MASTERING MACHANE CODE ON YOUR ON YOUR AMASTRAD A64/664/6128 Cive Gifford & Scott Vincent





MASTERING MACHINE CODE ON YOUR AMSTRAD 464/664/6128

CLIVE GIFFORD AND SCOTT VINCENT





MASTERING MACHINE CODE ON YOUR AMSTRAD 464/664/6128

Dedication: To all of our friends at Sheffield and Southampton Universities.

Acknowledgements:

To Sally and Peter for their enormous help with this book; to Andy, Doug and Luke for starting the decorating; to Catherine for her letters; and to summer opening hours for providing our main inspiration.

First published in the UK by: Interface Publications Ltd., 9-11 Kensington High Street, London W8 5NP

First published in Australia by: Interface Publications (Australia) Pty. Ltd., Chelsea House, 34 Camp Street, Chelsea, Victoria, 3196

Copyright[©] Gifford/Vincent, 1986 First printing March, 1986

ISBN 0 907563 91 0

The programs in this book have been included for their instructional value. They have been tested with care, but are not guaranteed for any particular purpose. Whilst every care has been taken, the publishers cannot be held responsible for any running mistakes which may occur.

ALL RIGHTS RESERVED

No use whatsoever may be made of the contents of this volume – programs and/or text – without the prior written permission of the copyright holder. Reproduction in any form or for any purpose beyond private study by the purchaser of this volume is forbidden.

Books published by Interface Publications are distributed in the UK by WHSD, St John's House, East Street, Leicester LE1 6NE (0533 551196) and in Australia and New Zealand by Pitman Publishing (Melbourne, Sydney, Brisbane and Wellington). Any queries regarding the contents of this volume should be directed by mail to Interface Publications at either the London or Chelsea, Australia, address.

Interface Publications has exclusive rights to the routines in this book, and to their names.

Typeset by Twickenham Printers Cover photograph – Tony Stone Associates

Contents

Pre-Packaged Machine Code Routines:

One – Read a Character
Two – Rotate Left
Three – Rotate Right
Four – Big Print
Five – Massive Print
Six – Screen Fill 100
Seven - Headerless Load/Save 103
Eight – Interrupt-Driven Music 108
Nine – Machine Code Monitor 128
Ten – Box Scrolls 139
Eleven – RSX Chords 169
Twelve – Screen Compactors 187
Thirteen – DEEK and DOKE 195
Fourteen – Games Writing Package 199

Appendices

A	_	Memory Map	215
В	_	Z8Ø Op Codes	216
С	_	Hexadecimal To Decimal Conversion	263

Foreword – Tim Hartnell

Now you have the chance to really master machine code programming on your Amstrad computer. No matter what anybody says, it is not an easy subject to come to grips with.

However, with the right guidance, even the most difficult terrain can be negotiated. I believe that Clive and Scott – two highly competent programmers, with a wealth of software and book writing experience behind them – are the ideal guides to help you understand the intricacies of Amstrad machine code programming.

You should work right through the book, skipping those sections which seem difficult the first time you read it. By the time you've finished the first reading, you should know enough to make sense of the sections whose meaning eluded you the first time through.

If you do this, you'll find the task is made much easier, and your progress won't be impeded because one particular section seems harder than the others.

And don't forget all the 'ready-to-run' machine code routines in this book (including a complete machine code games writing package) which you can use, in conjunction with your BASIC programs, even if you don't understand – at this stage – how they work.

Time now to start mastering machine code on your Amstrad.

Tim Hartnell, London, 1986

Tim Hartnell, head of Interface Publications, is one of the most widely-published computer authors in the world. Recent works include HARNESSING THE POWER OF YOUR ATARI ST, REPLICATING REALITY: EXPLORING COMPUTER SIMULATIONS and EXPLORING EXPERT SYSTEMS ON YOUR MICROCOMPUTER.

Authors' Introduction

We hope that this book helps you master programming in machine code. We've tried to write an easy-to-follow tutorial to give an insight into the workings of machine code, and its main commands, as well as providing a host of useful charts and tables. These take up about half of the book. The balance is a generous selection of machine code routines. These routines include an easy-to-use BASIC loader that can be typed in and run in a few minutes. And even if you don't want to learn the intricacies of machine code right now, this book will give you a library of useful and (in some cases) amazing routines.

It's been hard but enjoyable work writing this book. We hope you get as much value from the book as we did writing it.

Scott Vincent Clive Gifford

London, March 1986

Scott Vincent is a very experienced machine code programmer. Currently studying computer science at Sheffield University, Scott has had several machine code games published by well-known software houses. He has written many books which have been published by Interface and Virgin Books.

Clive Gifford has over twenty books in print and writes frequently for a number of computing magazines. His published titles include 'The Amstrad Programmer's Reference Guide', 'Using Computers In Education' and 'Using a Modem With Your Computer'.

The authors have worked previously on a number of projects together including several of the 'Games For Your...' series of books published by Virgin.

viii

Chapter One What is Machine Code?

Firstly, machine code is not some magical code that instantly creates programs of the same standard as commercial software. It is slow and laborious to write, very difficult to debug and often produces results that are only marginally better than a similar routine written in BASIC. It is not a wonder language and we feel that too much esteem is given to the terms '100% machine code program' and 'machine code programmer'.

To understand how machine code works, we need first to have a brief look into the internal workings of your computer. Your computer's processor consists of thousands of "gates", or switches which are either open (off) or closed (on). Binary numbers can be used to represent these settings easily with a one corresponding to on, and a zero for off. However, entering hundreds of ones and noughts just to get one command working is not an efficient use of time and so "high level" computer languages were created. BASIC is one such language. It replaces the masses of ones and zeros with easy-to-understand, English-like commands. It's not that difficult to ascertain what PRINT, CLEAR and END do. An *interpreter* is needed to translate BASIC commands into binary digits which the computer can understand, and this interpreting slows the speed of operation down greatly.

Machine code is the 'machine-level' series of instructions into which BASIC is converted. These low level instructions can then be acted upon directly by the processor. The 'brains' of your computer, where much of the action occurs, is the Central Processing Unit or CPU. Each design of CPU has its own collection of low-level machine instructions, known as the Instruction Set.

The CPU in your Amstrad is one of the most popular designs found in microcomputers, the Zilog Z8Ø CPU chip. So, naturally enough, this book deals with how to program in Z8Ø machine code and how to use the Z8Ø Instruction Set. (A number of other computers use the Z8Ø chip, the Spectrum and the MSX series being just two machines. But on each machine, there are certain conditions so one finds that while the Z8Ø chip is in a range of home computers, actually using it is a little different with each machine. A good example with the Amstrad is the non-standard way in which the screen is handled.)

The Instruction Set

Each machine code instruction forms part of the low level language. It may take half a dozen or more machine code instructions to emulate a single BASIC command.

There are a number of ways of representing machine code. One such way is Assembly Language which involves allocating a short descriptive name, known as a mnemonic, to each instruction. Another way is to use a decimal number. Binary and hexadecimal (base 16) numbers can also be used (more of that in the next chapter). For example, the instruction which makes the computer return from the machine code level to BASIC has a decimal value of 201, a binary value of 11001001, a hexadecimal value of C9 and a mnemonic of RET.

Now we have a very brief idea of what machine code is, let's look at the reasons for and against learning how to use it.

In Favour

Machine code, when written well, is usually much faster than a comparable BASIC routine.

Machine code allows you to perform actions that are impossible in BASIC, such as speech synthesis. As well machine code often takes far less memory than a similar BASIC program.

As well as this, there is a great deal of prestige in being able to program in machine code. Programs which work perfectly well in BASIC are often chastised in reviews because they are not written in machine code even when writing it in machine code would not produce any improvement whatsoever.

Arguments Against

Machine code often needs many instructions to emulate a simple action.

Machine code tends to be difficult to read, understand and debug. Writing machine code programs is often more difficult and timeconsuming than writing in a high level language such as BASIC.

It is very difficult to perform anything more than simple arithmetic in machine code, and the code can be very difficult to transfer from machine to machine. It is difficult enough if you are transferring a Z8Ø machine code program to another Z8Ø-based machine, but attempting to transfer a Z8Ø routine to, say, a Commodore 64 (which uses a 6502 CPU) is practically impossible.

Chapter Two Number Systems and Assembly Language

The three number systems that are often used when using machine code are decimal, binary and hexadecimal.

As you know, decimal numbers are built on a base of ten, binary numbers on a base of two, and hexadecimal numbers on a base of sixteen.

Imagine that, inside your Amstrad's CPU, are a whole stack of little pigeon holes, called addresses. Each address can hold eight binary digits (and a group of eight binary digits is known, for some arcane reason, as a **byte**). Obviously, if all eight digits are zeros, the number held in the address is zero. If, however, all eight digits are ones (i.e. 11111111), the address holds the number 255. Therefore, an address can hold any number between zero and 255.

The Amstrad command BIN\$ converts decimal numbers into their binary equivalents:

BIN\$(V,N)

In this statement, V is the value in decimal to be altered and N is the number of digits displayed. (The last parameter needn't be used, in which case the Amstrad will print the result without any leading zeros. PRINT BIN\$(4) will return an answer of 100, whereas if you specify that you want all of the eight digits you will get the answer 00000100.)

Hexadecimal (or "hex") is the correct name for the numbering system based on sixteen. The range of digits is \emptyset , 1, 2, 3, 4, 5, 6, 7, 8, 9, A, B, C, D, E and F. Hexadecimal is used in machine code because it is closely related to, and easy to convert to, binary, the computer's natural numbering system. Now, a single hexadecimal digit equals four binary digits (bits) or half a byte (a byte is eight bits). Half a byte is called a nibble! Thus \emptyset in hex equals $\emptyset\emptyset\emptyset\emptyset$ in binary, 2 equals $\emptyset\emptyset1\emptyset$ and F equals 1111. PRINT HEX\$(2 \emptyset 1) will produce the answer C9. This is obtained by 16 times the decimal value of the first hex digit (C=12) plus the decimal value of the second hex digit (9=9), or 192+9 which equals 2 \emptyset 1.

If you wish to include a hex number in a program, it should be prefixed with either "&" or "&H". A Binary number is prefixed with "&X".

Sixteen and Eight-Bit Numbers

So far, we have dealt with numbers in the decimal range of zero to 255. These are, as we said earlier, 8-bit numbers. However, to get numbers greater than 255, we need two byte (16-bit) numbers. These have a far larger range, between \emptyset and 65535. If you had a 16-bit number shown in hex as \emptyset 3C9 then the \emptyset 3 is known as the most significant or high byte while the other byte, C9, is the low or least significant byte. The total value of a 16-bit number is calculated as:

256 times high byte plus low byte

 $256 * \emptyset 3 + C9 = 256 \times 3 + 2\emptyset 1 = 969$

So we have seen how positive numbers between \emptyset and 65535 are obtained, but what about negative numbers? You have a most significant and least significant **bit** within a **byte**. The most significant bit is the **leftmost one**, and the furthest right is the least significant bit. Taking this **most significant bit** (which is normally worth either \emptyset or 128) we use **it** as the **sign** of the number. This effectively cuts down the maximum size of an 8-bit number to 127. However, it keeps the **range** of 256, as now you can have any value between -128 and +127. Setting this bit to zero makes the rest of the byte a **positive** value, while setting it to one will give a negative value. 16-bit numbers work in a similar way with the leftmost bit of the most significant or high byte becoming the **sign bit** as it is called.

However, creating negative binary numbers is unfortunately not quite as simple as that. The rule that adding the corresponding positive number to the negative value should result in the answer zero must be followed. To obtain this, we must use a number concept called **Two's Complement**. This is a two-part process. Firstly, you must swap the value of each bit of the byte. Therefore if a bit was equal to 1 then it should be made equal to Ø and vice versa. This is the actual complementing process. After this has been done, you add 1 to the result and this will equal the opposite-signed binary value of a number.

Hopefully, we can make this a little clearer with an example:

To change +7 to -7: +7 = 00000111 Complement the number = 11111000 Add one to the complement = 11111001 Therefore, 11111001 is the binary equivalent of -7.

(This works equally as well for conversions from minus to plus. You will notice that using this system, the signed bit is set correctly.

16-bit numbers are calculated in the same way and Amstrad's Basic holds all integer variables as 16-bit two's complement numbers. As with 8-bit numbers the range is kept the same but the maximum number is halved to 32767.)

Assembly Language

We hope you are still with us. Unfortunately, we need to go through this before getting to the interesting bits. Stick with it. It will be well worth it.

First, here's another coding system which you need to know about. Assembly language offers an alternative to using decimal and hex numbers to represent machine code routines and instructions. Whereas an instruction is shown as a number in decimal or hex, in assembly language it is shown by a mnemonic, a short name. A list of numbers would be very difficult to follow but a list of mnemonics such as RET (short for RETURN), JP (short for JUMP) and LD (short for LOAD) is far easier to understand.

In the back of the book there is an appendix which gives all of the Z8Ø instructions. These are displayed in both mnemonic and hexadecimal forms. As well, to make it easier to follow what is going on, all of the machine code routines in this book are supplied in both BASIC and assembly language forms. The BASIC listing is the one that you are most likely to actually use as this can be typed straight into your Amstrad and run. The assembly language listing needs a special program (called an **assembler**) to convert the mnemonics into code that can be acted upon by the Amstrad.

Buying an Assembler

There are many assemblers on the market. While they all offer varied additional features, all perform the one essential task of

turning assembly language listings into proper code. The value and importance of an assembler cannot be underestimated. You would find it extremely difficult to write your first machine code programs in decimal or hex, or converting the mnemonics into hex by hand, without an assembler. As you progress in the machine code world, you will find that longer routines are much simpler to write on an assembler. If you pick the right one, you'll also get a host of very useful features to aid writing, debugging and storing your machine code routines.

Despite this, you do not need an assembler to proceed with the material in this book. If you are serious about machine code, though, you will want to consider the purchase of one fairly soon.

If you are committed to machine code as something you wish to learn and use then you might as well buy a full-feature assembler, preferably one that comes with a monitor (a program allowing you to scan through memory) as well. Good ones available include **The Code Machine** by **Picturesque Software** costs a hefty £19.95 at the time of writing but well worth the money. We're not going to go into a long discussion of all the features save to say this is the assembler that we used in the creation of this book. Picturesque can be contacted at 6, Corkscrew Hill, West Wickham, Kent BR4 9BB.

If something a little cheaper is more to your taste then why not have a look at the Amstrad Assembly Language Course supplied by Glentop Publishers of Barnet, Herts. The assembler supplied with the course is pretty good, and as well you get a machine code course, all for $\pounds 12.95$ at the time of writing. Glentop have a freepost address, so it will cost you nothing to get some more information.

If you are an impecunious programmer, you can always turn to 'Computing With The Amstrad' magazine. Issue number 7 (July, 1985) included a reasonable assembler listing. Back issues of the magazines can be obtained from Freepost, Europa House, 68

Chester Road, Hazel Grove, Stockport SK7 5NY. Each back issue costs around $\pounds 1.25$ while a tape containing all the programs in the magazine costs $\pounds 3.75$ or so.

Directives

Moving on from assembler choice, we need to mention the various assembler **directives**. These are commands that you will find in the assembly language listing that are not actual machine code instructions but are commands produced by the assembler to aid the machine code programmer.

These vary slightly from assembler to assembler but usually include the main ones detailed below:

- **ORG** defines the address of the first piece of assembled code.
- END signifies the end of the program.
- EQU assigns a value to a label name.
- DEFL allows you to give a label a value as often as you like.
- **DEFB** assigns a value to a single byte at the current assembly address.
- **DEFW** assigns a value to the next two bytes at the current assembly address (a double-byte version of DEFB).
- **DEFS** creates a number of blank bytes from the current assembly address onwards.

- **DEFM** allows messages to be entered as a string of characters.
- **PRNT** allows printer to be turned on and off while the assembler is running through the listing attempting to assemble the code. The exact nature of this command varies depending on the assembler.
- ENT defines the entry point of the program.

Don't worry if you do not fully understand all of these directives. When you start using an assembler or are looking at the assembly language listings, you can always turn back to these pages.

.

Chapter Three Using Machine Code on the Amstrad

In this chapter, we'll look at the Amstrad's hardware, how to prepare the machine for a machine code program and how to save and load the program.

The Hardware

Let's have a brief look at the various parts that make up an Amstrad computer. These parts can be divided into input, output and processing. Input obviously includes the keyboard, but it can also take in such things as a joystick and any other input device such as a mouse.

Output includes the monitor and any interfaces which allow data to be output to a printer or another external device.

The internal processing parts of the Amstrad are most important for our discussion on machine code. As mentioned in the first chapter, the Amstrad is based around the Z8Ø CPU. The CPU is the central 'thinking' unit of the computer which co-ordinates the workings of all the other parts. Memory is perhaps the most essential of all the other parts.

Memory comes in two basic types, ROM and RAM. ROM is Read

Only Memory and cannot be altered through machine code. Its impregnability is needed as it holds important information – the operating system and the BASIC interpreter. As its name suggests, you can look at the ROM values, which can be very useful as we'll later see.

RAM stands for Random Access Memory and is the type of memory that you use to store your programs in. Not only can it be read, it can also be altered.

The Amstrad consists of 65536 RAM locations and 32768 ROM locations. You may be more familiar with the terms 64K and 32K (a K being a term for 1 \emptyset 24). Each location is capable of holding an 8-bit number which as we learned from binary in the previous chapter, means a number in the decimal range, \emptyset to 255. We also know that this 8-bit memory location is called a byte.

With that out of the way, let's turn to the other parts of the Amstrad.

The Sound

The PSG chip is the Programmable Sound Generator of the Amstrad. It is an AY-3-8912 which is found in many other home computers. The PPI is the Programmable Peripheral Interface and this important chip is the link between the CPU and the PSG, the printer port, the screen display and the cassette and/or disc drive. The CRTC is also concerned with the screen display (its full name is the Cathode Ray Tube Controller).

The final piece of the Amstrad's hardware to take a look at is the Gate Array. This is a nifty device, specific only to the Amstrad, which helps overlap the ROM and RAM as well as having a hand in the generation of the screen display. The overlapping of ROM and RAM will be discussed, along with the memory map, later.

Writing a Machine Code Program

Moving on from the Amstrad's hardware, we'll now look at the procedures involved in writing a machine code program.

Firstly, we must **reserve** an area of **memory** to place the machine code program in. Although your Amstrad has 64K (or more if you have the 6128 machine), not all of this memory is available for your programs. If you type PRINT FRE(\emptyset), you will be told that there is over 42K of memory at your disposal. However, you will only need a fraction of that for most machine code programs.

This spare user memory has an upper boundary above which lies the code for controlling the computer. This upper limit is called HIMEM and is a variable that can be printed and altered. If you switch your machine on and enter PRINT HIMEM, you should get a result of around 43K depending on which Amstrad model you have. HIMEM can be altered to a value below this original upper limit, via the MEMORY command. MEMORY 40000 lowers HIMEM to the address 40000 and gives you several K of RAM in which to place your machine code program or routine. We have done this with all the routines in the latter half of the book. The area above HIMEM is not affected by the command NEW.

Symbol After

There is another point concerning memory reservation that we must cover. The SYMBOL AFTER command reserves memory for using user-defined graphics. This memory is reserved immediately below HIMEM and stays there even if HIMEM is altered. If, for example, HIMEM was first 40000 and then altered to 36000, the memory for the UDG's would be above HIMEM and would not be accessed by the computer. This would cause many problems. If you turn to one of the many routines in the second half of the book, you will see that the first program line usually reads:

1Ø SYMBOL AFTER 256:MEMORY 39999:SYMBOL AFTER 24Ø

This line first gets rid of all memory saved for UDG's, alters HIMEM down to 39999 (meaning the first available address for a machine code routine is 40000, one above that) and then reserves memory for 16 UDG's starting from the new HIMEM position downwards. (For a more detailed explanation of SYMBOL AFTER and user-defined graphics, refer to the Amstrad manual.)

Once you have your machine code program entered into memory save it AT ONCE. Whatever you do, don't RUN the program without having a copy of it on tape or disc. Unlike BASIC programs with their friendly error commands, if a machine code program doesn't work properly, it is more than likely that it will 'crash', leaving you to switch your computer off and on again, thus losing all your hard work. We really cannot stress this point too strongly.

Saving It All

Saving a machine code program is just a matter of adding several extra parameters to the normal SAVE command. The new format is:

SAVE "PROGNAME", B, start address, length, entry point

The start address is the memory address from which program will be saved, while the entry point is the address from which the program actually runs.

The length of the program is the number of bytes that the routine occupies. This can be calculated by the following formula: End Address Of Routine – Start Address Of Routine + 1.

Loading It In

Loading a machine code program is simplicity itself. You just type LOAD "PROGNAME" or even LOAD "". You can re-locate the program by specifying a new start address, e.g. LOAD "SCREEN", 30000. However, you must bear in mind that the program may contain instructions that are dependent upon the original start address (this will become clearer in the following chapters, particularly the chapter concerned with the jump instructions).

Chapter Four Your First Machine Code Program

You now have a pretty good grounding in how to prepare for a machine code routine. As well, you now know what assembly language is. Now has come the time to use this in your first machine code program.

Recently, we read a letter in an Amstrad-specific magazine from a puzzled reader saying that they couldn't get their machine code working. They typed in hex pairs of numbers such as C9, and mnemonics such as RET and LD straight into the machine. Of course, all they got for their trouble was a series of 'syntax error' messages.

We already know that this process would not work as we know that (a) for using mnemonics we need an assembler and (b) code has to be actually POKEd into the memory of your Amstrad.

A Hex Loader

To do this job you need a small program known as a loader. Routines are often presented with an integral loader, such as all the routines in the second half of the book. For very short routines, such as the ones that we will be dealing with over the course of the next few chapters, it is better to use a special hexloading program. These are generally simple affairs (though we have seen some that add so many fancy features that they amount to pages of pages of listing just to perform the same function as a ten-line program).

Type in the following program and save it on disc or tape. You will need it a number of times before this book is out.

```
1 ' HEX LOADER
10 SYMBOL AFTER 256:MEMORY 39999:SYMBOL
AFTER 240
20 n=40000
30 INPUT a$
40 IF LEN(A$)/2<>INT(LEN(A$)/2) THEN PRI
NT "*ERROR* re-enter last line":GOTO 3
0
50 b$=LEFT$(a$,2)
60 POKE n,VAL("&"+b$)
70 n=n+1
80 a$=RIGHT$(a$,LEN(a$)-2)
90 IF a$="" GOTO 30 ELSE 50
```

Let's look at what the hexloader actually does. Line $1\emptyset$ simply sets up the memory reserve (as discussed in an earlier chapter). Line $3\emptyset$ accepts a string of hex digits and checks that an even number of characters have been given, i.e. if the string was nine characters long, then there would obviously be an error in the data entry.

Line 5Ø takes the first two characters of the string and POKEs them into the Amstrad's memory, starting at address 4ØØØØ. Line 8Ø chops off those two POKEd characters from the string, and if there are digits left, goes back to line $5\emptyset$ where the whole process is repeated again. If there are no more digits, the program goes back and asks for another input. When you have finished your data entry, press ESC twice to break the BASIC listing.

The Program

Now let's have our first magnificent, all-powerful machine code program. We will offer it in both its mnemonic form and its hexadecimal form...are you ready for this...

RET

C9

Not much to look at, is it? We have encountered this instruction before. It is the command for RETURNing from a subroutine and, if there are no subroutines, returning from BASIC. Therefore it is admirably suited for our purpose. If it works we will be in BASIC without any crashes and if it doesn't work then we have only lost a minute's typing. Make sure you save the Hexloader to tape before running it.

Run the hexloader, enter the hex pair and break the BASIC program. The routine has now been entered into your Amstrad's memory. To start it, you need to use the BASIC command, CALL. This must be followed by the memory location or address at which the routine starts. In this case, as with most of the routines in this book, it is 40000. CALL 40000 will execute the program and return you to BASIC without any problem. If you do have a crash, switch your machine off, load in the hexloader and check the program listing carefully. After that, run it and make sure that you enter C9, the hex for RET.

CALL is a powerful command on the Amstrad. You will find this is so when you read the chapter a little later on passing parameters to and from machine code programs.

