Transfer ABS m to A. |
|  |  | POP | BC | Restore the sign flag to B. |
|  |  | JP | C,31AD,REPORT-6 | Report the overflow now if |
|  |  | AND |  | ABS $m$ is greater than 255 or |
|  |  | JP | M,31AD,REPORT-6 | indeed greater than 127 (other values greater than about 39 will be detected later). |
|  |  | INC | B | Test the sign flag in B; '+' (i.e. +FF) will now set the zero flag. |
|  |  | $\begin{aligned} & \text { JR } \\ & \text { NEG } \end{aligned}$ | Z,2D18,E-FP-JUMP | Jump if sign of $m$ is ' + '. Negate $m$ if sign is ' - '. |
| 2D18 | E-FP-JUMP | JP | 2D4F,E-TOO-FP | Jump to assign to the 'last value' the result of $x^{\star} 10^{\wedge} m$. |

THE 'NUMERIC' SUBROUTINE
This subroutine returns with the carry flag reset if the present value of the A register denotes a valid digit.

| 2D1B NUMERIC | CP | +30 | C |
| :--- | :--- | :--- | :--- |
| RET | $+3 A$ | Test against 30 hex, the code <br> for '0'. <br> Return if not a valid character |  |
|  | CP |  | code. <br> Test against the upper limit. |
| CCF |  | Complement the carry flag. <br> RET |  |

## THE 'STK DIGIT' SUBROUTINE

This subroutine simply returns if the current value held in the A register does not represent a digit but if it does then the floating-point form for the digit becomes the 'last value' on the calculator stack.
2D22
STK-DIGIT CALL
2D1B,NUMERIC
RET
C
SUB
$+30$
s the character a digit?
Return if not in range.
Replace the code by the actual digit.

## THE 'STACK-A' SUBROUTINE

This subroutine gives the floating-point form for the absolute binary value currently held in the A register.

| 2D28 | STACK-A | LD | C, A | Transfer the value to the register. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | LD | $B,+00$ | Clear the B register |

## THE 'STACK-BC' SUBROUTINE

This subroutine gives the floating-point form for the absolute binary value currently held in the BC register pair.
The form used in this and hence in the two previous subroutines as well is the one reserved in the Spectrum for small integers n , where $-65535<=\mathrm{n}<=65535$. The first and fifth bytes are zero; the third and fourth bytes are the less significant and more significant bytes of the 16 bit integer n in two's complement form (if n is negative, these two bytes hold $65536+\mathrm{n}$ ); and the second byte is a sign byte, 00 for +' and FF for '-'


THE 'INTEGER TO FLOATING-POINT' SUBROUTINE
This subroutine returns a 'last value' on the calculator stack that is the result of converting an integer in a BASIC line, i.e. the integer part of the decimal number or the line number, to its floating-point form.

Repeated calls to CH-ADD+1 fetch each digit of the integer in turn. An exit is made when a code that does not represent a digit has been fetched.

| 2D3B INT-TO-FP | PUSH | AF | Save the first digit - in A. |
| :--- | :--- | :--- | :--- |
|  | RST | 0028,FP-CALC | Use the calculator. |
| DEFB | +A0,stk-zero | The 'last value' is now zero. |  |
|  | DEFB | +38, end-calc |  |
|  | POP | AF | Restore the first digit. |

Now a loop is set up. As long as the code represents a digit then the floating-point form is found and stacked under the 'last value'. The 'last value' is then multiplied by decimal 10 and added to the 'digit' to form a new 'last value' which is carried back to the start of the loop.

| NXT-DGT-2 | CALL | 2D22,STK-DIGIT |
| :--- | :--- | :--- |
|  | RET | C |

If the code represents a digit
then stack the floating-point
form.
Use the calculator.
'Digit' goes under 'last value'.
Define decimal 10.
'Last value' = 'last value' *10.
'Last value' = 'last value+
'digit'.
The next code goes into A.
Loop back with this code.

## THE ARITHMETIC ROUTINES

## THE 'E-FORMAT TO FLOATING-POINT' SUBROUTINE

(Offset 3C - see CALCULATE below: 'e-to-fp')
This subroutine gives a 'last value' on the top of the calculator stack that is the result of converting a number given in the form xEm where $m$ is a positive or negative integer. The subroutine is entered with $x$ at the top of the calculator stack and $m$ in the $A$ register.

The method used is to find the absolute value of $m$, say $p$, and to multiply or divide $\times$ by $10^{\wedge} p$ according to whether $m$ is positive or negative.

To achieve this, $p$ is shifted right until it is zero, and $x$ is multiplied or divided by $10^{\wedge}\left(2^{\wedge} n\right)$ for each set bit $b(n)$ of $p$. Since $p$ is never much more than decimal 39 , bits 6 and 7 of $p$ will not normally be set.

| 2D4F | E-TO-FP | RLCA RRCA |  | Test the sign of $m$ by rotating bit 7 of $A$ into the carry without changing $A$. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | JR | NC,2D55,E-SAVE | Jump if $m$ is positive. |
|  |  | CPL |  | Negate m in A without |
|  |  | INC | A | disturbing the carry flag. |
| 2D55 | E-SAVE | PUSH | AF | Save m in A briefly. |
|  |  | LD | HL,+5C92 | This is MEMBOT: a sign flag is |
|  |  | CALL | 350B,FP-0/1 | now stored in the first byte of mem-0, i.e. 0 for ' + ' and 1 for |
|  |  | RST | 0028,FP-CALC | The stack holds $x$. |
|  |  | DEFB | +A4,stk-ten | x,10 (decimal) |
|  |  | DEFB | +38, end-calc | x,10 |
|  |  | POP | AF | Restore min A. |
| 2D60 | E-LOOP | SRL | A | In the loop, shift out the next bit of $m$, modifying the carry and zero flags appropriately; |
|  |  | JR | NC,2D71,E-TST-END | jump if carry reset. |
|  |  | PUSH | AF | Save the rest of $m$ and the flags. |
|  |  | RST | 0028,FP-CALC | The stack holds $x^{\prime}$ and $10^{\wedge}\left(2^{\wedge} n\right)$, where $x^{\prime}$ is an interim stage in the multiplication of $x$ by $10^{\wedge} m$, and $n=$ $0,1,2,3,4$ or 5 . |
|  |  | DEFB | +C1,st-mem-1 | ( $10^{\wedge}\left(2^{\wedge} \mathrm{n}\right)$ is copied to mem-1). |
|  |  | DEFB | +E0,get-mem-0 | $\mathrm{x}^{\prime}, 10^{\wedge}\left(2^{\wedge} \mathrm{n}\right),(1 / 0)$ |
|  |  | DEFB | +00,jump-true | $\mathrm{x}^{\prime}, 10^{\wedge}\left(2^{\wedge} \mathrm{n}\right)$ |
|  |  | DEFB | +04, to E-DIVSN | $\mathrm{x}^{\prime}, 10^{\wedge}\left(2^{\wedge} \mathrm{n}\right)$ |
|  |  | DEFB | +04,multiply | $x^{\prime *} 10^{\wedge}\left(2^{\wedge} n\right)=x^{\prime \prime}$ |
|  |  | DEFB | +33,jump | x" |
|  |  | DEFB | +02, to E-FETCH | x" |
| 2D6D | E-DIVSN | DEFB | +05,division | $\begin{aligned} & x / 10^{\wedge}\left(2^{\wedge} n\right)=x^{\prime \prime}\left(x^{\prime \prime} \text { is } N^{\prime *} 10^{\wedge}\right. \\ & \left(2^{\wedge} n\right) \text { or } x^{\prime} / 10^{\wedge}\left(2^{\wedge} n\right) \end{aligned}$ |
|  | E-FETCH | DEFB | +E1,get-mem-1 |  |
| 2D6E |  | DEFB | +38,end-calc | $\mathrm{x}^{\prime \prime}, 10^{\wedge}\left(2^{\wedge} \mathrm{n}\right)$ |
|  |  | POP | AF | Restore the rest of $m$ in A, and the flags. |
| 2D71 | E-TST-END | JR | Z,2D7B,E-END | Jump if $m$ has been reduced to zero. |
|  |  | PUSH | AF | Save the rest of $m$ in $A$. |
|  |  | RST | 0028,FP-CALC | x", 10^(2^n) |
|  |  | DEFB | +31,duplicate | x", $10^{\wedge}\left(2^{\wedge} n\right), 10^{\wedge}\left(2^{\wedge} \mathrm{n}\right)$ |
|  |  | DEFB | +04,multiply | $\mathrm{x}^{\prime \prime}, 10^{\wedge}\left(2^{\wedge}(\mathrm{n}+1)\right)$ |
|  |  | DEFB | +38,end-calc | $\mathrm{x}^{\prime \prime}, 10^{\wedge}\left(2^{\wedge}(\mathrm{n}+1)\right.$ ) |
|  |  | POP | AF | Restore the rest of min A. |
|  |  | JR | 2D60,E-LOOP | Jump back for all bits of m. |
| 2D7B | E-END | RST | 0028,FP-CALC | Use the calculator to delete the |


| DEFB | +02, delete | final power of 10 reached, |
| :--- | :--- | :--- |
| DEFB | +28, end-calc | leaving the 'last value' $x^{* 1} 10^{\wedge} m$ |
| RET |  | on the stack |

## THE 'INT-FETCH' SUBROUTINE

This subroutine collects in DE a small integer $\mathrm{n}(-65535<=\mathrm{n}<=65535)$ from the location addressed by HL: i.e. n is normally the first (or second) number at the top of the calculator stack; but HL can alls access (by exchange with DE) a number which has been deleted from the stack. The subroutine does not itself delete the number from the stack or from memory; it returns HL pointing to the fourth byte of the number in its original position.

| 2D7F INT-FETCH | HL | Point to the sign byte of the <br> number. |  |
| :--- | :--- | :--- | :--- |
|  | LD | C,(HL) | Copy the sign byte to C. |

The following mechanism will twos complement the number if it is negative ( $C$ is $F F$ ) but leave it unaltered if it is positive ( $C$ is 00 )

| INC | HL | Point to the less significant <br> byte. |
| :--- | :--- | :--- |
| LD | A,(HL) | Collect the byte in A. <br> Ones complement it if negative |
| XOR | C | This adds 1 for negative <br> numbers; it sets the carry unless <br> the byte was 0. |
| LD | C | Less significant byte to E now. <br> Point to the more significant <br> byte. |
| INC | E,A | Collect it in A. |
| LD | A,(HL) | Finish two complementing in <br> the case of a negative number; <br> note that the carry is always |
| ADC | A,C | left reset. <br> More significant byte to D now. <br> Finished. |
| LD | D,A |  |

THE 'INT-STORE' SUBROUTINE
This subroutine stores a small integer $n(-65535<=n<=65535)$ in the location addressed by HL and the four following locations: i.e. $n$ replaces the first (or second) number at the top of the calculator stack. The subroutine returns HL pointing to the first byte of $n$ on the stack.

| 2D8C | P-INT-STO | LD | C, +00 | This entry point would store a <br> number known to be positive |
| :--- | :--- | :--- | :--- | :--- |
| 2D8E | INT-STORE | PUSH | HL | The pointer to the first location <br> is saved. |
|  |  | LD | (HL),+00 | The first byte is set to zero. |
|  |  | INC | HL | Point to the second location. |

The same mechanism is now used as in 'INT-FETCH' to twos complement negative numbers. This is needed e.g. before and after the multiplication of small integers. Addition is however performed without any further twos complementing before or afterwards.

| INC | HL | Point to the third location. <br> Collect the less significant |
| :--- | :--- | :--- |
| LD | A,E | byte. |
| XOR | Cwos complement it if the |  |
| SUB | C | number is negative |
| LD | (HL),A | Store the byte. |
| INC | HL | Point to the fourth location. |
| LD | A,D | Collect the more significant <br> byte. |
| ADC | A,C | Twos complement it if the |
| XOR | C | number is negative |


| LD | (HL),A | Store the byte. |
| :--- | :--- | :--- |
| INC | HL | Point to the fifth location. |
| LD | $(\mathrm{HL}),+00$ | The fifth byte is set to zero. |
| POP | HL | Return with HL pointing to the |
| RET | first byte on $n$ on the stack |  |

## THE 'FLOATING-POINT TO BC' SUBROUTINE

This subroutine is called from four different places for various purposes and is used to compress the floating-point 'last value' into the BC register pair. If the result is too large, i.e. greater than 65536 decimal, then the subroutine returns with the carry flag set. If the 'last value' is negative then the zero flag is reset. The low byte of the result is also copied to the A register.

| 2DA2 | FP-TO-BC | RST | 0028,FP-CALC | Use the calculator to make HL |
| :---: | :---: | :---: | :---: | :---: |
|  |  | DEFB | +38,end-calc | point to STKEND-5 |
|  |  | LD | A,(HL) | Collect the exponent byte of |
|  |  | AND | A | the 'last value'; jump if it is |
|  |  | JR | Z,2DAD,FP-DELETE zero, indicating a 'small integer'. |  |
|  |  | RST | 0028,FP-CALC | Now use the calculator to round |
|  |  | DEFB | +A2,stk-half | the 'last value' to the nearest |
|  |  | DEFB | +0F, addition | integer, which also changes it to |
|  |  | DEFB | +27, int | 'small integer' form on the |
|  |  | DEFB | +38,end-calc | calculator stack if that is possible, i.e. if -65535.5 <= x <65535.3 |
| 2DAD | FP-DELETE | RST | 0028,FP-CALC | Use the calculator to delete the |
|  |  | DEFB | +92,delete | integer from the stack; DE still |
|  |  | DEFB | +38, end-calc | points to it in memory (at STKEND). |
|  |  | PUSH | HL | Save both stack pointers. |
|  |  | PUSH | DE |  |
|  |  | EX | DE,HL | HL now points to the number. |
|  |  | LD | B,(HL) | Copy the first byte to B . |
|  |  | CALL | 2D7F,INT-FETCH | Copy bytes 2,3 and 4 to C, E and $D$. |
|  |  | XOR | A | Clear the A register. |
|  |  | SUB | B | This sets the carry unless B is zero. |
|  |  | BIT | 7,C | This sets the zero flag if the number is positive (NZ denotes negative). |
|  |  | LD | B, D | Copy the high byte to B . |
|  |  | LD | C, E | And the low byte to C. |
|  |  | LD | A, E | Copy the low byte to A too. |
|  |  | POP | DE | Restore the stack pointers. |
|  |  | POP | HL |  |
|  |  | RET |  | Finished. |

## THE 'LOG (2^A)' SUBROUTINE

This subroutine is called by the 'PRINT-FP' subroutine to calculate the approximate number of digits before the decimal in $x$, the number to be printed, or, if there are no digits before the decimal, then the approximate number of leading zeros after the decimal. It is entered with the A register containing e', the true exponent of $x$, or $e^{\prime}-2$, and calculates $z=\log$ to the base 10 of $\left(2^{\wedge} A\right)$. It then sets $A$ equal to ABS INT $(Z+0.5)$, as required, using FP-TO-A for this purpose.


| CALL | 2AB6,STK-STORE |  |
| :---: | :---: | :---: |
| RST | 0028,FP-CALC | The calculator is used |
| DEFB | +34,stk-data | Log 2 to the base 10 is now stacked. |
| DEFB | +EF, exponent + 7 F | The stack now holds a, log 2 . |
| DEFB | +1A,+20,+9A,+85 |  |
| DEFB | +04,multiply | $A^{*} \log 2$ i.e. $\log \left(2^{\wedge} \mathrm{A}\right)$ |
| DEFB | +27,int | INT $\log \left(2^{\wedge} \mathrm{A}\right)$ |
| DEFB | +38,end-calc |  |

The subroutine continues on into FP-TO-A to complete the calculation.

## THE 'FLOATING-POINT TO A' SUBROUTINE

This short but vital subroutine is called at least 8 times for various purposes. It uses the last but one subroutine, FP-TO-BC, to get the 'last value' into the A register where this is possible. It therefore tests whether the modulus of the number rounds to more than 255 and if it does the subroutine returns with the carry flag set. Otherwise it returns with the modulus of the number, rounded to the nearest integer, in the A register, and the zero flag set to imply that the number was positive, or reset to imply that it was negative.

| 2DD5 | FP-TO-A | CALL | 2DA2,FP-TO-BC | Compress the 'last value' into BC. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | RET | C | Return if out of range already. |
|  |  | PUSH | AF | Save the result and the flags. |
|  |  | DEC | B | Again it will be out of range |
|  |  | INC | B | if the $B$ register does not hold zero. |
|  |  | JR | Z,2DE1,FP-A-END | Jump if in range. |
|  |  | POP |  | Fetch the result and the flags |
|  |  | SCF |  | Signal the result is out of range. |
|  |  | RET |  | Finished - unsuccessful. |
| 2DE1 | FP-A-END | POP | AF | Fetch the result and the flags. Finished - successful. |

## THE 'PRINT A FLOATING-POINT NUMBER' SUBROUTINE

This subroutine is called by the PRINT command routine at 2039 and by STR $\$$ at 3630 , which converts to a string the number as it would be printed. The subroutine prints $x$, the 'last value' on the calculator stack. The print format never occupies more than 14 spaces The 8 most significant digits of $x$, correctly rounded, are stored in an ad hoc print buffer in mem-3 and mem-4. Small numbers, numerically less than 1 , and large numbers, numerically greater than $2^{\wedge} 27$, are dealt with separately. The former are multiplied by $10 \wedge$ n , where n is the approximate number of leading zeros after the decimal, while the latter are divided by $10 \wedge(\mathrm{n}-7)$, where n is the approximate number of digits before the decimal. This brings all numbers into the middle range, and the numbers of digits required before the decimal is built up in the second byte of mem-5. Finally the printing is done, using E-format if there are more than 8 digits before the decimal or, for small numbers, more than 4 leading zeros after the decimal.

The following program shows the range of print formats:
10 FOR $a=-11$ TO 12: PRINT SGN a*9^a,: NEXT a
i. First the sign of x is taken care of:

If X is negative, the subroutine jumps to PF-NEGATIVE, takes ABS x and prints the minus sign.
If x is zero, x is deleted from the calculator stack, a ' 0 ' is printed and a return is made from the subroutine.
If $x$ is positive, the subroutine just continues.

| 2DE3 | PRINT-FP | RST | 0028,FP-CALC | Use the calculator |
| :---: | :---: | :---: | :---: | :---: |
|  |  | DEFB | +31,duplicate |  |
|  |  | DEFB | +36,less-0 | $\mathrm{x},(1 / 0)$ Logical value of x . |
|  |  | DEFB | +00,jump-true | $x$ |
|  |  | DEFB | +0B,to PF-NEGTVE | x |
|  |  | DEFB | +31,duplicate | x, x |
|  |  | DEFB | +37,greater-0 | $\mathrm{x},(1 / 0)$ Logical value of X . |


|  |  | DEFB | +00,jump-true | X |
| :---: | :---: | :---: | :---: | :---: |
|  |  | DEFB | +0D, to PF-POSTVE | $x$ Hereafter $\mathrm{x}^{\prime}=A B S$ x. |
|  |  | DEFB | +02,delete | - |
|  |  | DEFB | +38,end-calc | - |
|  |  | LD | A, +30 | Enter the character code for ' 0 '. |
|  |  | RST | 0010,PRINT-A-1 | Print the ' 0 '. |
|  |  | RET |  | Finished as the 'last value' is zero. |
| 2DF2 | PF-NEGTVE | DEFB | +2A, abs | $x^{\prime} x^{\prime}=A B S$ x. |
|  |  | DEFB | +38, end-calc | $\mathrm{x}^{\prime}$ |
|  |  | LD | A,+2D | Enter the character code for '-'. |
|  |  | RST | 0010,PRINT-A-1 | Print the '-'. |
|  |  | RST | 0028,FP-CALC | Use the calculator again. |
| 2DF8 | PF-POSTVE | DEFB | +A0,stk-zero | The 15 bytes of mem-3, mem-4 |
|  |  | DEFB | +C3,st-mem-3 | and mem-5 are now initialised to |
|  |  | DEFB | +C4,st-mem-4 | zero to be used for a print |
|  |  | DEFB | +C5,st-mem-5 | buffer and two counters. |
|  |  | DEFB | +02, delete | The stack is cleared, except for $x^{\prime}$. |
|  |  | DEFB | +38, end-calc | $\mathrm{x}^{\prime}$ |
|  |  | EXX |  | H'L', which is used to hold |
|  |  | PUSH | HL | calculator offsets, (e.g. for |
|  |  | EXX |  | 'STR\$') is saved on the machine |

ii. This is the start of a loop which deals with large numbers. However every number x is first split into its integer part i and the fractional part $f$. If $i$ is a small integer, i.e. if $-65535<=i<=65535$, it is stored in D'E' for insertion into the print buffer.

| 2E01 | PF-LOOP | RST | 0028,FP-CALC | Use the calculator again. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | DEFB | +31,duplicate | $\mathrm{x}^{\prime} \times{ }^{\prime}$ |
|  |  | DEFB | +27, int | $\mathrm{x}^{\prime}$, INT ( $\mathrm{x}^{\prime}$ )=i |
|  |  | DEFB | +C2,st-mem-2 | (i is stored in mem-2). |
|  |  | DEFB | +03,subtract | $x^{\prime}-\mathrm{i}=\mathrm{f}$ |
|  |  | DEFB | +E2,get-mem-2 | f,i |
|  |  | DEFB | +01,exchange | i,f |
|  |  | DEFB | +C2,st-mem-2 | (f is stored in mem-2). |
|  |  | DEFB | +03, delete | i |
|  |  | DEFB | +38,end-calc | i |
|  |  | LD | A,(HL) | Is i a small integer (first byte |
|  |  | AND | A | zero) i.e. is $\mathrm{ABS} \mathrm{i}<=65535$ ? |
|  |  | JR | NZ,2E56,PF-LARGE | Jump if it is not |
|  |  | CALL | 2D7F,INT-FETCH | i is copied to DE ( i , like $\mathrm{x}^{\prime},>=0$ ). |
|  |  | LD | B,+10 | $B$ is set to count 16 bits. |
|  |  | LD | A, D | D is copied to A for testing: |
|  |  | AND | A | Is it zero? |
|  |  | JR | NZ,2E1E,PF-SAVE | Jump if it is not zero. |
|  |  | OR | E | Now test E. |
|  |  | JR | Z,2E24,PF-SMALL | Jump if DE zero: x is a pure fraction. |
|  |  | LD | D, E | Move E to D and set B for 8 |
|  |  | LD | B,+08 | bits: D was zero and E was not. |
| 2E1E | PF-SAVE | PUSH | DE | Transfer DE to D'E', via the |
|  |  | EXX |  | machine stack, to be moved |
|  |  | POP | DE | into the print buffer at |
|  |  | EXX |  | PF-BITS. |
|  |  | JR | 2E78,PF-BITS | Jump forward. |

iii. Pure fractions are multiplied by $10^{\wedge} \mathrm{n}$, where n is the approximate number of leading zeros after the decimal; and -n is added to the second byte of mem-5, which holds the number of digits needed before the decimals; a negative number here indicates leading zeros after the decimal;

| $2 E 24$ | PF-SMALL | RST | 0028,FP-CALC |
| :--- | :--- | :--- | :--- |
|  |  | DEFB | +E2,get-mem-2 |

Note that the stack is now unbalanced. An extra byte 'DEFB +02, delete' is needed at $2 E 25$, immediately after the RST 0028. Now an expression like " 2 " +STR $\$ 0.5$ is evaluated incorrectly as 0.5 ; the zero left on the stack displaces the " 2 " and is treated as a null string. Similarly all the string comparisons can yield incorrect values if the second string takes the form STR $\$ \mathrm{x}$ where x is numerically less than 1 ; e.g. the expression " 50 "<STR\$ 0.1 yields the logical value "true"; once again "" is used instead of " 50 ".

| LD | A,(HL) | The exponent byte e of $f$ is copied to A. |
| :---: | :---: | :---: |
| SUB | +7E | A becomes e-126 dec i.e. $e^{\prime}+2$, where $e^{\prime}$ is the true exponent of $f$. |
| CALL | 2DC1,LOG ( $2^{\wedge}$ A $)$ | The construction $\mathrm{A}=\mathrm{ABS}$ INT (LOG $\left(2^{\wedge} A\right)$ ) is performed (LOG is to base 10); i.e. A=n, |
| LD | D,A | say: n is copied from A to D . |
| LD | A,(mem-5-2nd) | The current count is collected |
| SUB | D | from the second byte of mem-5 |
| LD | (mem-5-2nd),A | and n is subtracted from it. |
| LD | A, D | n is copied from D to A . |
| CALL | 2D4F,E-TO-FP | $y=f^{\star} 10^{\wedge} n$ is formed and stacked. |
| RST | 0028,FP-CALC | i, y |
| DEFB | +31,duplicate | i, y, y |
| DEFB | +27,int | i, y, (INT (y) = i2 |
| DEFB | +C1,st-mem-1 | (i2 is copied to mem-1). |
| DEFB | +03, subtract | i, y-i2 |
| DEFB | +E1,get-mem-1 | i, y-i2, i2 |
| DEFB | +38,end-calc | $\mathrm{i}, \mathrm{f} 2, \mathrm{i} 2$ ( $\mathrm{f} 2=\mathrm{y}-\mathrm{i} 2$ ) |
| CALL | 2DD5,FP-TO-A | $i 2$ is transferred from the stack to A . |
| PUSH LD | HL (mem-3-1st),A | The pointer to $\uparrow 2$ is saved. i2 is stored in the first byte of |
|  |  | mem-3: a digit for printing. |
| DEC | A | i2 will not count as a digit for |
| RLA |  | printing if it is zero; $A$ is |
| SBC | A,A | manipulated so that zero will |
| INC | A | produce zero but a non-zero digit will produce 1. |
| LD | HL,+5CAB | The zero or one is inserted into |
| LD | (HL), A | the first byte of mem-5 (the no. |
| INC | HL | of digits for printing) and added |
| ADD | A,(HL) | to the second byte of mem-5 |
| LD | (HL), A | (the number of digits before the decimal). |
| POP | HL | The pointer to f2 is restored. |
| JP | 2ECF,PF-FRACTN | Jump to store f2 in buffer (HL |

iv. Numbers greater than $2^{\wedge} 27$ are similarly multiplied by $2^{\wedge}(-n+7)$, reducing the number of digits before the decimal to 8 , and the loop is re-entered at PF-LOOP.

| 2E56 | PF-LARGE | $\begin{aligned} & \text { SUB } \\ & \text { CP } \\ & \text { JR } \\ & \text { CALL } \\ & \text { SUB } \\ & \text { LD } \\ & \text { LD } \\ & \text { ADD } \\ & \text { LD } \end{aligned}$ | $\begin{aligned} & +80 \\ & +1 \mathrm{C} \\ & \text { C,2E6F,PF-MEDIUM } \\ & \text { 2DC1,LOG (2^A) } \\ & +07 \\ & \text { B,A } \\ & \text { HL,+5CAC } \\ & \text { A,(HL) } \\ & \text { (HL),A } \end{aligned}$ | e-80 hex = e', the true exponent of $i$. Is e' less than 28 decimal? Jump if it is less. n is formed in A . And reduced to $\mathrm{n}-7$. Then copied to B. $\mathrm{n}-7$ is added in to the second byte of mem-5, the number of digits required before the |
| :---: | :---: | :---: | :---: | :---: |


| LD | A,B | decimal in $x$. <br> Then $i$ is multiplied by |
| :--- | :--- | :--- |
| NEG | 10^(-n+7) |  |
| CALL | 2D4F,E-TO-FP | This will bring it into medium <br> range for printing. |
| JR | 2E01,PF-LOOP | Round the loop again to deal <br> with the now medium-sized <br> number. |

v . The integer part of x is now stored in the print buffer in mem-3 and mem-4.