Variables

BASIC programs could not survive without using variables. Machine code has its own variables, known as registers. Registers are very limited in comparison with the freedom of using BASIC variables.

There are a number of specialist registers but for the moment we will only consider the general purpose ones. There are six of them: B, C, D, E, H and L. Each register is rather like a single memory location. It can only hold a number between zero and 255. To overcome this, we pair the registers together getting BC, DE and HL. They are capable of dealing with 16-bit numbers in the range of zero to 65535.

We have to use these registers with the various machine code instructions. The most common of these is LD, which is short for LOAD. This is equivalent to a LET in BASIC. Therefore, LET B=10 or B=10 has the machine code equivalent of LD B, 10 (load B with 10). This form of loading is known as Immediate, Addressing.

You would find it most difficult to write a machine code routine without the LD instruction. It comes in a number of forms. As we have already seen, you can load a register with a numeric constant value. You can also LD a register with another register's value, as in LD B,C. Remember that it is always the first register that is being loaded with the value of the second register.

Accumulator

Before we go any further with the various LD's available, we need to look at the more specialised registers that exist. 'A' is known as the accumulator register. We will look at its features in chapter six dealing with simple arithmetic. There are a number of LDs specifically for the A register. It is the only 8-bit register capable of being loaded directly from a memory location. It can be used in the same way as any of the general purpose registers in the form LD A,C or LD A,19 \emptyset .

Let's look at these special LDs. LD A,(nn) would load A with the contents of the memory address nn. The brackets signify, 'with the contents of'. It is possible to reverse the process and load a specified memory location with the value of the A register. LD (nn),A would do this, again where nn is a memory location.

To demonstrate this a little further, let's look at an example routine. Follow the same procedure for this routine as for the earlier one, using the hex loader.

Assembler	Hexadecimal	
LD A,(41ØØØ)	3A 28 AØ	
LD (41ØØ1),A	32 29 AØ	
RET	С9	

Call 40000 will make this routine take the value of the address 41000 and load it into 41001, thus performing a simple machine code copy. To check to see if this works, type

PRINT PEEK(41000);" ";PEEK(41001)

and the numbers should

be the same.

Notice how easy it would have been to have fallen into a trap at this point. Many of the routines in this book start at address 40000. It would have been very easy to have used 40000 in this one, instead of 41000. However, by doing that, you would be placing a new value into part of the actual routine, with 40001 holding the same instruction as address 40000, LD A.

This should be avoided at all costs in machine code routines. This one isn't important but allowing the routine to tamper with itself when that routine is important or complex is obviously unwise.

Let's look at another couple of registers and how they are used with LD instructions. IX and IY are called **index registers**. They are 16-bit affairs and are particularly used for passing parameters from BASIC to machine code and vice versa, as discussed in the next chapter. LD IX,(addr) loads the IX register with the contents of an address.

F is the flags register, sometimes called the **status register**. From time to time this register cohabitates with A in the hope that no-one will notice it.

The bits of F each have a specific purpose.

They are:

Bit 7: The sign flag, or S.
Bit 6: The zero flag, or Z.
Bit 5: Not used.
Bit 4: The half carry flag, or H.
Bit 3: Not used.
Bit 2: The parity/overflow flag, or PV, or just P.
Bit 1: The subtract flag, or negate flag.
Bit 0: The carry flag, or C. (Not to be confused with register C).

The **half-carry** flag is set by arithmetic instructions if there is a carry from bit three into bit four or, in the case of register pairs, from bit eleven into bit twelve.

The subtract flag is set by any instruction which involves a subtraction, or reset by any instruction which involves an addition.

There are no instructions to directly alter the value of F. However, F can be made to hold a value of, say, xx using LD C,xx/PUSH BC/POP AF. Similarly, the state of the flags H and N may be read using PUSH AF/POP BC and then examining the bits of C.

Don't worry if this seems largely incomprehensible at the moment. It will become clearer in due course (we promise!), and when it does you can come back and reread this section with full understanding.

Chapter Five Passing Parameters

We already know that the BASIC CALL command is used to access a machine code routine. CALL can be followed by more than just the address of the routine to be CALLed. This section deals with using CALL to pass parameters not only from BASIC to the machine code, but also the other way.

There are a number of types of parameters that can be passed to and from machine code. Let's look at the three most common types.

First, there is the straight whole (integer) number which can be either an actual number or an integer variable. The value will be in the range of a 16-bit two's complement number, so 31092, 9, 220 and F% (where F% is equal to 1000) are all permitted.

When the machine code routine is called, the A register immediately holds the number of parameters following the CALL. This may or may not be of use to you. Each parameter is stored using the IX index register (discussed partly in the previous chapter and partly in the dictionary of machine code following shortly). The two byte parameter value is stored with the low byte first and the high byte second.

The register IX holds the address of the low byte of the last parameter. The high byte of the last parameter and the bytes of the other parameters, if there are any, are stored from address IX+1 onwards. Therefore if a CALL 40000, 258 was entered, (IX+0) would hold the low byte which is 2 and (IX+1) would hold the high

byte which is 1. (If you need your memory refreshing: 16-bit numbers are calculated by 256*high byte+low byte, in this case 256*1+2=258.)

Let's say that we wanted to enter a single parameter call at 30000 with the parameter being loaded into a two-byte normal register such as HL, and the parameter was a variable, B%. Your machine code routine would start with LD L, (IX+0) followed by LD H, (IX+1). If it was the third from last parameter of a CALL, for example, then it would become LD L, (IX+4) followed by LD H, (IX+5). In other words, the low byte always comes first and for every parameter, there are two bytes above the IX register.

The second type of parameter is still an integer variable but, unlike the previous type, it can be passed back from machine code to BASIC as well as travelling the other way. To achieve this in the CALL, an '@' sign is added before the variable name. (Naturally, this must be a variable rather than a number.) CALL 40000, @B% would make (IX+0) and (IX+1) hold the address of the value of B%.

Therefore, if you want to alter the value of B% then you first have to find the address where it is stored by examining $(IX+\emptyset)$ and (IX+1). Then you must alter the value of the two memory locations holding the value of B% before returning from machine code to BASIC.

For example, if you wanted to alter the value of B% to 258 then you would load the memory location given by $(IX+\emptyset)+256*(IX+1)$ with 2 (the low byte) and the memory location given by $(IX+\emptyset)+256*(IX+1)+1$ with 1 (the high byte).

Low Byte Change

Let's take a quick look at another small machine code program. This accepts your CALL with just one parameter and changes whatever the parameter's low byte value was, to 6. If you want to use the hex loader program and load in the equivalent hex code to check the routine out please do, but it is by no means compulsory.

Assembler	Hexadecimal	
LD L,(IX+Ø)	DD 6E ØØ	
LD H,(IX+1)	DD 66 Ø1	
LD (HL),6	36 Ø6	
RET	С9	

The first line takes the low byte of the parameter's address and puts it into register L. The second line takes the high byte of the parameter's address and puts it into the register H. With the address of the parameter now given by the HL register pair (by altering the memory pointed to by that register pair) we can alter the value of the variable at will. In the third line, we have changed it to 6. (In a moment I will show you a couple of other things to do with it.) In the fourth line, we have returned the computer to BASIC.

To use the routine type: $B\% = \emptyset$:CALL 40000, @B%. Printing B% after the routine has returned will show B% is valued at 6. Experimentation being the name of the game with machine code, alter the second hex pair on the third line (\emptyset 6 at present) to another value and run the routine again. B% will now be worth that new value. By looking up the instruction op codes at the back of the book, alter the LD (HL),n instruction to LD (HL),A. This will now return the number of parameters available.

Stringing Along

The third and final types of parameter that we are going to deal with at this point are strings. These work quite differently from numeric information. For a start you can only use string variables; no CALL 40000, "AMSTRAD" here. As well, a string variable must be preceded by an '@' character.

The IX register holds an address concerning the parameter, as in the previous numeric cases, but this time it isn't the address of where the parameter resides in memory but the address of another memory location – the start of the **string descriptor block**. The string descriptor block holds information describing the string, in terms of the string's start address and its length.

The first byte of the string descriptor block is the length of the string and is between \emptyset and 255. The second two bytes form the address of where the string starts in memory, and is in the usual form of low byte first followed by high byte. Thus, this address is the address of the first character in the string. You should always define a string variable beforehand, even if it is just as A="".

A String Routine

This tiny routine takes your string and tells you its length. To do this we need an additional parameter, an integer variable to hold the value to be returned.

Hexadecimal	
DD 6E ØØ	
DD 66 Ø1	
7 E	
DD 4E Ø2	
DD 46 Ø3	
Ø2	
С9	

Now execute a CALL $4\emptyset\emptyset\emptyset\emptyset, @V\%, @A$ \$ with both A\$ and V% already defined, to return a value in V% which equals the length of the string A\$.
Chapter Six Simple Arithmetic

We will now perform some simple arithmetic in machine code. To do this, we need some new instructions.

The simplest arithmetic operations in machine code are INC (increment) and DEC (decrement). INC adds one to the specified register which can be an 8-bit one, a 16-bit register pair or an index register. In addition, this instruction can alter the value of a memory location as in INC (HL) – this is known as **indirect addressing**.

INC is a single byte command. INC A would add 1 to the value of the A register, while INC IX would do the same to the index register, IX. If you are incrementing an 8-bit register and the register's value reaches 255, then it will return to zero. If a 16-bit register pair is incremented when its value reaches 65535 then it also returns to zero.

DEC works in much the same way as INC and all the formations of instructions that are valid for INC are also valid for DEC. These are:

INC or DEC r	(where r is an 8-bit register)
INC or DEC rr	(where rr is a 16-bit register
	pair)
INC or DEC IX	
INC or DEC IY	
INC or DEC (HL)	
INC or DEC (IX+d)	
INC or DEC (IY+d)	

If you decrement an 8-bit register when it has a value of zero it will cycle around to 255 – the opposite procedure to INC.

Whereas the 8-bit INC and DEC instructions affect most of the flags, the 16-bit registers (when used with INC or DEC) **do not** affect any of the flags. For the 8-bit instructions the **sign flag** is set (made equal to 1) if bit 7 of the result is set to 1. The **zero flag** is set if the result is zero.

The half carry is set if there is a carry or borrow from bit 4 of the result. The overflow flag is set if bit 7 of the result is altered. The subtract flag is set if the last operation performed was a subtraction instruction, such as DEC. Note that the carry flag is left unaltered.

Adding and Subtracting

ADD and SUB are examples of easy-to-understand mnemonics. It is pretty obvious that they stand for add and subtract.

ADD and SUB only work with one 8-bit register, the A register and with the HL register pair and the two index registers, IX and IY. Apart from their functions, which differ, the syntax and uses are almost the same.

It is a useful idea to show ADDs and SUBs emulating BASIC commands. This gives you a better, clearer idea of how you can use them within routines:

LET A=A+9	ADD A,9
LET A=A+C	ADD A,C
LET A=A-C	SUB A,C
LET A=A-178	SUB A,178

Only the A register can work directly with these commands. ADD B,9 or ADD C,D are not allowed. Therefore, if you wish to perform

arithmetic with the other registers, you have to use a procedure such as this one, which will subtract 8 from the register, E:

LD A,E	(Load A with E)
SUB 8	(Subtract 8 from A)
LD E,A	(Load E with A, the new value of E)

Whenever an addition or subtraction is made and 8-bit registers are used, you need to make sure that the result of the arithmetical operation can be stored within an 8-bit number.

For those operations that cannot, 16-bit addition or subtraction may work. Most 16-bit arithmetic is performed on the HL register pair. This can have other register pairs added to it and, similarly, subtracted from it as shown:

> ADD HL,BC ADD HL,DE ADD HL,HL ADD HL,SP

Note that the second register pair specified is not changed by these operations. The other 16-bit arithmetic operations work on the index registers as shown:

ADD IX,BC ADD IX,DE ADD IX,IX ADD IX,SP

IY is handled the same way. There is no instruction to add a 16-bit number to HL directly. One way of doing this would be to load DE with the desired number and then use ADD HL,DE but this corrupts the DE register pair. The carry flag is also affected by these operations. So, if the result obtained is larger than 65535 the carry flag is set, otherwise it is reset.

Carry

There are two more instructions dealing with addition and subtraction. These work in exactly the same way as ADD and SUB except that they make use of the **carry flag**. These are ADC (add with carry) and SBC (sub with carry). In the case of 8-bit arithmetic they only work on the A register as before. With 16-bit arithmetic they work on HL only, and not on IX or IY. ADC and SBC simply include the carry in the calculation, so the following operations are identical.

LET A=A+5+carry	ADC A,5
LET A=(A-12)-carry	SBC A,12

These instructions are very powerful as they allow you to perform arithmetic operations on, say, two 32-bit numbers. Sometimes you will need to make sure that the carry flag is reset before using them. One way of doing this is by using the instruction AND A which does not alter the value held in A but it does reset the carry flag as a carry can never occur.

No Division

You may well have noticed that there are no division or multiplication instructions. If your program requires many complex mathematical formulae and calculations, then it may be best to write those parts in BASIC. One disadvangage of machine code is its poor ability to handle complex mathematical functions.

However, it is possible to emulate multiplication very easily, using the ADD instruction. The A register will hold the original value at first and then later will hold the answer after the original value has been multiplied.

Multiplication by 2 is very simple: ADD A,A

In fact, any number that comes within the range of an 8-bit number and is a member of the family, such as 2x2=4x2=8x2=16x2=32, can be used as a multiplier very easily.

ADD A,A ADD A,A = multiplication by 4 and ADD A,A ADD A,A ADD A,A = multiplication by 8

You can continue this so that four ADDs are needed to multiply by 16, five for 32 and so on.

Other multipliers are not quite so easy to get but need the use of the LD command and in some cases the SUB command. Even so, they make up for only a few bytes of instructions and could be well worth using in your machine code routines of the future. Here are a couple of examples:

LD B,A ADD A,A ADD A,B	=	x 3
LD B,A ADD A,A ADD A,A ADD A,A SUB A,B	=	x 7

Note the last instruction in the times 7 routine. If you look at the Op Codes table you will see that it is not actually there. What happens is that SUB B is exactly the same as SUB A,B. You can leave the A out of the notation as it is only the A register that can legally be in that position. This applies to all the SUB operations that seem to only have one register.

Chapter Seven Stacking and Jumping

The Stack

There is an area of RAM that is set aside for storing various pieces of information to help the machine know what it's doing. It works as follows:

The word "stack" is something that the computer people have got straight out of a dictionary. It means exactly what it says. Imagine a stack of cardboard boxes. Each box is really a memory location, so each has an address – but if you want to know what's inside any particular cardboard box, the only one you can look at easily is the top one. If you tried to pull one of the boxes from somewhere in the middle then all the boxes above it would fall down. Conversely, to add a new box to the stack the only place you can put it easily is at the top.

The memory locations in the stack are just like that. You can put things on top of it, but only at the top, and you can take things from the top.

There are two special words which go with the stack – one which means "stacking a new number onto the top" and another which means "removing a number from the top"; the first word is PUSH, and the second word is POP. If you PUSH the number five onto the stack, and then you PUSH the number 9ABC, and then you PUSH, say, 8000h, then the first number you can POP is 8000h, since this number is now at the top (because it was PUSHed there last); the second number you can POP is 9ABC, and the third is five.

The stack is stored very high amongst the Amstrad's addresses, so that there is less chance of the BASIC program "colliding" with the stack as either one or the other is built up. The stack is actually very peculiar, because it's **upside down**! It turns out to be more efficient this way. So remember, the stack, or **machine stack** as it's sometimes called, is like a stack of cardboard boxes piled upon a shop floor, except that in a daring feat of defiance of Newton's laws this stack instead decides to reside on the ceiling and build up downwards! The top – the only part you can actually get at – is lower than the bottom!

The stack is so important to the computer that a special register is set aside just to store the position of the **top** of the stack (the part with the lowest address – the part we can get to). That register is called SP, which stands for Stack Pointer. It is actually a register pair because it stores two separate bytes, but unlike the other register pairs (BC, DE and HL) we **cannot** separate the two halves – they really are glued together quite solidly.

Here's how the instructions PUSH and POP work. I'll do this in hex because I think it's easier that way. Suppose HL contained a value of ABCD. This means that H contained AB and L contained CD. Now, the instruction PUSH HL would store the number ABCD at the top of the stack. It would do so by first of all stacking the **high** part (AB), and then stacking the **low** part (CD). It would alter the value of SP accordingly since two more bytes have been added to the stack, and the position of the top will therefore have moved (down) by two addresses.

It is unfortunately not possible to PUSH single registers onto the stack. You may only PUSH register **pairs**, so BC may be PUSHed, but B on its own may not. It is worth noting that PUSH BC will not in any way alter the value of BC, but will simply copy it without changing it. This, of course, goes for all PUSH instructions.

PUSH can be thought of in BASIC as three separate statements:

PUSH HL	SP=SP-2
	POKE SP+1,H
	POKE SP,L

POP, of course, works the other way round. POP HL will first of all remove L from the stack, and will then remove H. SP will be changed, since the top of the stack will have moved.

POP HL	L=PEEK SP
	H=PEEK (SP+1)
	SP=SP+2

Verifying, using the BASIC equivalents given, that PUSH HL followed by POP DE is the same thing as LD D,H followed by LD E,L.

PUSH

Here are the codes for the instruction PUSH. One of them will require a small degree of explanation.

F5	PUSH AF
C5	PUSH BC
D5	PUSH DE
E5	PUSH HL

The register pair AF, which normally cannot be used as such, is made up of the smaller single registers A and F, in the same way that BC is composed of B and C. A is a register which we've been using, but F is something completely different. The F stands for Flags. To understand the workings of F you have to look at it not in hex, but in **binary.** F could, for instance, have a value of 41h, which in binary is $\emptyset 1 \emptyset \emptyset \emptyset \emptyset \emptyset 1$. Each of the digits is either zero or one, and each of the different digits has a different meaning for they are all different flags. One of these flags, the carry flag, is actually stored as the rightmost binary digit of F.

POP

The codes for the POP instructions are very similar to the codes used for PUSH. They are:

F1	POP AF
C1	POP BC
D1	POP DE
E1	POP HL

One of the major uses of PUSH AF and POP AF is simply to PUSH and POP the value of A. The fact that F has been stacked alongside it can be conveniently forgotten about. PUSH AF will certainly stack the value of A until it's needed again, at which point it may be recovered by the use of POP AF. This could be useful if you have to use the A register to perform calculations of some kind that could not be performed in another register, but where the value of A will still be needed later on in the program.

For example, to add twenty-five to the value of B without altering the value of A (or any other register):

F5	PUSH AF
78	LD A,B
C619	ADD A,19h (=25d)
47	LD B,A
F1	POP AF

Why will only B and no other register be altered? (Not even the carry

flag!) See if you can work out precisely what the above routine is doing, before you read on.

Altering SP

We can actually use SP in much the same way that we use BC and DE. We can add and subtract it, and we can load it. The hex codes are:

F9	LD SP,HL
31????	LD SP,nn
ED7B????	LD SP,(nn)
ED73????	LD (nn),SP
39	ADD HL, SP
ED7A	ADC HL,SP
ED72	SBC HL,SP
33	INC SP
3B	DEC SP

This is very powerful, and very useful. Suppose you wanted to exchange the values of D and E without altering anything else. The following routine will do just that.

D5	PUSH DE
D5	PUSH DE
33	INC SP
D1	POP DE
33	INC SP

The final INC SP was necessary in order to restore the stack pointer to its original value. If this is not done you may get a pretty nasty crash.

SP is not the only specialised register in use. There is another two

byte register called PC, standing for Program Counter. Its job is to remember whereabouts we are in the program. Every time the Amstrad has to execute an instruction it will take a look at what PC says. If it says AØØ4 then it will execute the instruction at location AØØ4. It will then increment the value of PC by the number of bytes in that instruction, so that **next** time round it will be looking at the next instruction in sequence. For example, if AØØ4 contained the instruction LD A,B, then this would be carried out and PC would be incremented to AØØ5. If the instruction at AØØ5 was LD B,2 then once this was carried out, PC would be increased by **two**, since LD B,2 is a **two byte** instruction. PC would then be reading AØØ7 where the next instruction begins.

If you alter the value of PC then the effect is like a BASIC GO TO. The only difference is that machine code does not use line numbers, so you have to GO TO the right **address** rather than the right line number. The machine language instruction to do this is called JP which, of course, is short for JumP. JP AØØØ means GO TO address AØØØ and continue executing the machine code from there. Of course, all this instruction really does is to load the number AØØØ into PC (but without incrementing it at the end of the instruction) so that it thinks AØØØ is the next address in the program. It is far more useful for us human beings to think of it as a kind of GO TO though, because that's what we're used to.

Be Careful

Be careful with JP though. If you create an infinite loop in machine code then **tough**! You're stuck with it – you can never break out unless you actually reset the machine or switch it off. An example of an infinite loop would be:

77	AØØØ	LD (HL),A
23	AØØ1	INC HL
C3ØØFØ	AØØ2	ЈР АØØØ

I've written the actual addresses in the middle column. Usually this isn't done and important lines are marked with **labels**, or words which tell us which lines do what. These labels do not appear in the hex, and we only in fact write them out for our own convenience. If, for instance, we decided to call the first line START then our pretty bad program could be written as:

77	START	LD (HL),A
23		INC HL
C3ØØFØ		JP START

There is another instruction (similar to JP) called JR, or Jump Relative. It means jump forward a given number of bytes. In many ways it is better than JP because it is only two bytes long instead of three, and because a whole routine may be relocated without changing JP destinations all over the place. JR \emptyset has no effect whatsoever, and the next instruction will be executed in sequence; however, JR 1 will cause the next instruction (assuming it to be a single byte instruction) to be skipped. To skip over a two byte instruction, or two single byte instructions, you will need to use JR 2.

It is also possible to jump **backwards** using JR, since there is a convention that any hex number between 80 and FF will be taken to mean a negative number (actually 256d less than the number it would normally represent). Note that the number minus five, for example, is represented by FB, and so therefore it is possible to use the instruction JR - 5 - but note that because of this convention we are unable to say JR 129d, for instance, because 129d in hex is 81, which would be taken to mean - 127d and would be a jump backwards. The range we are limited to is, therefore, - 128d to 127d.

Doing Nothing

JR \emptyset , as has been said, does absolutely nothing. It will continue with the next instruction. It is important to remember that relative jumps

are counted from the next instruction. JR \emptyset means execute the next plus zero instruction. JR 1 means execute the next plus one instruction. Consequently, if we were to say JR - 2 then you must count backwards for two bytes starting at zero with the next instruction. You will find that two bytes leads you to precisely the instruction we have just executed – the instruction JR - 2. JR - 2 is therefore an infinite loop all by itself, and is thus not really a recommended instruction to use in a program.

The (rather silly) infinite loop program above can now be rewritten in one byte less using JR instead of JP.

77	START	LD (HL),A
23		INC HL
18FC		JR START

Notice how I wrote "JR START" instead of "JR – 4". This makes the program much easier to follow since all we have to do is look for the label START in order to know where the JR is taking us.

JR and JP are more or less useless on their own without conditions attached, in the same way that BASIC GO TO would be useless if the instruction IF/THEN GO TO were not around. We need some kind of **conditional** jump, so that we can say IF some condition is true THEN jump to a new address. Without this facility JP and JR can only really be used to create infinite loops. Although machine code does not have quite the same degree of flexibility as BASIC, it does allow us to check for four conditions on JR, or eight conditions on JP. These are (for JR):

18ee	JR e	Jump Relative by e
		bytes.
20ee	JR NZ,e	IF the last result calcu-
		lated was non zero
		THEN Jump Relative by
		e bytes.

	28ee	JR Z,e	IF the last result calcu- lated was zero THEN Jump Relative by e
	30ee	JR NC,e	IF CARRY = Ø THEN Jump Relative by e
	38ee	JR C,e	IF CARRY = 1 THEN Jump Relative by e bytes.
And for	JP:		
	С3qqpp С2qqpp	JP pq JP NZ,pq	Jump to address pq. IF the last result calcu- lated was non zero THEN jump to address
	САqqpp	JP Z,pq	IF the last result calcu- lated was zero THEN jump to address pg.
	D2qqpp	JP NC,pq	IF CARRY = \emptyset THEN jump to address pq.
	DAqqpp	JP C,pq	IF CARRY = 1 THEN jump to address pq .
	E2qqpp	JP PO,pq	See below.
	EAggpp	JP PE.pq	See below.
	F2qqpp	JP P,pq	IF the last result calcu- lated was positive (Plus) THEN jump to address
	FAqqpp	JP M,pq	IF the last result calcu- lated was negative (Minus) THEN jump to address pq.

Now although this is a far cry from IF A\$="HELLO" THEN PRINT "GOODBYE" as you're used to, you'll soon see that even this horrendous task may be carried out in machine code. First though, I'll explain about two of the instructions in the list above – JP PO and JP PE. Straightforwardly, JP PO means IF PV=ØTHEN jump to address pq, and JP PE means IF PV=1 THEN jump to address pq. But what exactly is PV?

Answer: PV is another flag – just like CARRY. It can only ever store one of two values – one or zero. The P stands for Parity and the V stands for oVerflow, because computer boffins, like all good mathematicians, can't spell properly. This usage is quite easy to explain:

JP PO can also be thought of as JP NV (Jump if No-oVerflow). JP PE can also be thought of as JP V (Jump if oVerflow).

Technically you shouldn't write JP NV or JP V because it's not standard convention, but it is certainly an aid to memory. I don't really think it matters if you're conventional or not as long as you know what it means (NV=PO and V=PE).

Now, to **overflows.** If we regard numbers from 80 to FF as negative and numbers from 00 to 7F as positive then strange things can happen in arithmetic if we try to cross that boundary. For instance, 41h (positive) plus 41h (positive) equals 82h (negative?). This sort of blunder is called an overflow, so JP V will jump if such an overflow has occured and JP NV will jump if such an overflow has not occurred. In simple terms, an overflow has occurred if the result of any arithmetic operation acting on numbers in positive/negative convention (also called "two's complement" convention) gives an answer with "the wrong sign".

To repeat once more: JPNV is a non-standard way of writing JPPO, and JPV is a way or writing JPPE. So NV=PO and V=PE. You should invent some way of remembering this.

All of these various tests, if combined with other instructions properly, can really check for any situation conceivable. In fact, there's only one other kind of instruction you need in order to make JP and JR as powerful as IF/THEN/GO TO – that instruction is called CP, or ComPare.

CP will compare the register A with any other register or with any numeric constant. For instance, we could have CP B (Compare A with B) or CP L (Compare A with L) or CP 3E (Compare A with the number 3E). What this instruction actually does can be thought of in BASIC as, for CP B, DUMMY = A - B.

In other words, a subtraction is performed, but the **result** of that subtraction is "forgotten", and the value of A never changes. The flags, however, will change and tell us all about the result. If A contains $\emptyset 5$ and B contains $\emptyset 6$ then after a CP B instruction JP NZ will work (since $\emptyset 5 - \emptyset 6$ is non zero), JP C will work (since $\emptyset 5$ is less than $\emptyset 6$), JP NV will work (there is no overflow since the result of $\emptyset 5 - \emptyset 6 = FF$ which is negative and which is **supposed** to be negative) and JP M will work (since FF is negative). Although the subtraction is actually carried out it should be emphasised that the result of the subtraction is discarded, and the value of A does not change.

You can do some useful tricks with CP:

If A=B THEN GO TO ...CP B / JR Z ...IF A<B THEN GO TO ...</td>CP B / JR C ...(This will only work if all numbers are assumed positive)IF A<B THEN GO TO ...</td>CP B / JP M ...

CALLing

Even in machine code we can have **subroutines.** The machine code equivalent of GO SUB is called CALL. We write CALL pq to mean GO SUB the subroutine starting at address pq. The equivalent

RETURN instruction is RET. I hope this looks familiar. RET has a dual purpose – at the end of a subroutine it means "return from that subroutine"; if there is no subroutine to return from, it means "return to BASIC". CALL and RET can both have conditions imposed on them as follows:

C9 RET	CALL pq	CDqqpp
CØ RET NZ	CALL NZ,pq	C4qqpp
C8 RET Z	CALL Z,pq	CCqqpp
DØ RET NC	CALL NC,pq	D4qqpp
D8 RET C	CALL C,pq	DCqqpp
EØ RET PØ	CALL PO,pq	E4qqpp
q" "RET NV"	"CALL NV,pq"	
E8 RET PE	CALL PE,pq	ECqqpp
" "RET V"	"CALL V,pq"	
FØ RET P	CALL P,pq	F4qqpp
F8 RET M	CALL M,pq	FCqqpp
DØ RET NC D8 RET C EØ RET PØ Pq" "RET N E8 RET PE " "RET V FØ RET P F8 RET P F8 RET M	CALL NC,pq CALL C,pq CALL PO,pq "CALL NV,pq" CALL PE,pq "CALL V,pq" CALL P,pq CALL M,pq	D4qqpp DCqqpp E4qqpp ECqqpp F4qqpp FCqqpp

As you'd probably guessed from the above, instructions like RET Z, etc., can also be used to conditionally return to BASIC, ie RET Z equals IF zero THEN RETURN to BASIC.

In BASIC there is no stack, so we never need to worry about it. In machine code, however, there is and so we **do** need to worry about it. There are two instructions other than PUSH and POP which alter the stack! These are CALL and RET!

CALL pq is equivalent to PUSH "the-address-of-the-next-instruction-to-be-executed", followed by JP pq.

RET is equivalent to POP DUMMY followed by "jump-to-address-DUMMY".

You should be able to see how this procedure causes CALL and RET to act as they do. Whilst this is certainly efficient, it does have some drawbacks which we must watch for.

The value of SP must not be altered during the course of a subroutine, since both CALL and RET rely on the stack. You can PUSH as many items as you like onto the stack during a subroutine, just so long as you POP an equal number before attempting to return. A clever trick you may like to know is how to actually **alter** the return address from a subroutine. Let's say you want to change the return address to BØØØ. Watch this:

E1	POP HL
21ØØEØ	LD HL,BØØØ
E5	PUSH HL

The first instruction deletes the original return address from the stack. The second and third instructions replace it by a new alternative return address. When, at any later stage, a RET instruction is reached, control will "return" to address BØØØ. Another useful trick to know is how to make absolutely sure that your subroutines will always exit with the stack "balanced". One way of doing this is to store the value of the stack pointer somewhere and then retrieve it at the end.

Chapter Eight A Dictionary of Machine Code Terms

This chapter should form a handy refresher after the tutorials of the last few chapters. This dictionary contains a short description of all the Z8Ø Op Codes, the registers and a type of command called a directive which is found in assembler listings. These summaries give you the chance to check up on particular commands or features that you may not have fully understood as well as providing a useful reference source for the future when you come across an obscure, infrequently-used command or register.

The dictionary also includes details of the Z8Ø Op Codes that we haven't yet covered. Reading about them here will give you a good understanding of how they work and are used.

Registers

In addition to the registers A,B,C,D,E,H and L, and the sign flags, PV and carry, there are a couple of other registers and flags, although not all of them are useful. Let's cover the registers first.

IX is a register pair; however, it may not be split into its component bytes like HL can. Any instruction in machine code which involves HL (not in brackets) may be written with IX instead of HL. (There are three exceptions to this rule: EX DE, HL, ADC HL, and SBC HL,.) The hex code of such an instruction is DD followed by the hex code of the corresponding instruction using HL. Any instruction in machine code which involves (HL) (with brackets) may be rewritten with (IX + dd) instead of (HL) - the dd represents any byte. For instance, since there is an instruction LD (HL),03 then there is also an instruction LD (IX + 2A),03. The displacement byte can be very useful. (There is an exception to this rule: JP (HL) may only be rewritten as JP (IX) – not JP (IX + dd).) The hex code of such an instruction, but with the displacement byte inserted into the third byte of the hex code. For example, the hex code for LD (HL),03 is 3603; therefore, the hex code of LD (IX + 2A),03 is DD362A03.

IY is also a register pair. It is used in exactly the same way as IX except that the hex codes of IY instructions use the byte FD rather than DD. Despite the fact that both of these register names begin with I they do **not** have the high part in common, and are both independent of each other.

SP is the stack pointer. It's a register pair like BC or DE, however, its two component bytes may not be separated. SP always points to the topmost item on the machine stack. If you change the value of SP then you automatically create a new stack at the given address. Addresses immediately below the current value of SP are likely to be overwritten without prior consent by something called the Z80's interrupt handling routine (see DI).

A is a register you should all know about by now. It's called the **accumulator** from time to time, since it may be used in one or two ways in which the other registers may not (eg ADD A, \emptyset 6).

B,C,D,E,H and L. These are registers which you should by now be totally familiar with.

The Alternate Register Set, featuring A', B', C', D', E', F', H' and L' is of little use except for copying the contents of your normal

registers to for safekeeping. However, Locomotive BASIC uses these registers and this plus Amstrad's warning not to use the alternate register set makes it advisable to leave them well alone.

F register holds all of the flags as discussed earlier.

All the Instructions

By now we've seen a fair number of Z8Ø instructions, so you'll be wanting to expand your vocabulary of these. Here now is a detailed list of every single instruction available to you. I shall cover them in alphabetical order so that you may use this chapter as a kind of dictionary of machine code instructions. For precisely that reason I shall also go over the ones we've already seen. You should, however, re-read them as this will provide a useful memory aid.

ADC Starting with ADC. It comes in two forms: ADC A,r and ADC HL,s. Here we are using r to stand for either A,B,C,D,E,H,L, a numerical constant, or the contents of an address (HL), (IX + d) or (IY + d). ADC A,r is a single byte instruction. It calculates the sum A plus r plus CARRY. The result is stored in A. ADC HL,s is a two byte instruction which evaluates HL plus s plus CARRY, and stores the result in HL. s here stands for either BC, DE, HL,IX or IY. Can you see why (ignoring flags) ADC A,A does precisely the same job as RLA?

ADD Very similar to ADC except that the carry flag is not used in the initial evaluation. It is, however, still altered by the final result. There are two important differences between ADC and ADD though. Firstly, the set of instructions ADD HL,s (where s means as it did in the description of ADC) are one byte instructions rather than two and secondly, it is permissible to use two further sets of instructions ADD IX,s and ADD IY,s. AND Only one form here – AND r. The value of the A register is altered one bit at a time. If such a bit is zero it will remain unaltered, otherwise it will take on the value of the corresponding bit of r. Thus, AND ØØ will always result in zero and AND FF will leave A unchanged. AND alters each of the flags. Specifically, the carry flag will always be reset to zero.

BIT The form of this instruction is BIT n,r where n is a number between zero and seven. The instruction alters the zero flag (only) according to the current value of the bit in question. If the bit is zero then the zero flag will be set, otherwise the zero flag will be reset. You can exploit this using instructions like JR Z (which will jump if the bit was zero) or RET NZ (which will return if the bit was not-zero). BIT does not alter the value of any of the registers, nor does it change the value of the carry flag. It is a two byte instruction. I tend to find it's not used all that often, but that when it is used it comes in very handy indeed.

CALL You've seen this one before – it's rather like GO SUB. Its exact function is as follows: PUSH the return address onto the stack, and then jump to the call address. Since the return address (now on the stack) is used by the instruction RET, it is vitally important that a subroutine should not alter the stack. You may only PUSH things onto the stack in a subroutine if you POP them off again before you attempt to return. CALL may also be used with conditions – for example, CALL Z,pq (pq is an absolute address) which means IF the zero flag is set then CALL pq, otherwise continue with the next instruction.

CCF Complement Carry Flag. If the carry flag was zero then set it to one. If it was one then reset it to zero.

CP In the form CP r, it will calculate the result of subtracting r from A, but will not store the result anywhere. The previous value of A (and, of course, r) remains unaltered. It will, on the other hand, alter

all of the flags, so conditional instructions like JP Z or JP C will still work. CP r followed by JR Z will jump if A equals r (since A minus r is zero) and so on.

CPD Imagine this as CP (HL) followed by DEC HL followed by DEC BC. The PV flag is reset if BC decrements to zero and set otherwise. The zero flag is set if CP part of the instruction found that A was equal to (HL), otherwise it is reset.

CPDR Basically this is the same as CPD except that the instruction is executed over and over again – a kind of automatic loop if you like. CPDR stands for ComPare with Decrement and Repeat. The loop will end in one of two cases: (i) either the comparison found A equal to (HL) (ie if now A = (HL + 1)) and the zero flag is set, or (ii) BC reaches zero, in which case the PV flag is reset.

CPI As CPD, except that HL is incremented instead of decremented.

CPIR As CPDR, except that HL is incremented instead of decremented.

CPL An abbreviation for ComPLement. The register A is altered bit by bit. If any particular bit starts off as zero then it is changed to one and **vice versa.** In other words, if A started off as 11010101(binary) then CPL would change it to 00101010 (binary). This is equivalent to subtracting A from FF. The flags are not affected by this instruction.

DAA Suppose you wanted to add sixteen to twenty-six without converting the numbers to hex. The following seems plausible: LD A,16 then ADD A,26. Unfortunately, because the machine works in hex the final value of A will be 3C not 42. The instruction DAA (Decimal Adjust Accumulator) will then change A from 3C to 42. How it works is rather complicated – it makes a note of what's been carried where and whether you've added or subtracted and so on; but it does always work. For instance, the sequence LD A,42 then SUB Ø6 will again leave A with 3C, but this time round DAA will change A to 36, since forty-two minus six is thirty-six. The instruction changes every flag appropriately.

DEC This is another one of those instructions which comes in two forms. It can be DEC r (a single register) or DEC s (a register pair). DEC r is very simple to understand – the value of the register r is decremented (decreased by one, or changed from $\emptyset\emptyset$ to FF), the carry flag is unaltered and the zero flag is changed as you'd expect it to be. DEC s, however, is a sneaky little instruction, for the zero flag is **not** altered! In fact, none of the flags at all are changed! Thus DEC BC/JR NZ, - 3 is either an infinite loop or has no effect! You must be very careful to remember this – a lot of our earlier programs crashed because we didn't!

DEFB Technically speaking this isn't really a machine code instruction – it's what's called a directive. The word DEFB must be followed by one or more bytes of data, each separated by a comma. Usually this data is in hex, but this isn't always necessary, eg DEFB 3A,C1,45d,11011110b,"f', yellow is valid. The data is inserted into the machine code program at the point it occurs, and in the order it's listed. Data which forms part of a machine code program should never be executed, since the Z80 cannot distinguish between data and program.

DEFM Similar to DEFB except that the data which follows the word DEFM must be a string of characters surrounded by quotes. Commas in the text will themselves be interpreted as data, not as separators. For instance, DEFM"SOLSTICE" will cause the bytes 53 4F 4C 53 54 49 43 45 to be inserted into the program. DEFM stands for DEFine Message, as opposed to DEFB which means DEFine Byte(s).

DEFS Yet another of those directives. This one means DEFine

Space(s). The word DEFS should be followed by one numerical constant. (Just one, mind.) The given number instructs the computer to insert **that number** of zeros. So, DEFS Ø8 will insert eight bytes into the program at that point. DEFS is used mainly to define "variables" in RAM; eg FRED DEFS Ø2 (FRED is a label) and then at some later point in the program LD (FRED),HL.

DEFW One last directive (for now). DEFW stands for DEFine Word. It is used in a very similar way to DEFB except that the items of data are two bytes long, not one, so that (for instance) DEFW 4000 is equivalent to DEFB 00,40. Note how the bytes are switched around. Labels and expressions are permitted in DEFW, so DEFW 70000, FRED, ERIC + 3 is quite valid.

DI This stands for Disable Interrupts and, although this sounds pretty confusing, its use is immensely simple. Fifty times a second a little pulse is sent down one of the pins of the Z8 \emptyset chip. There is a flag called IFF1 which stands for Interrupt Flip Flop One (no, I didn't make that up!), and the effect of DI can be thought of as RES IFF1. When the Z8 \emptyset receives one of these pulses it checks the value of the flag IFF1. If it is set then the computer acts as if a RST 38 instruction (or CALL $\emptyset\emptyset$ 38) has just been reached, with the return address being the next instruction in sequence. If IFF1 is **reset** however, then no action is taken and any machine code program will run as normal. the flag IFF1 must be set **before** any attempt is made to return to BASIC.

DJNZ Yet another abbreviation. This one stands for Decrement B and Jump if Not Zero. So if B is 7, DJNZ will reduce it to 6 and then jump to a new destination. If B is zero, DJNZ will reduce it to FF and again jump to a new destination. If B is one, however, then DJNZ will change it to zero and will **not** jump to a new destination. Instead, it will simply carry on with the next instruction. The form of the instruction is DJNZ e, where e is a single byte. If B is decremented to zero then the e is ignored. If not, then the e specifies how far to jump. The displacement calculated is as in JR.

EI Guess what? Another abbreviation. EI stands for Enable Interrupts and is the opposite of DI. This instruction is equivalent to SET IFF1. See DI for a more complete explanation.

EQU Short for EQUate. This is not a machine code instruction, but a directive. Each EQU directive must have a label attached to it, and the word EQU must be followed by a number (in the range $\emptyset\emptyset\emptyset\emptyset$ to FFFF) or an expression such as ERIC+2. When this directive is reached no action is taken, and no bytes are inserted into the program – therefore, it doesn't really matter whereabouts in the program an EQU appears although it is conventional to place all EQU directives at the very start of a program. What it does is to assign a numerical value (the one given) to the label attached. In other words, if you have ANNIE EQU 9 $\emptyset\emptyset\emptyset$ and then at some later stage LD HL,(ANNIE), then the instruction will be compiled as LD HL,(9 $\emptyset\emptyset\emptyset$).

EX Short for exchange. This instruction will swap the values contained by certain register pairs. There are five EX instructions – these are EX AF, AF'; EX DE, HL; EX (SP), HL; EX (SP), IX; and EX (SP), IY. They don't alter any of the flags, all they do is swap the values over – thus, EX DE, HL replaces DE by the value HL used to contain and HL by the value DE used to contain. The last three are rather interesting – the old value of HL (or IX or IY) is swapped with the topmost item on the stack, so LD BC,Ø123/PUSH BC/LD HL,4567/EX (SP), HL/POP BC will leave BC with 4567 and HL with Ø123. EX (SP), HL does not move the stack pointer and nor, of course, do EX (SP), IX or EX (SP), IY.

EXX You can imagine this as EX BC, B'C' followed by EX DE, D'E' followed by EX HL, H'L'. Basically, each of the common registers (except A) is exchanged with its corresponding alternative register.

HALT When a HALT instruction is reached, control will wait at that point in the program until the next interrupt occurs. When this happens the instruction RST 38 (CALLØØ38) will be executed and on return, control will then continue from the first instruction after HALT. Note that the flag IFF1 must be set if a HALT is to be executed, otherwise an interrupt will never occur. In this case, HALT will literally halt forever. There will be no way of breaking out except by pulling out the plug.

IM DANGER !!! Under no circumstances use this instruction.

IN This has two forms. The first is IN A,(n) where n is a numerical constant. This is equivalent to saying LET A=IN(256*A+n). The second form is IN r,(C) where r is a register. This is equivalent to saying LET r=IN(256*B+C). The argument of IN refers to a hardware device outside the Z80 chip – a different number for each different device. In the form IN A,(n) the flags are not altered; however, the flags are altered by IN r,(C).

INC Don't panic! At long last we're back to sensible instructions we can all understand. INC r increases the value of the register r by one, but without altering the carry flag. INC s increases the value of s by one but alters no flags at all.

IND IN with Decrement. IND can be thought of as IN(HL),(C) followed by DED HL followed by DEC B. The carry flag is unaltered, but the zero flag reflects the new result of B.

INDR As IND, except that the instruction re-executes over and over again, stopping only when B reaches zero.

INI As IND, except that HL is incremented instead of decremented.

INIR As INDR, except that HL is incremented instead of decremented.

JP If you can understand GO TO 10 then you can understand JP 7300. The destination is an address, not a line number, but the principle is exactly the same. JP is the machine language GO TO. We also have conditional jumps, for example JP NZ,7300 means IF non zero THEN jump to address 7300 (In other words if the zero flag is not set.) There is another form of JP which also has an analogy in BASIC – variable destinations. If you understand GO TO N then you'll understand JP (HL). JP (HL) means GO TO HL. In this form you may not have conditions: for instance, JP NC,(HL) is not allowed. Only one of three register pairs may be used as variable destinations – these are HL, IX and IY. Even so these are very powerful instructions – HL can be the result of a calculation, possibly even generated at random.

JR The same as JP only slightly less powerful, but one byte shorter. Only four of the eight conditions may be used: Z, NZ, C and NC. This means it is impossible to use (for instance) JR PO. It is also impossible to say JR (HL). JR does not use an absolute address – the R stands for Relative. You write the instruction as JR e (or JR Z, e or whatever) where the e is a single byte which specifies how far to jump. JR \emptyset has no effect, since it jumps forward by zero bytes. JR FE is an infinite loop because control will jump back to the JR FE instruction itself. The displacement byte starts counting from the instruction immediately after the JR e instruction. If the byte is between $\emptyset\emptyset$ and 7F then the jump is forward, and if the byte is between $\emptyset\emptyset$ and FF then the jump is backwards.

LD The most used instruction in the whole of machine language. All it does is to transfer data from one location to another. It has many, many forms: the simplest being LD r1,r2 – that is to transfer data from one register to another. Other forms are LD A,(BC), LD A,(DE), and LD A,(HL) – and in reverse, LD (BC), A, LD (DE), A and LD (HL), A. Remember that the brackets mean the **contents** of an address. Registers I and R may be loaded, in conjunction with A (but only A) the registers and the register pairs may all be loaded with numerical constants, the register pairs with the contents of any address and inversely any address may be loaded from a register pair (- note that register pairs store two bytes, not one, and these are transferred to and from the address pointed to and the address pointed to plus one). Also allowed are LD A,(pq) and LD (pq),A (where pq represents an address) and SP may be loaded from HL, IX, IY or (pq). ((pq) may be loaded from SP but HL, IX and IY may not.) In other words – there's a lot you can do and a lot you can't. You can't say LD HL,DE, for instance (- you must use LD H,D then LD L,E or vice versa). Fortunately, since LD is used so very, very often, it is extremely easy to become familiar with its many forms.

LDD LoaD with Decrement. Effectively, LD (DE),(HL) followed by DEC HL, DEC DE and DEC BC all in one go. The carry flag and zero flag are unaltered as is the sign flag, however the PV flag is reset to zero if and only if BC decrements to zero. Thus, JP PO will jump only if BC is zero after the instruction.

LDDR As LDD, but the instruction is executed repeatedly until BC reaches zero.

LDI As LDD except that DE and HL are both incremented instead of decremented. BC is still decremented as before.

LDIR As LDI, but the instruction is executed repeatedly until BC reaches zero.

NEG NEGates the accumulator (or A register). It works by performing the subtraction ØØ minus A and all the flags are changed accordingly. Thus, S reflects the sign of the result. Z will be set if and only if A is zero. P will be set if and only if A is 8Ø. C will always be set unless A is zero. NEG is equivalent to CPL followed by INC A (ignoring flags).

NOP This wondrous little beastie (whose name incidentally is short

for No OPeration) has a very simple purpose – to waste time. It has two major uses: (i) as a delay, or (ii) to overwrite previous machine code when debugging or editing. I suppose its nearest BASIC equivalent would be a blank REM statement.

OR In the form OR r this instruction is almost the opposite of AND r. Bit by bit the value of the A register is changed. If any given bit is **one** then it remains unaltered, otherwise it takes its new value from the corresponding bit of r. If A contains ØØ then (ignoring flags) OR r is the same as LD A,r. OR FF is effectively LD A,FF. All of the flags will change as you'd expect them to and the carry flag is reset to zero.

ORG ORG is a directive which must **not** have a label attached. The word ORG must be followed by a number in the range $\emptyset\emptyset\emptyset\emptyset$ to FFFF. It means all machine code from here on is to be written to the address given. Thus, ORG $7\emptyset\emptyset\emptyset$ followed by LD A, \emptyset 1 means that the instruction LD A, \emptyset 1 is to reside at address $7\emptyset\emptyset\emptyset$. Unless the next thing encountered is another ORG directive, then the next instruction will be at address $7\emptyset\emptyset2$ (since LD A, \emptyset 1 is a two byte instruction).