Note that the case where i us a small integer (less than 65536) re-enters here.

| 2E7B | PF-BITS | SLA | E | The mantissa of $i$ is now rotated |
| :---: | :---: | :---: | :---: | :---: |
|  |  | RL | D | left and all the bits of $i$ are thus |
|  |  | EXX |  | shifted into mem-4 and each |
|  |  | RL | E | byte of mem-4 is decimal adjusted at each shift. |
|  |  | RL | D | All four bytes of i. |
|  |  | EXX |  | Back to the main registers. |
|  |  | LD | HL, +5CAA | Address of fifth byte of mem-4 |
|  |  | LD | C, +05 | to HL; count of 5 bytes to C. |
| 2E8A | PF-BYTES | LD | A,(HL) | Get the byte of mem-4. |
|  |  | ADC | A, A | Shift it left, taking in the new bit. |
|  |  | DAA |  | Decimal adjust the byte. |
|  |  | LD | (HL), A | Restore it to mem-4. |
|  |  | DEC | HL | Point to next byte of mem-4. |
|  |  | DEC | C | Decrease the byte count by one. |
|  |  | JR | NZ,2E8A,PF-BYTES | Jump for each byte of mem-4. |
|  |  | DJNZ | 2E7B,PF-BITS | Jump for each bit of INT (x). |

Decimal adjusting each byte of mem-4 gave 2 decimal digits per byte, there being at most 9 digits. The digits will now be re-packed, one to a byte, in mem-3 and mem-4, using the instruction RLD.


|  |  | INC INC INC LD | DE <br> (mem-5-1st) <br> (mem-5-2nd) C,+00 |
| :---: | :---: | :---: | :---: |
| 2EB3 | PF-TEST-2 | $\begin{aligned} & \text { BIT } \\ & \text { JR } \\ & \text { INC } \end{aligned}$ | $\begin{aligned} & 0, \mathrm{~B} \\ & \text { Z,2EB8,PF,ALL-9 } \\ & \mathrm{HL} \end{aligned}$ |
| 2EB8 | PF-ALL-9 | $\begin{aligned} & \text { DJNZ } \\ & \text { LD } \\ & \text { SUB } \\ & \text { JR } \\ & \text { DEC } \end{aligned}$ | 2EA1,PF-DIGITS <br> A,(mem-5-1st) +09 <br> C,2ECB,PF-MORE (mem-5-1st) |
|  |  | $\begin{aligned} & \text { LD } \\ & \text { CP } \end{aligned}$ | $\begin{aligned} & \mathrm{A},+04 \\ & (\mathrm{mem}-4-4 \mathrm{th}) \end{aligned}$ |
| 2 ECB | PF-MORE | JR RST DEFB DEFB DEFB | 2FOC,PF-ROUND 0028,FP-CALC <br> +02, delete <br> +E2,get-mem-2 <br> +38 ,end-calc |

vi. The fractional part of $x$ is now stored in the print buffer.

| 2ECF | PF-FRACTN | EX <br> CALL | $\begin{aligned} & \text { DE,HL } \\ & \text { 2FBA,FETCH-TWO } \end{aligned}$ |
| :---: | :---: | :---: | :---: |
|  |  | EXX |  |
|  |  | LD | A, +80 |
|  |  | SUB | L |
|  |  | LD | L,+00 |
|  |  | SET | 7,D |
|  |  | EXX |  |
|  |  | CALL | 2FDD,SHIFT-FP |
| 2EDF | PF-FRN-LP | LP | A,(mem-5-1st) |
|  |  | CP | +08 |
|  |  | JR | C,2EEC,PR-FR-DGT |
|  |  | EXX |  |
|  |  | RL | D |
|  |  | EXX |  |
|  |  | JR | 2F0C,PF-ROUND |
| 2EEC | PF-FR-DGT | LD | BC,+0200 |
| 2EEF | PF-FR-EXX | LD | A, E |
|  |  | CALL | $2 \mathrm{~F} 8 \mathrm{~B}, \mathrm{CA}=10^{*} \mathrm{~A}+\mathrm{C}$ |
|  |  | LD | E,A |
|  |  | LD | A,D |
|  |  | CALL | $2 \mathrm{~F} 8 \mathrm{~B}, \mathrm{CA}=10^{*} \mathrm{~A}+\mathrm{C}$ |
|  |  | LD | D,A |
|  |  | PUSH | BC |
|  |  | EXX |  |
|  |  | POP | BC |
|  |  | DJNZ | 2EEF,PF-FR-EXX |
|  |  | LD | HL, +5CA1 |
|  |  | LD | A, C |
|  |  | LD | C,(mem-5-1st) |
|  |  | ADD | HL, BC |
|  |  | LD | (HL), A |

Point to next destination.
One more digit for printing, and one more before the decimal. Change the flag from leading zero to other zero.
The source pointer needs to be incremented on every second passage through the loop, when $B$ is odd.
Jump back for all 9 digits. Get counter: were there 9 digits excluding leading zeros? If not, jump to get more digits.
Prepare to round: reduce count to 8.
Compare 9th digit, byte 4 of mem-4, with 4 to set carry for rounding up.
Jump forward to round up. Use the calculator again. - (i is now deleted).
f
f

DE now points to $f$.
The mantissa of $f$ is now in D', E',D,E.
Get the exchange registers.
The exponent of $f$ is reduced to zero, by shifting the bits of $f 80$ hex - e places right, where L' contained e.
True numerical bit to bit 7 of D'.
Restore the main registers.
Now make the shift.
Get the digit count.
Are there already 8 digits?
If not, jump forward.
If 8 digits, just use f to round $i$
up, rotating $D$ ' left to set the carry.
Restore main registers and jump forward to round up
Initial zero to C, count of 2 to $B$. D'E'DE is multiplied by 10 in 2 stages, first DE then D'E', each byte by byte in 2 steps, and the integer part of the result is obtained in C to be passed into the print buffer.
The count and the result alternate between BC and $\mathrm{B}^{\prime} \mathrm{C}^{\prime}$.

Look back once through the exchange registers.
The start - 1st byte of mem-3.
Result to A for storing.
Count of digits so far in number to C .
Address the first empty byte.
Store the next digit.

| INC | (mem-5-1st) | Step up the count of digits. |
| :--- | :--- | :--- |
| JR | 2EDF,PF-FRN-LP | Loop back until there are 8 |
|  |  | digits. |

vii. The digits stored in the print buffer are rounded to a maximum of 8 digits for printing.

| 2F0C | PF-ROUND | PUSH | AF | Save the carry flag for the rounding. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | LD | HL, +5CA1 | Base address of number: mem-3, byte 1. |
|  |  | LD | C,(mem-5-1st) | Offset (number of digits in |
|  |  | LD | B,+00 | number) to BC. |
|  |  | ADD | HL, BC | Address the last byte of the number. |
|  |  | LD | B,C | Copy C to B as the counter. |
|  |  | POP | AF | Restore the carry flag. |
| 2F18 | PF-RND-LP | DEC | HL | This is the last byte of the number. |
|  |  | LD | A,(HL) | Get the byte into A. |
|  |  | ADC | A, +00 | Add in the carry i.e. round up. |
|  |  | LD | (HL), A | Store the rounded byte in the buffer. |
|  |  | AND | A | If the byte is 0 or 10, B will be |
|  |  | JR | Z,2F25,PF-R-BACK | decremented and the final zero |
|  |  | CP | +0A | (or the 10) will not be counted for printing. |
|  |  | CCF |  | Reset the carry for a valid digit. |
|  |  | JR | NC,2F2D,PF-COUNT | Jump if carry reset. |
| 2F25 | PF-R-BACK | DJNZ | 2F18,PF-RND-LP | Jump back for more rounding or more final zeros. |
|  |  | LD | (HL) , +01 | There is overflow to the left; |
|  |  | INC | B | an extra 1 is needed here. |
|  |  | INC | (mem-5-2nd) | It is also an extra digit before the decimal. |
| 2F2D | PF-COUNT | LD | (mem-5-1st),B | B now sets the count of the digits to be printed (final zeros will not be printed). |
|  |  | RST | 0028,FP-CALC | f is to be deleted. |
|  |  | DEFB | +02, delete | - |
|  |  | DEFB | +38,end-calc | - |
|  |  | $\begin{aligned} & \text { EXX } \\ & \text { POP } \end{aligned}$ | HL | The calculator offset saved on the stack is restored to H'L'. |
|  |  | EXX |  |  |

viii. The number can now be printed. First $C$ will be set to hold the number of digits to be printed, not counting final zeros, while $B$ will hold the number of digits required before the decimal.


The next entry point is also used to print the digits needed for E-format printing.
2F4A
PF-E-SBRN XOR
A
Start by setting A to zero.


## THE 'CA=10*A+C' SUBROUTINE'

This subroutine is called by the PRINT-FP subroutine to multiply each byte of D'E'DE by 10 and return the integer part of the result in the $C$ register. On entry, the A register contains the byte to be multiplied by 10 and the $C$ register contains the carry over from the previous byte. On return, the A register contains the resulting byte and the C register the carry forward to the next byte.
$\left.\begin{array}{llll}\mathrm{CA}=10^{*} \mathrm{~A}+\mathrm{C} & \mathrm{PUSH} & \mathrm{DE} & \begin{array}{l}\text { Save whichever DE pair is in use. } \\ \text { LD }\end{array} \\ & \text { LD } & \mathrm{H}, \mathrm{A} & \\ \text { Lopy the multiplicand from } \mathrm{A}\end{array}\right)$

## THE 'PREPARE TO ADD' SUBROUTINE.

This subroutine is the first of four subroutines that are used by the main arithmetic operation routines - SUBTRACTION, ADDITION, MULTIPLICATION and DIVISION
This particular subroutine prepares a floating-point number for addition, mainly by replacing the sign bit with a true numerical bit 1, and negating the number (two's complement) if it is negative. The exponent is returned in the A register and the first byte is set to Hex. 00 for a positive number and Hex.FF for a negative number.

| 2F9B | PREP-ADD | LD | A,(HL) | Transfer the exponent to A . |
| :---: | :---: | :---: | :---: | :---: |
|  |  | LD | (HL), +00 | Presume a positive number. |
|  |  | AND | A | If the number is zero then the |
|  |  | RET | Z | preparation is already finished. |
|  |  | INC | HL | Now point to the sign byte. |
|  |  | BIT | 7,(HL) | Set the zero flag for positive number. |
|  |  | SET | 7,(HL) | Restore the true numeric bit. |
|  |  | DEC | HL | Point to the first byte again. |
|  |  | RET | Z | Positive numbers have been prepared, but negative numbers need to be twos complemented. |
|  |  | PUSH | BC | Save any earlier exponent. |
|  |  | LD | BC, +0005 | There are 5 bytes to be handled. |
|  |  | ADD | HL, BC | Point one-past the last byte. |
|  |  | LD | B,C | Transfer the '5' to B. |
|  |  | LD | C,A | Save the exponent in C . |
|  |  | SCF |  | Set carry flag for negation. |
| 2FAF | NEG-BYTE | DEC | HL | Point to each byte in turn. |
|  |  | LD | A,(HL) | Get each byte. |
|  |  | CPL |  | One's complement the byte. |
|  |  | ADC | A, +00 | Add in carry for negation. |
|  |  | LD | (HL), A | Restore the byte. |
|  |  | DJNZ | 2FAF,NEG-BYTE | Loop the '5' times. |
|  |  | LD | A,C | Restore the exponent to A . |
|  |  | POP | BC | Restore any earlier exponent. |
|  |  | RET |  | Finished. |

## THE 'FETCH TWO NUMBERS' SUBROUTINE

This subroutine is called by ADDITION, MULTIPLICATION and DIVISION to get two numbers from the calculator stack and put them into the register, including the exchange registers.
On entry to the subroutine the HL register pair points to the first byte of the first number and the DE register pair points to the first byte of the second number.
When the subroutine is called from MULTIPLICATION or DIVISION the sign of the result is saved in the second byte of the first number.

| 2FBA FETCH-TWO | PUSH <br> PUSH | $\begin{aligned} & \mathrm{HL} \\ & \mathrm{AF} \end{aligned}$ | HL is preserved. $A F$ is preserved. |
| :---: | :---: | :---: | :---: |
| Call the five bytes of the and the second number | first num $-\mathrm{N} 1, \mathrm{~N}$ | $\begin{aligned} & \text { - M1, M2 } \\ & \text { N4 \& N5. } \end{aligned}$ |  |
|  | LD | C,(HL) | M1 to C. |
|  | INC | HL | Next. |
|  | LD | B,(HL) | M2 to B. |
|  | LD | (HL),A | Copy the sign of the result to (HL). |
|  | INC | HL | Next. |
|  | LD | A,C | M1 to A. |
|  | LD | C,(HL) | M3 to C. |
|  | PUSH | BC | Save M2 \& M3 on the machine stack. |
|  | INC | HL | Next. |
|  | LD | C,(HL) | M4 to C. |
|  | INC | HL | Next. |
|  | LD | B,(HL) | M5 to B. |
|  | EX | DE,HL | HL now points to N1. |
|  | LD | D,A | M1 to D. |
|  | LD | E,(HL) | N1 to E. |
|  | PUSH | DE | Save M1 \& N1 on the machine stack. |
|  | INC | HL | Next. |
|  | LD | D,(HL) | N2 to D. |
|  | INC | HL | Next. |
|  | LD | E,(HL) | N3 to E. |
|  | PUSH | DE | Save N2 \& N3 on the machine stack. |
|  | EXX |  | Get the exchange registers. |
|  | POP | DE | N2 to D' \& N3 to E'. |
|  | POP | HL | M1 to $\mathrm{H}^{\prime}$ \& N 1 to $\mathrm{L}^{\prime}$. |
|  | POP | BC | M2 to $\mathrm{B}^{\prime}$ \& M3 to $\mathrm{C}^{\prime}$. |
|  | EXX |  | Get the original set of registers. |
|  | INC | HL | Next. |
|  | LD | D,(HL) | N4 to D. |
|  | INC | HL | Next. |
|  | LD | E,(HL) | N5 to E. |
|  | POP | AF | Restore the original AF. |
|  | POP | HL | Restore the original HL. |
|  | RET |  | Finished. |
| Summary: | $\begin{aligned} & \text { M1 - M5 are in H', B', C', C, B. } \\ & \text { N1 - N5 are in: L', D', E', D, E. } \\ & \text { HL points to the first byte of the first number. } \end{aligned}$ |  |  |

## THE 'SHIFT ADDEND' SUBROUTINE

This subroutine shifts a floating-point number up to 32 decimal, Hex.20, places right to line it up properly for addition. The number with the smaller exponent has been put in the addend position before this subroutine is called. Any overflow to the right, into the carry, is added back into the number. If the exponent difference is greater than 32 decimal, or the carry ripples right back to the beginning of the number then the number is set to zero so that the addition will not alter the other number (the augend).

| 2FDD | SHIFT-FP | AND | A | If the exponent difference is |
| :---: | :---: | :---: | :---: | :---: |
|  |  | RET | Z | zero, the subroutine returns at |
|  |  | CP | +21 | once. If the difference is greater |
|  |  | JR | NC,2FF9,ADDEND-0 | than Hex.20, jump forward. |
|  |  | PUSH | BC | Save BC briefly. |
|  |  | LD | B,A | Transfer the exponent difference to B to count the shifts right. |
| 2FE5 | ONE-SHIFT | EXX |  | Arithmetic shift right for L', |
|  |  | SRA | L | preserving the sign marker bits. |



THE 'ADD-BACK' SUBROUTINE
This subroutine adds back into the number any carry which has overflowed to the right. In the extreme case, the carry ripples right back to the left of the number.
When this subroutine is called during addition, this ripple means that a mantissa of 0.5 was shifted a full 32 places right, and the addend will now be set to zero; when called from MULTIPLICATION, it means that the exponent must be incremented, and this may result in overflow.

| 3004 | ADD-BACK | INC | E | Add carry to rightmost byte. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | RET | NZ | Return if no overflow to left. |
|  |  | INC | D | Continue to the next byte. |
|  |  | RET | NZ | Return if no overflow to left. |
|  |  | EXX |  | Get the next byte. |
|  |  | INC | E | Increment it too. |
|  |  | JR | NZ,300D,ALL-ADDED | Jump if no overflow. |
|  |  | INC | D | Increment the last byte. |
| 300D | ALL-ADDED | EXX |  | Restore the original registers. |
|  |  | RET |  | Finished. |

THE 'SUBTRACTION' OPERATION
(Offset 03 - see CALCULATE below: 'subtract')
This subroutine simply changes the sign of the subtrahend and carried on into ADDITION.
Note that HL points to the minuend and DE points to the subtrahend. (See ADDITION for more details.)
300 F SUBTRACT EX Ex,HL Exange the pointers.
CALL 346E,NEGATE

Change the sign of the subtrahend.
Exchange the pointers back and continue into ADDITION.

## THE 'ADDITION' OPERATION

(Offset OF - see CALCULATE below: 'addition')
The first of three major arithmetical subroutines, this subroutine carries out the floating-point addition of two numbers, each with a 4byte mantissa and a 1-byte exponent. In these three subroutines, the two numbers at the top of the calculator stack are added/multiplied/divided to give one number at the top of the calculator stack, a 'last value'.

HL points to the second number from the top, the augend/multiplier/dividend. DE points to the number at the top of the calculator stack, the addend/multiplicand/divisor. Afterwards HL points to the resultant 'last value' whose address can also be considered to be STKEND - 5.

But the addition subroutine first tests whether the 2 numbers to be added are 'small integers'. If they are, it adds them quite simply in HL and BC, and puts the result directly on the stack. No twos complementing is needed before or after the addition, since such numbers are held on the stack in twos complement form, ready for addition.


Note that the number - 65536 decimal can arise here in the form 00 FF 000000 as the result of the addition of two smaller negative integers, e.g. -65000 and -536. It is simply stacked in this form. This is a mistake. The Spectrum system cannot handle this number.

Most functions treat it as zero, and it is printed as $-1 \mathrm{E}-38$, obtained by treating is as 'minus zero' in an illegitimate format.
One possible remedy would be to test for this number at about byte 3032 and, if it is present, to make the second byte 80 hex and the first byte 91 hex, so producing the full five byte floating-point form of the number, i.e. 9180000000 , which causes no problems. See also the remarks in 'truncate' below, before byte 3225, and the Appendix.

303C ADDN-OFLW DEC \begin{tabular}{llll}

HL \& DE \& | Restore the pointer to the first |
| :--- |
| number. | <br>

303E \& FULL-ADDN \& CALL \& 3293,RE-ST-TWO

 

Recond number.
\end{tabular}

The full ADDITION subroutine first calls PREP-ADD for each number, then gets the two numbers from the calculator stack and puts the one with the smaller exponent into the addend position. It then calls SHIFT-FP to shift the addend up to 32 decimal places right to line it up for addition. The actual addition is done in a few bytes, a single shift is made for carry (overflow to the left) if needed, the result is twos complemented if negative, and any arithmetic overflow is reported; otherwise the subroutine jumps to TEST-NORM to normalise the result and return it to the stack with the correct sign bit inserted into the second byte.


|  |  | $\begin{aligned} & \text { EX } \\ & \text { POP } \end{aligned}$ | $\begin{aligned} & \text { DE,HL } \\ & \mathrm{HL} \end{aligned}$ | The result is now in DED'E. Get the pointer to the exponent. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | RRA JR | NC,307C,TEST-NEG | The test for shift ( $\mathrm{H}^{\prime}$, L' were Hex. 00 for positive numbers and Hex.FF for negative numbers). |
|  |  | LD | A,+01 | A counts a single shift right. |
|  |  | CALL | 2FDD, SHIFT-FP | The shift is called. |
|  |  | INC | (HL) | Add 1 to the exponent; this |
|  |  | JR | Z,309F,ADD-REP-6 | may lead to arithmetic overflow. |
| 307C | TEST-NEG | EXX |  | Test for negative result: get |
|  |  | LD | A,L | sign bit of L' into A (this now |
|  |  | AND | +80 | correctly indicates the sign of |
|  |  | EXX |  | the result). |
|  |  | INC | HL | Store it in the second byte |
|  |  | LD | (HL), A | position of the result on |
|  |  | DEC | HL | the calculator stack. |
|  |  | JR | Z,30A5,GO-NC-MLT | If it is zero, then do not twos complement the result. |
|  |  | LD | A, E | Get the first byte. |
|  |  | NEG |  | Negate it. |
|  |  | CCF |  | Complement the carry for |
|  |  | LD | E,A | continued negation, and store byte. |
|  |  | LD | A, D | Get the next byte. |
|  |  | CPL |  | Ones complement it. |
|  |  | ADC | A, +00 | Add in the carry for negation. |
|  |  | LD | D,A | Store the byte. |
|  |  | EXX |  | Proceed to get next byte into |
|  |  | LD | A, E | the A register. |
|  |  | CPL |  | Ones complement it. |
|  |  | ADC | A, +00 | Add in the carry for negation. |
|  |  | LD | E,A | Store the byte. |
|  |  | LD | A, D | Get the last byte. |
|  |  | CPL |  | Ones complement it. |
|  |  | ADC | A, +00 | Add in the carry for negation. |
|  |  | JR | NC,30A3,END-COMPL | Done if no carry. |
|  |  | RRA |  | Else, get .5 into mantissa and |
|  |  | EXX |  | add 1 to the exponent; this will |
|  |  | INC | (HL) | be needed when two negative |
|  |  |  |  | numbers add to give an exact |
|  |  |  |  | power of 2, and it may lead to |
| 309F | ADD-REP-6 | JP | Z,31AD,REPORT-6 | Give the error if required. |
|  |  | EXX |  |  |
| 30A3 | END-COMPL | LD | D,A | Store the last byte. |
|  |  | EXX |  |  |
| 30A5 | GO-NC-MLT | XOR | A | Clear the carry flag. |
|  |  | JP | 3155,TEST-NORM | Exit via TEST-NORM. |

## THE 'HL=HL*DE' SUBROUTINE

This subroutine is called by 'GET-HL*DE' and by 'MULTIPLICATION' to perform the 16 -bit multiplication as stated. Any overflow of the 16 bits available is dealt with on return from the subroutine.