OUT The OUT instruction has two forms. The first is OUT (n), A – this is equivalent to saying OUT 256*A+n, A. The second form is OUT (C), r which is equivalent to OUT 256*B+C, r. OUT sends numbers out of the Z80 chip and into the hardware outside. It has absolutely no effect on the flags.

OUTD OUT with Decrement. Equivalent to OUT (C),(HL) followed by DEC HL followed by DEC B. The carry flag is unchanged, but the zero flag reflects the new value of B.

OTDR A slightly different spelling in no way alters the fact that this is still an OUT with Decrement and Repeat instruction – all it does is lead us to digress from alphabetical order in order to maintain consistency. Equivalent to OUTD repeated over and over again

until B reaches zero.

OUTI As OUTD, except that HL is incremented instead of decremented.

OTIR As OTDR, except that HL is incremented instead of decremented.

POP Removes two bytes of data from the top of the stack and loads them into a register pair. Register pairs BC, DE, HL, IX and IY may be used. In addition, the instruction POP AF may be used forming a pseudo "register pair" from the accumulator and the flags register. Specifically POP will remove the topmost byte from the stack into the low part of the register pair and the next byte into the high part. The stack pointer SP is updated automatically.

PUSH PUSH is the opposite of POP. It stores the contents of any register pair at the top of the stack. The high part is pushed first, then the low part. It "remembers" that a new item has been stacked by updating the value of SP. After a PUSH instruction, SP will always point to the low part of the topmost item on the stack.

RES With this instruction we can actually alter individual bits of any register. RES is short for RESet, which means "change to zero" in computing circles, so RES is the instruction which changes any required bit of a register to zero. For instance, to reset bit 3 of D you just have to say RES 3,D. RES has no effect on any of the flags.

RET RET is used to return from a subroutine. It works by POPping an address from the stack and then jumping to that address. It is possible to alter the address to which a subroutine will return by altering the value at the top of the stack. For example, POP HL/INC HL/PUSH HL will increase the return address by one. You could, for instance, store one byte of data immediately after the CALL instruction, then POP HL/LD A,(HL)/INC HL/PUSH HL will store that byte in A while at the same time ensuring that the subroutine will return to the address **after** that data. Another trick is to push an "artificial" return address onto the stack and then JP (or JR) to a subroutine instead of CALLing it. Now it will "return" to wherever you want it to go! Return may be used with conditions if needed. It does not alter the flags.

RL The form of this instruction is RL r. Each bit of the register specified is moved one position to the left. The leftmost bit is moved into the carry, and the rightmost bit takes on the previous value of the carry. Hence, Rotate Left. For example, if B contained 10010101 and the carry contained zero then RL B would leave B containing 00101010 and the carry containing one. RL alters all of the flags.

RLA Note that there is no space between the L and the A. RLA is a more efficient way of doing RL A! The instruction is one byte shorter and only the carry flag is affected by this instruction.

RLC Rotate Left without Carry. RLC r is almost the same as RL r, in fact, in the sense that each bit of the register in question is moved one position to the left. Here, however, the former leftmost bit becomes both the new value of CARRY and the new rightmost bit. The former value of CARRY does not enter into the process at all. All of the flags are changed.

RLCA In one byte instead of two, RLCA is just RLC A only quicker. Only the carry flag is changed by this operation.

RLD Now for a weird one. RLD is not to be confused with RL D for it is a completely different instruction which works as follows: write the value of A and the contents of address (HL) in hex. The second hex digit of (HL) is shifted left so that it becomes the new first digit. The former value of this first digit overwrites the second digit of A, which in turn becomes the second digit of (HL). Thus, if we start off with A containing 25 and (HL) containing A3 then RLD will change things to A=2A, (HL)=35. RLD incidentally, for some reason known only to the boffins above, is an abbreviation for Rotate Left Decimal.

RR As RL, except that the bits are moved right instead of left.

RRA As RLA, except that the bits are moved right instead of left.

RRC As RLC, except that the bits are moved right instead of left.

RRCA As RLCA, except that the bits are moved right instead of left.

RRD As RLD, except that the hex digits are moved right instead of left.

RST The same as CALL except that the instruction is but one byte long **altogether!** It is much less powerful though for two reasons: (i) you may **not** use conditions (eg RST 10 is legal but RST NZ, 10 is not); and (ii) only one of eight specific addresses may be called. These are 00,08, 10,18,20,28,30 and 38. Since the Amstrad begins executing the ROM from address 00000 onwards, RST 000 is the same as switching your machine on and off.

SBC SBC, like ADC, comes in two form. The first is SBC A,r, which will first of all subtract r from A, and will then subtract the carry. Similarly SBC HL,s will subtract both s and the carry flag from HL. SBC A,A is quite a useful instruction – it leaves the carry unchanged but alters A to ØØ if there is no carry or FF if there is a carry.

SCF Set the Carry Flag. All other flags are unchanged.

SET The opposite of RES. SET 4, H will set bit 4 of register H to one (for example). Any bit of any register may be set.

SLA Shift Left Arithmetic. The form of this instruction is SLA r. It is similar to RL r except that the rightmost bit is always replaced by zero. SLA r will multiply the register r by two.

SRA Shift Right Arithmetic. Any register may be shifted right using the format SRA r. The instruction is similar to RR r except that the leftmost bit remains unchanged. SRA r will divide the register r by two if that register contains a number which is to be regarded as being in two's complement form.

SRL Shift Right Logical. Again, similar to RR except that here the leftmost bit is replaced by zero always. SRL r will divide the register r by two if that register contains a number which is to be regarded as being in absolute value form.

SUB Written as SUB r (but sometimes as SUB A,r just to be confusing). This instruction will subtract r from the register A. Note that unlike with ADD there is no corresponding instruction SUB HL,s. If you wish to subtract s from HL you must first of all reset the carry flag (usually by the use of the instruction AND A) and then use SBC HL,s.

XOR XOR r changes the value of A bit by bit. If any given bit of A is identical to the corresponding bit of r then that bit of A is reset to zero, otherwise that bit of A will be set. XOR alters all the flags and, in particular, the carry flag is always reset. Note that XOR A is the same as LD A, $\emptyset\emptyset$ (ignoring flags) and that XOR FF is the same as CPL (also ignoring flags).

Chapter Nine Logical Operators and Manipulating Bits

In this chapter, we'll cover a fairly large number of instructions which all come under the headings of either being logical operators (AND, OR, XOR), or are instructions for manipulating bits within a byte (SET, RES, BIT, Rotations and Shifts).

Let's first look at the logical operators. They are also known as **Boolean Operations.** Each operation will be described along with its **Truth Table.** This is a table showing the result of the logical operation.

AND's truth table is shown below:

Α	В	A AND B
ø	ø	ø
1	ø	ø
ø	1	ø
1	1	1

There are a number of things that need to be explained after looking at this table. AND works with two values. The first must be the contents of the A register, while the other can be an 8-bit number, n, or any of the single registers (including A) or the value at address IX, IY or HL. We have used register B above as an example.

Because A must always be the first of the two values, the op code for
AND A,C is written as AND C for it is assumed that A comes before the C.

The truth table shows the result of a AND on the bits of two registers. This table applies to all eight bits of the registers. If bit \emptyset of A or bit \emptyset of B is zero then the resultant AND will make bit \emptyset equal zero. If both bits are zero then the resultant AND will be zero while if both bits had a value of 1 then the resultant AND would be 1.

The results from the AND are stored in the A register. The value of the B register (or whatever register is used) is left unchanged. This is very important to remember. Here is a short routine showing AND in action.

3E 66	LD A, 1Ø2	(102 in binary is 01100110)
Ø6 2B	LD B, 43	(43 in binary is ØØ1Ø1Ø11)
AØ	AND B	
32 28 AØ	LD (41000),A	

The routine AND's B and A (43 and 102) together which gives the below result:

Ø11ØØ11Ø ØØ1Ø1Ø11 ØØ1ØØØ1Ø

Only bits 1 and 5 were set (made equal to 1) as a result of the AND. The value of the A register is then placed into memory location 41000. To check that its value is indeed 00100010 (a decimal value of 34), type PRINT PEEK(41000).

OR's truth table is as follows:

Α	В	A OR B
ø	ø	ø
1	ø	1
ø	1	1
1	1	1

OR works with all of the same values as AND, they are:

AND r	(where r is an 8-bit register)
AND n	(where n is an 8-bit number)
AND (HL)	(the contents of memory location HL)
AND (IX+d)	(the contents of memory location IX or
AND (IY+d)	IY + the displacement)

XOR's truth table is below:

Α	В	A XOR B
ø	ø	ø
1	ø	1
ø	1	1
1	1	ø

This logical operator works with the same registers and addressing modes as OR and AND. XOR is short for Exclusive Oring and is a procedure in which a 1 is assigned as the result only if the two bits are different. Thus if both bits are \emptyset , or both bits are 1, then the result to be stored in the A register would be \emptyset .

We have discussed the logical operators. Now let's see them in action. They are capable of changing specific bits easily and the routine below makes use of this. You could call it a "sentence starter" or "capital letter creator". The routine takes the supplied string and changes the first character of that string to an upper case letter. How is this done? The upper case character set starts 32 ASCII codes below the lower case set. The program takes the character code of the first character and using the A register which contains a value of 223 and the instruction AND, effectively cancels the sixth bit, making it equal to zero thus subtracting 32 from the code value. The character code is then returned. Let's look at this in binary.

exact value of
nat it starts
it the vital 6th
ualling 32 in

The AND operation takes any bits that are both equal to 1 and makes the bit result 1. Any bits that do not equal this are reset (made equal to \emptyset). By making the sixth bit of the A register, zero, we are guaranteeing that the sixth bit will not be set and thus 32 is taken from the code value.

Assembler	Decimal
LD L,(IX+Ø)	DD 6E ØØ
LD H,(IX+1)	DD 66 01
INC HL	23
LD C,(HL)	4 E
LD A,223	3E DF
AND C	A1
LD A,C	79
LD (HL),A	77
RET	С9

Setting and Resetting

SET and RES (short for RESET) allow you to alter the value of specific bits within a register's 8-bit value. The format for SET and RES is:

SET/RES n, r

In this, r is any of the 8-bit general purpose registers or the 8-bit value at the location given by the HL, IX or IY registers. N is a number between \emptyset and 7 and this applies to which bit is being set or reset. (Set means being made equal to 1 while reset means making the bit equal to \emptyset .)

LD B,3	means that the first and second bits are
	set and if we add the instruction
SET 2,B	then the third bit is set (worth 4) and
	the B register is now worth in decimal
	terms 7, in binary ØØØØØ111.

A most important point to remember, not only with SET and RES but with many other instructions discussed here, is that the bit parameters always run from \emptyset to 7. Therefore if you are talking about the third bit, this is bit 2 (as is the case in the above example). Remember this, as many books and magazine articles do not make this clear, and confusion is often caused as a result.

If we had the two instructions:

LD C,7 RES Ø,C

then the value in the C register would become 6 because the set first bit (bit \emptyset) has been reset thus taking one off the C register's value.

The instruction BIT can be used to see what value a bit is. You must

specify the bit number, and the register, in exactly the same way as with SET and RES. Its answer is not put into a register but instead into one of the flags managed by the F register. The Z or Zero flag is made equal to 1 if the bit being tested by the BIT instruction is a zero. If the bit being examined is 1 then the Zero flag is made equal to zero (the opposite way round to what you might at first think).

Rotating and Shifting

Let's now move on to rotating and shifting. The principles involved with these operations are not difficult to understand but if explained badly do become very confusing. I have enlisted the help of some diagrams to explain the subject. Let's look at the first diagram on rotating:

```
*Indicates bit to watch
```



Note that the above equals a Rotate Left with Carry (RLC)

Rotate Right (RRC) is the exact opposite.

Rotating is an operation that involves a complete cycle and in this respect is different to shifting, as we will see in a moment. The end bit is taken off the byte and placed on the other end, thus moving all the other seven bits along one bit. Though the rotation shown in the diagram is to the left with all the bits moving up, this needn't always be the case.

The commands involved with this rotating are RLC and RRC. These stand for Rotate Left without Carry and Rotate Right without Carry. You can see from the diagram what this does.

There is another type of rotation and this is known as RL and RR (Rotate Left/Right with carry). What this does, in effect, is take the end bit (the starred one in the diagram) put it into the carry flag and take the old value from the carry flag and put that on the other end of the register. The second diagram shows this:



RL (Rotate Left with Carry)

These rotation instructions work with all 8-bit registers and the values given by IX, IY and HL.

Shifting is a similar process to rotation as the third diagram shows.

(1) Start with

* Ø Ø	1	1	1	ø	ø
-------	---	---	---	---	---

(2) Shift Operation



(3) End Result



Note that this constitutes a Shift Left (known as a Shift Left Arithmetic or SLA)

There are three major shift instructions: SLA, SRA and SRL. SLA stands for Shift Left Arithmetic and is the operation featured in the diagram. SRL stands for Shift Right Logical and works in a very similar way to SLA except it works from the right, hence bit 7 is reset and bit \emptyset (the rightmost bit) is pushed into carry while the rest of the bits are shifted down one bit. Thus a SRL instruction on 11100100 would equal 01110010.

The final shift instruction is Shift Right Arithmetic. This is similar to SRL except it shifts all the bits except bit 7 which is retained as it is. This instruction is used for when bit 7 holds the sign of the value (when used with signed integers between -127 and 127). All shifting instructions work with the same registers as the rotating instructions. What you may not have noticed, but is worth pointing out, is that SLA actually multiplies the number by 2 and SRL divides the number by 2. So to multiply a number by 8 you would use SLA three times.

Chapter Ten Screen and ROM Routines

The screen picture displayed by your Amstrad takes up a hefty 16K (16384 bytes) in all modes. There is nothing you can do about this. (You may wish to access the screen directly in some of your machine code programs rather than using the ROM routines available. This has the advantage of making programs run much faster in most cases. Unfortunately, associating a memory address with a position on the Amstrad's screen may seem an impossible task at first due to the way the screen is stored. As well, one pixel on the screen does not necessarily relate to one bit of a byte depending on which screen mode you are using, although this is the case in mode 2. Hence, there are only two colours available in this mode as a bit can be in one of two states, set or reset.)

We will now look at how the screen is stored, assuming that it is stored from address &CØØØ to address &FFFF. This will be the case unless a machine code program has moved it somewhere else. We will assume that the offset is equal to zero. This is **not** the case if you scroll the screen.

Two Hundred High

The screen is always 200 pixel lines high by 80 bytes across. When you set the screen mode you are telling the computer how to use these 80 bytes; whether you want a large number of pixels across the screen or whether you want a smaller number, using the extra bits to store a larger range of colour values.

The screen is stored as eight 2K blocks with each block containing information on one of the eight lines of every character position. Therefore, the first 8Ø bytes of block one relate to the top pixel line which is line 1 of the first character row, the second 8Ø bytes of block one relate to line 1 of the second character row, and so on.

This means that line 1 of the bottom character row (row 25) is stored in the first block from bytes 1921 (24*80+1) to 2000 (625*80). However, each block is 2K (2048 bytes) long and so the last 48 bytes of every block are not actually used. In the same way, the second block stores information on line 2 of every character row and block eight stores information on line 8 of every character row.

If you are a bit confused, type in the following BASIC program which POKEs every byte from &CØØØ to &FFFF with 255. This means that it is also altering the unused bytes but the computer does not mind you doing this. Note the order in which the lines appear on the screen.

10 MODE 1 20 FOR n=&C000 TO &FFFF 30 POKE n,255 40 NEXT

Change the MODE number in line $1\emptyset$ and run the program again. Notice that it takes the same time to fill the screen in all modes and that different colours appear on the screen for each mode.

There are a large number of ROM routines used by the interpreter for a whole host of tasks such as printing and initialising. (These are all documented in the firmware manual available from Amsoft.) Here are the addresses and names of a few of the more useful routines for the newer machine code programmer. Note all of the addresses are given in hexadecimal.

CALL Address	Description
BBØØ	Resets the keyboard buffer.
BB18	Waits for keypress.
BB24	Gets joystick value, number
	equals which bit is set:
	∮=up 1=down
	2=left 3=right
	4=fire2 5=fire1
	6=unused 7=reset
BB5A	Print character on the screen.
BBEA	Plot pixel with $DE = X$ co-ordinate
	and $HL = Y$ co-ordinate.
BBF6	Draw line from current cursor
	position to X, Y (held in the
	same register pairs as above).
BC38	Set border colours to B,C registers.
BC3E	Set flash periods H,L registers.
BC68	Set the cassette write speed.
BCAA	Add a sound to the sound queue.
BCB6	Halt all sounds.
BCD1	Introduce an RSX to the Firmware.
BD31	Send a character to the
	Centronics port (usually for the printer).
BB39	Set key repeat or not.
BB21	Get Caps Lock and Shift Lock states.

†L

Pre-Packaged Machine Code Routines

One – Read a Character Two – Rotate Left Three – Rotate Right Four – Big Print Five – Massive Print Six – Screen Fill Seven – Headerless Load/Save Eight – Interrupt-Driven Music Nine – Machine Code Monitor Ten – Box Scrolls Eleven – RSX Chords Twelve – Screen Compactors Thirteen – DEEK and DOKE Fourteen – Games Writing Package 9L

One – Read a Character

If many BASIC programmers were asked what one additional command they would like to see on the Amstrad 464, we wouldn't be surprised if the majority asked for a SCREEN\$ command. To the uninitiated, SCREEN\$ and similar commands tell the programmer what character is at a specified screen location. Thus if the programmer wanted to write a routine that involved having to check the centre of the screen to see if an asterisk ('*') was present, he may have a line such as this:

IF SCREEN\$(2Ø,12)=42 THEN...

The two numbers within brackets are the co-ordinates of a cursor position and the 42 is the ASCII code for an asterisk.

Unfortunately, the 464 does not have such a BASIC command and this makes life much more difficult for BASIC programmers wishing to write moving graphics games. The following routine solves this problem and will work in any of the Amstrad's three screen modes.

Type in this program and run it:

1 ' READ A CHARACTER
10 SYMBOL AFTER 256:MEMORY 39999:SYMBOL
AFTER 240
20 FOR n=40000 TO 40019
30 READ a\$:POKE n,VAL("&"+a\$)

40 NEXT 50 DATA DD,6E,02,DD,66,04,CD,75,BB,CD 60 DATA 60,BB,DD,6E,00,DD,66,01,77,C9

Type in the code and run. Then when you wish to check a specified screen position, use the following command:

CHAR%=Ø:CALL 40000, X, Y, @CHAR%

Here, X and Y are the co-ordinates of the position to be checked (in the same format as the BASIC LOCATE command). CHAR% is a specially created variable which, once the CALL has been executed, will contain the ASCII code of the character holding the specified position. If CHAR% is not defined in advance, an Improper Argument error will occur.

The CHAR variable must be an integer and can either be defined with a percentage sign following it, as above, or can be preceded by the command DEFINT C.

This is one of the shortest routines in the book and this is due to the Amstrad's ROM more than anything else. It contains a "read characters" routine which just can't be accessed by BASIC. Our program simply utilises this routine and passes the correct parameters to and from the routine.

This is what the routine looks like when disassembled:

	0001	;	READ	CHAR
	0002	;		
9C40	0010		ORG	40000
BB75	0020	CURSOR	DEFL	0BB75H

BB60		0030	RDCHAR	DEFL	0BB60H
9040	DD6E02	0040		LD	L,(IX+2)
9043	DD6604	0050		LD	H,(IX+4)
9046	CD75BB	0060		CALL	CURSOR
9049	CD60BB	0070		CALL	RDCHAR
9C4C	DD6E00	0080		LD	L,(IX+0)
9C4F	DD6601	0090		LD	H,(IX+1)
9052	77	0100		LD	(HL),A
9053	C9	0110		RET	
		0120		END	

Two – Rotate Left

This routine rotates a string of characters 90 degrees to the left, thus turning them on their side and printing them up the screen.

1 7 ROTATE LEFT 10 SYMBOL AFTER 256:MEMORY 39999:SYMBOL AFTER 240 20 FOR n=40000 TO 40091 30 READ a\$:POKE n.VAL("&"+a\$) 40 NEXT 50 DATA DD,6E,00,DD,66,01,7E,B7,C8,23 60 DATA 5E,23,56,DD,6E,02,DD.66,04.47 70 DATA C5,E5,CD,75,BB,1A,CD,64,9C,13 80 DATA E1,2D,C1,10,F1,C9,D5,CD,A5,BB 90 DATA EB,CD,AE,BB.06.07.23.10.FD.F5 100 DATA CD,06,89,0E,08,1A,13,D5,E5,16 110 DATA 80,1E,80,06,08,82,CB.16.2B.B3 120 DATA CB, 3A, CB, 2B, CB, F3, 10, F3, E1, D1 130 DATA 0D,20,E4,CD,09,B9,F1,CD,5A,BB 140 DATA D1.C9

The format for the command is: CALL 499999, X, Y, @A\$ X and Y are the co-ordinates of where the string is going to start being printed from and A\$ equals the string to be printed which must be preceded by an '@' character. This enables the string to be passed from BASIC to machine code. The A\$ is being used here to mean any string (in the same way that X can, in fact, be any numeric variable or value.) X\$ and pp\$ will work but 'HELLO' won't. In other words, it must be a string variable.

The X and Y co-ordinates are the same as in the BASIC command LOCATE; Y can be any row number from 1 to 25 and X can be any column number between 1 and 20 in mode \emptyset , $4\emptyset$ in mode 1, and $8\emptyset$ in mode 2.

The routine in effect takes the first character in the string, tips it on its side and places that character into the first UDG character available (in the BASIC loader, that is character number 240). It then prints the UDG character. This makes the first available UDG character not that suitable to use when incorporating this routine into a program. To overcome this, if needs be, you can simply define SYMBOL AFTER a little lower.

		0001	;	ROT	LEFT
		0002	;		
9C40		0010		ORG	40000
BB75		0020	CURSOR	DEFL	ØBB75H
9C40	DD6E00	0030		LD	L,(IX+0)
9043	DD6601	<i>0</i> 040		LD	H,(IX+1)
9046	7E	0050		LD	A,(HL)
9047	B7	0060		OR	А
9C48	C8	0070		RET	Z
9049	23	0080		INC	HL
9C4A	5E	0090		LD	E,(HL)

9C4B	23	0100		INC	HL
9C4C	56	0110		LD	D,(HL)
9C4D	DD6E02	0120		LD	L,(IX+2)
9050	DD6604	0130		LD	H,(IX+4)
9053	47	0140		LD	в,А
9054	C5	0150	NXTCHR	PUSH	BC
9C55	E5	0160		PUSH	HL
9056	CD75BB	0170		CALL	CURSOR
9059	1 A	0180		LD	A,(DE)
9C5A	CD649C	0190		CALL	ROTCHR
9C5D	13	0200		INC	DE
9C5E	E1	0210		POP	HL
9C5F	2D	0220		DEC	L
9060	C1	0230		POP	BC
9061	10F1	0240		DJNZ	NXTCHR
9063	C9	0250		RET	
9064	D5	0260	ROTCHR	PUSH	DE
BBA5		0270	MATRIX	DEFL	ØBBA5H
BBAE		0280	TABLE	DEFL	ØBBAEH
B906		0290	ROMENA	DEFL	0B906H
B909		0300	ROMDIS	DEFL	0B909H
BB5A		0310	тхтоит	DEFL	ØBB5AH
9065	CDA5BB	0320		CALL	MATRIX
9068	EB	0330		EX	DE,HL
9069	CDAEBB	0340		CALL	TABLE
9060	0607	0350		LD	в,7
9C6E	23	0360	ADD7	INC	HL

906F	10FD	0370		DJNZ	ADD7
9071	F5	0380		PUSH	AF
9072	CD06B9	0390		CALL	ROMENA
9075	ØEØ8	0400		LD	с,8
9077	1A	0410	NXTBYT	LD	A,(DE)
9078	13	0420		INC	DE
9079	D5	0430		PUSH	DE
9C7A	£5	<i>0</i> 440		PUSH	HL
9C7B	1680	0450		LD	D,128
9C7D	1E80	0460		LD	E,128
9C7F	0608	0470		LD	в,8
9081	82	0480	NXTBIT	ADD	A,D
9082	CB16	0490		RL	(HL)
9084	2B	0500		DEC	HL
9085	B3	0510		OR	E
9086	СВЗА	0520		SRL	D
9088	CB2B	0530		SRA	E
9C8A	CBF3	0540		SET	6,E
9080	10F3	0550		DJNZ	NXTBIT
9C8E	E1	0560		POP	HL
9C8F	D1	0570		POP	DE
9090	ØD	0580		DEC	С
9091	20E4	0590		JR	NZ,NXTBYT
9093	CD09B9	0600		CALL	ROMDIS
9096	F1	0610		POP	AF
9097	CD5ABB	0620		CALL	тхтоит

9C9A	D1	0630	POP	DE
9C9B	C9	0640	RET	
		0650	END	

Three – Rotate Right

This is more or less the opposite of the previous routine. With this one, the string will be printed down the screen, 90 degrees right of normal horizontal printing and 180 degrees right of Rotate Left printing.

```
1 7
               ROTATE RIGHT
10 SYMBOL AFTER 256:MEMORY 39999:SYMBOL
AFTER 240
20 FOR n=40000 TO 40086
30 READ a$:POKE n,VAL("&"+a$)
40 NEXT
50 DATA DD.6E.00.DD.66.01.7E.B7.C8.23
60 DATA 5E,23,56,DD,6E,02,DD,66,04,47
70 DATA C5,E5,CD,75,BB,1A,CD,64,9C,13
80 DATA E1,2C,C1,10,F1,C9,D5,CD,A5,BB
90 DATA EB.CD.AE.BB.F5.CD.06.B9.0E.08
100 DATA 1A,13,D5,E5,16,80,1E,80,06,08
110 DATA 82,CB,1E,23,B3,CB,3A,CB,2B,CB
120 DATA F3,10,F3,E1,D1,0D,20,E4,CD,09
130 DATA B9,F1,CD,5A,BB,D1,C9
```

The format for the call is the same as the previous one:

CALL 49999, X, Y, @A\$

Here X and Y are the co-ordinates for the string's starting position, and A\$ is the string that is to be printed down the screen.