## THE 'PREPARE TO MULTIPLY OR DIVIDE' SUBROUTINE

This subroutine prepares a floating-point number for multiplication or division, returning with carry set if the number is zero, getting the sign of the result into the A register, and replacing the sign bit in the number by the true numeric bit, 1.
30 CO PREP-M/D
If the number is zero, return

| PREP-M/D | CALL | $34 E 9$, |
| :--- | :--- | :--- |
|  | RET | C |
|  | INC | HL |
|  | XOR | $(\mathrm{HL})$ |
|  |  |  |
|  |  |  |
|  | SET | $7,(\mathrm{HL})$ |
|  | DEC | HL |
|  | RET |  |

with the carry flag set.
Point to the sign byte.
Get sign for result into A (like
signs give plus, unlike give minus); also reset the carry flag
Set the true numeric bit.
RET
Point to the exponent again
Return with carry flag reset.
THE 'MULTIPLICATION' OPERATION
(Offset 04 - see CALCULATE below: 'multiply')
This subroutine first tests whether the two numbers to be multiplied are 'small integers'. If they are, it uses INT-FETCH to get them from the stack, HL=HL*DE to multiply them and INT-STORE to return the result to the stack. Any overflow of this 'short multiplication' (i.e. if the result is not itself a 'small integer') causes a jump to multiplication in full five byte floating-point form (see below).

| 30CA | multiply | LD | A,(DE) | Test whether the first bytes of |
| :---: | :---: | :---: | :---: | :---: |
|  |  | OR | (HL) | both numbers are zero. |
|  |  | JR | NZ,30F0,MULT-LONG | If not, jump for 'long' multiplication. |
|  |  | PUSH | DE | Save the pointers: to the second number. |
|  |  | PUSH | HL | And to the first number. |
|  |  | PUSH | DE | And to the second number yet again. |
|  |  | CALL | 2D7F,INT-FETCH | Fetch sign in C, number in DE. |
|  |  | EX | DE,HL | Number to HL now. |
|  |  | EX | (SP), HL | Number to stack, second pointer to HL. |
|  |  | LD | B,C | Save first sign in B. |
|  |  | CALL | 2D7F,INT-FETCH | Fetch second sign in C , number in DE. |
|  |  | LD | A, B | Form sign of result in A: like |
|  |  | XOR | C | signs give plus (00), unlike give minus (FF). |
|  |  | LD | C, A | Store sign of result in C. |
|  |  | POP | HL | Restore the first number to HL. |
|  |  | CALL | 30A9,HL=HL*DE | Perform the actual multiplication. |
|  |  | EX | DE,HL | Store the result in DE. |
|  |  | POP | HL | Restore the pointer to the first number. |
|  |  | JR | C,30EF,MULT-OFLW | Jump on overflow to 'full' multiplication. |
| 30E5 |  | LD | A, D | These 5 bytes ensure that |


| OR | E <br> JR | NZ,30EA,MULT-RSLT <br> C,A | 00 FF 000000 is replaced by <br> zero; that they should not be <br> needed if this number were <br> excluded from the system (see |
| :--- | :--- | :--- | :--- |
| after 303B) above). |  |  |  |

The full MULTIPLICATION subroutine prepares the first number for multiplication by calling PREP-M/D, returning if it is zero; otherwise the second number is prepared by again calling PREP-M/D, and if it is zero the subroutine goes to set the result to zero. Next it fetches the two numbers from the calculator stack and multiplies their mantissas in the usual way, rotating the first number (treated as the multiplier) right and adding in the second number (the multiplicand) to the result whenever the multiplier bit is set. The exponents are then added together and checks are made for overflow and for underflow (giving the result zero). Finally, the result is normalised and returned to the calculator stack with the correct sign bit in the second byte.

| XOR | A | A is set to Hex. 00 so that the sign of the first number will go into A. |
| :---: | :---: | :---: |
| CALL | 30C0,PREP-M/D | Prepare the first number, and |
| RET | C | return if zero. (Result already zero.) |
| EXX |  | Exchange the registers. |
| PUSH | HL | Save the next literal address. |
| EXX |  | Exchange the registers. |
| PUSH | DE | Save the pointer to the multiplicand. |
| EX | DE,HL | Exchange the pointers. |
| CALL | 30C0,PREP-M/D | Prepare the 2nd number. |
| EX | DE,HL | Exchange the pointers again. |
| JR | C,315D,ZERO-RSLT | Jump forward if 2nd number is zero. |
| PUSH | HL | Save the pointer to the result. |
| CALL | 2FBA,FETCH-TWO | Get the two numbers from the stack. |
| LD | A, B | M5 to A (see FETCH-TWO). |
| AND | A | Prepare for a subtraction. |
| SBC | HL,HL | Initialise HL to zero for the result. |
| EXX |  | Exchange the registers. |
| PUSH | HL | Save M1 \& N1 (see FETCH-TWO). |
| SBC | HL,HL | Also initialise H'L' for the result. |
| EXX |  | Exchange the registers. |
| LD | B,+21 | B counts 33 decimal, Hex.21, shifts. |
| JR | 3125,STRT-MLT | Jump forward into the loop. |

Now enter the multiplier loop.

3114 MLT-LOOP JR \begin{tabular}{ll}

NC,311B,NO-ADD \& | Jump forward to NO-ADD if no |
| :--- |
| carry, i.e. the multiplier bit was |
| reset. | <br>

\& ADD <br>
EXX \& HL,DE
\end{tabular}

|  |  | EXX |  | H'L'HL. |
| :---: | :---: | :---: | :---: | :---: |
| 311B | NO-ADD | EXX |  | Whether multiplicand was added |
|  |  | RR | H | or not, shift result right in |
|  |  | RR | L | H'L'HL, i.e. the shift is done by |
|  |  | EXX |  | rotating each byte with carry, so |
|  |  | RR | H | that any bit that drops into the |
|  |  | RR | L | carry is picked up by the next |
|  |  |  |  | byte, and the shift continued |
| 3125 | STRT-MLT | EXX |  | Shift right the multiplier in |
|  |  | RR |  |  |
|  |  | RR | B | B'C'CA (see FETCH-TWO \& above). |
|  |  | RR | C | A final bit dropping into the |
|  |  | EXX |  | carry will trigger another add of |
|  |  | RR | C | the multiplicand to the result. |
|  |  | RRA |  |  |
|  |  | DJNZ | 3114,MLT-LOOP | Loop 33 times to get all the bits. |
|  |  | EX | DE,HL | Move the result from: |
|  |  | EXX |  |  |
|  |  | EX | DE,HL | H'L'HL to D'E'DE. |
|  |  | EXX |  |  |

Now add the exponents together.

|  |  | POP | BC | Restore the exponents - M1 \& N1. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | POP | HL | Restore the pointer to the exponent byte. |
|  |  | LD | A, B | Get the sum of the two |
|  |  | ADD | A,C | exponent bytes in A, and the correct carry. |
|  |  | JR <br> AND | NZ,313B,MAKE-EXPT A | If the sum equals zero then clear the carry; else leave it unchanged. |
| 313B | MAKE-EXPT | $\begin{aligned} & \text { DEC } \\ & \text { CCF } \end{aligned}$ | A | Prepare to increase the exponent by Hex. 80. |

The rest of the subroutine is common to both MULTIPLICATION and DIVISION.
$\left.\begin{array}{lll}\text { 313D } & \text { DIVN-EXPT } \\ & \begin{array}{ll}\text { RLA } \\ \text { CCF }\end{array} & \begin{array}{l}\text { These few bytes very cleverly } \\ \text { make the correct exponent } \\ \text { byte. }\end{array} \\ \text { Rotating left then right gets the } \\ \text { exponent byte (true exponent } \\ \text { plus Hex.80) into A. } \\ \text { If the sign flag is reset, no report } \\ \text { of arithmetic overflow needed. } \\ \text { Report the overflow if carry } \\ \text { reset. }\end{array}\right\}$

The remainder of the subroutine deals with normalisation and is common to all the arithmetic routines.

| 3155 | TEST-NORM | JR | NC,316C,NORMALISE | If no carry then normalise now. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | LD | A,(HL) | Else, deal with underflow (zero |
|  |  | AND | A | result) or near underflow |
| 3159 | NEAR-ZERO | LD | A, +80 | (result 2**-128): |
|  |  | JR | Z,315E,SKIP-ZERO | return exponent to $A$, test if $A$ |
| $\begin{aligned} & 315 \mathrm{D} \\ & 315 \mathrm{E} \end{aligned}$ | ZERO-RSLT <br> SKIP-ZERO | XOR | A | is zero (case $\left.2^{* *}-128\right)$ and if so |
|  |  | EXX |  | produce $2^{* *}-128$ if number is |
|  |  | AND | D | normal; otherwise produce zero |
|  |  | CALL | 2FFB,ZEROS-4/5 | The exponent must then be set |
|  |  | RLCA |  | to zero (for zero) or 1 (for 2**-128). |
|  |  | LD | (HL), A | Restore the exponent byte. |
|  |  | JR | C,3195,OFLOW-CLR | Jump if case $2^{* *}$-128. |
|  |  | INC | HL | Otherwise, put zero into second |
|  |  | LD | (HL), A | byte of result on the calculator |
|  |  | DEC | HL | stack. |
|  |  | JR | 3195,OFLOW-CLR | Jump forward to transfer the |

The actual normalisation operation.

| $\begin{aligned} & 316 \mathrm{C} \\ & 316 \mathrm{E} \end{aligned}$ | NORMALISE SHIFT-ONE | LD | B, +20 | Normalise the result by up to 32 |
| :---: | :---: | :---: | :---: | :---: |
|  |  | EXX |  | decimal, Hex.20, shifts left of |
|  |  | BIT | 7,D | D'E'DE (with A adjoined) until |
|  |  | EXX |  | bit 7 of $D^{\prime}$ is set. $A$ holds zero |
|  |  | JR | NZ,3186,NORML-NOW | after addition so no precision is |
|  |  | RLCA |  | gained or lost; A holds the fifth |
|  |  | RL | E | byte from B' after multiplica- |
|  |  | RL | D | tion or division; but as only |
|  |  | EXX |  | about 32 bits can be correct, no |
|  |  | RL | E | precision is lost. Note that $A$ is |
|  |  | RL | D |  |
|  |  | EXX |  | at carry .... eventually a random process. |
|  |  | DEC | (HL) | The exponent is decremented on each shift. |
|  |  | JR | Z,3159,NEAR-ZERO | If the exponent becomes zero, then number from $2^{* *}$-129 are rounded up to $2^{* *}-128$. |
|  |  | DJNZ | 316E,SHIFT-ONE | Loop back, up to 32 times. |
|  |  | JR | 315D,ZERO-RSLT | If bit 7 never became 1 then the whole result is to be zero. |

Finish the normalisation by considering the 'carry'.


The final part of the subroutine involves passing the result to the bytes reserved for it on the calculator stack and resetting the pointers.

3195 OFLOW-CLR PUSH \begin{tabular}{lll}

INC \& HL \& | Save the result pointer. |
| :--- |
| Point to the sign byte in the |
| result. | <br>

\& EXX \& <br>
The result is moved from its
\end{tabular}

| POP | BC |  |
| :---: | :---: | :---: |
| LD | A, B | The sign bit is retrieved from |
| RLA |  | its temporary store and trans- |
| RL | (HL) | ferred to its correct position of |
| RRA |  | bit 7 of the first byte of the mantissa. |
| LD | (HL), A | The first byte is stored. |
| INC | HL | Next. |
| LD | (HL), C | The second byte is stored. |
| INC | HL | Next. |
| LD | (HL), D | The third byte is stored. |
| INC | HL | Next. |
| LD | (HL), E | The fourth byte is stored. |
| POP | HL | Restore the pointer to the result. |
| POP | DE | Restore the pointer to second number. |
| EXX |  | Exchange the register. |
| POP | HL | Restore the next literal address. |
| EXX |  | Exchange the registers. |
| RET |  | Finished. |

Report 6 - Arithmetic overflow

| 31 AD REPORT-6 | RST <br> DEFB | $0008, E R R O R-1$ <br> +05 | Call the error handling <br> routine. |
| :--- | :--- | :--- | :--- |

## THE 'DIVISION' OPERATION

(Offset 05 - see CALCULATE below: 'division')
This subroutine first prepared the divisor by calling PREP-M/D, reporting arithmetic overflow if it is zero; then it prepares the dividend again calling PREP-M/D, returning if it is zero. Next fetches the two numbers from the calculator stack and divides their mantissa by means of the usual restoring division, trial subtracting the divisor from the dividend and restoring if there is carry, otherwise adding 1 to the quotient. The maximum precision is obtained for a 4-byte division, and after subtracting the exponents the subroutine exits by joining the later part of MULTIPLICATION.


| XOR | A | Clear A and reset the carry flag. |
| :--- | :--- | :--- |
| LD | B,+DF | B will count upwards from -33 <br> to -1, twos complement, |
|  |  | He. DF to FF, looping on |
| minus and will jump again on |  |  |
|  |  | zero for extra precision. <br> Jump forward into the division <br> loop for the first trial <br> subtraction. |

Now enter the division loop.

| 31D2 | DIV-LOOP | RLA <br> RL <br> EXX <br> RL <br> RL <br> EXX | $\begin{aligned} & \mathrm{C} \\ & \mathrm{C} \\ & \mathrm{~B} \end{aligned}$ | Shift the result left into $B^{\prime} C^{\prime} C A$, shifting out the bits already there, picking up 1 from the carry whenever it is set, and rotating left each byte with carry to achieve the 32 bit shift. |
| :---: | :---: | :---: | :---: | :---: |
| 31 DB | DIV-34TH | $\begin{aligned} & \text { ADD } \\ & \text { EXX } \\ & \text { ADC } \\ & \text { EXX } \end{aligned}$ | HL,HL HL,HL | Move what remains of the dividend left in H'L'HL before the next trial subtraction; if a bit drops into the carry, force no restore and a bit for the quotient, thus retrieving the lost |
|  |  | JR | C,31F2,SUBN-ONLY | bit and allowing a full 32-bit divisor. |
| 31 E 2 | DIV-START | $\begin{aligned} & \text { SBC } \\ & \text { EXX } \end{aligned}$ | HL,DE | Trial subtract divisor in D'E'DE from rest of dividend in H'L'HL; |
|  |  | SBC EXX | HL,DE | there is no initial carry (see previous step). |
|  |  | JR | NC,31F9,NO-RSTORE | Jump forward if there is no carry. |
|  |  | $\begin{aligned} & \text { ADD } \\ & \text { EXX } \end{aligned}$ | HL,DE | Otherwise restore, i.e. add back the divisor. Then clear the carry |
|  |  | ADC | HL,DE | so that there will be no bit for the quotient (the divisor 'did |
|  |  | AND | A | not go'). |
|  |  | JR | 31FA,COUNT-ONE | Jump forward to the counter. |
| 31F2 | SUBN-ONLY | AND | A | Just subtract with no restore |
|  |  | $\begin{aligned} & \text { SBC } \\ & \text { EXX } \end{aligned}$ | HL, DE | and go on to set the carry flag because the lost bit of the divi- |
|  |  | $\begin{aligned} & \text { SBC } \\ & \text { EXX } \end{aligned}$ | HL,DE | dend is to be retrieved and used for the quotient. |
| $\begin{aligned} & 31 \mathrm{F9} \\ & 31 \mathrm{FA} \end{aligned}$ | NO-RSTORE COUNT-ONE | SCF | One for the quotient in $\mathrm{B}^{\prime} \mathrm{C}^{\prime} \mathrm{CA}$. |  |
|  |  | INC | B | Step the loop count up by one. Loop 32 times for all bits. Save any 33rd bit for extra precision (the present carry). Trial subtract yet again for any 34th bit; the PUSH AF above saves this bit too. |
|  |  | JP | M,31D2,DIV-LOOP |  |
|  |  | PUSH | AF |  |
|  |  | JR | Z,31E2,DIV-START |  |

Note: This jump is made to the wrong place. No 34th bit will ever be obtained without first shifting the dividend. Hence important results like $1 / 10$ and $1 / 1000$ are not rounded up as they should be. Rounding up never occurs when it depends on the 34th bit. The jump should have been to 31DB DIV-34TH above: i.e. byte 3200 hex in the ROM should read DA hex ( 128 decimal) instead of E1 hex (225 decimal).

| LD | E,A | Now move the four bytes that <br> form the mantissa bytes of the |
| :--- | :--- | :--- |
| LD | D,C | result from $B^{\prime} C^{\prime} C A$ to D'E'DE. |
| EXX | E,C |  |
| LD | D,B | Then put the 34th and 33rd bits |


| RR | B | into 'B' to be picked up on |
| :--- | :--- | :--- |
| POP | AF | normalisation. |
| RR | B |  |
| EXX | BC | Restore the exponent bytes, M1 <br> POP |
| \& N1. |  |  |

## THE 'INTEGER TRUNCATION TOWARDS ZERO' SUBROUTINE

(Offset 3A - see CALCULATE below: 'truncate')
This subroutine (say $\mathrm{I}(\mathrm{x})$ ) returns the result of integer truncation of x , the 'last value', towards zero. Thus $\mathrm{I}(2.4)$ is 2 and $\mathrm{I}(-2.4)$ is -2 . The subroutine returns at once if $x$ is in the form of a 'short integer'. It returns zero if the exponent byte of $x$ if less than 81 hex (ABS $x$ is less than 1 ). If $\mathrm{I}(\mathrm{x})$ is a 'short integer' the subroutine returns it in that form. It returns x if the exponent byte is A 0 hex or greater ( x has no significant non-integral part). Otherwise the correct number of bytes of $x$ are set to zero and, if needed, one more byte is split with a mask.

| 3214 | truncate | LD | A,(HL) | Get the exponent byte of $X$ into A. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | AND | A | If $A$ is zero, return since $x$ is |
|  |  | RET | Z | already a small integer. |
|  |  | CP | +81 | Compare e, the exponent, to 81 hex. |
|  |  | JR | NC,3221,T-GR-ZERO | Jump if e is greater than 80 hex. |
|  |  | LD | (HL), +00 | Else, set the exponent to zero; |
|  |  | LD | A, +20 | enter 32 decimal, 20 hex, into A |
|  |  | JR | 3272,NIL-BYTES | and jump forward to NIL- <br> BYTES to make all the bits of $x$ be zero. |
| 3221 | T-GR-ZERO | CP | +91 | Compare e to 91 hex, 145 decimal. |
| 3223 |  | JR | NZ,323F,T-SMALL | Jump if e not 91 hex. |

The next 26 bytes seem designed to test whether $x$ is in fact -65536 decimal, i.e. 9180000000 , and if it is, to set it to 00 FF 000000 . This is a mistake. As already stated at byte 303B above, the Spectrum system cannot handle this number. The result here is simply to make INT (-65536) return the value -1. This is a pity, since the number would have been perfectly all right if left alone. The remedy would seem to be simply to omit the 28 bytes from 3223 above to 323 E inclusive from the program.

| 3225 |  | INC | HL | HL is pointed at the fourth byte |
| :---: | :---: | :---: | :---: | :---: |
|  |  | INC | HL | of $x$, where the 17 bits of the |
|  |  | INC | HL | integer part of $x$ end after the first bit. |
|  |  | LD | A, +80 | The first bit is obtained in A. |
|  |  | AND | (HL) | using 80 hex as a mask. |
|  |  | DEC | HL | That bit and the previous 8 bits |
|  |  | OR | (HL) | are tested together for zero. |
|  |  | DEC | HL | $H L$ is pointed at the second byte of $x$. |
|  |  | JR | NZ,3233,T-FIRST | If already non-zero, the test can end. |
|  |  | LD | A, +80 | Otherwise, the test for -65536 is |
|  |  | XOR | (HL) | now completed: 9180000000 will leave the zero flag set now. |
| 3233 | T-FIRST | DEC | HL | HL is pointed at the first byte of |
|  |  |  |  | x. <br> If zero reset the jump is made |
|  |  | LR | (HL), A | If zero reset, the jump is made. The first byte is set to zero. |
|  |  | INC | HL | HL points to the second byte. |


| LD | $(\mathrm{HL}),+$ FF | The second byte is set to FF. <br> DEC |
| :--- | :--- | :--- |
| HL | HL again points to the first <br> byte. |  |
| LD | A,+18 | The last 24 bits are to be zero. |
| JR | 3272, NIL-BYTES | The jump to NIL-BYTES <br> completes the number 00 FF |
|  |  | 000000. |

If the exponent byte of x is between 81 and 90 hex ( 129 and 144 decimal) inclusive, $\mathrm{I}(\mathrm{x})$ is a 'small integer', and will be compressed into one or two bytes. But first a test is made to see whether $x$ is, after all, large.

| 323F | T-SMALL | JR | NC,326D,X-LARGE |
| :--- | :--- | :--- | :--- | | Jump with exponent byte 92 |
| :--- |
| or more (it would be better to |
| jump with 91 too). |

Large values of $x$ remains to be considered.

| 326C | T-EXPNENT | LD | A,(HL) |
| :--- | :--- | :--- | :--- |
| 326D | X-LARGE | SUB | + A0 |
|  |  | RET | $P$ |

Get the exponent byte of $x$ into A.

Subtract 160 decimal, A0 hex, from e.
Return on plus - $x$ has no significant non-integral part. (If the true exponent were reduced to zero, the 'binary point' would come at or after the end of the four bytes of the mantissa).
Else, negate the remainder; this gives the number of bits to become zero (the number of bits after the 'binary point').

Now the bits of the mantissa can be cleared.

| 3272 | NIL-BYTES | PUSH | DE | Save the current value of DE (STKEND). |
| :---: | :---: | :---: | :---: | :---: |
|  |  | EX | DE,HL | Make HL point one past the fifth byte. |
|  |  | DEC | HL | HL now points to the fifth byte of $x$. |
|  |  | LD | B,A | Get the number of bits to be set |
|  |  | SRL | B | to zero in B and divide it by B |
|  |  | SRL | B | to give the number of whole |
|  |  | SRL | B | bytes implied. |
|  |  | JR | Z,3283,BITS-ZERO | Jump forward if the result is zero. |
| 327E | BYTE-ZERO | LD | (HL) , +00 | Else, set the bytes to zero; |
|  |  | DEC | HL | $B$ counts them. |
|  |  | DJNZ | 327E,BYTE-ZERO |  |
| 3283 | BITS-ZERO | AND | +07 | Get $\mathrm{A}(\bmod 8)$; this is the number of bits still to be set to zero. |
|  |  | JR | Z,3290,IX-END | Jump to the end if nothing more to do. |
|  |  | LD | B,A | B will count the bits now. |
|  |  | LD | A, +FF | Prepare the mask. |
| 328A | LESS-MASK | SLA | A | With each loop a zero enters the |
|  |  | DJNZ | 328A,LESS-MASK | mask from the right and thereby a mask of the correct length is produced. |
|  |  | AND | (HL) | The unwanted bits of ( HL ) are |
|  |  | LD | (HL), A | lost as the masking is performed. |
| 3290 | IX-END | EX | DE,HL | Return the pointer to HL. |
|  |  | POP | DE | Return STKEND to DE. |
|  |  | RET |  | Finished. |

THE 'RE-STACK TWO' SUBROUTINE
This subroutine is called to re-stack two 'small integers' in full five byte floating-point form for the binary operations of addition, multiplication and division. It does so by calling the following subroutine twice.

3293 RE-ST-TWO CALL \begin{tabular}{lll}

3269,RESTK-SUB \& | Call the subroutine, and then |
| :--- |
| continue into it for the second |
| call. | <br>

3296 \& RESTK-SUB EX \& DE,HL

 

Exchange the pointers at each <br>
call.
\end{tabular}

## THE 'RE-STACK TWO' SUBROUTINE

(Offset 3D - see CALCULATE below: 're-stack')
This subroutine is called to re-stack one number (which could be a 'small integer') in full five byte floating-point form. It is used for a single number by ARCTAN and also, through the calculator offset, by EXP, LN and 'get-argt'.

3297 RE-STACK |  | LD | A,(HL) | If the first byte is not zero, |
| :--- | :--- | :--- | :--- |
|  | AND | A | return - the number cannot be |
|  | RET | NZ | a 'small integer'. |
|  | PUSH | DE | Save the 'other' pointer in DE. |
|  | CALL | 2D7F,INT-FETCH | Fetch the sign in C and the |
|  |  |  | number in DE. |
|  | XOR | A | Clear the A register. |
|  | INC | HL | Point to the fifth location. |
|  | LD | (HL),A | Set the fifth byte to zero. |
|  | DEC | HL | Point to the fourth location. |
|  | LD | (HL),A | Set the fourth byte to zero: |
|  |  |  | bytes 2 and 3 will hold the man- |
|  |  | tissa. |  |



## THE FLOATING-POINT CALCULATOR

## THE TABLE OF CONSTANTS

This first table holds the five useful and frequently needed numbers zero, one, a half, a half of pi and ten. The numbers are held in a condensed form which is expanded by the STACK LITERALS subroutine, see below, to give the required floating-point form.

|  |  | data: | constant | when expanded gives: exp. mantissa: (Hex.) |
| :---: | :---: | :---: | :---: | :---: |
| 32 C 5 | stk-zero | $\begin{aligned} & \text { DEFB +00 } \\ & \text { DEFB +B0 } \\ & \text { DEFB +00 } \end{aligned}$ | zero | 0000000000 |
| 32 C 8 | stk-one | $\begin{aligned} & \text { DEFB }+40 \\ & \text { DEFB }+B 0 \\ & \text { DEFB }+00 \\ & \text { DEFB }+01 \end{aligned}$ | one | 0000010000 |
| 32CC | stk-half | $\begin{aligned} & \text { DEFB }+30 \\ & \text { DEFB }+00 \end{aligned}$ | a half | 8000000000 |
| 32CE | stk-pi/2 | $\begin{aligned} & \text { DEFB +F1 } \\ & \text { DEFB + 49 } \\ & \text { DEFB +0F } \\ & \text { DEFB +DA } \\ & \text { DEFB +A2 } \end{aligned}$ | a half of pi | 8149 OF DA A2 |
| 32D3 | stk-ten | $\begin{aligned} & \text { DEFB +40 } \\ & \text { DEFB +B0 } \\ & \text { DEFB }+00 \\ & \text { DEFB }+0 A \end{aligned}$ | ten | 0000 OA 0000 |

## THE TABLE OF ADDRESSES

This second table is a look-up table of the addresses of the sixty-six operational subroutines of the calculator. The offsets used to index into the table are derived either from the operation codes used in SCANNING, see 2734, etc., or from the literals that follow a RST 0028 instruction.

| 32 D 7 | offset 00 | label jump-true | address <br> 8F $36$ | 3319 | offset 21 | label <br> tan | address DA |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  |  |  |  |  |  |
|  |  |  |  |  |  |  | 37 |
| 32D9 | 01 | exchange | 3C | 331B | 22 | asn | 33 |
|  |  |  | 34 |  |  |  | 38 |
| 32DB | 02 | delete | A1 | 331D | 23 | acs | 43 |
|  |  |  | 33 |  |  |  | 38 |
| 32DD | 03 | subtract | 0F | 331F | 24 | atn | E2 |
|  |  |  | 30 |  |  |  | 37 |
| 32DF | 04 | multiply | CA | 3321 | 25 | In | 13 |
|  |  |  | 30 |  |  |  | 37 |
| 32E1 | 05 | division | AF | 3323 | 26 | exp | C4 |
|  |  |  | 31 |  |  |  | 36 |
| 32E3 | 06 | to-power | 51 | 3325 | 27 | int | AF |
|  |  |  | 38 |  |  |  | 36 |
| 32 E 5 | 07 | or | 1B | 3327 | 28 | sqr | 4A |
|  |  |  | 35 |  |  |  | 38 |
| 32 E 7 | 08 | no-\&-no | 24 | 3329 | 29 | sgn | 92 |
|  |  |  | 35 |  |  |  | 34 |
| 32E9 | 09 | no-l-eql | 3B | 332B | 2A | abs | 6A |
|  |  |  | 35 |  |  |  | 34 |


| 32EB | 0A | no-gr-eq | 3B | 332D | 2 B | peek | AC |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  | 35 |  |  |  | 34 |
| 32ED | OB | nos-neql | 3B | 332F | 2 C | in | A5 |
|  |  |  | 35 |  |  |  | 34 |
| 32EF | OC | no-grtr | 3B | 3331 | 2D | usr-no | B3 |
|  |  |  | 35 |  |  |  | 34 |
| 32F1 | OD | no-less | 3B | 3333 | 2E | str\$ | 1F |
|  |  |  | 35 |  |  |  | 36 |
| 32F3 | OE | nos-eql | 3B | 3335 | 2 F | chr\$ | C9 |
|  |  |  | 35 |  |  |  | 35 |
| 32F5 | OF | addition | 14 | 3337 | 30 | not | 01 |
|  |  |  | 30 |  |  |  | 35 |
| 32F7 | 10 | str-\&-no | 2D | 3339 | 31 | duplicate | C0 |
|  |  |  | 35 |  |  |  | 33 |
| 32F9 | 11 | str-l-eql | 3B | 333B | 32 | n-mod-m | A0 |
|  |  |  | 35 |  |  |  | 36 |
| 32FB | 12 | str-gr-eq | 3B | 333D | 33 | jump | 86 |
|  |  |  | 35 |  |  |  | 36 |
| 32FD | 13 | strs-neql | 3B | 333F | 34 | stk-data | C6 |
|  |  |  | 35 |  |  |  | 33 |
| 32FF | 14 | str-grtr | 3B | 3341 | 35 | dec-jr-nz | 7A |
|  |  |  | 35 |  |  |  | 36 |
| 3301 | 15 | str-less | 3B | 3343 | 36 | less-0 | 06 |
|  |  |  | 35 |  |  |  | 35 |
| 3303 | 16 | strs-eql | 3B | 3345 | 37 | greater-0 | F9 |
|  |  |  | 35 |  |  |  | 34 |
| 3305 | 17 | strs-add | 9C | 3347 | 38 | end-calc | 9B |
|  |  |  | 35 |  |  |  | 36 |
| 3307 | 18 | val\$ | DE | 3349 | 39 | get-argt | 83 |
|  |  |  | 35 |  |  |  | 37 |
| 3309 | 19 | usr-\$ | BC | 334B | 3A | truncate | 14 |
|  |  |  | 34 |  |  |  | 32 |
| 330B | 1A | read-in | 45 | 334D | 3B | fp-calc-2 | A2 |
|  |  |  | 36 |  |  |  | 33 |
| 330D | 1B | negate | 6E | 334F | 3C | e-to-fp | 4F |
|  |  |  | 34 |  |  |  | 2D |
| 330F | 1C | code | 69 | 3351 | 3D | re-stack | 97 |
|  |  |  | 36 |  |  |  | 32 |
| 3311 | 1D | val | DE | 3353 | 3E | series-06 | 49 |
|  |  |  | 45 |  |  | etc. | 34 |
| 3313 | 1E | len | 74 | 3355 | 3F | stk-zero | 1B |
|  |  |  | 36 |  |  | etc. | 34 |
| 3315 | 1F | $\sin$ | B5 | 3357 | 40 | st-mem-0 | 2D |
|  |  |  | 37 |  |  | etc. | 34 |
| 3317 | 20 | cos | AA | 3359 | 41 | get-mem-0 | 0F |
|  |  |  | 37 |  |  | etc. | 34 |

Note: The last four subroutines are multi-purpose subroutines and are entered with a parameter that is a copy of the right hand five bits of the original literal. The full set follows:
Offset 3E: series-06, series-08, \& series-0C; literals 86,88 \& 8C.
Offset 3F: stk-zero, stk-one, stk-half, stk-pi/2 \& stk-ten; literals A0 to A4.
Offset 40: st-mem-0, st-mem-1, st-mem-2, st-mem-3, st-mem-4 \& st-mem-5; literals C0 to C5.
Offset 41: get-mem-0, get-mem-1, get-mem-2, get-mem-3, get-mem-4 \& get-mem-5; literals E0 to E5.

## THE 'CALCULATE' SUBROUTINE

This subroutine is used to perform floating-point calculations. These can be considered to be of three types:
I. Binary operations, e.g. addition, where two numbers in floating-point form are added together to give one 'last value'.
II. Unary operations, e.g. sin, where the 'last value' is changed to give the appropriate function result as a new 'last value'.
III. Manipulatory operations, e.g. st-mem-0, where the 'last value' is copied to the first five bytes of the calculator's memory area.

The operations to be performed are specified as a series of data-bytes, the literals, that follow an RST 0028 instruction that calls this subroutine. The last literal in the list is always ' 38 ' which leads to an end to the whole operation.

In the case of a single operation needing to be performed, the operation offset can be passed to the CALCULATOR in the $B$ register, and operation '3B', the SINGLE CALCULATION operation, performed.

It is also possible to call this subroutine recursively, i.e. from within itself, and in such a case it is possible to use the system variable BREG as a counter that controls how many operations are performed before returning.

The first part of this subroutine is complicated but essentially it performs the two tasks of setting the registers to hold their required values, and to produce an offset, and possibly a parameter, from the literal that is currently being considered.

The offset is used to index into the calculator's table of addresses, see above, to find the required subroutine address.
The parameter is used when the multi-purpose subroutines are called.

Note: A floating-point number may in reality be a set of string parameters.

335B CALCULATE CALL \begin{tabular}{ll}

35BF,STK-PNTRS \& | Presume a unary operation and |
| :--- |
| therefore set HL to point to the |
| start of the 'last value' on the |
| calculator stack and DE one- |
| past this floating-point number |
| (STKEND). | <br>

Either, transfer a single <br>
operation offset to BREG <br>
temporarily, or, when using the <br>
subroutine recursively pass the <br>
parameter to BREG to be used <br>
as a counter.
\end{tabular}

|  |  | INC | HL | Make H'L' point to the next literal. |
| :---: | :---: | :---: | :---: | :---: |
| 336C | SCAN-ENT | PUSH | HL | This pointer is saved briefly on the machine stack. SCAN-ENT is used by the SINGLE CALCULATION subroutine to find the subroutine that is required. |
|  |  | AND | A | Test the A register. |
|  |  | JP | P,3380,FIRST-3D | Separate the simple literals from the multi-purpose literals. Jump with literals 00 -3D. |
|  |  | LD | D,A | Save the literal in D. |
|  |  | AND | +60 | Continue only with bits 5 \& 6 . Four right shifts make them now bits $1 \& 2$. |
|  |  | RRCA |  |  |
|  |  | RRCA |  |  |
|  |  | RRCA |  |  |
|  |  | RRCA |  |  |
|  |  | ADD | A, +7 C | The offsets required are $3 \mathrm{E}-41$. and $L$ will now hold double the required offset. |
|  |  | LD | L,A |  |
|  |  | LD | A, D | Now produce the parameter by |
|  |  | AND | +1F | taking bits $0,1,2,3 \& 4$ of the literal; keep the parameter in A. |
|  |  | JR | 338E,ENT-TABLE | Jump forward to find the address of the required subroutine. |
| 3380 | FIRST-3D | CP | +18 | Jump forward if performing a unary operation. <br> All of the subroutines that perform binary operations require that HL points to the first operand and DE points to the second operand (the 'last value') as they appear on the calculator stack. |
|  |  | JR | NC,338C,DOUBLE-A |  |
|  |  | EXX |  |  |
|  |  | LD | BC,+FFFB |  |
|  |  | LD | D,H |  |
|  |  | LD | E,L |  |
|  |  | ADD | HL, BC |  |
|  |  | EXX |  |  |
| 338C | DOUBLE-A | $\begin{aligned} & \text { RLCA } \\ & \text { LD } \end{aligned}$ |  |  |
|  |  |  | L,A | addresses takes up two bytes the offset produced is doubled. |
| 338E | ENT-TABLE | LD | $\begin{aligned} & \mathrm{DE},+32 \mathrm{D} 7 \\ & \mathrm{H},+00 \\ & \mathrm{HL}, \mathrm{DE} \\ & \mathrm{E},(\mathrm{HL}) \\ & \mathrm{HL} \\ & \mathrm{D},(\mathrm{HL}) \\ & \mathrm{HL},+3365 \\ & \text { (SP),HL } \\ & \text { DE } \end{aligned}$ | The base address of the table. |
|  |  | LD |  | The address of the required |
|  |  | ADD |  | table entry is formed in HL; and |
|  |  | LD |  | the required subroutine address |
|  |  | INC |  | is loaded into the DE register |
|  |  | LD |  | pair. |
|  |  | LD |  | The RE-ENTRY address of 3365 |
|  |  | EX |  | is put on the machine stack |
|  |  | PUSH |  | underneath the subroutine address. |
|  |  | EXX |  | Return to the main set of registers. |
|  |  | LD | BC,(STKEND-hi) | The current value of BREG is transferred to the B register thereby returning the single operation offset. <br> (See COMPARISON at 353B) |
| 33A1 | delete | RET |  | An indirect jump to the required subroutine. |

## THE 'DELETE' SUBROUTINE

(Offset 02: 'delete)
This subroutine contains only the single RET instruction at 33A1, above. The literal '02' results in this subroutine being considered as a binary operation that is to be entered with a first number addressed by the HL register pair and a second number addressed by the DE register pair, and the result produced again addressed by the HL register pair.
The single RET instruction thereby leads to the first number being considered as the resulting 'last value' and the second number considered as being deleted. Of course the number has not been deleted from the memory but remains inactive and will probably soon be overwritten.

## THE 'SINGLE OPERATION' SUBROUTINE

(Offset 3B: 'fp-calc-2')
This subroutine is only called from SCANNING at 2757 hex and is used to perform a single arithmetic operation. The offset that specifies which operation is to be performed is supplied to the calculator in the $B$ register and subsequently transferred to the system variable BREG

The effect of calling this subroutine is essentially to make a jump to the appropriate subroutine for the single operation.

| 33A2 | fp-calc-2 | POP <br> LD <br> EXX <br> JR | AF <br> A,(BREG) 336C,SCAN-ENT | Discard the RE-ENTRY address Transfer the offset to $A$. Enter the alternate register set. Jump back to find the required address; stack the RE-ENTRY address and jump to the subroutine for the operation. |
| :---: | :---: | :---: | :---: | :---: |

## THE 'TEST 5-SPACES' SUBROUTINE

This subroutine tests whether there is sufficient room in memory for another 5-byte floating-point number to be added to the calculator stack.

| $33 \mathrm{A9}$ | TEST-5-SP | PUSH | DE | Save DE briefly. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | PUSH | HL | Save HL briefly. |
|  |  | LD | BC,+0005 | Specify the test is for 5 bytes. |
|  |  | CALL | 1F05,TEST-ROOM | Make the test. |
|  |  | POP | HL | Restore HL. |
|  |  | POP | DE | Restore DE. |
|  |  | RET |  | Finished. |

## THE 'STACK NUMBER' SUBROUTINE

This subroutine is called by BEEP and SCANNING twice to copy STKEND to DE, move a floating-point number to the calculator stack, and reset STKEND from DE. It calls 'MOVE-FP' to do the actual move.

| 33B4 | STACK-NUM LD | DE,(STKEND) | Copy STKEND to DE as <br> destination address. |
| :--- | :--- | :--- | :--- |
|  | CALL | 33C0,MOVE-FP | Move the number. <br> Reset STKEND from DE. |
|  | LD | (STKEND),DE | Finished. |

## THE 'MOVE A FLOATING-POINT NUMBER' SUBROUTINE

(Offset 31: 'duplicate')
This subroutine moves a floating-point number to the top of the calculator stack ( 3 cases) or from the top of the stack to the calculator's memory area (1 case). It is also called through the calculator when it simply duplicates the number at the top of the calculator stack, the 'last value', thereby extending the stack by five bytes.

| $33 C 0$ | MOVE-FP | CALL |
| :--- | :--- | :--- |
|  | LDIR | 33A9,TEST-5-SP |$\quad$| A test is made for room. |
| :--- |
| RET |

## THE 'STACK LITERALS' SUBROUTINE <br> (Offset 34: 'stk-data')

This subroutine places on the calculator stack, as a 'last value', the floating-point number supplied to it as $2,3,4$ or 5 literals. When called by using offset ' 34 ' the literals follow the ' 34 ' in the list of literals; when called by the SERIES GENERATOR, see below, the literals are supplied by the sub-routine that called for a series to be generated; and when called by SKIP CONSTANTS \& STACK A CONSTANT the literals are obtained from the calculator's table of constants (32C5-32D6)
In each case, the first literal supplied is divided by Hex.40, and the integer quotient plus 1 determines whether 1, 2, 3 or 4 further literals will be taken from the source to form the mantissa of the number. Any unfilled bytes of the five bytes that go to form a 5-byte floatingpoint number are set to zero. The first literal is also used to determine the exponent, after reducing mod Hex.40, unless the remainder is zero, in which case the second literal is used, as it stands, without reducing mod Hex.40. In either case, Hex. 50 is added to the literal, giving the augmented exponent byte, e (the true exponent e' plus Hex.80). The rest of the 5 bytes are stacked, including any zeros needed, and the subroutine returns.

| 33C6 | STK-DATA | $\begin{aligned} & \text { LD } \\ & \text { LD } \end{aligned}$ | $\begin{aligned} & \mathrm{H}, \mathrm{D} \\ & \mathrm{~L}, \mathrm{E} \end{aligned}$ | This subroutine performs the manipulatory operation of adding a 'last value' to the calculator stack; hence HL is set to point one-past the present 'last value' and hence point to the result. |
| :---: | :---: | :---: | :---: | :---: |
| 33 C 8 | STK-CONST | CALL | 33A9,TEST-5-SP | Now test that there is indeed room. |
|  |  | $\begin{aligned} & \text { EXX } \\ & \text { PUSH } \\ & \text { EXX } \end{aligned}$ | HL | Go to the alternate register set and stack the pointer to the next literal. |
|  |  | EX | (SP), HL | Switch over the result pointer and the next literal pointer. |
|  |  | PUSH | BC | Save BC briefly. |
|  |  | LD | A,(HL) | The first literal is put into $A$ |
|  |  | AND | +C0 | and divided by Hex. 40 to give |
|  |  | RLCA <br> RLCA |  | the integer values 0,1 , 2 or 3 . |
|  |  | LD | C, A | The integer value is transferred |
|  |  | INC | C | to $C$ and incremented, thereby giving the range $1,2,3$ or 4 for the number of literals that will be needed. |
|  |  | LD | A,(HL) | The literal is fetch anew, |
|  |  | AND | +3F | reduced mod Hex. 40 and dis- |
|  |  | JR | NZ,33DE,FORM-EXP | carded as inappropriate if the |
|  |  | INC | HL | remainder if zero; in which case |


|  |  | LD | A,(HL) | the next literal is fetched and used unreduced. |
| :---: | :---: | :---: | :---: | :---: |
| 33DE | FORM-EXP | ADD <br> LD | $\begin{aligned} & \mathrm{A},+50 \\ & \text { (DE), } \mathrm{A} \end{aligned}$ | The exponent, e, is formed by the addition of Hex. 50 and passed to the calculator stack as the first of the five bytes of the result. |
|  |  | LD | A, +05 | The number of literals specified |
|  |  | SUB | C | in C are taken from the source |
|  |  | INC | HL | and entered into the bytes of |
|  |  | INC | DE | the result. |
|  |  | LD | B, +00 |  |
|  |  | LDIR |  |  |
|  |  | POP | BC | Restore BC. |
|  |  | EX | (SP), HL | Return the result pointer to HL |
|  |  | EXX |  | and the next literal pointer to |
|  |  | POP | HL | its usual position in $\mathrm{H}^{\prime}$ \& L'. |
|  |  | EXX |  |  |
|  |  | LD | B,A | The number of zero bytes |
|  |  | XOR | A | required at this stage is given by |
| 33F1 | STK-ZEROS | DEC | B | $5-\mathrm{C}-1$; and this number of zeros |
|  |  | RET | Z | is added to the result to make |
|  |  | LD | (DE), A | up the required five bytes. |
|  |  | INC | DE |  |
|  |  | JR | 33F1, STK-ZEROS |  |

## THE 'SKIP CONSTANTS' SUBROUTINE

This subroutine is entered with the HL register pair holding the base address of the calculator's table of constants and the A register holding a parameter that shows which of the five constants is being requested.
The subroutine performs the null operations of loading the five bytes of each unwanted constant into the locations 0000, 0001, 0002, 0003 and 0004 at the beginning of the ROM until the requested constant is reached.
The subroutine returns with the HL register pair holding the base address of the requested constant within the table of constants.

| 33F7 | SKIP-CONS | AND | A | The subroutine returns if the |
| :---: | :---: | :---: | :---: | :---: |
| 33F8 | SKIP-NEXT | RET | Z | parameter is zero, or when the |
|  |  |  |  | requested constant has been reached. |
|  |  | PUSH | AF | Save the parameter. |
|  |  | PUSH | DE | Save the result pointer. |
|  |  | LD | DE,+0000 | The dummy address. |
|  |  | CALL | 33C8,STK-CONST | Perform imaginary stacking of an expanded constant. |
|  |  | POP | DE | Restore the result pointer. |
|  |  | POP | AF | Restore the parameter. |
|  |  | DEC | A | Count the loops. |
|  |  | JR | 33F8,SKIP-NEXT | Jump back to consider the value of the counter. |

## THE 'MEMORY LOCATION' SUBROUTINE

This subroutine finds the base address for each five byte portion of the calculator's memory area to or from which a floating-point number is to be moved from or to the calculator stack. It does this operation by adding five times the parameter supplied to the base address for the area which is held in the HL register pair.
Note that when a FOR-NEXT variable is being handled then the pointers are changed so that the variable is treated as if it were the calculator's memory area (see address 1D20).
3406
LOC-MEM

| LD | C,A |
| :--- | :--- |
| RLCA |  |
| RLCA |  |
| ADD | A,C |

Copy the parameter to C .
Double the parameter.
Double the result.
Add the value of the parameter
to give five times the original

|  |  | value. |
| :--- | :--- | :--- |
| LD | C,A | This result is wanted in the |
| LD | B,+00 | BC register pair. |
| ADD | HL,BC | Produce the new base address. |
| RET |  | Finished. |

## THE 'GET FROM MEMORY AREA' SUBROUTINE

(Offsets E0 to E5: 'get-mem-0' to 'get-mem-5')
This subroutine is called using the literals E0 to E5 and the parameter derived from these literals is held in the A register. The subroutine calls MEMORY LOCATION to put the required source address into the HL register pair and MOVE A FLOATING-POINT NUMBER to copy the five bytes involved from the calculator's memory area to the top of the calculator stack to form a new 'last value'.

| 340F get-mem-0 | PUSH | DE | Save the result pointer. <br> etc. |
| :--- | :--- | :--- | :--- |
|  | LD | HL,(MEM) | Fetch the pointer to the current <br> memory area (see above). |
|  | CALL | 3406,LOC-MEM | The base address is found. |
|  | CALL | 33C0,MOVE-FP | The five bytes are moved. |

## THE 'STACK A CONSTANT' SUBROUTINE

(offsets A0 to A4: 'stk-zero','stk-one','stk-half','stk-pi/2' \& 'stk-ten')
This subroutine uses SKIP CONSTANTS to find the base address of the requested constants from the calculator's table of constants and then calls STACK LITERALS, entering at STK-CONST, to make the expanded form of the constant the 'last value' on the calculator stack.

| stk-zero <br> etc. | LD | H,D | Set HL to hold the result <br> pointer. |
| :--- | :--- | :--- | :--- |
|  | LD | L,E | Go to the alternate register set <br> and save the next literal pointer. |
|  | PUSH | HL | The base address of the calcul- <br> ator's table of constants. |
|  | EXX | HL,+32C5 | Back to the main set of registers. |
|  | Cind the requested base address. |  |  |

## THE 'STORE IN MEMORY AREA' SUBROUTINE

(Offsets C0 to C5: 'st-mem-0' to 'st-mem-5')

This subroutine is called using the literals C0 to C5 and the parameter derived from these literals is held in the A register. This subroutine is very similar to the GET FROM MEMORY subroutine but the source and destination pointers are exchanged.


Note that the pointers HL and DE remain as they were, pointing to STKEND-5 and STKEND respectively, so that the 'last value' remains on the calculator stack. If required it can be removed by using 'delete'.

## THE 'EXCHANGE' SUBROUTINE

(Offset 01: 'exchange')
This binary operation 'exchanges' the first number with the second number, i.e. the topmost two numbers on the calculator stack are exchanged.

| $343 C$ | EXCHANGE | LD | B, +05 |
| :--- | :--- | :--- | :--- |
| $343 E$ | SWAP-BYTE | LD | A,(DE) | | There are five bytes involved. |
| :--- |
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## THE 'SERIES GENERATOR' SUBROUTINE

(Offsets 86,88 \& 8C: 'series-06','series-08' \& 'series-0C')
This important subroutine generates the series of Chebyshev polynomials which are used to approximate to SIN, ATN, LN and EXP and hence to derive the other arithmetic functions which depend on these (COS, TAN, ASN, ACS, ** and SQR).
The polynomials are generated, for $n=1,2, \ldots$, by the recurrence relation:
$T_{n+1}(z)=2 z T_{n}(z)-T_{n-1}(z)$, where $T_{n}(z)$ is the $n$th Chebyshev polynomial in $z$.
The series in fact generates:
$T_{0}, 2 T_{1}, 2 T 2, \ldots .2 T_{n-1}$, where $n$ is 6 for SIN, 8 for EXP and 12 decimal, for LN and ATN.
The coefficients of the powers of $z$ in these polynomials may be found in the Handbook of Mathematical Functions by M. Abramowitz and I.A. Stegun (Dover 1965), page 795.
BASIC programs showing the generation of each of the four functions are given here in the Appendix.
In simple terms this subroutine is called with the 'last value' on the calculator stack, say $Z$, being a number that bears a simple relationship to the argument, say X, when the task is to evaluate, for instance, SIN X. The calling subroutine also supplies the list of constants that are to be required (six constants for SIN). The SERIES GENERATOR then manipulates its data and returns to the calling routine a 'last value' that bears a simple relationship to the requested function, for instance, SIN X.
This subroutine can be considered to have four major parts:
i. The setting of the loop counter:

The calling subroutine passes its parameters in the A register for use as a counter. The calculator is entered at GEN-ENT-1 so that the counter can be set.

| 3449 | series-06 | LD | B,A |
| :--- | :--- | :--- | :--- |$\quad$| Move the parameter to B. |
| :--- |
| etc. |$\quad$ CALL $\quad$ 335E,GEN-ENT-1 $\quad$| In effect a RST 0028 |
| :--- |
| instruction but sets the counter. |

ii. The handling of the 'last value', Z :

The loop of the generator requires $2^{*} Z$ to be placed in mem-0, zero to be placed in mem-2 and the 'last value' to be zero.

|  |  | calculator stack |  |
| :--- | :--- | :--- | :--- |
| DEFB | +31, duplicate | Z,Z |  |
| DEFB | $+0 F$, addition | $2^{\star} Z$ |  |
| DEFB | $+C 0$, st-mem-0 | $2^{\star} Z$ | mem-0 holds $2^{*} Z$ |
| DEFB | +02, delete | - |  |


| DEFB | + A0,stk-zero | 0 |  |
| :--- | :--- | :--- | :--- |
| DEFB | + C2,st-mem-2 | 0 | mem-2 holds 0 |

iii. The main loop:

The series is generated by looping, using BREG as a counter; the constants in the calling subroutine are stacked in turn by calling STK-DATA; the calculator is re-entered at GEN-ENT-2 so as not to disturb the value of BREG; and the series is built up in the form: $B(R)=2^{*} Z^{*} B(R-1)-B(R-2)+A(R)$, for $R=1,2, \ldots, N$, where $A(1), A(2), \ldots, A(N)$ are the constants supplied by the calling subroutine (SIN, ATN, LN and EXP) and $B(0)=0=B(-1)$
The (R+1)th loop starts with $B(R)$ on the stack and with $2 * Z, B(R-2)$ and $B(R-1)$ in mem -0 , mem-1 and mem- 2 respectively.

| 3453 | G-LOOP | DEFB | +31,duplicate | $B(R), B(R)$ |
| :---: | :---: | :---: | :---: | :---: |
|  |  | DEFB | +E0,get-mem-0 | $B(R), B(R), 2 * Z$ |
|  |  | DEFB | +04,multiply | $B(R), 2^{*} B(R) * Z$ |
|  |  | DEFB | +E2,get-mem-2 | $B(R), 2^{*} B(R) * Z, B(R-1)$ |
|  |  | DEFB | +C1,st-mem-1 | mem-1 holds $B(\mathrm{R}-1)$ |
| DEFB | +38,end-calc | DEFB | +03,subtract | $B(R), 2 * B(R) * Z-B(R-1)$ |

The next constant is placed on the calculator stack.

$$
\text { CALL } 33 C 6, S T K-D A T A \quad B(R), 2 * B(R) * Z-B(R-1), A(R+1)
$$

The Calculator is re-entered without disturbing BREG.

| CALL | 3362,GEN-ENT-2 |  |
| :---: | :---: | :---: |
| DEFB | +0F,addition | $B(R), 2^{*} B(R) * Z-B(R-1)+A(R+1)$ |
| DEFB | +01,exchange | $2^{*} B(R) * Z-B(R-1)+A(R+1), B(R)$ |
| DEFB | +C2,st-mem-2 | mem-2 holds $B(R)$ |
| DEFB | +02, delete | $\begin{aligned} & 2^{*} B(R)^{\star} Z-B(R-1)+A(R!1)= \\ & B(R!1) \end{aligned}$ |
| DEFB | +35,dec-jr-nz | $B(R+1)$ |
| DEFB | +EE,to 3453,G-LOOP |  |

iv. The subtraction of $\mathrm{B}(\mathrm{N}-2)$ :

The loop above leaves $B(N)$ on the stack and the required result is given by $B(N)-B(N-2)$.

| DEFB | + E1, get-mem-1 | $B(N), B(N-2)$ |
| :--- | :--- | :--- |
| DEFB | +03, subtract | $B(N)-B(N-2)$ |
| DEFB | +38, end-calc |  |
| RET |  | Finished |

## THE 'ABSOLUTE MAGNITUDE' FUNCTION

(Offset 2A: 'abs')
This subroutine performs its unary operation by ensuring that the sign bit of a floating-point number is reset 'Small integers' have to be treated separately. Most of the work is shared with the 'unary minus' operation.

| 346A abs | LD | B,+FF | B is set to FF hex. |
| :--- | :--- | :--- | :--- |
|  | JR | 3474, NEG-TEST | The jump is made into 'unary |

## THE 'UNARY MINUS' OPERATION

(Offset 1B: 'negate')
This subroutine performs its unary operation by changing the sign of the 'last value' on the calculator stack.
Zero is simply returned unchanged. Full five byte floating-point numbers have their sign bit manipulated so that it ends up reset (for abs') or changed (for 'negate'). 'Small integers' have their sign byte set to zero (for 'abs') or changed (for 'negate')
346E
NEGATE CALL
34E9,TEST-ZERO
If the number is zero, the
subroutine returns leaving
0000000000 unchanged.
LD B,$+00 \quad B$ is set to +00 hex for 'negate'.
'ABS' enters here.