The routines have much in common as one can see from the assembler listing. The UDGs are affected in the same way with Rotate Right as they were with Rotate Left.

		0001	;	ROT	RIGHT
		0002	;		
9040		0010		ORG	40000
BB75		0020	CURSOR	DEFL	ØBB75H
9040	DD6E00	0030		LD	L,(IX+0)
9043	DD6601	<i>0</i> 040		LD	H,(IX+1)
9046	7E	0050		LD	A,(HL)
9047	B7	<i>0060</i>		OR	А
9048	C8	0070		RET	Z
9049	23	0080		INC	HL
9C4A	5E	0090		LD	E,(HL)
9C4B	23	0100		INC	HL
9C4C	56	0110		LD	D,(HL)
9C4D	DD6E02	0120		LD	L,(IX+2)
9050	DD6604	0130		LD	H,(IX+4)
9053	47	0140		LD	В,А
9054	C5	0150	NXTCHR	PUSH	BC
9055	E5	0160		PUSH	HL
9056	CD75BB	0170		CALL	CURSOR
9059	1A	0180		LD	A,(DE)
9C5A	CD649C	0190		CALL	ROTCHR

9C5D	13	0200		INC	DE
9C5E	E1	0210		POP	HL
9C5F	2C	0220		INC	L
9060	C1	0230		POP	BC
9061	10F1	0240		DJNZ	NXTCHR
9063	C9	0250		RET	
9064	D5	0260	ROTCHR	PUSH	DE
BBA5		0270	MATRIX	DEFL	ØBBA5H
BBAE		0280	TABLE	DEFL	ØBBAEH
B906		0290	ROMENA	DEFL	0B906H
B909		0300	ROMDIS	DEFL	0B909H
BB5A		0310	тхтоит	DEFL	ØBB5AH
9065	CDA5BB	0320		CALL	MATRIX
9068	EB	0330		ΕX	DE,HL
9069	CDAEBB	0340		CALL	TABLE
9C6C	F5	0350		PUSH	AF
9C6D	CD06B9	0360		CALL	ROMENA
9070	0E08	0370		LD	с,8
9072	1 A	0380	NXTBYT	LD	A,(DE)
9073	13	0390		INC	DE
9074	D5	0400		PUSH	DE
9075	E5	0410		PUSH	HL
9076	1680	0420		LD	D,128
9078	1E80	0430		LD	E,128
9C7A	0608	0440		LD	в,8
9070	82	0450	NXTBIT	ADD	A,D
9C7D	CB1E	<i>0</i> 460		RR	(HL)

9C7F	23	0470	INC	HL
9080	B3	0480	OR	E
9081	СВЗА	0490	SRL	D
9083	CB2B	0500	SRA	E
9085	CBF3	0510	SET	6,E
9087	10F3	0520	DJNZ	NXTBIT
9089	E1	0530	POP	HL
9C8A	D1	0540	POP	DE
9C8B	ØD	0550	DEC	с
9C8C	20E4	0560	JR	NZ,NXTBYT
9C8E	CD09B9	0570	CALL	ROMDIS
9091	F1	0580	POP	AF
9092	CD5ABB	0590	CALL	TXTOUT
9095	D1	0600	POP	DE
9096	C9	0610	RET	
		0620	END	

Four – Big Print

Double-size characters are now possible with this routine. There are many obvious uses for such a routine, one being for titles of programs or screen displays.

The BASIC loader should be typed in as normal, saved onto tape or cassette and then run.

1 7 BIG PRINT 10 SYMBOL AFTER 256:MEMORY 39999:SYMBOL AFTER 240 20 FOR n=40000 TO 40159 30 READ a\$:POKE n.VAL("&"+a\$) **40 NEXT** 50 DATA CD,93,BB,F5,DD,6E,04,DD,66,05 60 DATA 46,23,5E,23,56,DD,6E,06,DD,66 70 DATA 08,C5,D5,E5,1A,47,CD,06,B9,78 80 DATA CD,9D,9C,47,CD,09,B9,E1,5D,54 90 DATA CD,75,88,DD,7E,02,CD,90,88,78 100 DATA CD,5A,BB,3C,CD,5A,BB,6B,62,2C 110 DATA CD.75,BB.DD.7E.00,CD.90,BB.78 120 DATA 3C, 3C, CD, 5A, BB, 3C, CD, 5A, BB, 6B 130 DATA 62,24,24,D1,13,C1,10,BD,F1,CD 140 DATA 90,BB,C9,CD,A5,BB,EB,CD,AE,BB

150 DATA F5,0E,02,06,04,C5,1A,0F,0F,0F 160 DATA 0F,06,04,1F,CB,1E,CB,2E,10,F9 170 DATA 7E,23,77,06,07,23,10,FD,1A,06 180 DATA 04,1F,CB,1E,CB,2E,10,F9,7E,23 190 DATA 77,06,07,2B,10,FD,13,C1,10,D3 200 DATA 06,08,23,10,FD,0D,20,C9,F1,C9

To use the routine, you need the following CALL command:

CALL 40000, X, Y, @A\$, P1, P2

By now you should be able to work out the first two or three parameters for yourself. In case you don't know, or you've been very naughty and flicked through the book, reading at odd points, then I'll tell you. X and Y are the top left co-ordinates. The double size characters are printed from here. A\$ is the string containing the characters that are to be printed. The particular variables used in the example need not be the ones that you use in your program; any numeric or string variable may be used and for the co-ordinates, ordinary numbers can be utilised.

This is also the case for the two final parameters which I've given the names P1 and P2. They can use either numeric variables or ordinary numbers and they set the colour of the characters. P1 controls the top half of the character and P2 the bottom half. These pen variables can be set to **any** number. This is because, if the parameter is out of the range of the usual values (shown in the table below), then the routine will mask off enough bits to make them fit within the normal pen colour range.

Normal Values:	Mode	Range
	ø	Ø-15
	1	Ø-3
	2	Ø-1

The routine sets the current pen colours back to what they were before the routine was called. This routine utilises the first four UDG characters to create the double-size characters. Therefore make sure that there are at least four UDG characters available. (If you are not using any UDGs then leave the SYMBOL AFTER 24Ø command, found in the BASIC loader, in your program.)

.

		0001	;	BIG	PRINT
		0002	;		
9040		0010		ORG	40000
B906		0020	ROMENA	DEFL	0B906H
B909		0030	ROMDIS	DEFL	0B909H
BBA5		0040	MATRIX	DEFL	ØBBA5H
BBAE		0050	TABLE	DEFL	ØBBAEH
BB75		0060	CURSOR	DEFL	0BB75H
BB93		0070	GETPEN	DEFL	0BB93H
BB90		0080	SETPEN	DEFL	0BB90H
BB5A		0090	TXTOUT	DEFL	0BB5AH
9040	CD93BB	0100		CALL	GETPEN
9043	F5	0110		PUSH	AF
9044	DD6EØ4	0120		LD	L,(IX+4)
9047	DD6605	0130		LD	H,(IX+5)
9C4A	46	0140		LD	B,(HL)
9C4B	23	0150		INC	HL
9040	5E	0160		LD	E,(HL)
9C4D	23	0170		INC	HL
9C4E	56	0180		LD	D,(HL)
9C4F	DD6E06	0190		LD	$L_{1}(IX+6)$

9052	DD6608	0200		LD	H,(IX+8)
9055	C5	0210	NXTCHR	PUSH	BC
9056	D5	0220		PUSH	DE
9057	E5	0230		PUSH	HL
9058	1 A	0240		LD	A,(DE)
9059	47	0250		LD	в,А
9C5A	CD06B9	0260		CALL	ROMENA
9C5D	78	0270		LD	А,В
9C5E	CD9D9C	0280		CALL	BIGCHR
9061	47	0290		LD	в,А
9062	CD09B9	0300		CALL	ROMDIS
9065	E 1	0310		POP	HL
9066	5D	0320		LD	E,L
9067	54	0330		LD	D,H
9068	CD75BB	0340		CALL	CURSOR
9C6B	DD7E02	0350		LD	A,(IX+2)
9C6E	CD90BB	0360		CALL	SETPEN
9071	78	0370		LD	А,В
9072	CD5ABB	0380		CALL	TXTOUT
9075	3C	0390		INC	Α
9076	CD5ABB	0400		CALL	TXTOUT
9079	6B	0410		LD	L,E
9C7A	62	0420		LD	H,D
9C7B	20	0430		INC	L
9070	CD75BB	0440		CALL	CURSOR
9C7F	DD7E00	0450		LD	A,(IX+0)
9082	CD90BB	0460		CALL	SETPEN

9085	78	0470		LD	А,В
9086	3C	0480		INC	Α
9087	3C	0490		INC	Α
9088	CD5ABB	0500		CALL	тхтоит
9C8B	3C	0510		INC	A
9080	CD5ABB	0520		CALL	TXTOUT
9C8F	6B	0530		LD	L,E
9090	62	0540		LD	H,D
9091	24	0550		INC	н
9092	24	0560		INC	Н
9093	D1	0570		POP	DE
9094	13	0580		INC	DE
9095	C1	0590		POP	BC
9096	10BD	0600		DJNZ	NXTCHR
9098	F1	0610		POP	AF
9099	CD90BB	0620		CALL	SETPEN
9090	C9	0630		RET	
9C9D	CDA5BB	0640	BIGCHR	CALL	MATRIX
9CA0	EB	0650		ΕX	DE,HL
9CA1	CDAEBB	0660		CALL	TABLE
9CA4	F5	0670		PUSH	AF
9CA5	0E02	0680		LD	С,2
9CA7	0604	0690	NXTSET	LD	в,4
9CA9	C5	0700	NXTROW	PUSH	BC
9CAA	1 A	0710		LD	A,(DE)
9CAB	0F	0720		RRCA	
9CAC	ØF	0730		RRCA	

9CAD	ØF	0740		RRCA	
9CAE	0F	0750		RRCA	
9CAF	0604	0760		LD	в,4
9CB1	1F	0770	LFTBYT	RRA	
9CB2	CB1E	0780		RR	(HL)
9CB4	CB2E	0790		SRA	(HL)
9CB6	10F9	0800		DJNZ	LFTBYT
9CB8	7E	0810		LD	A,(HL)
9CB9	23	0820		INC	HL
9CBA	77	0830		LD	(HL),A
9CBB	0607	0840	,	LD	в,7
9CBD	23	0850	NXT1	INC	HL
9CBE	10FD	0860		DJNZ	NXT1
9000	1A	0870		LD	A,(DE)
9001	0604	0880		LD	В,4
9003	1F	0890	RGTBYT	RRA	
9004	CB1E	0900		RR	(HL)
9006	CB2E	0910		SRA	(HL)
9008	10F9	0920		DJNZ	RGTBYT
90CA	7E	0930		LD	A,(HL)
9CCB	23	0940		INC	HL
9000	77	0950		LD	(HL),A
9CCD	0607	0960		LD	в,7
9CCF	2B	0970	NXT2	DEC	HL
9CD0	10FD	0980		DJNZ	NXT2
9CD2	13	0990		INC	DE
9CD3	C1	1000		POP	BC

9CD4	10D3	1010		DJNZ	NXTROW
9CD6	0608	1020		LD	в,8
9CD8	23	1030	NXT3	INC	HL
9CD9	10FD	1040		DJNZ	NXT3
9CDB	0D	1050		DEC	С
9CDC	2009	1060		JR	NZ,NXTSET
9CDE	F1	1070		POP	AF
9CDF	C9	1080		RET	
		1090		END	

.

•

Five – Massive Print

Massive Print performs character enlargement but on a far bigger scale than **Big Print**. Each character when reproduced on the screen is 16 times bigger than normal.

If you glance at the two routines, **Big Print** and **Massive Print**, you will see many differences. In fact, the two routines are completely different. **Massive Print** was far easier to write than **Big Print**. The **Massive Print** routine assigns a two by two block of CHR\$(143)'s (solid blocks) for each pixel of a normal character whereas the **Big Print** routine cannot use this time-saving method of character enlargement and must rely on a more complex and longer routine to work.

1	MASSIVE PRINT
10	SYMBOL AFTER 256:MEMORY 39999:SYMBOL
AF	ER 240
20	FOR n=40000 TO 40085
30	READ a\$:POKE n,VAL("&"+a\$)
40	NEXT
50	DATA CD,93,BB,F5,CD,06,B9,DD,7E,04
60	DATA CD,A5,BB,EB,DD,6E,06,DD,66,08
70	DATA 06,08,C5,0E,80,06,02,C5,E5,CD
80	DATA 75,BB,CB,40,28,05,DD,7E,00,18
90	DATA 03,DD,7E,02,CD,90,BB,E1,06,08

100 DATA 1A,A1,28,04,3E,8F,18,02,3E,20 110 DATA CD,5A,BB,CD,5A,BB,CB,39,10,EC 120 DATA 2C,C1,10,D1,13,C1,10,C8,CD,09 130 DATA B9,F1,CD,90,BB,C9

The format for the CALL is:

CALL 40000, X, Y, CHAR, P1, P2

The CALL parameters are very similar to the previous large printing routine with one exception. Whereas **Big Print** will print a string of characters, **Massive Print** with its far larger size asks for the ASCII code of the one character which it will enlarge. If you wish to disply more than one enlarged character on the screen, call the routine once for every character.

The routine will work in all three of the screen modes and uses the two pen colour parameters in the same way as **Big Print.** If, of course, you want single colour characters, make P1 and P2 the same number.

	0001	;	MSVE	PRINT
	0002	;		
9C40	0010		ORG	40000
BB93	0020	GETPEN	DEFL	0BB93H
BB90	0030	SETPEN	DEFL	0BB90H
BBA5	<i>0</i> 040	MATRIX	DEFL	ØBBA5H
BB75	0050	CURSOR	DEFL	0BB75H
BB5A	0060	TXTOUT	DEFL	ØBB5AH

B906		0070	ROMENA	DEFL	0B906H
B909		0080	ROMDIS	DEFL	0B909H
9040	CD93BB	0090		CALL	GETPEN
9043	F5	0100		PUSH	AF
9044	CD06B9	0110		CALL	ROMENA
9047	DD7E04	0120		LD	A,(IX+4)
9C4A	CDA5BB	0130		CALL	MATRIX
9C4D	EB	0140		ЕX	DE,HL
9C4E	DD6E06	0150		LD	L,(IX+6)
9051	DD6608	0160		LD	H,(IX+8)
9054	0608	0170		LD	в,8
9056	C5	0180	NXTROW	PUSH	BC
9057	0E80	0190		LD	C,128
9059	0602	0200		LD	B,2
9C5B	C5	0210	AGAIN	PUSH	BC
9050	E5	0220		PUSH	HL
9C5D	CD75BB	0230		CALL	CURSOR
9060	CB40	0240		BIT	Ø,B
9062	2805	0250		JR	Z,COL2
9064	DD7E00	0260		ĹD	A,(IX+0)
9067	1803	0270		JR	SETCOL
9069	DD7E02	0280	COL2	LD	A,(IX+2)
9060	CD90BB	0290	SETCOL	CALL	SETPEN
906F	E1	0300		POP	HL
9070	0608	0310		LD	B,8
9072	1A	0320	NXTCOL	LD	A,(DE)
9073	A1	0330		AND	С
9074	2804	0340		JR	Z,SPACE
------	--------	------	-------	------	---------
9076	3E8F	0350		LD	A,143
9078	1802	0360		JR	PRINT
9C7A	3E20	0370	SPACE	LD	A,32
9C7C	CD5ABB	0380	PRINT	CALL	TXTOUT
9C7F	CD5ABB	0390		CALL	тхтоот
9082	CB39	0400		SRL	С
9084	10EC	0410		DJNZ	NXTCOL
9086	20	0420		INC	L
9087	C1	0430		POP	BC
9088	10D1	0440		DJNZ	AGAIN
9C8A	13	0450		INC	DE
9C8B	C1	Ø46Ø		POP	BC
9080	10C8	0470		DJNZ	NXTROW
9C8E	CD09B9	0480		CALL	ROMDIS
9091	F1	0490		POP	AF
9092	CD90BB	0500		CALL	SETPEN
9095	C9	0510		RET	
		0520		END	

Six – Screen Fill

This simple routine fills the Amstrad's screen with whatever ASCII characters you require.

1 ' SCREEN FILL 10 SYMBOL AFTER 256:MEMORY 39999:SYMBOL AFTER 240 20 FOR n=40000 TO 40021 30 READ a\$:POKE n,VAL("&"+a\$) 40 NEXT 50 DATA DD,7E,00,01,E8,03,F5,C5,CD,5A 60 DATA BB,C1,0B,78,B1,28,03,F1,18,F2 70 DATA F1,C9

The format for the call is straightforward:

CALL 40000,A

A is simply the ASCII code of the character you wish to print onscreen and A can be a variable or can be replaced by a normal number. Hence CALL 40000,66 will fill the screen with B's. The routine, at present, will only work with the MODE 1 screen.

Here is an ideal opportunity for you to delve into a little alteration work. If you look at the assembler listing of this routine you can see that line 4 \emptyset has the instruction LD BC, 1000. This loading of the register BC with the value 1000 is telling the routine how many screen spaces there are to fill with a character. In mode 1, this is 40 x 25 = 1000. Mode \emptyset is only 500 (20 x 25) and mode 2 twice as many (80 x 25 = 2000). Therefore altering BC to 500 will enable it to print in mode \emptyset just filling one complete screen.

		0001	Ţ	SCRE	EN FILL
		0002	;		
9040		0010		ORG	40000
BB5A		0020	PRNT	DEFL	ØBB5AH
9040	DD7E00	0030		LD	A,(IX+0)
9043	01E803	0040		LD	BC,1000
9046	F5	0050	LOOP	PUSH	AF
9047	C5	0060		PUSH	BC
9048	CD5ABB	0070		CALL	PRNT
9C4B	C1	0080		POP	BC
9C4C	ØB	0090		DEC	BC
9C4D	78	0100		LD	А,В
9C4E	B1	0110		OR	С
9C4F	2803	0120		JR	Z,END
9051	F1	0130		POP	AF
9052	18F2	0140		JR	LOOP
9054	F1	0150	END	POP	AF
9055	C9	0160		RET	
		0170		END	

If you have an assembler, it is just a matter of editing line $4\emptyset$ and altering it to LD BC, $5\emptyset\emptyset$ but if you don't then it's a fraction more

complicated. The three hex pairs to the left of the line number 40 are the equivalent of the op code LD BC, 1000. We wish to alter the last two pairs which make up the number, 10000. Using the HEX\$ command, PRINT HEX\$(5000) to give hexadecimal value, F4 01. Returning to the BASIC listing, we alter the appropriate data (the fifth and sixth numbers) to F4 01 and then re-run the program, after first saving it to tape or disc. The program should now work perfectly in mode 0. From this, it can easily be altered to work in mode 2 or to just print a few lines of a certain character in a certain mode.

Seven – Headerless Load/Save

If you have a 464, and you've watched it loading in a long program, you would have noticed that it loads in a large number of blocks. Each block is proceeded by a header which contains information such as the start address and the length of the data.

This program scraps the headers, making the program load and save in just one big block and cutting down greatly on the saving and loading time needed. When called, the routine does not issue prompt messages.

```
1 ' HEADERLESS
10 SYMBOL AFTER 256:MEMORY 39999:SYMBOL
AFTER 240
20 FOR n=40000 TO 40166
30 READ a$:POKE n,VAL("&"+a$)
40 NEXT
50 DATA FE,03,28,16,21,C0,9C,FE,02,20
60 DATA 5F,AF,DD,5E,00,DD,56,01,DD,6E
70 DATA 02,DD,66,03,18,0F,DD,7E,00,DD
80 DATA 5E,02,DD,56,03,DD,6E,04,DD,66
90 DATA 05,CD,9E,BC,30,36,C9,FE,03,28
100 DATA 16,21,C0,9C,FE,02,20,30,AF,DD
110 DATA 5E,00,DD,56,01,DD,6E,02,DD,66
```

120 DATA 03,18,0F,DD,7E,00,DD,5E,02,DD 130 DATA 56,03,DD,6E,04,DD,66,05,CD,A1 140 DATA BC,D8,B7,28,05,21,D3,9C,18,06 150 DATA CD,03,BB,21,DF,9C,7E,23,CB,7F 160 DATA 20,05,CD,5A,BB,18,F5,CB,BF,CD 170 DATA 5A,BB,3E,07,CD,5A,BB,C9,2A,50 180 DATA 41,52,41,4D,45,54,45,52,20,4D 190 DATA 49,53,53,49,4E,47,AA,2A,52,45 200 DATA 41,44,20,45,52,52,4F,52,AA,2A 210 DATA 45,53,43,41,50,45,AA

The format for this CALL is:

CALL 40000, S, L, SYNC

S is the start address and L the length of the program as there are no headers to carry this information. The SYNC variable is used by your Amstrad to distinguish the header (character 44) from the data (represented by a SYNC value of 22). As we have no headers, the value included is not important, except that the value which is used to SAVE a program, must be used to LOAD it.

	0001	;		HEADERLESS
	0002	;		
9040	0010		ORG	40000
	0020	;	SAVE	ROUTINE
BC9E	0030	SAVE	DEFL	ØBC9EH
BCA1	0040	LOAD	DEFL	ØBCA1H
BB5A	0050	тхтоит	DEFL	ØBB5AH

BB03		0060	KRESET	DEFL	0BB03H
9040	FE03	0070		CP	3
9042	2816	0080		JR	Z,SVSYNC
9044	210090	0090		LD	HL,PARAM
9047	FE02	0100		СР	2
9049	205F	0110		JR	NZ,ERROR
9C4B	AF	0120		XOR	А
9C4C	DD5E00	0130		LD	E,(IX+0)
904F	DD5601	0140		LD	$D,\langle IX+1\rangle$
9052	DD6E02	0150		LD	L,(IX+2)
9055	DD6603	0160		LD	H,(IX+3)
9058	180F	0170		JR	SVESET
9C5A	DD7E00	0180	SVSYNC	LD	A,(IX+0)
9C5D	DD5E02	0190		LD	E,(IX+2)
9060	DD5603	0200		LD	D,(IX+3)
9063	DD6E04	0210		LD	L,(IX+4)
9066	DD6605	0220		LD	H,(IX+5)
9069	CD9EBC	0230	SVESET	CALL	SAVE
9C6C	3036	0240		JR	NC,ESC
9C6E	C9	0250		RET	
		0260	;		
		0270	;	LOAD	ROUTINE
9C6F	FE03	0280		СР	3
9071	2816	0290		JR	Z,LDSYNC
9073	210090	0300		LD	HL, PARAM
9076	FE02	0310		СР	2
9078	2030	0320		JR	NZ,ERROR

.