3474 NEG-TEST | LD | A,(HL) | If the first byte is zero, the |
| :--- | :--- | :--- |
|  | AND | A |
| JR | Z,3483,INT-CASE | jump is made to deal with a |
| INC | HL | Small integer'. |

The 'integer case' does a similar operation with the sign byte.

| 3483 | INT-CASE | PUSH | DE | Save STKEND in DE. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | PUSH | HL | Save pointer to the number in HL. |
|  |  | CALL | 2D7F,INT-FETCH | Fetch the sign in C , the number in DE. |
|  |  | POP | HL | Restore the pointer to the number in HL. |
|  |  | LD | A,B | Get +FF for 'abs', +00 for 'negate'. |
|  |  | OR | C | Now +FF for 'abs', no change for 'negate' |
|  |  | $\begin{aligned} & \text { CPL } \\ & \text { LD } \end{aligned}$ | C,A | Now +00 for 'abs', and a changed byte for 'negate': store it in C. |
|  |  | CALL | 2D8E,INT-STORE | Store result on the stack. |
|  |  | $\mathrm{POP}$ RET | DE | Return STKEND to DE. |

## THE 'SIGNUM' FUNCTION

(Offset 29: 'sgn')
This subroutine handles the function SGN $X$ and therefore returns a 'last value' of 1 if $X$ is positive, zero if $X$ is zero and -1 if $X$ is negative.
3492 sgn

| CALL 34E9, | TEST-ZERO |
| :--- | :--- |
| RET | C |
| PUSH | DE |
| LD | DE, +0001 |
| INC | HL |
| RL | $(\mathrm{HL})$ |
| DEC | HL |
| SBC | $\mathrm{A}, \mathrm{A}$ |
| LD | $\mathrm{C}, \mathrm{A}$ |
| CALL | 2D8E,INT-STORE |
| POP | DE |
|  |  |

[^1]
## THE 'IN' FUNCTION

(Offset 2C: 'in')
This subroutine handles the function $I N X$. It inputs at processor level from port $X$, loading $B C$ with $X$ and performing the instruction $I N$ A,(C).

| 34A5 in | CALL | 1E99,FIND-INT2 |
| :--- | :--- | :--- |
|  | The 'last value', X, is |  |
|  | AN (C) | compressed into BC. |
|  | TR | The signal is received. |

## THE 'PEEK' FUNCTION

Offset 2B: 'peek')
This subroutine handles the function PEEK X. The 'last value' is unstacked by calling FIND-INT2 and replaced by the value of the contents of the required location.

| 34AC peek CALL | 1E99,FIND-INT2 | Evaluate the 'last value', <br> rounded to the nearest integer; <br> test that it is in range and return |
| :--- | :--- | :--- | :--- |
| it in BC. |  |  |

## THE 'USR' FUNCTION

(Offset 2D: 'usr-no')
This subroutine ('USR number' as distinct from 'USR string') handles the function USR X , where X is a number. The value of X is obtained in BC, a return address is stacked and the machine code is executed from location X .

34B3
usr-no

| CALL | 1E99,FIND-INT |
| :--- | :--- |
|  |  |
|  |  |
| LD | $H L,+2 D 2 B$ |
| PUSH | $H L$ |
| PUSH | $B C$ |
| RET |  |

Evaluate the 'last value', rounded to the nearest integer; test that it is in range and return it in BC.
Make the return address be that of the subroutine STACK-BC Make an indirect jump to the required location.

Note: It is interesting that the IY register pair is re-initialised when the return to STACK-BC has been made, but the important H'L' that holds the next literal pointer is not restored should it have been disturbed. For a successful return to BASIC, H'L' must on exit from the machine code contain the address in SCANNING of the 'end-calc' instruction, 2758 hex (10072 decimal).

## THE 'USR-STRING' FUNCTION

(Offset 19: 'usr-\$')
This subroutine handles the function USR $X \$$, where $X \$$ is a string. The subroutine returns in BC the address of the bit pattern for the user-defined graphic corresponding to $X \$$. It reports error $A$ if $X \$$ is not a single letter between a and u or a user-defined graphic.

| usr-\$ | CALL | 2BF1,STK-FETCH |
| :--- | :--- | :--- |
|  | DEC | BC |
|  |  |  |
|  | LD | A,B |
|  | OR | C |
|  | LR | NZ,34E7,REPORT-A |
|  |  | A,(DE) |
|  | CALL | 2C8D,ALPHA |
|  | JR | C,34D3,USR-RANGE |
|  | SUB | +90 |
|  |  |  |
|  | JR | C,34E7,REPORT-A |
|  | CP | +15 |
|  | JR | NC,34E7,REPORT-A |
|  | INC | A |

Fetch the parameters of the string $\mathrm{X} \$$.
Decrease the length by 1 to test it. If the length was not 1 , then jump to give error report A.

Fetch the single code of the string.
Does it denote a letter?
If so, jump to gets its address. Reduce range for actual userdefined graphics to 0-20 decimal. Give report A if out of range. Test the range again. Give report A if out of range. Make range of user-defined graphics 1 to 21 decimal, as for a to $u$.

| 34D3 | USR-RANGE | A | Now make the range 0 to 20 <br> decimal in each case. |
| :--- | :--- | :--- | :--- |
|  | ADD | A,A | Multiply by 8 to get an offset <br> for the address. |
|  | ADD | A,A | Test the range of the offset. |

REPORT A - Invalid argument.
34E7 REPORT-A RST
DEFB
0008,ERROR-1
Call the error handling routine.

## THE 'TEST-ZERO' SUBROUTINE

This subroutine is called at least nine times to test whether a floating-point number is zero. This test requires that the first four bytes of the number should each be zero. The subroutine returns with the carry flag set if the number was in fact zero.

| 34E9 | TEST-ZERO | PUSH | HL | Save HL on the stack. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | PUSH | BC | Save BC on the stack. |
|  |  | LD | B,A | Save the value of $A$ in $B$. |
|  |  | LD | A,(HL) | Get the first byte. |
|  |  | INC | HL | Point to the second byte. |
|  |  | OR | (HL) | OR first byte with second. |
|  |  | INC | HL | Point to the third byte. |
|  |  | OR | (HL) | OR the result with the third byte. |
|  |  | INC | HL | Point to the fourth byte. |
|  |  | OR | (HL) | OR the result with the fourth byte. |
|  |  | LD | A, B | Restore the original value of A . |
|  |  | POP | BC | And of BC. |
|  |  | POP | HL | Restore the pointer to the number to HL. |
|  |  | RET | NZ | Return with carry reset if any of the four bytes was non-zero. |
|  |  | SCF |  | Set the carry flag to indicate |
|  |  | RET |  | that the number was zero, and return. |

## THE 'GREATER THAN ZERO' OPERATION <br> (Offset 37: 'greater-0')

This subroutine returns a 'last value' of one if the present 'last value' is greater than zero and zero otherwise. It is also used by other subroutines to 'jump on plus'.

34F9 GREATER-0 CALL
34E9,TEST-ZERO Is the 'last-value' zero?
,
GREATER-0 CALL
C
LD A,+FF
JR 3507,SIGN-TO-C

If so, return.
Jump forward to LESS THAN
ZERO but signal the opposite action is needed.

## THE 'NOT' FUNCTION

(Offset 30: 'not')
This subroutine returns a 'last value' of one if the present 'last value' is zero and zero otherwise. It is also used by other subroutines to 'jump on zero'.

The carry flag will be set only if the 'last value' is zero; this gives the correct result.
Jump forward.
THE 'LESS THAN ZERO' OPERATION
(Offset 36: 'less-0')
This subroutine returns a 'last value' of one if the present 'last value' is less than zero and zero otherwise. It is also used by other subroutines to 'jump on minus'.

| 3506 | less-0 | XOR | A |
| :--- | :--- | :--- | :--- |
| 3507 | SIGN-TO-C | INC | HL |
|  |  | XOR | (HL) |
|  |  | DEC | HL |
|  |  |  | RLCA |

Clear the A register.
Point to the sign byte.
The carry is reset for a positive number and set for a negative number; when entered from GREATER-0 the opposite sign goes to the carry.

## THE 'ZERO OR ONE' SUBROUTINE

This subroutine sets the 'last value' to zero if the carry flag is reset and to one if it is set. When called from 'E-TO-FP' however it creates the zero or one not on the stack but in mem-0.

| 350B | FP-0/1 | PUSH | HL | Save the result pointer. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | LD | A, +00 | Clear A without disturbing the carry. |
|  |  | LD | (HL),A | Set the first byte to zero. |
|  |  | INC | HL | Point to the second byte. |
|  |  | LD | (HL), A | Set the second byte to zero. |
|  |  | INC | HL | Point to the third byte. |
|  |  | RLA |  | Rotate the carry into A, making |
|  |  |  |  | A one if the carry was set, but zero if the carry was reset. |
|  |  | LD | (HL), A | Set the third byte to one or zero. |
|  |  | RRA |  | Ensure that A is zero again. |
|  |  | INC | HL | Point to the fourth byte. |
|  |  | LD | (HL), A | Set the fourth byte to zero. |
|  |  | INC | HL | Point to the fifth byte. |
|  |  | LD | (HL), A |  |
|  |  | POP | HL | Restore the result pointer. |
|  |  | RET |  |  |

## THE 'OR' OPERATION

(Offset 07: 'or')
This subroutine performs the binary operation ' X OR $\mathrm{Y}^{\prime}$ and returns X if Y is zero and the value 1 otherwise.

| 351B | or | EX | DE,HL | Point HL at Y , the second number. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | CALL | 34E9,TEST-ZERO | Test whether Y is zero. |
|  |  | EX | DE,HL | Restore the pointers. |
|  |  | RET | C | Return if Y was zero; X is now the 'last value'. |
|  |  | $\begin{aligned} & \text { SCF } \\ & \text { JR } \end{aligned}$ | 350B,FP-0/1 | Set the carry flag and jump back to set the 'last value' to 1 . |

## THE 'NUMBER AND NUMBER' OPERATION

(Offset 08: 'no-\&-no')
This subroutine performs the binary operation ' X AND $Y$ ' and returns X if Y is non-zero and the value zero otherwise

Point HL at Y , DE at X .
CALL 34E9,TEST-ZERO
EX DE,HL
RET NC

AND A
JR 350B,FP-0/1
Test whether $Y$ is zero.
Swap the pointers back.
Return with X as the 'last value'
if $Y$ was non-zero.
Reset the carry flag and jump
back to set the 'last value' to
zero.

## THE 'STRING AND NUMBER' OPERATION

(Offset 10: 'str-\&-no')
This subroutine performs the binary operation ' $X \$$ AND $Y$ ' and returns $X \$$ if $Y$ is non-zero and a null string otherwise

| 352D | str-\&-no | EX | DE,HL | Point HL at Y, DE at X\$ |
| :---: | :---: | :---: | :---: | :---: |
|  |  | CALL | 34E9,TEST-ZERO | Test whether $Y$ is zero. |
|  |  | EX | DE,HL | Swap the pointers back. |
|  |  | RET | NC | Return with $\mathrm{X} \$$ as the 'last value' if $Y$ was non-zero. |
|  |  | PUSH | DE | Save the pointer to the number. |
|  |  | DEC | DE | Point to the fifth byte of the string parameters i.e. lengthhigh. |
|  |  | XOR | A | Clear the A register. |
|  |  | LD | (DE), A | Length-high is now set to zero. |
|  |  | DEC | DE | Point to length-low. |
|  |  | LD | (DE), A | Length-low is now set to zero. |
|  |  | POP | DE | Restore the pointer. |
|  |  | RET |  | Return with the string |
|  |  |  |  | parameters being the 'last value |

## THE 'COMPARISON' OPERATIONS

(Offsets 09 to 0E \& 11 to 16: 'no-l-eql', 'no-gr-eq', 'nos-neql', 'no-grtr', 'no-less', 'nos-eql', 'str-l-eql', 'str-gr-eq', 'strs-neql', 'str-grtr', 'strless' \& 'strs-eql')

This subroutine is used to perform the twelve possible comparison operations. The single operation offset is present in the B register at the start of the subroutine.

| 353B | no-l-eql etc. | LD | A, B | The single offset goes to the A register. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | SUB | +08 | The range is now 01-06 \& 09-0E. |
|  |  | BIT | 2,A | This range is changed to: |
|  |  | JR | NZ,3543,EX-OR-NOT |  |
|  |  | DEC | A | 0C-0E. |
| 3543 | EX-OR-NOT | RRCA |  | Then reduced to 00-07 with carry set for 'greater than or equal to' \& 'less than'; the operations with carry set are |
|  |  | JR | NC,354E,NU-OR-STR | then treated as their |
|  |  | PUSH | AF | complementary operation once |
|  |  | PUSH | HL | their values have been exchanged |
|  |  | CALL | 343C,EXCHANGE |  |
|  |  | POP | DE |  |
|  |  | EX | DE,HL |  |
|  |  | POP | AF |  |
| 354E | NU-OR-STR | BIT | 2,A | The numerical comparisons are |
|  |  | JR | NZ,3559,STRINGS | now separated from the string comparisons by testing bit 2. |
|  |  | RRCA |  | The numerical operations now have the range 00-01 with carry set for 'equal' and 'not equal'. |
|  |  | PUSH | AF | Save the offset. |
|  |  | CALL | 300F,SUBTRACT | The numbers are subtracted for |


| 3559 | STRINGS | $\begin{aligned} & \text { JR } \\ & \text { RRCA } \end{aligned}$ | 358C,END-TESTS | the final tests. <br> The string comparisons now have the range 02-03 with carry set for 'equal' and 'not equal'. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | PUSH | AF | Save the offset. |
|  |  | CALL | 2BF1,STK-FETCH | The lengths and starting |
|  |  | PUSH | DE | addresses of the strings are |
|  |  | PUSH | BC | fetched from the calculator |
|  |  | CALL | 2BF1,STK-FETCH | stack. |
|  |  | POP | HL | The length of the second string. |
| 3564 | BYTE-COMP | LD | A, H |  |
|  |  | OR | L |  |
|  |  | EX | (SP), HL |  |
|  |  | LD | A, B |  |
|  |  | JR | NZ,3575,SEC-PLUS | Jump unless the second string |
|  |  | OR | C | is null. |
| 356B | SECND-LOW | POP | BC | Here the second string is either null or less than the first. |
|  |  | JR | Z,3572,BOTH-NULL |  |
|  |  | POP | AF |  |
|  |  | CCF |  | The carry is complemented to |
|  |  | JR | 3588,STR-TEST | give the correct test results. |
| 3572 | BOTH-NULL | POP | AF | Here the carry is used as it |
|  |  | JR | 3588,STR-TEST | stands. |
| 3575 | SEC-PLUS | OR | C |  |
|  |  | JR | Z,3585,FRST-LESS | The first string is now null, the second not. |
|  |  | LD | A,(DE) | Neither string is null, so their |
|  |  | SUB | (HL) | next bytes are compared. |
|  |  | JR | C,3585,FRST-LESS | The first byte is less. |
|  |  | JR | NZ,356B,SECND-LOW | The second byte is less. |
|  |  | DEC | BC | The bytes are equal; so the |
|  |  | INC | DE | lengths are decremented and a |
|  |  | INC | HL | jump is made to BYTE-COMP |
|  |  | EX | (SP), HL | to compare the next bytes of |
|  |  | DEC | HL | the reduced strings. |
|  |  | JR | 3564,BYTE-COMP |  |
| 3585 | FRST-LESS | POP | AF |  |
|  |  | AND | A | The carry is cleared here for the correct test results. |
| 3588 | STR-TEST | PUSH |  | For the string tests, a zero is |
|  |  | RST | 0028,FP-CALC | put on to the calculator stack. |
|  |  | DEFB | +A0,stk-zero |  |
|  |  | DEFB | +38,end-calc |  |
| 358C | END-TESTS | POP | AF | These three tests, called as |
|  |  | PUSH | AF | needed, give the correct results |
|  |  | CALL | C,3501,NOT | for all twelve comparisons. The |
|  |  | POP | AF | initial carry is set for 'not equal' |
|  |  | PUSH | AF | and 'equal', and the final carry |
|  |  | CALL | NC,34F9,GREATER-0 | is set for 'greater than', 'less |
|  |  | POP | AF | than' and 'equal'. |
|  |  | RRCA |  |  |
|  |  | CALL | NC,3501,NOT |  |
|  |  | RET |  | Finished. |

## THE 'STRING CONCATENATION' OPERATION

(Offset 17: 'strs-add')
This subroutine performs the binary operation ' $\mathrm{A} \$+\mathrm{B} \$$. The parameters for these strings are fetched and the total length found. Sufficient room to hold both the strings is made available in the work space and the strings are copied over. The result of this subroutine is therefore to produce a temporary variable $A \$+B \$$ that resides in the work space.

| 359C | strs-add | CALL | 2BF1,STK-FETCH | The parameters of the second |
| :---: | :---: | :---: | :---: | :---: |
|  |  | PUSH | DE | string are fetched and saved. |
|  |  | PUSH | BC |  |
|  |  | CALL | 2BF1,STK-FETCH | The parameters of the first string are fetched. |
|  |  | POP | HL |  |
|  |  | PUSH | HL | The lengths are now in HL and $B C$. |
|  |  | PUSH | DE | The parameters of the first |
|  |  | PUSH | BC | string are saved. |
|  |  | ADD | HL, BC | The total length of the two |
|  |  | LD | B, H | strings is calculated and passed |
|  |  | LD | C,L | to BC. |
|  |  | RST | 0030,BC-SPACES | Sufficient room is made available. |
|  |  | CALL | 2AB2,STK-STORE | The parameters of the new string are passed to the calculator stack. |
|  |  | POP | BC | The parameters of the first |
|  |  | POP | HL | string are retrieved and the |
|  |  | LD | A, B | string copied to the work space |
|  |  | OR | C | as long as it is not a null string. |
|  |  | JR <br> LDIR | Z,35B7,OTHER-STR |  |
| 35B7 | OTHER-STR | POP | BC | Exactly the same procedure is |
|  |  | POP | HL | followed for the second string |
|  |  | LD | A, B | thereby giving 'A\$+B\$'. |
|  |  | OR | C |  |
|  |  | JR | Z,35BF,STK-PNTRS |  |
|  |  | LDIR |  |  |

## THE 'STK-PNTRS' SUBROUTINE

This subroutine resets the HL register pair to point to the first byte of the 'last value', i.e. STKEND-5, and the DE register pair to point one-past the 'last value', i.e. STKEND.

35BF
STK-PNTRS LD

| LD | DE,+FFFB |
| :--- | :--- |
| PUSH | HL |
| ADD | HL,DE |
| POP | DE |
| RET |  |

Fetch the current value of STKEND. Set DE to -5, twos complement.
Stack the value for STKEND.
Calculate STKEND-5.
DE now holds STKEND and HL

## THE 'CHR\$' FUNCTION

(Offset 2F: 'chrs')
This subroutine handles the function CHR\$ X and creates a single character string in the work space.

| 35C9 | chrs | CALL 2DD5,FP-TO-A |  | The 'last value' is compressed |
| :---: | :---: | :---: | :---: | :---: |
|  |  |  |  | into the A register. |
|  |  | JR | C,35DC,REPORT-B | Give the error report if $X$ was greater than 255 decimal, or |
|  |  | JR | NZ,35DC,REPORT-B | X was a negative number. |
|  |  | PUSH | AF | Save the compressed value of $X$. |
|  |  | LD | BC, +0001 | Make one space available in the |
|  |  | POP | AF | Fetch the value. |
|  |  | LD | (DE), A | Copy the value to the work space. |
|  |  | CALL | 2AB2,STK-STORE | Pass the parameters of the new string to the calculator stack. |
|  |  | EX | DE,HL | Reset the pointers. |
|  |  | RET |  | Finished. |

REPORT-B - Integer out of range
35DC REPORT-B RST

0008,ERROR-
+0A

Call the error handling routine.

## THE 'VAL' AND 'VAL\$' FUNCTION

(Offsets 1D: 'val' and 18: 'val\$')
This subroutine handles the functions VAL $X \$$ and VAL\$ $X \$$. When handling VAL $X \$$, it return a 'last value' that is the result of evaluating the string (without its bounding quotes) as a numerical expression. when handling VAL $\$ \times \$$, it evaluates $\mathrm{X} \$$ (without its bounding quotes) as a string expression, and returns the parameters of that string expression as a 'last value' on the calculator stack.

| 35DE | val (also val\$) | $\begin{aligned} & \text { LD } \\ & \text { PUSH } \\ & \text { LD } \end{aligned}$ | $\begin{aligned} & \mathrm{HL},(\mathrm{CH}-\mathrm{ADD}) \\ & \mathrm{HL} \\ & \mathrm{~A}, \mathrm{~B} \end{aligned}$ | The current value of CH-ADD is preserved on the machine stack. The 'offset' for 'val' or 'val\$' must be in the $B$ register; it is now copied to $A$. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | ADD | A, +E3 | Produce +00 and carry set for 'val', +FB and carry reset for 'val\$'. |
|  |  | SBC | A,A | Produce +FF (bit 6 therefore set) for 'val', but +00 (bit 6 reset) for 'val\$'. |
|  |  | PUSH | AF | Save this 'flag' on the machine stack. |
|  |  | CALL | 2BF1,STK-FETCH | The parameters of the string are |
|  |  | PUSH | DE | fetched; the starting address is |
|  |  | INC | BC | saved; one byte is added to the |
|  |  | RST | 0030,BC-SPACES | length and room made available for the string (+1) in the work space. |
|  |  | POP | HL | The starting address of the string goes to HL as a source address. |
|  |  | LD | (CH-ADD),DE | The pointer to the first new |
|  |  | PUSH | DE | space goes to CH-ADD and to the machine stack. |
|  |  | LDIR |  | The string is copied to the work space, together with an extra byte. |
|  |  | EX | DE,HL | Switch the pointers. |
|  |  | DEC | HL | The extra byte is replaced by a |
|  |  | LD | (HL), +0D | 'carriage return' character. |
|  |  | RES | 7,(FLAGS) | The syntax flag is reset and the |
|  |  | CALL | 24FB,SCANNING | string is scanned for correct syntax. |
|  |  | RST | 0018,GET-CHAR | The character after the string is fetched. |
|  |  | CP | +0D | A check is made that the end of the expression has been reached. |
|  |  | JR | NZ,360C,V-RPORT-C | If not, the error is reported. |
|  |  | POP | HL | The starting address of the string is fetched. |
|  |  | POP | AF | The 'flag' for 'val/val\$' is |
|  |  | XOR | (FLAGS) | fetched and bit 6 is compared |
|  |  | AND | +40 | with bit 6 of the result of the syntax scan. |
| 360 C | V-RPORT-C | JP | NZ,1C8A,REPORT-C | Report the error if they do not match. |
|  |  | LD | (CH-ADD), HL | Start address to CH-ADD again. |
|  |  | SET | 7,(FLAGS) | The flag is set for line execution. |


| CALL | 24FB,SCANNING | The string is treated as a 'next <br> expression' and a 'last value' <br> produced. |
| :--- | :--- | :--- |
| POP | HL | The original value of CH-ADD is <br> restored. |
| JR | (CH-ADD),HL | The subroutine exits via STK- |
|  | 35BF,STK-PNTRS | PNTRS which resets the pointers. |

## THE 'STR\$' FUNCTION

(Offset 2E: 'str\$')
This subroutine handles the function STR\$ $X$ and returns a 'last value' which is a set of parameters that define a string containing what would appear on the screen if $X$ were displayed by a PRINT command.

| 361F | $\operatorname{str}$ \$ | LD | $\mathrm{BC},+0001$ | One space is made in the work |
| :---: | :---: | :---: | :---: | :---: |
|  |  | RST | 0030,BC-SPACES | space and its address is copied |
|  |  | LD | (K-CUR), HL | to K-CUR, the address of the cursor. |
|  |  | PUSH | HL | This address is saved on the stack too. |
|  |  | LD | HL,(CURCHL) | The current channel address is |
|  |  | PUSH | HL | saved on the machine stack. |
|  |  | LD | A,+FF | Channel 'R' is opened, allowing |
|  |  | CALL | 1601,CHAN-OPEN | the string to be 'printed' out into the work space. |
|  |  | CALL | 2DE3,PRINT-FP | The 'last value', X , is now printed out in the work space and the work space is expanded with each character. |
|  |  | POP | HL | Restore CURCHL to HL and |
|  |  | CALL | 1615,CHAN-FLAG | restore the flags that are appropriate to it. |
|  |  | POP | DE | Restore the start address of the string. |
|  |  | LD | HL,(K-CUR) | Now the cursor address is one |
|  |  | AND | A | past the end of the string and |
|  |  | SBC | HL, DE | hence the difference is the length. |
|  |  | LD | B, H | Transfer the length to BC. |
|  |  | LD | C,L |  |
|  |  | CALL | 2AB2,STK-STO-\$ | Pass the parameters of the new string to the calculator stack. |
|  |  | $\begin{aligned} & \text { EX } \\ & \text { RET } \end{aligned}$ | DE,HL | Reset the pointers. Finished. |

Note: See PRINT-FP for an explanation of the 'PRINT "A"+STR\$ 0.1' error

## THE 'READ-IN' SUBROUTINE

(Offset 1A: 'read-in')
This subroutine is called via the calculator offset through the first line of the S-INKEY\$ routine in SCANNING. It appears to provide for the reading in of data through different streams from those available on the standard Spectrum. Like INKEY\$ the subroutine returns a string.

3645 read-in |  | CALL | 1E94,FIND-INT1 |
| :--- | :--- | :--- |
|  |  |  |
|  | CP | +10 |
|  | JP | NC,1E9F,REPORT-B |
|  | LD | $\mathrm{HL},(\mathrm{CURCHL})$ |
|  | PUSH | HL |
|  | CALL | 1601, CHAN-OPEN |

The numerical parameter is compressed into the A register Is it smaller than 16 decimal? If not, report the error. The current channel address is saved on the machine stack. The channel specified by the parameter is opened.


## THE 'CODE' FUNCTION

(Offset 1C: 'code')
This subroutine handles the function CODE $A \$$ and returns the Spectrum code of the first character in $A \$$, or zero if $A \$$ should be null.

| 3669 | code | CALL | 2BF1,STK-FETCH | The parameters of the string are fetched. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | LD | A, B | The length is tested and the A |
|  |  | OR | C | register holding zero is carried |
|  |  | JR | Z,3671,STK-CODE | forward is $\mathrm{A} \$$ is a null string. |
|  |  | LD | A,(DE) | The code of the first character is put into A otherwise. |
| 3671 | STK-CODE | JP | 2D28,STACK-A | The subroutine exits via |
|  |  |  |  | STACK-A which gives the |

## THE 'LEN' FUNCTION

(Offset 1E: 'len')
This subroutine handles the function LEN A\$ and returns a 'last value' that is equal to the length of the string

| 3674 len | 2BF1,STK-FETCH | The parameters of the string are <br> fetched. |
| :--- | :--- | :--- |
| JP | 2D2B,STACK-BC | The subroutine exits via <br> STACK-BC which gives the <br> correct 'last value'. |

## THE 'DECREASE THE COUNTER' SUBROUTINE <br> (Offset 35: 'dec-jr-nz')

This subroutine is only called by the SERIES GENERATOR subroutine and in effect is a 'DJNZ' operation but the counter is the system variable, BREG, rather than the $B$ register.

| 367A dec-jr-nz | EXX <br> PUSH | HL |
| :--- | :--- | :--- |
|  |  | Go to the alternative register set <br> and save the next literal pointer |
| on the machine stack. |  |  |

THE 'JUMP' SUBROUTINE
(Offset 33: 'jump')
This subroutine executes an unconditional jump when called by the literal '33'. It is also used by the subroutines DECREASE THE COUNTER and JUMP ON TRUE

| 3686 | JUMP | EXX |  |
| :--- | :--- | :--- | :--- |
| 3687 | JUMP-2 | LD | E,(HL) |
|  |  |  |  |
|  |  | LD | A,E |
|  |  | RLA |  |
|  |  | SBC | A,A |
|  |  | LD | D,A |
|  |  | ADD | HL,DE |
|  |  | EXX |  |
|  |  | RET |  |

Go to the next alternate register set.
The next literal (jump length) is put in the E' register.
The number 00 hex or FF hex is formed in A according as $\mathrm{E}^{\prime}$ is positive or negative, and is then copied to $\mathrm{D}^{\prime}$. The registers H' \& L' now hold the next literal pointer. Finished.

## THE 'JUMP ON TRUE' SUBROUTINE

(Offset 00: 'jump-true')
This subroutine executes a conditional jump if the 'last value' on the calculator stack, or more precisely the number addressed currently by the DE register pair, is true.

| 368F | jump-true | INC | DE | Point to the third byte, which is |
| :---: | :---: | :---: | :---: | :---: |
|  |  | INC | DE | zero or one. |
|  |  | LD | A, (DE) | Collect this byte in the A register. |
|  |  | DEC | DE | Point to the first byte once |
|  |  | DEC | DE | again. |
|  |  | AND | A | Test the third byte: is it zero? |
|  |  | JR | NZ,3686,JUMP | Make the jump if the byte is non-zero, i.e. if the number is not-false. |
|  |  | EXX |  | Go to the alternate register set. |
|  |  | INC | HL | Pass over the jump length. |
|  |  | EXX |  | Back to the main set of registers. |
|  |  | RET |  | Finished. |

## THE 'END-CALC' SUBROUTINE

(Offset 38: 'end-calc')
This subroutine ends a RST 0028 operation.

| 369B end-calc | POP | The return address to the <br> calculator ('RE-ENTRY') is <br> discarded. |
| :--- | :--- | :--- |
| EXX |  |  |
| EX | (SP).HL | Instead, the address in H'L' is <br> put on the machine stack and <br> an indirect jump is made to it. |
| H'L' will now hold any earlier |  |  |
| address in the calculator chain |  |  |
| of addresses. |  |  |
| Finished. |  |  |

## THE 'MODULUS' SUBROUTINE

(Offset 32: 'n-mod-m')
This subroutine calculates M (mod M ), where M is a positive integer held at the top of the calculator stack, the 'last value', and N is the integer held on the stack beneath M .
The subroutine returns the integer quotient INT (N/M) at the top of the calculator stack, the 'last value', and the remainder N-INT (N/M) in the second place on the stack.
This subroutine is called during the calculation of a random number to reduce N mod 65537 decimal.


| DEFB | +E0,get-mem-0 | N, INT (N/M),M |
| :---: | :---: | :---: |
| DEFB | +01, exchange | N, M, INT (N/M) |
| DEFB | +C0,st-mem-0 | N, M, INT (N/M) mem-0 holds INT (N/M) |
| DEFB | +04,multiply | N, M*INT (N/M) |
| DEFB | +03,subtract | $\mathrm{n}-\mathrm{M}^{*} \mathrm{INT}(\mathrm{N} / \mathrm{M})$ |
| DEFB | +E0,get-mem-0 | $\mathrm{n}-\mathrm{M} *$ INT ( $\mathrm{N} / \mathrm{M}$ ), INT ( $\mathrm{N} / \mathrm{M}$ ) |
| DEFB | +38,end-calc |  |
| RET |  | Finished. |

## THE 'INT' FUNCTION

(Offset 27: 'int')
This subroutine handles the function INT X and returns a 'last value' that is the 'integer part' of the value supplied. Thus INT 2.4 gives 2 but as the subroutine always rounds the result down INT -2.4 gives -3 .
The subroutine uses the INTEGER TRUNCATION TOWARDS ZERO subroutine at 3214 to produce $I(X)$ such that $I(2.4)$ gives 2 and $I$ $(-2.4)$ gives -2. Thus, INT $X$ is gives by $\mathrm{I}(\mathrm{X})$ for values of X that are greater than or equal to zero, and $\mathrm{I}(\mathrm{X})-1$ for negative values of X that are not already integers, when the result is, of course, $I(X)$.

| 36AF int | RST | 0028,FP-CALC | X |
| :--- | :--- | :--- | :--- |
|  |  | DEFB | +31,duplicate |
|  | DEFB | +36,less-0 | $\times, X$ |
|  |  | DEFB | +00,jump-true |
|  | DEFB | +04, to $36 B 7$, X-NEG | X |
|  |  |  |  |

For values of $X$ that have been shown to be greater than or equal to zero there is no jump and $I(X)$ is readily found.

| DEFB | $+3 A$, truncate | $I(X)$ |
| :--- | :--- | :--- |
| DEFB | +38, end-calc | Finished. |
| RET |  |  |


| 36B7 X-NEG | DEFB | +31,duplicate | $\mathrm{X}, \mathrm{X}$ |
| :---: | :---: | :---: | :---: |
|  | DEFB | +3A,truncate | X, I (X) |
|  | DEFB | +C0,st-mem-0 | X, I (X) mem-0 holds I (X) |
|  | DEFB | +03,subtract | X-I (X) |
|  | DEFB | +E0,get-mem-0 | X-I (X), I (X) |
|  | DEFB | +01, exchange | $\mathrm{I}(\mathrm{X}), \mathrm{X}-\mathrm{I}(\mathrm{X})$ |
|  | DEFB | +30,not | $1(X),(1 / 0)$ |
|  | DEFB | +00,jump-true | 1 (X) |
|  | DEFB | +03,to 36C2,EXIT | 1 (X) |

The jump is made for values of $X$ that are negative integers, otherwise there is no jump and $I(X)-1$ is calculated.

| DEFB | + A1 1, stk-one | $I(X), 1$ |
| :--- | :--- | :--- |
| DEFB | +03, subtract | $I(X)-1$ |

In either case the subroutine finishes with;


## THE 'EXPONENTIAL' FUNCTION

(Offset 26: 'exp')
This subroutine handles the function EXP $X$ and is the first of four routines that use SERIES GENERATOR to produce Chebyshev polynomials.
The approximation to EXP X is found as follows:
i. X is divided by LN 2 to give Y , so that 2 to the power Y is now the required result.
ii. The value N is found, such that $\mathrm{N}=\mathrm{INT} \mathrm{Y}$.
iii. The value W is found, such that $\mathrm{W}=\mathrm{Y}-\mathrm{N}$, where $0<=\mathrm{W}<=1$, as required for the series to converge.
$i v$. The argument $Z$ if formed, such that $Z=2^{*} w-1$.
v. The SERIES GENERATOR is used to return $2^{* *}$ W.
vi. Finally N is added to the exponent, giving $2^{* *}(\mathrm{~N}+\mathrm{W})$, which is $2^{* *} \mathrm{Y}$ and therefore the required answer for EXP X .

The method is illustrated using a BASIC program in the Appendix.

| $36 C 4 \quad$ EXP | RST | 0028,FP-CALC | X |
| :--- | :--- | :--- | :--- |
| Perform step i. |  |  |  |
|  | DEFB | +3D,re-stack | X (in full floating-point form) |
|  | DEFB | +34,stk-data <br> +F1,exponent+81 | X, 1/LN 2 |
|  | DEFB | +38,+AA,+3B,+29 |  |
|  | DEFB | +04,multiply | X/LN 2 $=$ Y |

Perform step ii.

| DEFB | +31, duplicate | $Y, Y$ |
| :--- | :--- | :--- |
| DEFB | +27, int,1C46 | $Y$, INT Y $=N$ |
| DEFB | $+C 3$, st-mem-3 | $Y, N \quad$ mem- -3 holds $N$ |

Perform step iii.
DEFB $\quad+03$, subtract $\quad \mathrm{Y}-\mathrm{N}=\mathrm{W}$
Perform step iv.

| DEFB | +31, duplicate | $W, W$ |
| :--- | :--- | :--- |
| DEFB | $+0 F$, addition | $2^{*} W$ |
| DEFB | $+A 1$, stk-one | $2^{*} W, 1$ |
| DEFB | +03, subtract | $2^{*} W-1=Z$ |

Perform step v, passing to the SERIES GENERATOR the parameter ' 8 ' and the eight constants required.

|  | DEFB | +88,series-08 | Z |
| :---: | :---: | :---: | :---: |
| 1. | DEFB | +13,exponent+63 |  |
|  | DEFB | +36,(+00,+00,+00) |  |
| 2. | DEFB | +58, exponent+68 |  |
|  | DEFB | +65,+66,(+00,+00) |  |
| 3. | DEFB | +9D, exponent+6D |  |
|  | DEFB | +78,+65,+40,(+00) |  |
| 4. | DEFB | +A2,exponent+72 |  |
|  | DEFB | +60,+32,+C9,(+00) |  |
| 5. | DEFB | +E7, exponent+77 |  |
|  | DEFB | +21,+F7,+AF,+24 |  |
| 6. | DEFB | +EB, exponent+7B |  |
|  | DEFB | +2F,+B0,+B0,+14 |  |
| 7. | DEFB | +EE, exponent +7E |  |
|  | DEFB | +7E,+BB,+94,+58 |  |
| 8. | DEFB | +F1,exponent+81 |  |
|  | DEFB | +3A,+7E, +F8,+CF |  |

At the end of the last loop the 'last value' is $2^{* *} W$.
Perform step vi.

| DEFB | +E3,get-mem-3 | $2 *$ W, N |
| :--- | :--- | :--- |
| DEFB | +38,end-calc |  |
| CALL | 2DD5,FP-TO-A | The absolute value of N mod <br> 256 decimal, is put into the A <br> register. |
| JR | NZ,3705,N-NEGTV | Jump forward if N was negative. |
| JR | C,3703,REPORT-6 | Error if ABS N greater than 255 <br> dec. |
| ADD | A,(HL) | Now add ABS N to the <br> exponent. <br> Jump unless e greater than 255 <br> dec. |

Report 6 - Number too big

| 3703 | REPORT-6 | $\begin{aligned} & \text { RST } \\ & \text { DEFB } \end{aligned}$ | $\begin{aligned} & \text { 0008,ERROR-1 } \\ & +05 \end{aligned}$ | Call the error handling routine. |
| :---: | :---: | :---: | :---: | :---: |
| 3705 | N-NEGTV | JR | C,370E,RSLT-ZERO | The result is to be zero if N is less than -255 decimal. |
|  |  | SUB | (HL) | Subtract ABS N from the exponent as N was negative. |
|  |  | $\begin{aligned} & \text { JR } \\ & \text { NEG } \end{aligned}$ | NC,370E,RSLT-ZERO | Zero result if e less than zero. Minus e is changed to $e$. |
| 370C | RESULT-OK | $\begin{aligned} & \text { LD } \\ & \text { RET } \end{aligned}$ | (HL), A | The exponent, e, is entered. |
| 370E | RSLT-ZERO | RST | 0028,FP-CALC | Use the calculator to make the |
|  |  | DEFB | +02,delete | 'last value' zero. |
|  |  | DEFB | +A0,stk-zero |  |
|  |  | DEFB | +38,end-calc |  |
|  |  | RET |  | Finished, with EXP $\mathrm{X}=0$. |

## THE 'NATURAL LOGARITHM' FUNCTION

(Offset 25: 'In')
This subroutine handles the function LN X and is the second of the four routines that use SERIES GENERATOR to produce Chebyshev polynomials.
The approximation to LN X is found as follows:
I. $X$ is tested and report $A$ is given if $X$ is not positive.
II. $X$ is then split into its true exponent, $e^{\prime}$, and its mantissa $X^{\prime}=X /\left(2^{* *} e^{\prime}\right)$, where $X^{\prime}$ is greater than, or equal to, 0.5 but still less than 1.
III. The required value $Y 1$ or $Y 2$ is formed. If $X^{\prime}$ is greater than 0.8 then $Y 1=e^{\prime *} L N 2$ and if otherwise $Y 2=\left(e^{\prime}-1\right)^{*} L N 2$.
IV. If $X^{\prime}$ is greater than 0.8 then the quantity $X^{\prime}-1$ is stacked; otherwise $2^{*} X^{\prime}-1$ is stacked.
V. Now the argument $Z$ is formed, being if $X^{\prime}$ is greater than $0.8, Z=2.5^{*} X^{\prime}-3$; otherwise $Z=5^{\star} X^{\prime}-3$. In each case, $-1<=Z<=1$, as required for the series to converge.
VI. The SERIES GENERATOR is used to produce the required function.
VII. Finally a simply multiplication and addition leads to $\mathrm{LN} X$ being returned as the 'last value'.
3713 In
RST
0028,FP-CALC
X

Perform step i.

| DEFB | +3D,re-stack | $X$ (in full floating-point form) |
| :--- | :--- | :--- |
| DEFB | +31,duplicate | $X, X$ |
| DEFB | +37, greater-0 | $X,(1 / 0)$ |
| DEFB | +00, jump-true | $X$ |
| DEFB | +04,to 371C, VALID | $X$ |
| DEFB | +38, end-calc | $X$ |

Report A - Invalid argument

| 371A REPORT-A | RST | 0008, ERROR |
| :--- | :--- | :--- |
|  | DEFB | +09 |

Call the error handling routine.

Perform step ii.

3 VALID | 371C | DEFB | +A0,stk-zero |
| :--- | :--- | :--- |
|  | DEFB | +02,delete |
|  | DEFB | +38,end-calc |
|  | LD | A,(HL) |
|  | LD | (HL),+80 |
|  | CALL | 2D28,STACK-A |
|  | RST | 0028, FP-CALC |
|  | DEFB | +34,stk-data |
|  | DEFB | +38,exponent+88 |

$\mathrm{X}, 0 \quad$ The deleted 1 is
X
X
The exponent, e, goes into A .
X is reduced to $\mathrm{X}^{\prime}$.
The stack holds: $\mathrm{X}^{\prime}$, e.
$\mathrm{X}^{\prime}, \mathrm{e}$
$\mathrm{X}^{\prime}, \mathrm{e}, 128$ (decimal)

| DEFB | $+00,(+00,+00,+00)$ |  |
| :--- | :--- | :--- |
| DEFB | +03, subtract | $X^{\prime}, e^{\prime}$ |

Perform step iii

|  |  | DEFB | +01, exchange | e', X' |
| :---: | :---: | :---: | :---: | :---: |
|  |  | DEFB | +31, duplicate | e', X', X' |
|  |  | DEFB | +34,stk-data | e', X', X',0.8 (decimal) |
|  |  | DEFB | +F0, exponent+80 |  |
|  |  | DEFB | +4C,+CC,+CC,+CD |  |
|  |  | DEFB | +03, subtract | e', X', X'-0.8 |
|  |  | DEFB | +37, greater-0 | $e^{\prime}, X^{\prime},(1 / 0)$ |
|  |  | DEFB | +00,jump-true | $\mathrm{e}^{\prime}, \mathrm{X}^{\prime}$ |
|  |  | DEFB | +08,to 373D, GRE. 8 | e', X' |
|  |  | DEFB | +01, exchange | $\mathrm{X}^{\prime}, \mathrm{e}^{\prime}$ |
|  |  | DEFB | +A1,stk-one | X', e', 1 |
|  |  | DEFB | +03,subtract | $X^{\prime}, e^{\prime}-1$ |
|  |  | DEFB | +01, exchange | $e^{\prime}-1, X^{\prime}$ |
|  |  | DEFB | +38,end-calc | $e^{\prime}-1, X^{\prime}$ |
|  |  | INC | (HL) | Double $\mathrm{X}^{\prime}$ to give $2^{*} \mathrm{X}^{\prime}$. |
|  |  | RST | 0028,FP-CALC | $e^{\prime}-1,2^{*} \mathrm{X}^{\prime}$ |
| 373D | GRE. 8 | DEFB | +01,exchange | X', e' - X' large. |
|  |  |  |  | $2^{*} \mathrm{X}^{\prime}, \mathrm{e}^{\prime}-1$ - $\mathrm{X}^{\prime}$ small. |
|  |  | DEFB | +34,stk-data | X',e',LN 2 |
|  |  | DEFB | +F0, exponent+80 | 2* ${ }^{\prime}$, e' ${ }^{\prime}$-1, LN 2 |
|  |  | DEFB | +31,+72,+17,+F8 |  |
|  |  | DEFB | +04,multiply | $\mathrm{X}^{\prime}, \mathrm{e}^{\prime *} \mathrm{LN} 2=\mathrm{Y} 1$ |
|  |  |  |  | $2^{*} X^{\prime}$, (e'-1)* ${ }^{*}$ N 2 = Y2 |

Perform step iv.

| DEFB | +01,exchange | $\mathrm{Y} 1, \mathrm{X}^{\prime} \quad-\mathrm{X}^{\prime}$ large. |
| :---: | :---: | :---: |
|  |  | Y2, 2* $\mathrm{X}^{\prime}$ - $\mathrm{X}^{\prime}$ small. |
| DEFB | +A2,stk-half | Y1, X', . 5 (decimal) |
|  |  | Y2, 2*X', . 5 |
| DEFB | +03,subtract | Y1, X'-. 5 |
|  |  | Y2, 2* ${ }^{\prime}$-. 5 |
| DEFB | +A2,stk-half | Y1, X'-.5, . 5 |
|  |  | Y2, 2*X'-.5, . 5 |
| DEFB | +03,subtract | Y1, X'-1 |
|  |  | Y2, 2* $\mathrm{X}^{\prime}-1$ |

Perform step v.

| DEFB | +31, duplicate | $Y, X^{\prime}-1, X^{\prime}-1$ |
| :--- | :--- | :--- |
|  |  | $Y 2,2^{*} X^{\prime}-1,2^{*} X^{\prime}-1$ |
| DEFB | +34, stk-data | $Y 1, X^{\prime}-1, X^{\prime}-1,2.5$ (decimal) |
| DEFB | +32, exponent +82 | $Y 2,2^{*} X^{\prime}-1,2^{*} X^{\prime}-1,2.5$ |
| DEFB | $+20,(+00,+00,+00)$ |  |
| DEFB | +04, multiply | $Y 1, X^{\prime}-1,2.5^{*} X^{\prime}-3=Z$ |
|  |  | $Y 2,2^{*} X^{\prime}-1,5^{*} X^{\prime}-3=Z$ |

Perform step vi, passing to the SERIES GENERATOR the parameter ' 12 ' decimal, and the twelve constant required.

| DEFB | $+8 C$, series $-0 C$ | $Y 1, X^{\prime}-1, Z$ or $Y 2,2^{*} X^{\prime}-1, Z$ |
| :--- | :--- | :--- |
| DEFB | +11, exponent+61 |  |
| DEFB | + AC,(+00,+00,+00) |  |
| DEFB | +14, exponent +64 |  |
| DEFB | $+09,(+00,+00,+00)$ |  |
| DEFB | +56, exponent+66 |  |
| DEFB | $+D A,+$ A5,(+00,+00) |  |


| 4. | DEFB | +59,exponent+69 |
| :---: | :---: | :---: |
|  | DEFB | +30,+C5,(+00,+00) |
| 5. | DEFB | +5C,exponent+6C |
|  | DEFB | +90,+AA, (+00,+00) |
| 6. | DEFB | +9E,exponent+6E |
|  | DEFB | +70,+6F,+61,(+00) |
| 7. | DEFB | +A1, exponent+71 |
|  | DEFB | +CB,+DA,+96,(+00) |
| 8. | DEFB | +A4, exponent+74 |
|  | DEFB | +31,+9F,+B4,(+00) |
| 9. | DEFB | +E7,exponent+77 |
|  | DEFB | +A0,+FE, +5C,+FC |
| 10. | DEFB | +EA, exponent+7A |
|  | DEFB | +1B,+43,+CA, +36 |
| 11. | DEFB | +ED, exponent+7D |
|  | DEFB | +A7,+9C,+7E,+5E |
| 12. | DEFB | +F0, exponent+80 |
|  | DEFB | +6E,+23,+80,+93 |