9C7A	AF	0330		XOR	А
9C7B	DD5E00	0340		LD	E,(IX+0)
9C7E	DD5601	0350		LD	D,(IX+1)
9081	DD6E02	0360		LD	L,(IX+2)
9084	DD6603	0370		LD	H,(IX+3)
9087	130F	0380		JR	LDSET
9089	DD7E00	0390	LDSYNC	LD	A,(IX+0)
9080	DD5E02	0400		LD	$E_{i}(IX+2)$
908F	DD5603	0410		LD	D,(IX+3)
9092	DD6E04	0420		LD	L,(IX+4)
9095	DD6605	0430		LD	H,(IX+5)
9098	CDA1BC	0440	LDSET	CALL	LOAD
9C9B	D8	0450		RET	С
909C	B7	0460		OR	А
9C9D	2805	0470		JR	Z,ESC
9C9F	21D39C	0480		LD	HL, TAPERR
9CA2	1806	0490		JR	ERROR
9CA4	CD03BB	0500	ESC	CALL	KRESET
9CA7	21DF9C	0510		LD	HL,ESCAPE
90AA	7E	0520	ERROR	LD	A,(HL)
9CAB	23	0530		INC	HL
9CAC	CB7F	0540		BIT	7,A
9CAE	2005	0550		JR	NZ, ENDERR
9CB0	CD5ABB	0560		CALL	TXTOUT
9CB3	18F5	0570		JR	ERROR
9CB5	CBBF	0580	ENDERR	RES	7,A

9CB7	CD5ABB	0590		CALL	TXTOUT
9CBA	3E07	0600		LD	Α,7
9CBC	CD5ABB	0610		CALL	TXTOUT
9CBF	C9	0620		RET	
9000		0630	PARAM	DEFM	"*PARAMETER"
90CA		0640		DEFM	" MISSING"
9CD2	AA	0650		DEFB	" * "+80H
9CD3		0660	TAPERR	DEFM	"*READ"
9CD8		0670		DEFM	" ERROR"
9CDE	AA	0680		DEFB	"*"+80H
9CDF		0690	ESCAPE	DEFM	"*ESCAPE"
9CE6	AA	0700		DEFB	" * "+80H
		0710		END	

Eight – Interrupt-Driven Music

Many commercial arcade games allow music to play constantly in the background while the game is played. We wanted to produce a similar effect that you could include within your games. The result is this program.

The program is one of the largest in the book. The storing of data for the tune is a byte-consuming task. Even so, the whole routine (including the tune) takes up only 590 bytes leaving you still with around 43K for your arcade game. The machine code routine needs to be used in conjunction with a small piece of BASIC:

> 10 CALL 40000,0 20 EVERY 6 GOSUB 13000

... Rest Of Game ...

13000 CALL 40000:RETURN

The first CALL needs to be followed by any parameter (we've chosen \emptyset here; it doesn't really matter which value you use). This initialises the routine and sets the music data up ready to play. The second line utilises the BASIC command EVERY for controlling system interrupts. The 6 is approximately 6/50ths of a second. If you are not familiar with the command, EVERY or system interrupts then I suggest you look in the Amstrad manual, chapter 9.3 and chapter 10. EVERY 6/50ths of a second the program will jump to the subroutine at line 130000 (though you can locate this subroutine at any place you wish). The CALL here simply plays the next note of

the tune and then returns to the normal program.

The 6/50ths time can be altered to whatever the user requires, this is just an example. The sound uses only channel 1 and the ENV volume envelope 15. This envelope can be re-defined through BASIC to create some strange effects. The tune can be stopped by simply disabling the system interrupts. Fortunately, we have a BASIC command to take the hard work out of this, it is simply called DI (Disable Interrupts). The opposite of this command is EI (Enable Interrupts).

1 1 INTERRUPT MUSIC 10 SYMBOL AFTER 256:MEMORY 39999:SYMBOL AFTER 240 20 FOR n=40000 TO 40590 30 READ a\$:POKE n.VAL("&"+a\$) 40 NEXT 50 DATA B7,28,14,21,95,9C,22,5D,9C,3E 60 DATA 01,32,84,9C,3E,0F,21,85,9C,CD 70 DATA BC, BC, C9, 21, 84, 9C, 35, C0, 11, 95 80 DATA 9C,1A,B7,C8,13,77,3C,20,09,1A 90 DATA 77,21,95,9C,22,5D,9C,C9,1A,6F 100 DATA 13,1A,67,13,ED,53,5D,9C,22,8F 110 DATA 9C,21,8C,9C,CD,AA,BC,C9,00,02 120 DATA 03,05,01,0F,FF,0A,81,0F,00,00 130 DATA 00,00,00,00,00 140 REM MUSIC DATA 150 DATA 02,AA,01,02,92 160 DATA 01,01,7B,01,02,EF,00,01,7B,01 170 DATA 02,EF,00,01,7B,01,06,EF,00,01

180	DATA	EF,00,01,D5,00,01,C9,00,01,BE
190	DATA	00,01,EF,00,01,D5,00,02,BE,00
200	DATA	01,FD,00,02,D5,00,06,EF,00,02
210	DATA	AA,01,92,92,01,01,7B,01,02,EF
220	DATA	00,01,7B,01,02,5F,00,01,7B,01
230	DATA	06,EF,00,01,1C,01,01,3F,01,01
240	DATA	52,01,01,1C,01,01,EF,00,02,BE
250	DATA	00,01,D5,00,01,EF,00,01,1C,01
260	DATA	06,D5,00,02,AA,01,02,92,01,01
270	DATA	7B,01,02,EF,00,01,7B,01,02,EF
280	DATA	00,01,7B,01,06,EF,00,01,EF,00
290	DATA	01,D5,00,01,C9,00,01,BE,00,01
300	DATA	EF,00,01,D5,00,02,BE,00,01,FD
310	DATA	00,02,D5,00,06,EF,00,01,EF,00
320	DATA	01,D5,00,01,BE,00,01,EF,00,01
330	DATA	D5,00,02,BE,00,01,EF,00,01,D5
340	DATA	00,01,EF,00,01,BE,00,01,EF,00
350	DATA	01,D5,00,02,BE,00,01,EF,00,01
360	DATA	D5,00,01,EF,00,01,BE,00,01,EF
370	DATA	00,01,D5,00,02,BE,00,01,FD,00
380	DATA	02,D5,00,06,EF,00,01,7B,01,01
390	DATA	66,01,01,52,01,02,3F,01,01,1C
400	DATA	01,02,3F,01,01,7B,01,01,66,01
410	DATA	01,52,01,02,3F,01,01,1C,01,02
420	DATA	3F,01,01,BE,00,01,EF,00,01,3F
430	DATA	01,01,1C,01,01,FD,00,01,EF,00

440 DATA 01.D5.00.01.BE.00.01.D5.00.01 450 DATA EF,00,01,D5,00,05,3F,01,01,3F 460 DATA 01.01.7B.01.01.66.01.02.3F.01 470 DATA 01,1C,01,02,3F,01,01,7B,01,01 480 DATA 66.01.01.52.01.02.3F.01.01.1C 490 DATA 01,02,3F,01,01,3F,01,01,1C.01 500 DATA 01,0C,01,01,FD,00,02,FD,00,02 510 DATA FD.00.01.1C.01.01.52.01.01.AA 520 DATA 01,05,3F,01,01,7B,01,01,66,01 530 DATA 01.52.01.02.3F.01.01.1C.01.02 540 DATA 3F,01,01,7B,01,01,66,01,01,52 550 DATA 01.02.3F.01.01.1C.01.02.3F.01 560 DATA 01.BE.00.01.EF.00.01.3F.01.01 570 DATA 1C,01,01,FD,00,01,EF,00,01,D5 580 DATA 00,01,BE,00,01,D5,00,01,EF,00 590 DATA 01,D5,00,05,EF,00,01,3F,01,01 600 DATA 52,01,01,3F,01,02,EF,00,01,1C 610 DATA 01.02.EF.00.01.1C.01.01.EF.00 620 DATA 01,1C,01,01,3F,01,01,EF,00,01 630 DATA BE,00,02,9F,00,01,BE,00,01,EF 640 DATA 00,01,3F,01,02,1C,01,02,EF,00 650 DATA 01,BE,00,03,D5,00,06,EF,00,FF 660 DATA 01

Each note of the tune consists of three bytes of data. The first byte deals with the length of the note:

1 = a semi-quaver
2 = a quaver
3 = a dotted quaver
4 = a crochet
6 = a dotted crochet
8 = a minim
12 = a dotted minim

The pitch of your note can theoretically be between 1 and 4000, though in practice you will probably use a range of between 50 and 1300. This means that we have to use two bytes in which to hold the value. Byte two holds the low byte of the tone and byte three holds the high byte. These correspond to the formula:

Pitch = Byte Two + 256 * Byte Three

The music data starts at line 150 of the BASIC loader. It's simpler to use an assembler to enter the music data, but if you don't have one you can work out the hexadecimal values for each byte, and enter them into the BASIC loader. Keep in mind that the low byte precedes the high byte of the pitch.

Two values for the length of the note have special effects. Zero constitutes the end of the data and will not play any more even if the BASIC EVERY command is still calling the subroutine. The second special value 255,n tells the program to repeat the tune again after a wait of n, where n is comparable to one of the note values (shown earlier in a table).

		0001	;	INT	MUSIC
		0002	;		
9040		0010		ORG	40000
BCBC		0020	AMPENV	DEFL	ØBCBCH
BCAA		0030	ADDSND	DEFL	ØBCAAH
9040	B7	0040		OR	А
9041	2814	0050		JR	Z,NXTSND
9043	219590	0060		LD	HL,MUSIC
9046	225D9C	0070		LD	(DATADR+1),HL
9049	3EØ1	0080		LD	A,1
9C4B	328490	0090		LD	(TIME),A
9C4E	3EØF	0100		LD	A,15
9050	218590	0110		LD	HL,AMPL
9053	CDBCBC	0120		CALL	AMPENV
9056	C9	0130		RET	
9057	218490	0140	NXTSND	LD	HL,TIME
9C5A	35	0150		DEC	(HL)
9C5B	CØ	0160		RET	NZ
9050	119590	0170	DATADR	LD	DE,MUSIC
9C5F	1 A	0180		LD	A,(DE)
9060	B7	0190		OR	A
9061	C8	0200		RET	Z
9062	13	0210		INC	DE
9063	77	0220		LD	(HL),A
9064	3C	0230		INC	Α
9065	2009	0240		JR	NZ,CONT
9067	1A	0250		LD	A,(DE)

9068	77	0260	LD	(HL),A
9069	219590	0270	LD	HL,MUSIC
9060	225D9C	0280	LD	(DATADR+1),HL
906F	C9	0290	RET	
9070	1A	0300 CONT	LD	A,(DE)
9071	6F	0310	LD	L,A
9072	13	0320	INC	DE
9073	1 A	0330	LD	A,(DE)
9074	67	0340	LD	Н,А
9075	13	0350	INC	DE
9076	ED535D9C	0360	LD	(DATADR+1),DE
9C7A	228F9C	0370	LD	(TONE),HL
9C7D	218090	0380	LD	HL, SOUND
9080	CDAABC	0390	CALL	ADDSND
9083	C9	0400	RET	
9084	00	0410 TIME	DEFB	Ø
9085	02	0420 AMPL	DEFB	2,3,5,1
	03 05 01			
9089	0F	0430	DEFB	15,255,10
	FF ØA			
9080	81	0440 SOUND	DEFB	129,15,0
	0F 00			
9C8F	00	0450 TONE	DEFB	0,0,0,0,0,0
	00 00 00	00 00		
9095	02	0460 MUSIC	DEFB	2
9096	AA01	0470	DEFW	426

*

9098	02	0480	DEFB	2
9099	9201	0490	DEFW	402
9C9B	Ø1	0500	DEFB	1
9090	7801	0510	DEFW	379
9C9E	02	0520	DEFB	2
9C9F	EF00	0530	DEFW	239
9CA1	01	0540	DEFB	1
9CA2	7BØ1	0550	DEFW	379
9CA4	02	0560	DEFB	2
9CA5	EF00	0570	DEFW	239
9CA7	Ø1	0580	DEFB	1
9CA8	7BØ1	0590	DEFW	379
9CAA	06	0600	DEFB	6
9CAB	EF00	0610	DEFW	239
9CAD	01	0620	DEFB	1
9CAE	EF00	0630	DEFW	239
9CB0	01	0640	DEFB	1
9CB1	D500	0650	DEFW	213
9CB3	01	0660	DEFB	1
9CB4	C900	0670	DEFW	201
9CB6	Ø1	0680	DEFB	1
9CB7	BE00	0690	DEFW	190
9CB9	Ø1	0700	DEFB	1
9CBA	EF00	0710	DEFW	239
9CBC	Ø1	0720	DEFB	1
9CBD	D500	0730	DEFW	213
9CBF	02	0740	DEFB	2

9000	BE00	0750	DEFW 190
9002	Ø1	0760	DEFB 1
9003	FD00	0770	DEFW 253
9005	02	0780	DEFB 2
9006	D500	0790	DEFW 213
9008	Ø6	0800	DEFB 6
9009	EF00	0810	DEFW 239
9CCB	02	0820	DEFB 2
9000	AAØ1	0830	DEFW 426
9CCE	02	0840	DEFB 2
9CCF	9201	0850	DEFW 402
9CD1	01	0860	DEFB 1
9CD2	7BØ1	0870	DEFW 379
9CD4	02	0880	DEFB 2
9CD5	EF00	0890	DEFW 239
9CD7	01	0900	DEFB 1
9CD8	7801	0910	DEFW 379
9CDA	02	0920	DEFB 2
9CDB	EF00	0930	DEFW 239
9CDD	01	0940	DEFB 1
9CDE	7BØ1	0950	DEFW 379
9CEØ	Ø6	0960	DEFB 6
9CE1	EF00	0970	DEFW 239
9CE3	Ø1	0980	DEFB 1
9CE4	1C01	0990	DEFW 284
9CE6	Ø1	1000	DEFB 1
9CE7	3FØ1	1010	DEFW 319

9CE9	Ø1	1020	DEFB	1
9CEA	5201	1030	DEFW	338
9CEC	Ø1	1040	DEFB	1
9CED	1CØ1	1050	DEFW	284
9CEF	Ø1	1060	DEFB	1
9CF0	EF00	1070	DEFW	239
9CF2	02	1080	DEFB	2
9CF3	BE00	1090	DEFW	190
9CF5	Ø1	1100	DEFB	1
9CF6	D500	1110	DEFW	213
9CF8	Ø1	1120	DEFB	1
9CF9	EF00	1130	DEFW	239
9CFB	01	1140	DEFB	1
9CFC	1CØ1	1150	DEFW	284
9CFE	Ø6	1160	DEFB	6
9CFF	D500	1170	DEFW	213
9DØ1	02	1180	DEFB	2
9D02	AAØ1	1190	DEFW	426
9D04	02	1200	DEFB	2
9D05	9201	1210	DEFW	402
9D07	01	1220	DEFB	1
9D08	7BØ1	1230	DEFW	379
9DØA	02	1240	DEFB	2
9D0B	EF00	1250	DEFW	239
9D0D	01	1260	DEFB	1
9D0E	7801	1270	DEFW	379 _.

9D10	<i>0</i> 2	1280	DEFB	2
9D11	EF00	1290	DEFW	239
9D13	01	1300	DEFB	1
9D14	7801	1310	DEFW	379
9D16	Ø6	1320	DEFB	6
9D17	EF00	1330	DEFW	239
9D19	Ø1	1340	DEFB	1
9D1A	EF00	1350	DEFW	239
9D1C	Ø1	1360	DEFB	1
9D1D	D500	1370	DEFW	213
9D1F	01	1380	DEFB	1
9D20	C900	1390	DEFW	201
9D22	01	1400	DEFB	1
9D23	BE00	1410	DEFW	190
9D25	Ø1	1420	DEFB	1
9D26	EF00	1430	DEFW	239
9D28	01	1440	DEFB	1
9D29	D500	1450	DEFW	213
9D2B	02	1460	DEFB	2
9D2C	BE00	1470	DEFW	190
9D2E	01	1480	DEFB	1
9D2F	FD00	1490	DEFW	253
9D31	02	1500	DEFB	2
9D32	D500	1510	DEFW	213
9D34	06	1520	DEFB	6
9D35	EF00	1530	DEFW	239

9D37	Ø1	1540	DEFB	1
9D38	EF00	1550	DEFW	239
9D3A	01	1560	DEFB	1
9D3B	D500	1570	DEFW	213
9D3D	01	1580	DEFB	1
9D3E	BE00	1590	DEFW	190
9D40	01	1600	DEFB	1
9D41	EF00	1610	DEFW	239
9D43	Ø1	1620	DEFB	1
9D44	D500	1630	DEFW	213
9D46	02	1640	DEFB	2
9D47	BE00	1650	DEFW	190
9D49	01	1660	DEFB	1
9D4A	EF00	1670	DEFW	239
9D4C	Ø1	1680	DEFB	1
9D4D	D500	1690	DEFW	213
9D4F	Ø1	1700	DEFB	1
9D50	EF00	1710	DEFW	239
9D52	Ø1	1720	DEFB	1
9D53	BE00	1730	DEFW	190
9D55	Ø1	1740	DEFB	1
9D56	EF00	1750	DEFW	239
9D58	Ø1	1760	DEFB	1
9D59	D500	1770	DEFW	213
9D5B	02	1780	DEFB	2
9D5C	BE00	1790	DEFW	190
9D5E	Ø1	1800	DEFB	1

•

9D5F	EF00	1810	DEFW 239
9D61	01	1820	DEFB 1
9D62	D500	1830	DEFW 213
9D64	01	1840	DEFB 1
9D65	EF00	1850	DEFW 239
9D67	01	1860	DEFB 1
9 D 68	BE00	1870	DEFW 190
9D6A	01	1880	DEFB 1
9D6B	EF00	1890	DEFW 239
9D6D	Ø1	1900	DEFB 1
9D6E	D500	1910	DEFW 213
9D70	02	1920	DEFB 2
9D71	BE00	1930	DEFW 190
9D73	01	1940	DEFB 1
9D74	FD00	1950	DEFW 253
9D76	02	1960	DEFB 2
9D77	D500	1970	DEFW 213
9D79	06	1980	DEFB 6
9D7A	EF00	1990	DEFW 239
9D7C	01	2000	DEFB 1
9D7D	7B01	2010	DEFW 379
9D7F	01	2020	DEFB 1
9D80	6601	2030	DEFW 358
9D82	01	2040	DEFB 1
9D83	5201	2050	DEFW 338
9D85	02	2060	DEFB 2

9D86	3FØ1	2070	DEFW	319
9D88	Ø1	2080	DEFB	1
9D89	1CØ1	2090	DEFW	284
9D8B	02	2100	DEFB	2
9D8C	3F01	2110	DEFW	319
9D8E	Ø1	2120	DEFB	1
9D8F	7BØ1	2130	DEFW	379
9D91	Ø1	2140	DEFB	1
9 D 92	6601	2150	DEFW	358
9D94	01	2160	DEFB	1
9D95	5201	2170	DEFW	338
9D97	02	2180	DEFB	2
9 D 98	3F01	2190	DEFW	319
9D9A	01	2200	DEFB	1
9D9B	1CØ1	2210	DEFW	284
9D9D	02	2220	DEFB	2
9D9E	3FØ1	2230	DEFW	319
9DA0	Ø1	2240	DEFB	1
9DA1	BE00	2250	DEFW	190
9DA3	01	2260	DEFB	1
9DA4	EF00	2270	DEFW	239
9DA6	01	2280	DEFB	1
9DA7	3F01	2290	DEFW	319
9DA9	01	2300	DEFB	1
9DAA	1CØ1	2310	DEFW	284
9DAC	Ø1	2320	DEFB	1

9DAD	FD00	2330	DEFW	253
9DAF	01	2340	DEFB	1
9DB0	EF00	2350	DEFW	239
9DB2	Ø1	2360	DEFB	ڊ ل
9DB3	D500	2370	DEFW	213
9DB5	Ø1	2380	DEFB	1
9DB6	BE00	2390	DEFW	190
9DB8	Ø1	2400	DEFB	1
9DB9	D500	2410	DEFW	213
9DBB	Ø1	2420	DEFB	1
9DBC	EF00	2430	DEFW	239
9DBE	Ø1	2440	DEFB	1
9DBF	D500	2450	DEFW	213
9DC1	05	2460	DEFB	5
9DC2	3FØ1	2470	DEFW	319
9DC4	Ø1	2480	DEFB	1
9DC5	3FØ1	2490	DEFW	319
9DC7	01	2500	DEFB	1
9DC8	7BØ1	2510	DEFW	379
9DCA	01	2520	DEFB	1
9DCB	6601	2530	DEFW	358
9DCD	02	2540	DEFB	2
9DCE	3FØ1	2550	DEFW	319
9DD0	01	2560	DEFB	1
9DD1	1001	2570	DEFW	284
9DD3	02	2580	DEFB	2
9DD4	3F01	2590	DEFW	319

9DD6	Ø1	2600	DEFB	1
9DD7	7BØ1	2610	DEFW	379
9DD9	01	2620	DEFB	1
9DDA	6601	2630	DEFW	358
9DDC	01	2640	DEFB	1
9DDD	5201	2650	DEFW	338
9DDF	02	2660	DEFB	2
9DEØ	3FØ1	2670	DEFW	319
9DE2	01	2680	DEFB	1
9DE3	1CØ1	2690	DEFW	284
9DE5	02	2700	DEFB	2
9DE6	3FØ1	2710	DEFW	319
9DE8	01	2720	DEFB	1
9DE9	3FØ1	2730	DEFW	319
9DEB	01	2740	DEFB	1
9DEC	1CØ1	2750	DEFW	284
9DEE	01	2760	DEFB	1
9DEF	0C01	2770	DEFW	268
9DF 1	Ø1	2780	DEFB	1
9DF2	FD00	2790	DEFW	253
9DF4	02	2800	DEFB	2
9DF5	FD00	2810	DEFW	253
9DF7	02	2820	DEFB	2
9DF8	FD00	2830	DEFW	253
9DFA	01	2840	DEFB	1
9DFB	1CØ1	2850	DEFW	284
9DFD	01	2860	DEFB	1
9DFE	5201	2870	DEFW	338
9E00	01	2880	DEFB	1

9EØ1	AAØ1	2890	DEFW 426
9E03	05	2900	DEFB 5
9E04	3F01	2910	DEFW 319
9E06	01	2920	DEFB 1
9E07	7BØ1	2930	DEFW 379
9E09	01	2940	DEFB 1
9E0A	6601	2950	DEFW 358
9E0C	01	2960	DEFB 1
9EØD	5201	2970	DEFW 338
9EØF	02	2980	DEFB 2
9E10	3F01	2990	DEFW 319
9E12	01	3000	DEFB 1
9E13	1CØ1	3010	DEFW 284
9E15	02	3020	DEFB 2
9E16	3F01	3030	DEFW 319
9E18	01	3040	DEFB 1
9E19	7B01	3050	DEFW 379
9E1B	01	3060	DEFB 1
9E1C	6601	3070	DEFW 358
9E1E	01	3080	DEFB 1
9E1F.	5201	3090	DEFW 338
9E21	02	3100	DEFB 2
9E22	3F01	3110	DEFW 319
9E24	01	3120	DEFB 1
9E25	1CØ1	3130	DEFW 284
9E27	02	3140	DEFB 2

9E28	3FØ1	3150	DEFW	319
9E2A	01	3160	DEFB	1
9E2B	BE00	3170	DEFW	190
9E2D	01	3180	DEFB	1
9E2E	EF00	3190	DEFW	239
9E30	Ø1	3200	DEFB	1
9E31	3F01	3210	DEFW	319
9E33	01	3220	DEFB	1
9E34	1CØ1	3230	DEFW	284
9E36	Ø1	3240	DEFB	1
9E37	FD00	3250	DEFW	253
9E39	Ø1	3260	DEFB	1
9E3A	EF00	3270	DEFW	239
9E3C	Ø1	3280	DEFB	1
9E3D	D500	3290	DEFW	213
9E3F	Ø1	3300	DEFB	1
9E40	BE00	3310	DEFW	190
9E42	Ø1	3320	DEFB	1
9E43	D500	3330	DEFW	213
9E45	Ø1	3340	DEFB	1
9E46	EF00	3350	DEFW	239
9E48	01	3360	DEFB	1
9E49	D500	3370	DEFW	213
9E4B	05	3380	DEFB	5
9E4C-	EF00	3390	DEFW	239
9E4E	Ø1	3400	DEFB	1

9E4F	3FØ1	3410	DEFW	319
9E51	01	3420	DEFB	1
9E52	5201	3430	DEFW	338
9E54	01	3440	DEFB	1
9E55	3F01	3450	DEFW	319
9E57	02	3460	DEFB	2
9E58	EF00	3470	DEFW	239
9E5A	01	3480	DEFB	1
9E5B	1CØ1	3490	DEFW	284
9E5D	02	3500	DEFB	2
9E5E	EF00	3510	DEFW	239
9E60	Ø1	3520	DEFB	1
9E61	1CØ1	3530	DEFW	284
9E63	Ø1	3540	DEFB	1
9E64	EF00	3550	DEFW	239
9E66	Ø1	3560	DEFB	1
9E67	1001	3570	DEFW	284
9E69	Ø1	3580	DEFB	1
9E6A	3F01	3590	DEFW	319
9E6C	01	3600	DEFB	1
9E6D	EF00	3610	DEFW	239
9E6F	01	3620	DEFB	1
9E70	BE00	3630	DEFW	190
9E72	02	3640	DEFB	2
9E73	9F00	3650	DEFW	159
9E75	01	3660	DEFB	1
9E76	BE00	3670	DEFW	190

9E78	Ø1	3680	DEFB 1
9E79	EFØØ	3690	DEFW 239
9E7B	01	3700	DEFB 1
9E7C	3F01	3710	DEFW 319
9E7E	02	3720	DEFB 2
9E7F	1CØ1	3730	DEFW 284
9E81	02	3740	DEFB 2
9E82	EF00	3750	DEFW 239
9E84	Ø1	3760	DEFB 1
9E85	BE00	3770	DEFW 190
9E87	03	3780	DEFB 3
9E88	D500	3790	DEFW 213
9E8A	06	3800	DEFB 6
9E8B	EF00	3810	DEFW 239
9E8D	FF	3820	DEFB 255,1
	01		
		3830	END

Nine – Machine Code Monitor

This program allows you to examine the contents of your Amstrad's memory. You can flick through screenfuls of ROM and RAM with it. In addition, you can alter the **contents** of memory locations, so making the program more than just a "memory viewer".

The available commands are:

'3' = Fast forwards through memory.
'E' = Slow forwards through memory.
'D' = Slow backwards through memory.
'X' = Fast backwards through memory.
'Q' = Quit. (You can't use ESC to break the program).
'ENTER' key to enter a number into an address.

There are several points to keep in mind when using the ENTER command. The address that can be altered when the ENTER key is pressed is that address at the very top of the screen display. Secondly, be careful where you attempt to alter address values, particularly between addresses 40000 and 40276, as this is where the monitor program is stored.

```
10 SYMBOL AFTER 256:MEMORY 39999:SYMBOL
AFTER 240
20 CLS
30 monitor=40000:scroll=40034
40 LOCATE 11,2:PRINT"Machine Code Monito
r"
```

```
50 IF PEEK(40000)=&DD AND PEEK(40001)=&6
E GOTO 90
60 FOR n=40000 TO 40276
70 READ a$:POKE n.VAL("&"+a$)
80 NEXT
90 LOCATE 11.6:INPUT"Start address";st
100 LOCATE 11.8:PRINT"For addresses &000
0-&4000"
110 LOCATE 11.9:PRINT"and &C000-&FFFF ex
amine"
120 LOCATE 11,10:PRINT"(A) ROM or (B) RA
M?"
130 rom%=0
140 IF INKEY(69)=0 THEN rom%=1:GOTO 160
150 IF INKEY(54)<>0 GOTO 140
160 \text{ start%}=INT(UNT(st))-25
170 POKE scroll,25
180 CALL monitor,@rom%,@start%
190 IF rom%=99 THEN CLS:END
200 LOCATE 12,1:PRINT"> <":LOCATE 13,1
210 GOSUB 260:h$=a$
220 GOSUB 260:1$=a$
230 IF INKEY(57)*INKEY(58)*INKEY(61)=0 G
OTO 230
240 POKE start%,VAL("&"+h$+1$)
250 POKE scroll,1:GOTO 180
```

```
260 a$=UPPER$(INKEY$):IF a$<"0" OR a$>"F
" OR a$>"9" AND a$<"A" GOTO 260
270 PRINT a$:
280 RETURN
290 DATA DD.6E.00.DD.66.01.5E.23.56.DD
300 DATA 6E,02,DD,66,03,7E,B7,28,08,CD
310 DATA 00, B9, CD, 06, B9, 18, 06, CD, 03, B9
320 DATA CD.09.B9.06.19.C5.3E.19.CD.D1
330 DATA 9C,C1,10,F7,3E,39,CD,1E,BB,28
340 DATA 07,3E,19,CD,D1,9C,18,F2,3E,3F
350 DATA CD,1E,BB,28,07,3E,01,CD,D1,9C
360 DATA 18,E4,3E,3A,CD,1E,BB,28,07,3E
370 DATA 19,CD,C7,9C,18,D6,3E,3D,CD,1E
380 DATA BB,28,07,3E,01,CD,C7,9C,18,C8
390 DATA 3E,12,CD,1E,BB,20,0F,3E,43,CD
400 DATA 1E.BB.28.BA.DD.6E.02.DD.66.03
410 DATA 36,63,DD,6E,00,DD,66,01,73,23
420 DATA 72,CD,03,BB,C9,F5,01,00,50,0B
430 DATA 78, B1, 20, FB, F1, F5, D5, 47, CB, 38
440 DATA CB, 38, CB, 38, CB, 38, AF, CD, 4D, BC
450 DATA 3E,1F,CD,5A,BB,3E,05,CD,5A,BB
460 DATA D1,F1,F5,CD,5A,BB,FE,01,28,07
470 DATA 13,21,18,00,19,18,03,18,62,68
480 DATA 44,CD,34,9D,45,CD,34,9D,3E,1F
490 DATA CD, 5A, BB, 3E, 0D, CD, 5A, BB, F1, F5
500 DATA CD,5A,BB,46,CD,34,9D,C1,7E,FE
```

510 DATA 21,F8,FE,7F,F0,3E,1F,CD,5A,BB 520 DATA 3E,13,CD,5A,BB,78,CD,5A,BB,7E 530 DATA CD,5A,BB,C9,78,E6,0F,4F,CB,38 540 DATA CB,38,CB,38,CB,38,78,06,02,FE 550 DATA 0A,F2,4C,9D,C6,30,18,02,C6,37 560 DATA CD,5A,BB,79,10,EF,C9

Note that the first 64 bytes of the lower ROM are identical to the first 64 bytes of RAM. This is because those first bytes constitute the low kernel jump block and are copied from ROM to RAM when the machine is switched on.

You will be asked for a starting address when the program runs. Enter this in decimal or hex. If you use hex, the address must be preceded by an '&' character. As you can see from the dump of the program in action, all addresses and values are displayed in hexadecimal.

> Machine Code Monitor Start address? &742 For addresses &0000-&4000 and &C000-&FFFF examine (A) ROM or (B) RAM? 0742 53 S 0743 63 c 0744 68 h

0745	6E	n
0746	65	е
0747	69	i
0748	64	đ
0749	65	e
074A	72	r
074B	00	
074C	ØA	
074D	20	
074E	41	А
074F	77	W
0750	61	a
0751	00	
0752	ØA	
0753	20	
0754	53	S
0755	6F	o
0756	6C	1
0757	61	а
0758	76	v
0759	6F	o
075A	78	x

		0001	;	MC	MONITOR
		0002	;		
9040		0010		ORG	40000
B900		0020	UROMON	DEFL	0B900H
B903		0030	UROMOF	DEFL	0B903H
B906		0040	LROMON	DEFL	0B906H
B909		0050	LROMOF	DEFL	0B909H
BB1E		0060	TSTKEY	DEFL	ØBB1EH
BB03		0070	KRESET	DEFL	0BB03H
BC4D		0080	SCROLL	DEFL	ØBC4DH
BB5A		0090	тхтоит	DEFL	ØBB5AH
9040	DD6E00	0100		LD	L,(IX+0)
9043	DD6601	0110		LD	H,(IX+1)
9046	5E	0120		LD	E,(HL)
9047	23	0130		INC	HL
9048	56	0140		LD	D,(HL)
9049	DD6E02	0150		LD	L,(IX+2)
9C4C	DD6603	0160		LD	H,(IX+3)
9C4F	7E	0170		LD	A,(HL)
9050	B7	0180		OR	А
9051	2808	0190		JR	Z,ROMOFF
9053	CD00B9	0200		CALL	UROMON
9056	CD06B9	0210		CALL	LROMON
9059	1806	0220		JR	ROMSET
9C5B	CD03B9	0230	ROMOFF	CALL	UROMOF
9C5E	CD09B9	0240		CALL	LROMOF
9061	0619	0250	ROMSET	LD	B,25

9063	C5	0260	SETUP	PUSH	BC
9064	3E19	0270		LD	A,25
9066	CDD19C	0280		CALL	NOWAIT
9069	C1	0290		POP	BC
9C6A	10F7	0300		DJNZ	SETUP
9C6C	3E39	0310	MAIN	LD	A,57
9C6E	CD1EBB	0320		CALL	TSTKEY
9071	2807	0330		JR	z,NOT3
9073	3E19	0340		LD	A,25
9075	CDD19C	0350		CALL	NOWAIT
9078	18F2	0360		JR	MAIN
9C7A	3E3F	0370	NOT3	LD	A,63
9070	CD1EBB	0380		CALL	TSTKEY
9C7F	2807	0390		JR	Z,NOTX
9081	3EØ1	0400		LD	Α,1
9083	CDD19C	0410		CALL	NOWAIT
9086	18E4	0420		JR	MAIN
9088	3E3A	0430	NOTX	LD	A,58
9C8A	CD1EBB	0440		CALL	TSTKEY
9C8D	2807	0450		JR	Z,NOTE
9C8F	3E19	0460		LD	A,25
9091	CDC79C	0470		CALL	PAUSE
9094	18D6	0480		JR	MAIN
9096	3E3D	0490	NOTE	LD	A,61
9098	CD1EBB	0500		CALL	TSTKEY
9C9B	2807	0510		JR	Z,NOTD
9C9D	3EØ1	0520		LD	Α,1
9C9F	CDC79C	0530		CALL	PAUSE
------	--------	------	--------	------	----------
9CA2	1808	0540		JR	MAIN
9CA4	3E12	0550	NOTD	LD	A,18
9CA6	CD1EBB	0560		CALL	TSTKEY
9CA9	200F	0570		JR	NZ,ENTER
9CAB	3E43	0580		LD	A,67
9CAD	CD1EBB	0590		CALL	TSTKEY
9CB0	28BA	0600		JR	Z,MAIN
9CB2	DD6E02	0610		LD	L,(IX+2)
9CB5	DD6603	0620		LD	H,(IX+3)
9CB8	3663	0630		LD	(HL),99
9CBA	DD6E00	0640	ENTER	LD	L,(IX+0)
9CBD	DD6601	0650		LD	H,(IX+1)
9000	73	0660		LD	(HL),E
9001	23	0670		INC	HL
9002	72	0680		LD	(HL),D
9003	CD03BB	0690		CALL	KRESET
9006	C9	0700		RET	
9007	F5	0710	PAUSE	PUSH	AF
9008	010050	0720		LD	вс,5000Н
9CCB	0B	0730	DELAY	DEC	BC
9000	78	0740		LD	А,В
9CCD	B1	0750		OR	С
9CCE	20FB	0760		JR	NZ,DELAY
9CD0	F1	0770		POP	AF
9CD1	F5	0780	NOWAIT	PUSH	AF
9CD2	D5	0790		PUSH	DE

.

9CD3	47	0800		LD	в,А
9CD4	CB38	0810		SRL	В
9CD6	CB38	0820		SRL	в
9CD8	CB38	0830		SRL	В
9CDA	CB38	0840		SRL	в
9CDC	AF	0850		XOR	А
9CDD	CD4DBC	0860		CALL	SCROLL
9CE0	3E1F	0870		LD	A,31
9CE2	CD5ABB	0880		CALL	тхтоит
9CE5	3E05	0890		LD	Α,5
9CE7	CD5ABB	0900		CALL	тхтоит
9CEA	D1	0910		POP	DE
9CEB	F1	0920		POP	AF
9CEC	F5	0930		PUSH	AF
9CED	CD5ABB	0940		CALL	TXTOUT
9CFØ	FEØ1	0950		СР	1
9CF2	2807	0960		JR	Z,DOWN
9CF4	13	0970		INC	DE
9CF5	211800	0980		LD	HL,24
9058	19	0990		ADD	HL,DE
9CF9	1803	1000		JR	PRINT
9CFB	1B	1010	DOWN	DEC	DE
9CFC	62	1020		LD	H,D
9CFD	6B	1030		LD	L,E
9CFE	44	1040	PRINT	LD	в,н
9CFF	CD349D	1050		CALL	HEXOUT
9D02	45	1060		LD	B,L

136

9D03	CD349D	1070		CALL	HEXOUT
9D06	3E1F	1080		LD	A,31
9D08	CD5ABB	1090		CALL	тхтоит
9D0B	3E0D	1100		LD	A,13
9D0D	CD5ABB	1110		CALL	тхтоит
9D10	F1	1120		POP	AF
9D11	F5	1130		PUSH	AF
9D12	CD5ABB	1140		CALL	TXTOUT
9D15	46	1150		LD	B,(HL)
9D16	CD349D	1160		CALL	HEXOUT
9D19	C1	1170		POP	BC
9D1A	7E	1180		LD	A,(HL)
9D1B	FE21	1190		СР	33
9D1D	F8	1200		RET	м
9D1E	FE7F	1210		СР	127
9D20	FØ	1220		RET	Р
9D21	3E1F	1230		LD	A,31
9D23	CD5ABB	1240		CALL	TXTOUT
9 D 26	3E13	1250		LD	A,19
9D28	CD5ABB	1260		CALL	TXTOUT
9D2B	78	1270		LD	А,В
9D2C	CD5ABB	1280		CALL	тхтоит
9D2F	7E	1290		LD	A,(HL)
9D30	CD5ABB	1300		CALL	тхтоот
9D33	C9	1310		RET	
9D34	78	1320	HEXOUT	LD	А,В
9D35	E60F	1330		AND	15

9D37	4F	1340		LD	С,А
9D38	CB38	1350		SRL	В
9D3A	CB38	1360		SRL	В
9D3C	CB38	1370		SRL	В
9D3E	CB38	1380		SRL	в
9D40	78	1390		LD	А,В
9D41	0602	1400		LD	В,2
9D43	FEØA	1410	DIGIT	СР	10
9D45	F24C9D	1420		JP	P,LETTER
9D48	C630	1430		ADD	A,48
9D4A	1802	1440		JR	PRNTCH
9D4C	C637	1450	LETTER	ADD	A,55
9D4E	CD5ABB	1460	PRNTCH	CALL	тхтоот
9D51	79	1470		LD	А,С
9D52	10EF	1480		DJNZ	DIGIT
9D54	C9	1490		RET	
		1500		END	

Ten – Box Scrolls

We are quite proud of these routines. You should find them very handy for any programs which require moving graphics.

Scroll Box Up

Every machine code book worth its salt includes some sort of scroll routine. This is usually restricted to scrolling a row, a column or the whole screen. Our routine for **Scroll Box Up** (and the three other direction routines, down, left and right) allows you to specify a "box", or window on-screen, and make that box scroll smoothly.

A smooth scroll is a pixel by pixel scroll. Even in machine code it is not dramatically fast, but should be rapid enough for most requirements. Its great asset is the smoothness of the scroll, compared with a character by character scroll.

The box is defined in a similar way to the WINDOW command. The format is:

CALL 40000, L, R, U, D

In this, L stands for the leftmost co-ordinate of the box, R for the rightmost co-ordinate, U for the top, and D for the bottom co-ordinate. Therefore, the box can only be square or rectangular. Obviously, L must be equal to, or smaller than R; and U must be equal to, or smaller than D. Note that L, R, U and D can be either numeric variables or numbers.

```
1 7
              SCROLL BOX UP
10 SYMBOL AFTER 256:MEMORY 39999:SYMBOL
AFTER 240
20 FOR n=40000 TO 40145
30 READ a$:POKE n.VAL("&"+a$)
40 NEXT
50 DATA DD.6E.02.DD.66.06.2D.25.DD.7E
60 DATA 04,94,4F,DD,7E,00,95,57,CD,1A
70 DATA BC.3E.00.81.10.FD.5F.18.38.D5
80 DATA 43,5D,7C,C6.38,57,7D,C6.50.6F
90 DATA 30.0A.24.7C.E6.07.20.04.7C.D6
100 DATA 08,67,E5,7E,12,1C,20,0A,14,7A
110 DATA E6,07,20,04,7A,D6,08,57,2C,20
120 DATA 0A,24,7C,E6,07,20,04,7C,D6,08
130 DATA 67,10,E2,E1,D1,D5,E5,4B,E5,06
140 DATA 07,5D,54,7C,C6,08.67.7E.12.10
150 DATA F6,E1,2C.20.0A.24.7C.E6.07.20
160 DATA 04,7C,D6,08,67,0D,20,E2,E1,D1
170 DATA 15,20,A2,7C,C6,38,67,43,36,00
180 DATA 2C,20,0A,24,7C,E6,07,20,04,7C
190 DATA D6,08,67,10,EF,C9
```

When you make the CALL, the specified box would scroll up one single pixel with the top line disappearing and an extra blank line appearing at the bottom of the box. One single scroll is not enough for most uses. To scroll a number of times, we use a FOR/NEXT loop.

There are few limitations in using this program. It works in any

screen mode, and several boxes can be scrolled at the same time. Of course, the more boxes you have on-screen, as well as the larger the boxes you have, the slower the scroll procedure works. Therefore, it is best to restrict the size and number of boxes to that which is really necessary. For a weird effect, try scrolling two boxes that overlap each other.

		0010	;SCRL	BOX	UP
		0020	;		
9040		0030		ORG	40000
BC1A		0040	CHRPOS	DEFL	ØBC1AH
9040	DD6E02	0050		LD	$L_{,(IX+2)}$
9043	DD6606	<i>0060</i>		LD	H,(IX+6)
9046	2D	0070		DEC	L
9047	25	0080		DEC	Н
9048	DD7E04	0090		LD	A,(IX+4)
9C4B	94	0100		SUB	н
9C4C	4F	0110		LD	С,А
9C4D	DD7E00	0120		LD	A,(IX+0)
9050	95	0130		SUB	L
9051	57	0140		LD	D,A
9052	CD1ABC	0150		CALL	CHRPOS
9055	3E00	0160		LD	Α,0
9057	81	0170	CHRWID	ADD	A,C
9058	10FD	0180		DJNZ	CHRWID
9C5A	5F	0190		LD	Е,А
9C5B	1838	0200		JR	START
9C5D	D5	0210	NXTROW	PUSH	DE
9C5E	43	0220		LD	B,E

9C5F	5D	0230		LD	E,L
9060	7C	0240		LD	А,Н
9061	C638	0250		ADD	A,56
9063	57	0260		LD	D,A
9064	7D	0270		LD	A,L
9065	C650	0280		ADD	A,80
9067	6F	0290		LD	L,A
9068	300A	0300		JR	NC, NEWL IN
9C6A	24	0310		INC	н
9C6B	7C	0320		LD	А,Н
9060	E607	0330		AND	7
9C6E	2004	0340		JR	NZ,NEWLIN
9070	7C	0350		LD	А,Н
9071	D608	0360		SUB	8
9073	67	0370		LD	H,A
9074	E5	0380	NEWLIN	PUSH	HL
9075	7E	0390	LASTLN	LD	A,(HL)
9076	12	0400		LD	(DE),A
9077	1C	0410		INC	Ε
9078	200A	0420		JR	NZ,DEOK
9C7A	14	0430		INC	D
9C7B	7A	0440		LD	A,D
9070	E607	0450		AND	7
9C7E	2004	0460		JR	NZ,DEOK
9080	7A	0470		LD	A,D
9081	D608	0480		SUB	8

9083	57	0490		LD	D,A
9084	2C	0500	DEOK	INC	L
9085	200A	0510		JR	NZ,HLOK
9087	24	0520		INC	н
9088	7C	0530		LD	А,Н
9089	E607	0540		AND	7
9C8B	2004	0550		JR	NZ,HLOK
9C8D	7C	0560		LD	А,Н
9C8E	D608	0570		SUB	8
9090	67	0580		LD	H,A
9091	10E2	0590	HLOK	DJNZ	LASTLN
9093	E1	0600		POP	HL
9094	D1	0610		POP	DE
9095	D5	0620	START	PUSH	DE
9096	E5	0630		PUSH	HL
9097	4B	0640		LD	C,E
9098	E5	0650	NXTCHR	PUSH	HL
9099	0607	0660		LD	в,7
9C9B	5D	0670	NXTLIN	LD	E,L
9090	54	0680		LD	р,н
9C9D	7C	0690		LD	А,Н
9C9E	C608	0700		ADD	А,8
9CA0	67	0710		LD	H,A
9CA1	7E	0720		LD	A,(HL)
9CA2	12	0730		LD	(DE),A
9CA3	10F6	0740		DJNZ	NXTLIN

9CA5	E1	0750		POP	HL
9CA6	2C	0760		INC	L
9CA7	200A	0770		JR	NZ,OK
9CA9	24	0780		INC	н
9CAA	7C	0790		LD	А,Н
9CAB	E607	0800		AND	7
9CAD	2004	0810		JR	NZ,OK
9CAF	7C	0820		LD	А,Н
9CB0	D608	0830		SUB	8
9CB2	67	0840		LD	H,A
9CB3	ØD	0850	ок	DEC	С
9CB4	20E2	0860		JR	NZ,NXTCHR
9CB6	E1	0870		POP	HL
9CB7	D1	0880		POP	DE
9CB8	15	0890		DEC	D
9CB9	20A2	0900		JR	NZ,NXTROW
9CBB	7C	0910		LD	А,Н
9CBC	C638	0920		ADD	A,56
9CBE	67	0930		LD	H,A
9CBF	43	0940		LD	B,E
9000	3600	0950 1	BLANK	LD	(HL),0
9002	2C	0960		INC	L
9003	200A	0970		JR	NZ,CONT
9005	24	0980		INC	н
9006	7C	0990		LD	А,Н
9007	E607	1000		AND	7
9009	2004	1010		JR	NZ,CONT

9CCB	7C	1020		LD	А,Н
9000	D608	1030		SUB	8
9CCE	67	1040		LD	H,A
9CCF	10EF	1050	CONT	DJNZ	BLANK
9CD1	C9	1060		RET	
		1070		END	

Scroll Box Down

This routine is similar in design and use to the previous routine used for scrolling a box up the screen. It works in all three screen modes as does Scroll Box Up and the format to use it is exactly the same:

CALL 40000, L R, U, D

```
1 ' SCROLL BOX DOWN
10 SYMBOL AFTER 256:MEMORY 39999:SYMBOL
AFTER 240
20 FOR n=40000 TO 40151
30 READ a$:POKE n,VAL("&"+a$)
40 NEXT
50 DATA DD,6E,00,DD,66,06,2D,25,DD,7E
60 DATA 04,94,4F,7D,C6,02,DD,96,02,57
70 DATA CD,1A,BC,7C,C6,08,67,3E,00,81
80 DATA 10,FD,5F,18,38,D5,43,5D,7C,D6
90 DATA 38,57,7D,D6,50,6F,30,0A,7C,25
100 DATA E6,07,20,04,7C,C6,08,67,E5,7E
110 DATA 12,1C,20,0A,14,7A,E6,07,20,04
```

120	DATA	7A,D6,08,57,2C,20,0A,24,7C,E6
130	DATA	07,20,04,7C,D6,08,67,10,E2,E1
140	DATA	D1,D5,E5,4B,E5,06,07,5D,54,7C
150	DATA	D6,08,67,7E,12,10,F6,E1,2C,20
160	DATA	ØA,24,7C,E6,07,20,04,7C,D6,08
170	DATA	67,0D,20,E2,E1,D1,15,20,A2,7C
180	DATA	D6,38,67,43,36,00,2C,20,0A,24
190	DATA	7C,E6,07,20,04,7C,D6,08,67,10
200	DATA	EF,C9

Once again, L is left, R is right, U is up and D is down. As with the previous routine, a single CALL will result in a one pixel-line scroll so a loop will be necessary if you intend to do a lot of scrolling.