At the end of the last loop the 'last value' is:

| either | LN X |
| :--- | :--- |
| or | $\mathrm{LN}\left(2^{\star} \mathrm{X}^{\prime}\right) /\left(2^{\star}-\mathrm{X}^{\prime}-1\right)$ for the targer values of $\mathrm{X}^{\prime}$ |
|  |  |

Perform step vii.

| DEFB | +04, multiply | $\mathrm{Y} 1=\mathrm{LN}\left(2^{\star *} e^{\prime}\right)$, LN X' |  |
| :--- | :--- | :--- | :--- |
|  |  | $\mathrm{Y} 2=\mathrm{LN}\left(2^{* *}\left(e^{\prime}-1\right)\right)$, LN $\left(2^{*} \mathrm{X}^{\prime}\right)$ |  |
| DEFB | $+0 F$, addition | LN $\left.\left(2^{* *} e^{\prime}\right)^{\star} X^{\prime}\right)$ | $=$ LN X |
|  |  | LN $\left(2^{* *}\left(e^{\prime}-1\right)^{*} 2^{*} X^{\prime}\right)$ | $=$ LN X |
| DEFB | +38, end-calc | LN X |  |
| RET |  | Finished: 'last value' is LN X. |  |

## THE 'REDUCE ARGUMENT' SUBROUTINE

(Offset 39: 'get-argt')
This subroutine transforms the argument X of $\operatorname{SIN} \mathrm{X}$ or COS X into a value V .
The subroutine first finds a value Y such that:
$\left.\mathrm{Y}=\mathrm{X} /\left(2^{*} \mathrm{PI}\right)-\mathrm{INT}\left(\mathrm{X} / 2^{*} \mathrm{PI}\right)+0.5\right)$, where Y is greater than, or equal to, -.5 but less than +.5 .
The subroutine returns with:
$\mathrm{V}=4^{*} \mathrm{Y}$ if $-1<=4^{*} \mathrm{Y}<=1$ - case i .
or, $V=2-4^{*} Y \quad$ if $1<4^{*} Y<2 \quad$ - case ii.
or, $V=-4^{*} Y-2 \quad$ if $-2<=4^{*} Y<-1$. - case iii.
In each case, $-1<=\mathrm{V}<=1$ and $\operatorname{SIN}(\mathrm{PI} * \mathrm{~V} / 2)=\mathrm{SIN} \mathrm{X}$


Note: Adding 0.5 and taking INT rounds the result to the nearest integer.

| DEFB | +31, duplicate | Y, Y |
| :---: | :---: | :---: |
| DEFB | +0F, addition | 2*Y |
| DEFB | +31, duplicate | 2*Y, 2*Y |
| DEFB | +0F, addition | 4*Y |
| DEFB | +31, duplicate | 4*Y, 4*Y |
| DEFB | +2A, abs | 4*Y, ABS (4*Y) |
| DEFB | +A1,stk-one | 4*Y, ABS (4*Y), 1 |
| DEFB | +03, subtract | $4^{*} \mathrm{Y}, \mathrm{ABS}\left(4^{*} \mathrm{Y}\right)-1=\mathrm{Z}$ |
| DEFB | +31, duplicate | $4^{*} \mathrm{Y}, \mathrm{Z}, \mathrm{Z}$ |
| DEFB | +37, greater-0 | 4*Y, Z, (1/0) |
| DEFB | +C0,st-mem-0 | Mem-0 holds the result of the test. |
| DEFB | +00,jump-true | 4*Y, Z |
| DEFB | +04, to 37A1,ZPLUS | 4*Y, Z |
| DEFB | +02, delete | 4*Y |
| DEFB | +38,end-calc | $4^{*} \mathrm{Y}=\mathrm{V}$ - case i . |
| RET |  | Finished. |

If the jump was made then continue.

| 37A1 | ZPLUS | DEFB | +A1,stk-one | $4^{*} \mathrm{Y}, \mathrm{Z}, 1$ |
| :---: | :---: | :---: | :---: | :---: |
|  |  | DEFB | +03,subtract | 4*Y, Z-1 |
|  |  | DEFB | +01,exchange | Z-1,4*Y |
|  |  | DEFB | +36,less-0 | Z-1,(1/0) |
|  |  | DEFB | +00,jump-true | Z-1 |
|  |  | DEFB | +02,to 37A8,YNEG | Z-1 |
|  |  | DEFB | +1B, negate | 1-Z |
| 37A8 | YNEG | DEFB | +38,end-calc | $1-\bar{Z}=V-\text { case ii. }$ $\mathrm{Z}-1=\mathrm{V} \text { - case iii. }$ |
|  |  | RET |  | Finished. |

## THE 'COSINE' FUNCTION

(Offset 20: 'cos')
This subroutine handles the function $\operatorname{COS} \mathrm{X}$ and returns a 'last value' 'that is an approximation to COS X .
The subroutine uses the expression:
$\operatorname{COS} X=\operatorname{SIN}\left(P I^{*} \mathrm{~W} / 2\right)$, where $-1<=\mathrm{W}<=1$.
In deriving $W$ for $X$ the subroutine uses the test result obtained in the previous subroutine and stored for this purpose in mem- 0 . It then jumps to the SINE, subroutine, entering at C-ENT, to produce a 'last value' of COS X.

| 37AA | cos | RST <br> DEFB <br> DEFB <br> DEFB <br> DEFB <br> DEFB <br> DEFB <br> DEFB | $\begin{aligned} & \text { 0028,FP-CALC. } \\ & \text { +39,get-argt } \\ & \text { +2A,abs } \\ & \text { +A1,stk-one } \\ & \text { +03,subtract } \\ & \text { +E0,get-mem-0 } \\ & \text { +00,jump-true } \\ & \text { +06, to 37B7,C-ENT } \end{aligned}$ | X <br> V <br> ABS V <br> ABS V, 1 <br> ABS V-1 <br> ABS V-1, (1/0) <br> ABS V-1 <br> ABS V-1 $=W$ |
| :---: | :---: | :---: | :---: | :---: |
| If the jump was not made then continue. |  |  |  |  |
|  |  | DEFB | +1B, negate | 1-ABS V |
|  |  | DEFB | +33,jump | 1-ABS V |
|  |  | DEFB | +03, to 37B7,C-ENT | $1-A B S V=W$ |

## THE 'SINE' FUNCTION

(Offset 1F: 'sin')
This subroutine handles the function SIN X and is the third of the four routines that use SERIES GENERATOR to produce Chebyshev polynomials.
The approximation to SIN X is found as follows:
i. The argument X is reduced and in this case $\mathrm{W}=\mathrm{V}$ directly.

Note that $-1<=\mathrm{W}<=1$, as required for the series to converge.
ii. The argument $Z$ is formed, such that $Z=2^{*} W * W-1$.
iii. The SERIES GENERATOR is used to return (SIN (PI*W/2))/W
iv. Finally a simple multiplication gives SIN X.
37B5 $\sin \quad$ RST 0028 FP-CALC X

Perform step i.
DEFB +39, get-argt W

Perform step ii. The subroutine from now on is common to both the SINE and COSINE functions.

| 37B7 | C-ENT | DEFB | +31, duplicate | W, W |
| :---: | :---: | :---: | :---: | :---: |
|  |  | DEFB | +31, duplicate | W, W, W |
|  |  | DEFB | +04, multiply | W, W*W |
|  |  | DEFB | +31, duplicate | W, W*W, W*W |
|  |  | DEFB | +0F,addition | W, 2*W*W |
|  |  | DEFB | +A1,stk-one | W, 2*W*W, 1 |
|  |  | DEFB | +03,subtract | W, 2*W*W-1 = Z |

Perform step iii, passing to the SERIES GENERATOR the parameter ' 6 ' and the six constants required.

|  | DEFB | +86, series-06 | W, Z |
| :---: | :---: | :---: | :---: |
| 1. | DEFB | +14,exponent+64 |  |
|  | DEFB | +E6,(+00,+00,+00) |  |
| 2. | DEFB | +5C,exponent+6C |  |
|  | DEFB | +1F,+0B,(+00,+00) |  |
| 3. | DEFB | +A3,exponent+73 |  |
|  | DEFB | +8F,+38,+EE,(+00) |  |
| 4. | DEFB | +E9,exponent+79 |  |
|  | DEFB | +15,+63,+BB,+23 |  |
| 5. | DEFB | +EE, exponent+7E |  |
|  | DEFB | +92,+0D,+CD, +ED |  |
| 6. | DEFB | +F1,exponent+81 |  |
|  | DEFB | +23,+5D,+1B,+EA |  |

At the end of the last loop the 'last value' is (SIN ( $\mathrm{PI*} \mathrm{~W} / 2$ ) )/W.
Perform step v.

| DEFB | +04, multiply | SIN $\left(\mathrm{PI}^{*} \mathrm{~W} / 2\right)=$ SIN X $($ or $=$ |
| :--- | :--- | :--- |
| COFB | +38, end-calc |  |
| RET |  | Finished: 'last value' $=$ SIN X <br> or ('last value' $=\operatorname{COS} X)$ |
|  |  |  |

## THE 'TAN' FUNCTION

(Offset 21: 'tan')
This subroutine handles the function TAN $X$. The subroutine simply returns SIN X/COS $X$, with arithmetic overflow if COS $X=0$.

| 37DA | tan | RST | 0028,FP-CALC | X |
| :---: | :---: | :---: | :---: | :---: |
|  |  | DEFB | +31, duplicate | X, X |
|  |  | DEFB | +1F,sin | X, SIN X |
|  |  | DEFB | +01, exchange | SIN X, X |
|  |  | DEFB | +20, cos | SIN X,COS X |
|  |  | DEFB | +05,division | SIN X/COS X = TAN X |
|  |  |  |  | Report arithmetic overflow if needed. |
|  |  | DEFB | +38,end-calc | TAN X |
|  |  | RET |  | Finished: 'last value' = TAN X. |

## THE 'ARCTAN' FUNCTION

(Offset 24: 'atn')
This subroutine handles the function ATN $X$ and is the last of the four routines that use SERIES GENERATOR to produce Chebyshev polynomials. It returns a real number between $-\mathrm{Pl} / 2$ and $\mathrm{PI} / 2$, which is equal to the value in radians of the angle whose tan is X . The approximation to ATN $X$ is found as follows:
i. The values $W$ and $Y$ are found for three cases of $X$, such that:

| if $-1<X<1$ | then $W=0$ | $\& Y=X$ |
| :--- | :--- | :--- |
| if $-1<=X$ then $W=P I / 2 \& Y=-1 / X$ | - case ii. |  |
| if $X<=-1$ then $W=-\mathrm{PI} / 2 \& Y=-1 / X$ | - case iii. |  |

In each case, $-1<=Y<=1$, as required for the series to converge.
ii. The argument $Z$ is formed, such that:
if $-1<X<1 \quad$ then $Z=2^{\star} Y^{\star} Y-1=2^{*} X^{\star} X-1 \quad$ - case $i$.
if $1<X$ then $Z=2^{*} Y^{*} Y-1=2 /\left(X^{*} X\right)-1$ - case ii.
if $X<=-1$ then $Z=2^{*} Y^{*} Y-1=2 /\left(X^{*} X\right)-1 \quad$ - case iii.
iii. The SERIES GENERATOR is used to produce the required function
iv. Finally a simple multiplication and addition give ATN X.

Perform stage i.

| 37E2 | atn | CALL | 3297,RE-STACK | Use the full floating-point form of $X$. |
| :---: | :---: | :---: | :---: | :---: |
|  |  | LD | A,(HL) | Fetch the exponent of $X$. |
|  |  | CP | +81 |  |
|  |  | JR | C,37F8,SMALL | Jump forward for case $\mathrm{i}: \mathrm{Y}=\mathrm{X}$. |
|  |  | RST | 0028,FP-CALC | X |
|  |  | DEFB | +A1,stk-one | X, 1 |
|  |  | DEFB | +1B,negate | X,-1 |
|  |  | DEFB | +01, exchange | -1, X |
|  |  | DEFB | +05,division | -1/X |
|  |  | DEFB | +31,duplicate | -1/X, -1/X |
|  |  | DEFB | +36, less-0 | -1/X, (1/0) |
|  |  | DEFB | +A3,stk-pi/2 | -1/X, (1/0), PI/2 |
|  |  | DEFB | +01, exchange | -1/X, PI/2, (1/0) |
|  |  | DEFB | +00,jump-true | -1/X, PI/2 |
|  |  | DEFB | +06, to 37FA,CASES | Jump forward for case ii: $Y=-1 / X \quad W=P I / 2$ |
|  |  | DEFB | +1B, negate | -1/X, -PI/2 |
|  |  | DEFB | +33,jump | -1/X, -PI/2 |
|  |  | DEFB | +03,to 37FA,CASES | Jump forward for case iii: $Y=-1 / X \quad W=-P I / 2$ |
| 37F8 | SMALL | RST | 0028,FP-CALC | Y |
|  |  | DEFB | +A0,stk-zero | Y, 0 |
|  |  |  |  | Continue for case i: W = 0 |

Perform step ii.

| 37FA | CASES | DEFB | +01,exchange | W, Y |
| :---: | :---: | :---: | :---: | :---: |
|  |  | DEFB | +31,duplicate | W, Y, Y |
|  |  | DEFB | +31,duplicate | W, Y, Y, Y |
|  |  | DEFB | +04, multiply | W, Y, Y*Y |
|  |  | DEFB | +31,duplicate | $W, Y, Y^{*} Y, Y^{*} Y$ |
|  |  | DEFB | +0F, addition | W, Y, 2* ${ }^{*} \mathrm{Y}$ |
|  |  | DEFB | +A1,stk-one | W, Y, $2^{*} Y^{*} \mathrm{Y}, 1$ |
|  |  | DEFB | +03,subtract | $W, Y, 2^{*} Y^{*} Y-1=Z$ |

Perform step iii, passing to the SERIES GENERATOR the parameter ' 12 ' decimal, and the twelve constants required.

| DEFB | +8 C,series-0C | $\mathrm{W}, \mathrm{Y}, \mathrm{Z}$ |
| :--- | :--- | :--- |
| DEFB | +10, exponent +60 |  |
| DEFB | $+B 2,(+00,+00,+00)$ |  |


| 2. | DEFB | +13,exponent+63 |
| :---: | :---: | :---: |
|  | DEFB | +0E,(+00,+00,+00) |
| 3. | DEFB | +55, exponent+65 |
|  | DEFB | +E4,+8D,(+00,+00) |
| 4. | DEFB | +58,exponent+68 |
|  | DEFB | +39,+BC,(+00,+00) |
| 5. | DEFB | +5B,exponent+6B |
|  | DEFB | +98,+FD, (+00,+00) |
| 6. | DEFB | +9E, exponent+6E |
|  | DEFB | +00,+36,+75,(+00) |
| 7. | DEFB | +A0,exponent+70 |
|  | DEFB | +DB,+E8,+B4,(+00) |
| 8. | DEFB | +63,exponent+73 |
|  | DEFB | +42,+C4,(+00,+00) |
| 9. | DEFB | +E6,exponent+76 |
|  | DEFB | +B5,+09,+36,+BE |
| 10. | DEFB | +E9, exponent+79 |
|  | DEFB | +36,+73, +1B, +5D |
| 11. | DEFB | +EC,exponent+7C |
|  | DEFB | +D8,+DE, +63,+BE |
| 12. | DEFB | +F0,exponent+80 |
|  | DEFB | $+61,+\mathrm{A} 1,+\mathrm{B} 3,+0 \mathrm{C}$ |

At the end of the last loop the 'last value' is:
ATN X/X - case i.
ATN $(-1 / X) /(-1 / X) \quad$ - case ii.
ATN $(-1 / X) /(-1 / X) \quad$ - case iii.
Perform step iv.

| DEFB | +04, multiply | W, ATN X- case i. |
| :--- | :--- | :--- |
|  |  | W, ATN $(-1 / X)$ |
| WEFB | +0 case ii. |  |
| DE, addition | WTN $(-1 / X) \quad-$ case iii. |  |
| DEFB | +38, end-calc | ATN X - all cases now. |
| RET |  |  |
|  |  |  |

## THE 'ARCSIN' FUNCTION

(Offset 22: 'asn')
This subroutine handles the function ASN $X$ and return a real real number from $-\mathrm{PI} / 2$ to $\mathrm{PI} / 2$ inclusive which is equal to the value in radians of the angle whose sine is $X$. Thereby if $Y=A S N X$ then $X=S I N Y$.
This subroutine uses the trigonometric identity:
$\operatorname{TAN}(\mathrm{Y} / 2)=\operatorname{SIN~Y/1}(1+\operatorname{COS} \mathrm{Y})$
to obtain TAN (Y/2) and hence (using ATN) Y/2 and finally Y.

| 3833 | asn | RST | 0028,FP-CALC | $X$ |
| :---: | :---: | :---: | :---: | :---: |
|  |  | DEFB | +31, duplicate | X, X |
|  |  | DEFB | +31, duplicate | X, X, X |
|  |  | DEFB | +04,multiply | X, ${ }^{*}$ X |
|  |  | DEFB | +A1,stk-one | X, $\mathrm{X}^{*}$ X, 1 |
|  |  | DEFB | +03,subtract | X, $\mathrm{X}^{*} \mathrm{X}$-1 |
|  |  | DEFB | +1B, negate | X,1-X*X |
|  |  | DEFB | +28,sqr | X,SQR (1-X* ${ }^{\text {a }}$ ) |
|  |  | DEFB | +A1,stk-one | X,SQR (1-X*X), 1 |
|  |  | DEFB | +0F,addition | X, $1+$ SQR ( $\left.1-\mathrm{X}^{*} \mathrm{X}\right)$ |
|  |  | DEFB | +05, division | $\begin{aligned} & X /\left(1+S Q R\left(1-X^{*} X\right)\right)=\text { TAN } \\ & (Y / 2) \end{aligned}$ |
|  |  | DEFB | +24,atn | Y/2 |
|  |  | DEFB | +31, duplicate | Y/2, Y/2 |
|  |  | DEFB | +0F,addition | $Y=A S N X$ |
|  |  | $\begin{aligned} & \text { DEFB } \\ & \text { RET } \end{aligned}$ | +38,end-calc | Finished: 'last value' = ASN X. |

## THE 'ARCCOS' FUNCTION

(Offset 23: 'acs')
This subroutine handles the function ACS X and returns a real number from zero to PI inclusive which is equal to the value in radians of the angle whose cosine is $X$.

This subroutine uses the relation:

$$
\mathrm{ACS} X=\mathrm{Pl} / 2-\mathrm{ASN} \mathrm{X}
$$

3843 acs

| RST | 0028, FP-CALC |
| :--- | :--- |
| DEFB | +22, asn |
| DEFB | + +A3,stk-pi/2 |
| DEFN | +03, subtract |
| DEFB | +1 B,negate |
| DEFB | +38, end-calc |
| RET |  |

X
ASN X
ASN X,PI/2
ASN X-PI/2
PI/2-ASN X = ACS X

Finished: 'last value' = ACS X.

## THE 'SQUARE ROOT' FUNCTION

(Offset 28: 'sqr')
This subroutine handles the function SQR $X$ and returns the positive square root of the real number $X$ if $X$ is positive, and zero if $X$ is zero. A negative value of $X$ gives rise to report $A$ - invalid argument (via $\ln$ in the EXPONENTIATION subroutine).
This subroutine treats the square root operation as being $X^{* *} .5$ and therefore stacks the value .5 and proceeds directly into the EXPONENTIATION subroutine.

| RST | 0028, FP-CALC |
| :--- | :--- |
| DEFB | +31,duplicate |
| DEFB | +30, not |
| DEFB | +00 ,jump-true |
| DEFB | $+1 E$, to 386C,LAST |

$X$
$X, X$
$X,(1 / 0)$
$X$
$X$

The jump is made if $X=0$, otherwise continue with:

| DEFB | + A2,stk-half | $X, .5$ |
| :--- | :--- | :--- |
| DEFB | +38, end-calc |  |

and then find the result of $X^{* *} .5$.

## THE 'EXPONENTIATION' OPERATION

(Offset 06: 'to-power')
This subroutine performs the binary operation of raising the first number, X , to the power of the second number, Y
The subroutine treats the result $\mathrm{X}^{* *} \mathrm{Y}$ as being equivalent to $\mathrm{EXP}\left(\mathrm{Y}^{*} \mathrm{LN} X\right)$. It returns this value unless X is zero, in which case it returns 1 if $Y$ is also zero $\left(0^{* *} 0=1\right)$, returns zero if $Y$ is positive and reports arithmetic overflow if $Y$ is negative.

3851 to-power

| RST | 0028,FP-CALC |
| :--- | :--- |
| DEFB | +01, exchange |
| DEFB | +31, duplicate |
| DEFB | +30, not |
| DEFB | +00, jump-true |
| DEFB | +07, to 385D,XIS0 |

$X, Y$
$Y, X$
$Y, X, X$
$Y, X,(1 / 0)$
$Y, X$
$Y, X$

The jump is made if $X=0$, otherwise EXP $\left(Y^{*} L N X\right)$ is formed.
DEFB $\quad+25, \ln$
Y,LN X
Giving report A if X is negative.

| DEFB | +04, multiply | $Y^{*}$ LN X |
| :--- | :--- | :--- |
| DEFB | +38, end-calc |  |
| JP | 36C4,EXP | Exit via EXP to form EXP |
|  |  | $(Y *$ LN X). |

The value of X is zero so consider the three possible cases involved.

| 385D XIS0 | DEFB | +02,delete | Y |
| :--- | :--- | :--- | :--- |
|  | DEFB | +31,duplicate | Y,Y |
|  | DEFB | +30,not | Y,(1/0) |
|  | DEFB | +00, jump-true | $Y$ |

DEFB +09,to 386A,ONE Y
The jump is made if $\mathrm{X}=0$ and $\mathrm{Y}=0$, otherwise proceed.

| DEFB | +A0,stk-zero | Y,0 |
| :--- | :--- | :--- |
| DEFB | +01,exchange | $0, Y$ |
| DEFB | +37,greater-0 | $0,(1 / 0)$ |
| DEFB | +00,jump-true | 0 |
| DEFB | +06,to 386 C,LAST | 0 |

The jump is made if $X=0$ and $Y$ is positive, otherwise proceed.

| DEFB | +A1,stk-one | 0,1 |
| :--- | :--- | :--- |
| DEFB | +01,exchange | 1,0 |
| DEFB | +05,division | Exit via 'division' as dividing by |
|  |  | zero gives 'arithmetic overflow'. |

The result is to be 1 for the operation.
386A ONE DEFB +02,delete
DEFB +A1,stk-one
1
Now return with the 'last value' on the stack being $0^{* *}$ Y.
386C LAST
DEFB +38 ,end-calc
RET
(1/0)
Finished: 'last value' is 0 or 1 .
386E - 3CFF These locations are 'spare'. They all hold +FF.
3D00-3FFF These locations hold the 'character set'. There are 8 byte representations for all the characters with codes +20 (space) to +7 F (©).
e.g. the letter ' $A$ ' has the representation $003 C 4242$ 7E 424200 and thereby the form:

00000000
00111100
01000010
01000010
01111110
01000010
01000010
00000000

## APPENDIX

## BASIC PROGRAMS FOR THE MAIN SERIES

The following BASIC programs have been included as they give a good illustration of how Chebyshev polynomials are used to produce the approximations to the functions SIN, EXP, LN and ATN.

The series generator:
This subroutine is called by all the 'function' programs.

```
5 0 0 ~ R E M ~ S E R I E S ~ G E N E R A T O R , ~ E N T E R ~
5 1 0 \text { REM USING THE COUNTER BREG}
5 2 0 ~ R E M ~ A N D ~ A R R A Y - A ~ H O L D I N G ~ T H E ~
530 REM CONSTANTS
540 REM FIRST VALUE IN Z.
550 LET M0=2*Z
500 LET M2=0
570 LET T=0
580 FOR I=BREG TO 1 STEP -1
590 LET M1=M2
600 LET U=T*M0-M2+A(BREG+1-I)
6 1 0 ~ L E T ~ M 2 = T ~ T
6 2 0 ~ L E T ~ T = U ~
6 3 0 ~ N E X T ~ I ~
6 4 0 ~ L E T ~ T = T - M 1 ~
6 5 0 ~ R E T U R N
660 REM LAST VALUE IN T.
```

In the above subroutine the variable are:

| $Z-$ | the entry value. |
| :--- | :--- |
| T - | the exit value. |
| M0 - | mem-0 |
| M1- | mem-1 |
| M2 - | mem-2 |
| $1-$ | the counter for BREG. |
| $\mathrm{U}-$ | a temporary variable for T. |
| $\mathrm{A}(1)$ to |  |
| A(BREG) - | the constants. |
| BREG - | the number of constants to be used. |

To see how the Chebyshev polynomials are generated, record on paper the values of $U$ M1, M2 and T through the lines 550 to 630, passing, say, 6 times through the loop, and keeping the algebraic expressions for $A(1)$ to $A(6)$ without substituting numerical values. Then record T-M1. The multipliers of the constants $A(1)$ to $A(6)$ will then be the required Chebyshev polynomials. More precisely, the multiplier of $A(1)$ will be $2^{*} T 5(Z)$, for $A(2)$ it will be $2^{*} T 4(Z)$ and so on to $2^{*} T 1(Z)$ for $A(5)$ and finally $T 0(Z)$ for $A(6)$.

Note that $T 0(Z)=1, T 1(Z)=Z$ and, for $n>=2, T n(Z)=2^{*} Z^{*} T n-1(Z)-T n-2(Z)$.

## SIN X

```
10 REM DEMONSTRATION FOR SIN X
20 REM USING THE 'SERIES GENERATOR'
30 DIM A(6)
40 LET A(1)=-. 000000003
50 LET A(2)=0.000000592
60 LET A(3)=-. 000068294
70 LET A(4) \(=0.004559008\)
80 LET A(5)=-. 142630785
90 LET \(A(6)=1.276278962\)
100 PRINT
110 PRINT "ENTER START VALUE IN DEGREES"
120 INPUT C
130 CLS
140 LET C=C-10
150 PRINT "BASIC PROGRAM","ROM PROGRAM"
160 PRINT
```

$\qquad$

```
170 PRINT
180 FOR J=1 TO 4
190 LET C=C+10
200 LET Y=C/360-INT (C/360+.5)
210 LET W=4*Y
220 IF \(W>1\) THEN LET W=2-W
230 IF W <-1 THEN LET W=-W-2
240 LET Z=2*W*W-1
250 LET BREG=6
260 REM USE 'SERIES GENERATOR'
270 GO SUB 550
280 PRINT TAB 6; "SIN ";C;" DEGREES"
290 PRINT
300 PRINT T*W,SIN (PI*C/180)
310 PRINT
320 NEXT J
330 GO TO 100
```


## NOTES:

I. When $C$ is entered this program calculates and prints SIN C degrees, SIN ( $C+10$ ) degrees, $\operatorname{SIN}(C+20)$ degrees and SIN (C+30) degrees. It also prints the values obtained by using the ROM program. For a specimen of results, try entering these values in degrees: 0; 5; 100; -80; -260; 3600; -7200.
II. The constants $A(1)$ to $A(6)$ in lines 40 to 90 are given (apart from a factor of $1 / 2$ ) in Abramowitz and Stegun Handbook of Mathematical Functions (Dover 1965) page 76. They can be checked by integrating (SIN ( $\mathrm{PI}{ }^{*} \mathrm{X} / 2$ ) )/X over the interval $\mathrm{U}=0$ to PI , after first multiplying by $\operatorname{COS}\left(N^{*} U\right)$ for each constant (i.e. $\left.N=1,2, \ldots, 6\right)$ and substituting $\operatorname{COS} U=2^{*} X^{*} X-1$. Each result should then be divided by PI. (This integration can be performed by approximate methods e.g. using Simpson's Rule if there is a reasonable computer or programmable calculator to hand.)

## EXP X

| 10 | REM DEMONSTRATION FOR EXP X |
| :---: | :---: |
| 20 | REM USING THE 'SERIES GENERATOR' |
| 30 | LET $\mathrm{T}=0 \quad$ (This makes T the first variable.) |
| 40 | DIM A(8) |
| 50 | LET A(1)=0.000000001 |
| 60 | LET A(2)=0.000000053 |
| 70 | LET A(3)=0.000001851 |
| 80 | LET A(4)=0.000053453 |
| 90 | LET A(5)=0.001235714 |
| 100 | LET A(6)=0.021446556 |
| 110 | LET A(7)=0.248762434 |
| 120 | LET A(8)=1.456999875 |
| 130 | PRINT |
| 140 | PRINT "ENTER START VALUE" |
| 150 | INPUT C |
| 160 | CLS |
| 170 | LET C=C-10 |
| 180 | PRINT "BASIC PROGRAM", "ROM PROGRAM" |
| 190 | PRINT "------------", "-----------" |
| 200 | PRINT |
| 210 | FOR J=1 TO 4 |
| 220 | LET C=C+10 |
| 230 | LET D=C*1.442695041 ( $\mathrm{D}=\mathrm{C}^{*}(1 / \mathrm{LN} 2)$;EXP C=2**D). |
| 240 | LET $\mathrm{N}=$ INT D |
| 250 | LET Z=D-N ( $2^{* *}(\mathrm{~N}+\mathrm{Z})$ is now required $)$. |
| 260 | LET Z=2*Z-1 |
| 270 | LET BREG=8 |
| 280 | REM USE "SERIES GENERATOR" |
| 290 | GO SUB 550 |
| 300 | LET V=PEEK 23627+256*PEEK 23628+1 (V=(VARS)+1) |
| 310 | LET N=N+PEEK V |
| 320 | IF $\mathrm{N}>255$ THEN STOP (STOP with arithmetic overflow). |
| 330 | IF N < 0 THEN GO TO 360 |
| 340 | POKE V,N |
| 350 | GO TO 370 |
| 360 | LET T=0 |
| 370 | PRINT TAB 11;"EXP ";C |
| 380 | PRINT |
| 390 | PRINT T,EXP C |
| 400 | PRINT |
| 410 | NEXT J |
|  | GO TO 130 |

## NOTES:

I. When C is entered this program calculates and prints EXP C, EXP $(\mathrm{C}+10)$, EXP $(\mathrm{C}+20)$ and EXP $(\mathrm{C}+30)$. It also prints the values obtained by using the ROM program. For a specimen of results, try entering these values: $0 ; 15 ; 65$ (with overflow at the end); -100; -40.
II. The exponent is tested for overflow and for a zero result in lines 320 and 330 . These tests are simpler in BASIC than in machine code, since the variable N , unlike the A register, is not confined to one byte.
III. The constants $\mathrm{A}(1)$ to $\mathrm{A}(8)$ in lines 50 to 120 can be obtained by integrating $2^{* *} \mathrm{X}$ over the interval $\mathrm{U}=0$ to PI , after first multiplying the $\operatorname{COS}\left(N^{*} U\right)$ for each constant (i.e. for $\left.N=1,2, \ldots, 8\right)$ and substituting $\operatorname{COS} U=2^{*} X-1$. Each result should then be divided by PI.

## LN X:

| 10 | REM DEMONSTRATION FOR LN X |
| :---: | :---: |
| 20 | REM USING THE 'SERIES GENERATOR' |
| 30 | LET $\mathrm{D}=0 \quad$ (This makes D the first variable). |
| 40 | DIM A(12) |
| 50 | LET A(1)= -. 0000000003 |
| 60 | LET A(2)=0.0000000020 |
| 70 | LET A(3) = -.0000000127 |
| 80 | LET A(4)=-0.0000000823 |
| 90 | LET A(5) $=-.0000005389$ |
| 100 | LET $A(6)=0.0000035828$ |
| 110 | LET A(7)= .0000243013 |
| 120 | LET $A(8)=0.0001693953$ |
| 130 | LET A(9)= .0012282837 |
| 140 | LET $\mathrm{A}(10)=0.0094766116$ |
| 150 | LET A(11) = .0818414567 |
| 160 | LET A(12)=0.9302292213 |
| 170 | PRINT |
| 180 | PRINT "ENTER START VALUE" |
| 190 | INPUT C |
| 200 | CLS |
| 210 | PRINT "BASIC PROGRAM", "ROM PROGRAM" |
| 220 | PRINT "-------------", "----------" |
| 230 | PRINT |
| 240 | LET C=SQR C |
| 250 | FOR J=1 TO 4 |
| 260 | LET C=C*C |
| 270 | IF C=0 THEN STOP (STOP with 'invalid argument'.) |
| 280 | LET D=C |
| 290 | LET V=PEEK 23627+256*PEEK 23628+1 |
| 300 | LET N=PEEK V-128 ( N holds $\mathrm{e}^{\prime}$ ). |
| 310 | POKE V,128 |
| 320 | IF D<=0.8 THEN GO TO 360 (D holds X'). |
| 330 | LET S=D-1 |
| 340 | LET $\mathrm{Z}=2.5^{*} \mathrm{D}-3$ |
| 350 | GO TO 390 |
| 360 | LET $\mathrm{N}=\mathrm{N}-1$ |
| 370 | LET S=2*D-1 |
| 380 | LET $\mathrm{Z}=5 \times \mathrm{D}-3$ |
| 390 | LET R=N*0.6931471806 (R holds N*LN 2). |
| 400 | LET BREG=12 |
| 410 | REM USE 'SERIES GENERATOR' |
| 420 | GO SUB 550 |
| 430 | PRINT TAB 8;"LN ";C |
| 440 | PRINT |
| 450 | PRINT S*T+R,LN C |
| 460 | PRINT |
| 470 | NEXT J |
| 480 | GO TO 170 |

## NOTES:

I. When $C$ is entered this program calculates and prints $L N C, L N\left(C^{* *} 2\right), L N\left(C^{* *} 4\right)$ and $L N\left(C^{* *} 8\right)$. It also prints the values obtained by using the ROM program. For a specimen of results, try entering these values: 1.1; 0.9; 300; 0.004; 1E5 (for overflow) and 1E-5 (STOP as 'invalid argument')
II. The constants $A(1)$ to $A(12)$ in lines 50 to 160 can be obtained by integrating $5^{*} L N\left(4^{*}(X+1) / 5\right) /\left(4^{*} X-1\right)$ over the interval $U=0$ to PI , after first multiplying by COS ( $\mathrm{N}^{*} \mathrm{U}$ ) for each constant (i.e. for $\mathrm{N}=1,2, \ldots, 12$ ) and substituting COS $\mathrm{U}=2^{*} \mathrm{X}-1$. Each result should then be divided by PI.

## ATN X:

| 10 | REM DEMONSTRATION FOR ATN $X$ |
| :---: | :---: |
| 20 | REM USING THE 'SERIES GENERATOR' |
| 30 | DIM A(12) |
| 40 | LET A(1) $=-.0000000002$ |
| 50 | LET A(2)=0.0000000010 |
| 60 | LET $\mathrm{A}(3)=-.0000000066$ |
| 70 | LET A(4)=0.0000000432 |
| 80 | LET A(5) $=-.0000002850$ |
| 90 | LET A(6)=0.0000019105 |
| 100 | LET A(7) $=-.0000131076$ |
| 110 | LET A (8)=0.0000928715 |
| 120 | LET A $(9)=-.0006905975$ |
| 130 | LET A(10) $=0.0055679210$ |
| 140 | LET A(11)= -.0529464623 |
| 150 | LET A(12)=0.8813735870 |
| 160 | PRINT |
| 170 | PRINT "ENTER START VALUE" |
| 180 | INPUT C |
| 190 | CLS |
| 200 | PRINT "BASIC PROGRAM", "ROM PROGRAM" |
| 210 | PRINT "-------------", "----------" |
| 220 | PRINT |
| 230 | FOR J=1 TO 4 |
| 240 | LET B=J*C |
| 250 | LET D=B |
| 260 | IF ABS B>=1 THEN LET $\mathrm{D}=-1 / \mathrm{B}$ |
| 270 | LET $\mathrm{Z}=2^{*} \mathrm{D}^{*} \mathrm{D}-1$ |
| 280 | LET BREG=12 |
| 290 | REM USE "SERIES GENERATOR" |
| 300 | GO SUB 550 |
| 310 | LET T=D*T |
| 320 | IF B > = 1 THEN LET T=T+P//2 |
| 330 | IF B < =-1 THEN LET T=T-P//2 |
| 340 | PRINT TAB 8;"ATN ";B |
| 350 | PRINT |
| 360 | PRINT T,ATN B (or PRINT T*180/PI,ATN B*180/PI |
| 370 | PRINT to obtain the answers in degrees) |
| 380 | NEXT J |
| 390 | GO TO 160 |

## NOTES:

I. When C is entered this program calculates and prints ATN C, ATN (C*2), ATN (C*3) and ATN (C*4). For a specimen of results, try entering these values: $0.2 ;-1 ; 10$ and -100 . The results may be found more interesting if converted to yield degrees by multiplying the answers in line 360 by 180/PI.
II. The constants $A(1)$ to $A(12)$ in lines 40 to 150 are given (apart from a factor of $1 / 2$ ) in Abramowitz and Stegun Handbook of Mathematical Functions (Dover 1965) page 82. They can be checked by integrating ATN X/X over the interval U=0 to PI, after first multiplying by $\operatorname{COS}\left(\mathrm{N}^{*} \mathrm{U}\right)$ for each parameter (i.e. for $\mathrm{n}=1,2, \ldots, 12$ ) and substituting $\operatorname{COS} \mathrm{U}=2^{*} \mathrm{X}^{*} \mathrm{X}-1$. Each result should then be divided by PI.

An alternative subroutine for SIN X:
It is straightforward to produce the full expansion of the Chebyshev polynomials and this can be written in BASIC as follows:

```
550 LET T =(32*Z*Z*Z*Z*Z-40*Z*Z*Z+10*Z)*A(1)
    +(16*Z*Z*Z*Z-16*Z*Z+2)*A(2)
    +(8*Z*Z*Z-6*Z)*A(3)
    +(4*Z*Z-2)*A(4)
    +2*Z *A(5)
    +A(6)
5 6 0 ~ R E T U R N
```

This subroutine is called instead of the SERIES GENERATOR and can be seen to be of a similar accuracy.
An alternative subroutine for EXP X:
The full expansion for EXP X is:

```
550 LET T =(128*Z*Z*Z*Z*Z*Z*Z-224*Z*Z*Z*Z*Z+112*Z*Z*Z-14*Z)*A(1)
    +64*Z*Z*Z*Z*Z*Z-96*Z*Z*Z*Z+36*Z*Z-2)*A(2)
    +(32*\mp@subsup{Z}{}{*}\mp@subsup{Z}{}{*}\mp@subsup{Z}{}{*}\mp@subsup{Z}{}{*}Z-40*\mp@subsup{Z}{}{*}\mp@subsup{Z}{}{*}Z+10*Z)*
    +(16*Z*Z*Z*Z-16*Z*Z+2)*A(4)
    +(8*Z*Z*Z-6*Z)*A(5)
    +(4*Z*Z-2)*A(6)
    +2*Z*A(7)
    +A(8)
5 6 0 \text { RETURN}
```

The expansion for LN X and A TN X, given algebraically, will be:

```
(2048z 11-5632z 9}+5632\mp@subsup{z}{}{7}-2464\mp@subsup{z}{}{5}+440\mp@subsup{z}{}{3}-2\mp@subsup{2}{}{z}) * A (1)
1024z 10-2560z}\mp@subsup{}{8}{+}+2240\mp@subsup{z}{}{6}-800\mp@subsup{z}{}{4}+100\mp@subsup{z}{}{2}-2) * A(2
(512z9-1152z7+864z5-240z3+18z) * A(3)
(256z8-512z6+320z4-64z2+2) * A(4)
(128z7-224z5+112z3-14z) * A(5)
(64z6-96z4+36z2-2) * A(6)
(32z5-40z3+10z) * A(7)
(16z4-16z2+2) * A(8)
(8z3-6z) * A(9)
(4z2-2) * A(10)
(2z) * A(11)
A(12)
```


## THE 'DRAW' ALGORITHM

The following BASIC program illustrates the essential parts of the DRAW operation when being used to produce a straight line. The program in its present form only allows for lines where $\mathrm{X}>\mathrm{Y}$.

```
10 REM DRAW 255,175 PROGRAM
20 REM SET ORIGIN
30 LET PLOTx=0: LET PLOTy=0
4 0 ~ R E M ~ S E T ~ L I M I T S ~
50 LET X=255: LET Y=175
60 REM SET INCREMENT,i
70 LET i=X/2
80 REM ENTER LOOP
90 FOR B=X TO 1 STEP -1
100 LET A=Y+i
110 IF X> A THEN GO TO }16
120 REM UP A PIXEL ON THIS PASS
130 LET A=A-X
140 LET PLOTy=PLOTy+1
150 REM RESET INCREMENT,i
160 LET i=A
170 REM ALWAYS ALONG ONE PIXEL
180 LET PLOTx=PLOTx+1
190 REM NOW MAKE A PLOT
200 PLOT PLOTx,PLOTy
2 1 0 ~ N E X T ~ B ~
```

A complete algorithm is to found in the following program, as a subroutine that will 'DRAW A LINE' from the last position to $\mathrm{X}, \mathrm{Y}$

## THE 'CIRCLE' ALGORITHM

The following BASIC program illustrates how the CIRCLE command produces its circles.
Initially the number of arcs required is calculated. Then a set of parameters is prepared in the 'memory area' and the 'calculator stack'.

The arcs are then drawn by repeated calls to the line drawing subroutine that on each call draws a single line from the 'last position' to the position ' $\mathrm{X}, \mathrm{Y}$ '.

Note: In the ROM program there is a final 'closing' line but this feature has not been included here.

```
10 REM A CIRCLE PROGRAM
20 LET X=127: LET Y=87: LET Z=87
30 REM How many arcs?
4 0 ~ L E T ~ A r c s = 4 * I N T ~ ( I N T ~ ( A B S ~ ( P I * S Q R ~ Z ) + 0 . 5 ) / 4 ) + 4 ~
50 REM Set up memory area; M0-M5
60 LET M0=X+Z
LET M1=0
80 LET M2=2*Z*SIN (PI/Arcs)
90 LET M3=1-2*(SIN (PI/Arcs)) ^ 2
100 LET M4=SIN (2*PI/Arcs)
110 LET M5=2*PI
120 REM Set up stack; Sa-Sd
130 LET Sa=X+Z
140 LET Sb=Y-Z*SIN (Pl/Arcs)
150 LET Sc=Sa
160 LET Sd-Sb
170 REM Initialise COORDS
180 POKE 23677,Sa: POKE 23678,Sb
190 LET M0=Sd
200 REM 'DRAW THE ARCS'
210 LET M0=M0+M2
220 LET Sc=Sc+M1
230 LET X=Sc-PEEK }2367
240 LET Y=MO-PEEK 23678
```

```
2 5 0 \text { GO SUB 510}
260 LET Arcs=Arcs-1: IF Arcs=0 THEN STOP
270 LET MM1=M1
280 LET M1=M1*M3-M2*M4
290 LET M2=MM1*M4+M2*M3
300 GO TO 210
500 REM 'DRAW A LINE' from last position to X,Y
510 LET PLOTx=PEEK 23677: LET PLOTy=PEEK }2367
520 LET dx=SGN X: LET dy=SGN Y
530 LET X=ABS X: LET Y=ABS Y
540 IF X> = Y THEN GO TO 580
50 LET L=X: LET B=Y
500 LET ddx=0: LET ddy=dy
5 7 0 \text { GO TO } 6 1 0
580 IF X+Y=0 THEN STOP
590 LET L=Y: LET B=X
600 LET ddx=dx: LET ddy=0
6 1 0 \text { LET H=B}
6 2 0 ~ L E T ~ i = I N T ~ ( B / 2 ) ~
6 3 0 ~ F O R ~ N = B ~ T O ~ 1 ~ S T E P ~ - 1 ~
6 4 0 ~ L E T ~ i = i + L
6 5 0 ~ I F ~ i ~ < ~ H ~ T H E N ~ G O ~ T O ~ 6 9 0 ~
600 LET i=i-H
670 LET ix=dx: LET iy=dy
6 8 0 \text { GO TO 700}
690 LET ix=ddx: LET iy=ddy
700 LET PLOTy=PLOTy+iy
710 IF PLOTy <0 OR PLOTy > 175 THEN STOP
720 LET PLOTx=PLOTx+ix
730 IF PLOTx <0 OR PLOTx > 255 THEN STOP
7 4 0 ~ P L O T ~ P L O T x , P L O T y ~
7 5 0 ~ N E X T ~ N ~
760 RETURN
```


## NOTE ON SMALL INTEGERS AND -65536.

1. Small integers n are those for which -65535 is less than or equal to n which is less than or equal to 65535 . The form in which they are held is described in 'STACK-BC'. Note that the manual is inaccurate when it says that the third and fourth bytes hold $n$ plus 131072 if $n$ is negative. Since the range of $n$ is then -1 to -65535 , the two bytes can only hold $n$ plus 131072 if it is taken mod 65536; i.e. they hold $n$ plus 65536. The manual is fudging the issue. The fact is that this is not a true twos complement form (as the form n plus 131072 , in other circumstances, could be). Here the same number can stand for two different numbers according to the sign byte: e.g. 0001 stands for 1 if the sign byte is 00 and for -65535 if the sign byte is FF; similarly FF FF stands for 65535 if the sign byte is 00 and for -1 if the sign byte is FF.
2. Accepting that negative numbers are given a special 'twos complement' form, the main feature about this method of holding numbers is that they are ready for 'short addition' without any further twos complementing. They are simply fetched and stored direct by the addition subroutine. But for multiplication they need to be fetched by INT-FETCH and stored afterwards by INT-STORE. These subroutines twos complement the number when fetching or storing it. The calls to INT-STORE are from 'multiply' (after 'short multiplication'), from 'truncate' (after forming a 'small integer' between -65535 and 65535 inclusive), from 'negate'/'abs' for the 'integer case' and from 'sgn' to store 1 or -1. The calls to INT-FETCH are from PRINT-FP to fetch the integer part of the number when it is 'small', from 'multiply' twice to fetch two 'small integers', from 'RE-STACK' to fetch a 'small integer' for re-stacking, from 'negate'/'abs' to fetch a 'small integer' for manipulation and from FP-TO-BC to fetch the integer for transfer to BC.

## The Number -65536.

3. The number -65536 can fit into the 'small integer' format as 00 FF 000000 . It is then the 'limiting number', the one which when twos complemented overflows (cf. 80 hex in a simple one byte or 7 bit system, i.e. - 128 decimal, which when twos complemented still gives 80 hex i.e. -128 decimal since the positive number 128 decimal does not fit into the system).
4. Some awareness of this may have inspired the abortive attempt to create 00 FF 000000 in 'truncate'. It is abortive since it does not even survive the INT routine of which 'truncate' is a part. It just leads to the mistake INT (-65536) equals -1 .
5. But the main error is that this number has been allowed to arise from 'short addition' of two smaller negative integers and then simply put on the stack as 00 FF 000000 . The system cannot cope with this number. The solution proposed in 'addition' is to form the full five byte floating-point form at once; i.e. test for the number first, at about byte 3032, as follows:

| 3032 |  | PUSH | AF | Save the sign byte in A . |
| :---: | :---: | :---: | :---: | :---: |
| 3033 |  | INC | A | Make any FF in A into 00. |
| 3034 |  | OR | E | Test all 3 bytes now for zero. |
| 3035 |  | OR | D |  |
| 3036 |  | JR | NZ,3040,ADD-STORE | Jump if not -65536. |
| 3038 |  | POP | AF | Clear the stack. |
| 3039 |  | LD | (HL), +80 | Enter 80 hex into second byte. |
| 303B |  | DEC | HL | Point to the first byte. |
| 303C |  | LD | (HL), +91 | Enter 91 hex into the first byte. |
| 303E |  | JR | 3049,ADD-RSTOR | Jump to set the pointer and exit. |
| 3040 | ADD-STORE | POP | AF | Restore the sign byte in A. |
| 3041 |  | LD | (HL), A | Store it on the stack. |
| 3042 |  | INC | HL | Point to the next location. |
| 3043 |  | LD | (HL), E | Store the low byte of the result. |
| 3044 |  | INC | HL | Point to the next location. |
| 3045 |  | LD | (HL), D | Store the high byte of the result. |
| 3046 |  | DEC | HL | Move the pointer back to |
| 3047 |  | DEC | HL | address the first byte of the |
| 3048 |  | DEC | HL | result. |
| 3049 | ADD-RSTOR | POP | DE | Restore STKEND to DE. |
| 304A |  | RET |  | Finished. |

6. The above amendment (i.e. 15 extra bytes) with the omission of bytes 3223 to 323 E inclusive from 'truncate' should solve the problems. It would be nice to be able to test this. The calls of INT-STORE should not lead to 00 FF 000000 being stacked. In 'multiply' the number will lead to overflow if it occurs, since 65536 will set the carry flag; so 'long' multiplication will be used. As noted at 30E5, the 5 bytes starting there could probably be omitted if the above amendments were made. 'Negate' avoids stacking 00 FF 000000 by treating zero separately and returning it unaltered. Truncate deals separately with -65536 , as noted above. SGN stores only 1 and -1 .

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## Spectrum ROM Disassembly

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The Sinclair Spectrum is a complex microcomputer whose normal operation is controlled by the 16 K ROM program that is inside every Spectrurh, ? In this book, Dr. Iantogan and Dr. Frank O'Hara examine this program and explain exactly what it is that makes the Spectrum operate in the way that it does, Every routine in the ROM has been disassembled and has full comments os what its function is and how it relates to the other functions in the ROM.
Each aspect of the Spectrum operation is discussed in detail:
The input/output routines:
These cover the keyboard routines, the loudspeaker, the cassette handling routines, and the screen and printer handling routines.
BASIC line and command interpretation:
This part of the ROM considers each BASIC statement as a set of commands. For each command there is a 'command routine', and it is the execution of the machine code in the appropriate 'command routine' that affect's the 'interpretation'.
Expression evaluation:
The Spectrum has a most comprehensive expression evaluator allowing for a wide range of variable types, functions and operators.
The arithmetic routines and the floating point calculator:
This part of the ROM handles all the numbers in a unique five byte floating point form, as well as all the mathematical functions.
Overall, the 16 K ROM program offers an extremely wide range of BASIC commands and functions. This book makes all the functions and entry points available for use in your own programs or for modification into your own special routines.
The COMPLETE SPECTRUM ROM DISASSEMBLY is a must for all serious programmers of the Spectrum.


[^0]:    Save the base address and the number of the line. It is called '176' times. Fetch the line number and the base address. The base address is updated by '256' locations for each line of pixels.

[^1]:    If X is zero, just return with
    zero as the 'last value'.
    Save the pointer to STKEND.
    Store 1 in DE.
    Point to the second byte of $X$.
    Rotate bit 7 into the carry flag.
    Point to the destination again.
    Set C to zero for positive X and
    to FF hex for negative $X$.
    Stack 1 or -1 as required.
    Restore the pointer to
    STKEND'
    Finished.