		0001	;	SCRL	DOWN
		0002	;		
9040		0010		ORG	40000
BC1A		0020	CHRPOS	DEFL	ØBC1AH
9040	DD6E00	0030		LD	L,(IX+0)
9043	DD6606	<i>0040</i>		LD	H,(IX+6)
9046	2D	0050		DEC	L
9047	25	<i>0060</i>		DEC	н
9048	DD7EØ4	<i>0</i> 070		LD	A,(IX+4)
9C4B	94	0080		SUB	н
9C4C	4F	0090		LD	С,А
9C4D	7D	0100		LD	A,L
9C4E	C602	0110		ADD	Α,2
9050	DD9602	0120		SUB	(IX+2)

9053	57	0130		LD	D,A
9054	CD1ABC	0140		CALL	CHRPOS
9057	7C	0150		LD	А,Н
9058	C638	0160		ADD	A,56
9C5A	67	0170		LD	H,A
9C5B	3E00	0180		LD	Α,Ø
9C5D	81	0190	CHRWID	ADD	A,C
9C5E	10FD	0200		DJNZ	CHRWID
9060	5F	0210		LD	E,A
9061	1838	0220		JR	START
9063	D5	0230	NXTROW	PUSH	DE
9064	43	0240		LD	в,Е
9065	5D	0250		LD	E,L
9066	7C	0260		LD	А,Н
9067	D638	0270		SUB	56
9069	57	0280		LD	D,A
9C6A	7D	0290		LD	A,L
9C6B	D650	0300		SUB	80
9C6D	6F	0310		LD	L,A
9C6E	300A	0320		JR	NC,NEWLIN
9070	7C	0330		LD	А,Н
9071	25	0340		DEC	н
9072	E607	0350		AND	7
9074	2004	0360		JR	NZ,NEWLIN
9076	7C	0370		LD	А,Н
9077	C608	0380		ADD	A,8
9079	67	0390		LD	Н,А

9C7A	E5	0400	NEWLIN	PUSH	HL
9C7B	7E	0410	LASTLN	LD	A,(HL)
9C7C	12	0420		LD	(DE),A
9C7D	1C	0430		INC	Ε
9C7E	200A	0440		JR	NZ,DEOK
9080	14	0450		INC	D
9081	7A	0460		LD	A,D
9082	E607	0470		AND	7
9084	2004	0480		JR	NZ,DEOK
9086	7A	0490		LD	A,D
9C87	D608	0500		SUB	8
9089	57	0510		LD	D,A
9C8A	2C	0520	DEOK	INC	L
9C8B	200A	0530		JR	NZ,HLOK
9C8D	24	0540		INC	н
9C8E	7C	0550		LD	А,Н
9C8F	E607	0560		AND	7
9091	2004	0570		JR	NZ,HLOK
9093	7C	0580		LD	А,Н
9094	D608	0590		SUB	8
9096	67	0600		LD	Н,А
9097	10E2	0610	HLOK	DJNZ	LASTLN
9099	E1	0620		POP	HL
9C9A	D1	0630		POP	DE
9C9B	D5	0640	START	PUSH	DE
9090	E5	0650		PUSH	HL
9C9D	4B	0660		LD	C,E

9C9E	E5	0670	NXTCHR	PUSH	HL
9C9F	0607	0680		LD	в,7
9CA1	5D	0690	NXTLIN	LD	E,L
9CA2	54	0700		LD	р,н
9CA3	7C	0710		LD	А,Н
9CA4	D608	0720		SUB	8
9CA6	67	0730		LD	H,A
9CA7	7E	0740		LD	A,(HL)
9CA8	12	0750		LD	(DE),A
9CA9	10F6	0760		DJNZ	NXTLIN
9CAB	E 1	0770		POP	HL
9CAC	2C	0780		INC	L
9CAD	200A	0790		JR	NZ,OK
9CAF	24	0800		INC	н
9CB0	7C	0810		LD	А,Н
9CB1	E607	0820		AND	7.
9CB3	2004	0830		JR	NZ,OK
9CB5	7C	0840		LD	А,Н
9CB6	D608	0850		SUB	8
9CB8	67	0860		LD	H,A
9CB9	ØD	0870	0K	DEC	С
9CBA	20E2	0880		JR	NZ,NXTCHR
9CBC	E1	0890		POP	HL
9CBD	D1	0900		POP	DE
9CBE	15	0910		DEC	D
9CBF	20A2	0920		JR	NZ,NXTROW

9001	7C	0930		LD	А,Н
9002	D638	0940		SUB	56
9004	67	0950		LD	н,А
9005	43	0960		LD	B,E
9006	3600	0970	BLANK	LD	(HL),0
9008	2C	0980		INC	L
9009	200A	0990		JR	NZ,CONT
9CCB	24	1000		INC	н
9000	7C	1010		LD	А,Н
9CCD	E607	1020		AND	7
9CCF	2004	1030		JR	NZ,CONT
9CD1	7C	1040		LD	А,Н
9CD2	D608	1050		SUB	8
9CD4	67	1060		LD	Н,А
9CD5	10EF	1070	CONT	DJNZ	BLANK
9CD7	C9	1080		RET	
		1090		END	

Scroll Box Left

Though scroll left (and right) work in exactly the same way as scroll up and down, the **routine** to create the effect has some differences. This is mainly due to the horizontal scrolling movement as opposed to the vertical movement of the previous two routines.

The format to use it is exactly the same:

CALL 40000, L, R, U, D

1 7 SCROLL BOX LEFT 10 SYMBOL AFTER 256:MEMORY 39999:SYMBOL AFTER 240 20 FOR n=40000 TO 40271 30 READ a\$:POKE n.VAL("&"+a\$) 40 NEXT 50 DATA DD,6E,02,DD,66,06,2D,25,DD,7E 60 DATA 04,94,4F,DD,7E,00,95,57,CD,1A 70 DATA BC.3E.00.81.10.FD.47.CD.11.BC 80 DATA DA.FA.9C.28.48.DD.6E.02.DD.66 90 DATA 04,2D,25,C5,CD,1A,BC,C1,0E,08 100 DATA C5,E5,B7,CB,16,F5,7D,2D,B7,20 110 DATA 0A.7C.25.E6.07.20.04.7C.C6.08 120 DATA 67,F1,10,EB,E1,7C,C6,08,67,C1 130 DATA 0D,20,DF,7C,D6,40,67,7D,C6,50 140 DATA 6F, 30, 0A, 24, 7C, E6, 07, 20, 04, 7C 150 DATA D6,08,67,15,20,C6,C9,05,D5,0E 160 DATA 08,C5,E5,CB,26,5D,54,2C,20,0A 170 DATA 24,7C,E6,07,20,04,7C,D6,08,67 180 DATA 1A,CB,26,38,04,CB,A7,18,02,CB 190 DATA E7,CB,66,28,02,CB,C7,12,10,DD 200 DATA CB,A6,E1,7C,C6,08,67,C1,0D,20 210 DATA CE,7C,D6,40,67,7D,C6,50,6F,30 220 DATA 0A,24,7C,E6,07,20,04,7C,D6,08 230 DATA 67,D1,15,20,B3,C9,05,D5,0E,08 240 DATA C5,E5,C5,5D,54,2C,20,0A,24,7C 250 DATA E6,07,20,04,7C,D6,08,67,7E,1F

260	DATA	4F,1A,17,17,06,04,17,CB,21,CB
270	DATA	21,17,10,F8,12,C1,10,DC,7E,17
280	DATA	06,04,07,CB,27,10,FB,77,E1,7C
290	DATA	C6,08,67,C1,0D,20,C7,7C,D6,40
300	DATA	67,7D,C6,50,6F,30,0A,24,7C,E6
310	DATA	07,20,04,7C,D6,08,67,D1,15,20
320	DATA	AC,C9

As usual, L is left, R is right, U is up and D is down. As with the other routines, a single CALL will result in a one pixel-line scroll so a loop will be necessary if you intend to do a lot of scrolling. This scroll also works in all three screen modes as does Scroll Right.

		0001	;	SCRL	LEFT
		0002	;		
9040		0010		ORG	40000
BC1A		0020	CHRPOS	DEFL	ØBC1AH
BC11		0030	SCMODE	DEFL	ØBC11H
9040	DD6E02	0040		LD	L,(IX+2)
9043	DD6606	0050		LD	H,(IX+6)
9046	2D	0060		DEC	L
9047	25	0070		DEC	н
9048	DD7E04	0080		LD	A,(IX+4)
9C4B	94	0090		SUB	н
9C4C	4F	0100		LD	С,А
9C4D	DD7E00	0110		LD	A,(IX+0)
9050	95	0120		SUB	L
9051	57	0130		LD	D,A

9052	CD1ABC	0140		CALL	CHRPOS
9055	3E00	0150		LD	Α,0
9057	81	0160	CHRWID	ADD	A,C
9058	10FD	0170		DJNZ	CHRWID
9C5A	47	0180		LD	в,А
9C5B	CD11BC	0190		CALL	SCMODE
9C5E	DAFA9C	0200		JP	C,MODE0
9061	2848	0210		JR	Z,MODE1
9063	DD6E02	0220		LD	L,(IX+2)
9066	DD6604	0230		LD	H,(IX+4)
9069	2D	0240		DEC	L
9C6A	25	0250		DEC	Н
9C6B	C5	0260		PUSH	BC
9C6C	CD1ABC	0270		CALL	CHRPOS
9C6F	C1	0280		POP	BC
9C70	0E08	0290	ROW2	LD	С,8
9072	C5	0300	LINE2	PUSH	BC
9073	E5	0310		PUSH	HL
9074	B7	0320		OR	Α
9075	CB16	0330	BYTE2	RL	(HL)
9077	F5	0340		PUSH	AF
9078	7D	0350		LD	A,L
9079	2D	0360		DEC	L
9C7A	B7	0370		OR	A
9C7B	200A	0380		JR	NZ,OK2
9C7D	7C	0390		LD	А,Н
9C7E	25	0400		DEC	н

9C7F	E607	0410	AND	7
9081	2004	0420	JR	NZ,OK2
9083	7C	0430	LD	А,Н
9084	C608	0440	ADD	Α,8
9086	67	0450	LD	Н,А
9087	F1	0460 OK2	POP	AF
9088	10EB	0470	DJNZ	BYTE2
9C8A	E 1	0480	POP	HL
9C8B	7C	0490	LD	А,Н
9080	C608	0500	ADD	А,8
9C8E	67	0510	LD	Н,А
9C8F	C1	0520	POP	BC
9090	ØD	0530	DEC	С
9091	20DF	0540	JR	NZ,LINE2
9093	7C	0550	LD	А,Н
9094	D640	0560	SUB	64
9096	67	0570	LD	Н,А
9097	7D	0580	LD	A,L
9098	C650	0590	ADD	A,80
9C9A	6F	0600	LD	L,A
9C9B	300A	0610	JR	NC,DONE2
9C9D	24	0620	INC	н
9C9E	7C	0630	LD	А,Н
9C9F	E607	0640	AND	7
9CA1	2004	0650	JR	NZ,DONE2
9CA3	7C	0660	LD	А,Н
9CA4	D608	0670	SUB	8

9CA6	67	0680		LD	H,A
9CA7	15	0690	DONE2	DEC	D
9CA8	2006	0700		JR	NZ,ROW2
9CAA	C9	0710		RET	
9CAB	Ø5	0720	MODE 1	DEC	В
9CAC	D5	0730	ROW1	PUSH	DE
9CAD	0E08	0740		LD	с,8
9CAF	C5	0750	LINE1	PUSH	BC
9CB0	E5	0760		PUSH	HL
9CB1	CB26	0770		SLA	(HL)
9CB3	5D	0780	BYTE1	LD	E,L
9CB4	54	0790		LD	D,H
9CB5	20	0800		INC	L
9CB6	200A	0810		JR	NZ,0K1
9CB8	24	0820		INC	н
9CB9	7C	0830		LD	А,Н
9CBA	E607	0840		AND	7
9CBC	2004	0850		JR	NZ,0K1
9CBE	7C	0860		LD	А,Н
9CBF	D608	0870		SUB	8
9001	67	0880		LD	H,A
9002	1 A	0890	0K1	LD	A,(DE)
9003	CB26	0900		SLA	(HL)
9005	3804	0910		JR	C,SETBIT
9007	CBA7	0920		RES	4,A
9009	1802	0930		JR	BITOK
9CCB	CBE7	0940	SETBIT	SET	4,A

9CCD	CB66	0950	BITOK	BIT	4,(HL)
9CCF	2802	0960		JR	Z,BITSET
9CD1	CBC7	0970		SET	0,A
9CD3	12	0980	BITSET	LD	(DE),A
9CD4	10DD	0990		DJNZ	BYTE1
9CD6	CBA6	1000		RES	4,(HL)
9CD8	E1	1010		POP	HL
9CD9	7C	1020		LD	А,Н
9CDA	C608	1030		ADD	Α,8
9CDC	67	1040		LD	Н,А
9CDD	C1	1050		POP	BC
9CDE	ØD	1060		DEC	С
9CDF	20CE	1070		JR	NZ,LINE1
9CE1	7C	1080		LD	А,Н
9CE2	D640	1090		SUB	64
9CE4	67	1100		LD	Н,А
9CE5	7D	1110		LD	A,L
9CE6	C650	1120		ADD	A,80
9CE8	6F	1130		LD	L,A
9CE9	300A	1140		JR	NC, DONE 1
9CEB	24	1150		INC	н
9CEC	7C	1160		LD	А,Н
9CED	E607	1170		AND	7
9CEF	2004	1180		JR	NZ,DONE1
9CF1	7C	1190		LD	А,Н
9CF2	D608	1200		SUB	8

.

9CF4	67	1210		LD	H,A
9CF5	D1	1220	DONE 1	POP	DE
9CF6	15	1230		DEC	D
9CF7	20B3	1240		JR	NZ,ROW1
9CF9	C9	1250		RET	
9CFA	05	1260	MODEØ	DEC	B
9CFB	D5	1270	ROWØ	PUSH	DE
9CFC	ØEØ8	1280		LD	с,8
9CFE	C5	1290	LINEØ	PUSH	BC
9CFF	E5	1300		PUSH	HL
9D00	C5	1310	ΒΥΤΕ Ø	PUSH	BC
9DØ1	5D	1320		LD	E,L
9D02	54	1330		LD	D,H
9D03	2C	1340		INC	L
9D04	200A	1350		JR	NZ,OKØ
9D06	24	1360		INC	н
9D07	7C	1370		LD	А,Н
9D08	E607	1380		AND	7
9D0A	2004	1390		JR	NZ,OKØ
9DØC	7C	1400		LD	А,Н
9D0D	D608	1410		SUB	8
9DØF	67	1420		LD	Н,А
9D10	7E	1430	OKØ	LD	A,(HL)
9D11	1 F	1440		RRA	
9D12	4F	1450		LD	с,А
9D13	1A	1460		LD	A,(DE)
9D14	17	1470		RLA	

9D15	17	1480		RLA	
9D16	0604	1490		LD	в,4
9D18	17	1500	NXTROT	RLA	
9D19	CB21	1510		SLA	С
9D1B	CB21	1520		SLA	С
9D1D	17	1530		RLA	
9D1E	10F8	1540		DJNZ	NXTROT
9D20	12	1550		LD	(DE),A
9D21	C1	1560		POP	BC
9D22	10DC	1570		DJNZ	BYTE0
9D24	7E	1580		LD	A,(HL)
9D25	17	1590		RLA	
9D26	0604	1600		LD	В,4
9D28	07	1610	LSTBYT	RLCA	
9D29	CB27	1620		SLA	Α
9D2B	10FB	1630		DJNZ	LSTBYT
9D2D	77	1640		LD	(HL),A
9D2E	E1	1650		POP	HL
9D2F	7C	1660		LD	А,Н
9D30	C608	1670		ADD	Α,8
9D32	67	1680		LD	Н,А
9D33	C1	1690		POP	BC
9D34	0D	1700		DEC	С
9D35	2007	1710		JR	NZ,LINE0
9D37	7C	1720		LD	А,Н
9D38	D640	1730		SUB	64
9D3A	67	1740		LD	Н,А

9D3B	7D	1750		LD	A,L
9D3C	C650	1760		ADD	A,80
9D3E	6F	1770		LD	L,A
9D3F	300A	1780		JR	NC,DONE0
9D41	24	1790		INC	н
9D42	7C	1800		LD	А,Н
9D43	E607	1810		AND	7
9D45	2004	1820		JR	NZ,DONEØ
9D47	7C	1830		LD	А,Н
9D48	D608	1840		SUB	8
9D4A	67	1850		LD	н,А
9D4B	D1	1860 DONI	EØ	POP	DE
9D4C	15	1870		DEC	D
9D4D	20AC	1880		JR	NZ,ROWØ
9D4F	C9	1890		RET	
		1900		END	

Scroll Box Right

With scroll up, scroll down and scroll left, it's not difficult to work out what this routine will do.

Once again, the format to use it is exactly the same as the other routines:

٠

CALL 40000, L, R, U, D

As with the other routines, a single CALL will result in a one pixelline scroll so a loop will be necessary if you intend to do a lot of scrolling.

1 ' SCROLL BOX RIGHT 10 SYMBOL AFTER 256:MEMORY 39999:SYMBOL AFTER 240 20 FOR n=40000 TO 40277 30 READ a\$:POKE n.VAL("&"+a\$) 40 NEXT 50 DATA DD.6E.02,DD.66,04,2D,25,7C,C6 60 DATA 02.DD.96.06.4F.DD.7E.00.95.57 70 DATA CD, 1A, BC, 2B, 3E, 00, 81, 23, 10, FC 80 DATA 47,CD.11.BC.DA.FE.9C.28.46.DD 90 DATA 6E,02,DD,66,06,2D,25,C5,CD,1A 100 DATA BC,C1,0E,08,C5,E5,B7,CB,1E,F5 110 DATA 2C,20,0A,24,7C,E6,07,20,04,7C 120 DATA D6,08,67,F1,10,ED,E1,7C,C6,08 130 DATA 67,C1,0D,20,E1,7C,D6,40,67,7D 140 DATA C6.50.6F.30.0A.24.7C.E6.07.20 150 DATA 04,7C,D6,08,67,15,20,C8,C9,05 160 DATA D5,0E,08,C5,E5,CB,3E,5D,54,7D 170 DATA 2D, B7, 20, 0A, 7C, 25, E6, 07, 20, 04 180 DATA 7C,C6,08,67,1A,CB,3E,38,04,CB 190 DATA 9F,18,02,CB,DF,CB,5E,28,02,CB 200 DATA FF,12,10,DB,CB,9E,E1,7C,C6,08 210 DATA 67,C1,0D,20,CC,7C,D6,40,67,7D 220 DATA C6,50,6F,30,0A,24,7C,E6,07,20 230 DATA 04,7C,D6,08,67,D1,15,20,B1,C9 240 DATA 05,D5,0E,08,C5,E5,C5,5D,54,7D 250 DATA 2D, B7, 20, 0A, 7C, 25, E6, 07, 20, 04

260 DATA 7C,C6,08,67,7E,17,4F,1A,1F,1F 270 DATA 06,04,1F,CB,39,CB,39,1F,10,F8 280 DATA 12,C1,10,DA,7E,1F,06,04,0F,CB 290 DATA 3F,10,FB,77,E1,7C,C6,08,67,C1 300 DATA 0D,20,C5,7C,D6,40,67,7D,C6,50 310 DATA 6F,30,0A,24,7C,E6,07,20,04,7C 320 DATA D6,08,67,D1,15,20,AA,C9

		0001	;	SCRL	RIGHT
		0002	;		
9040		0010		ORG	40000
BC1A		0020	CHRPOS	DEFL	ØBC1AH
BC11		0030	SCMODE	DEFL	ØBC11H
9040	DD6E02	0040		LD	L,(IX+2)
9043	DD6604	0050		LD	H,(IX+4)
9046	2D	0060		DEC	L
9047	25	0070		DEC	н
9048	7C	0080		LD	А,Н
9049	C602	0090		ADD	Α,2
9C4B	DD9606	0100		SUB	(IX+6)
9C4E	4F	0110		LD	с,А
9C4F	DD7E00	0120		LD	A,(IX+0)
9052	95	0130		SUB	L
9053	57	0140		LD	D,A

9054	CD1ABC	0150		CALL	CHRPOS
9057	2B	0160		DEC	HL
9058	3E00	0170		LD	Α,0
9C5A	81	0180	CHRWID	ADD	A,C
9C5B	23	0190		INC	HL
9050	10FC	0200		DJNZ	CHRWID
905E	47	0210		LD	в,А
9C5F	CD11BC	0220		CALL	SCMODE
9062	DAFE9C	0230		JP	C,MODE0
9065	2846	0240		JR	Z,MODE1
9067	DD6E02	0250		LD	L,(IX+2)
9C6A	DD6606	0260		LD	H,(IX+6)
9C6D	2D	0270		DEC	L
9C6E	25	0280		DEC	н
9C6F	C5	0290		PUSH	BC
9070	CD1ABC	0300		CALL	CHRPOS
9073	C1	0310		POP	BC
9074	0E08	0320	ROW2	LD	с,8
9076	C5	0330	LINE2	PUSH	BC
9077	E5	0340		PUSH	HL
9078	B7	0350		OR	А
9079	CB1E	0360	BYTE2	RR	(HL)
9C7B	F5	0370		PUSH	AF
9C7C	2C	0380		INC	L
9C7D	200A	0390		JR	NZ,0K2
9C7F	24	0400		INC	н

9080	7C	0410	LD	А,Н
9081	E607	0420	AND	7
9083	2004	0430	JR	NZ,0K2
9085	7C	0440	LD	А,Н
9086	D608	0450	SUB	8
9088	67	0460	LD	H,A
9089	F1	0470 OK2	POP	AF
9C8A	10ED	0480	DJNZ	BYTE2
9080	E1	0490	POP	HL
9C8D	7C	0500	LD	А,Н
9C8E	C608	0510	ADD	А,8
9090	67	0520	LD	н,А
9091	C1	0530	POP	BC
9092	ØD	0540	DEC	С
9093	20E1	0550	JR	NZ,LINE2
9095	7C	0560	LD	А,Н
9096	D640	0570	SUB	64
9098	67	0580	LD	н,А
9099	7D	0590	LD	A,L
9C9A	C650	0600	ADD	A,80
909 0	6F	0610	LD	L,A
9C9D	300A	0620	JR	NC,DONE2
9C9F	24	0630	INC	Н
9CAØ	7C	0640	LD	А,Н
9CA1	E607	0650	AND	7
9CA3	2004	0660	JR	NZ,DONE2
9CA5	7C	0670	LD	А,Н

9CA6	D608	0680		SUB	8
9CA8	67	0690		LD	Н,А
9CA9	15	0700	DONE2	DEC	D
90AA	20C8	0710		JR	NZ,ROW2
9CAC	C9	0720		RET	
9CAD	05	0730	MODE 1	DEC	В
9CAE	D5	0740	ROW1	PUSH	DE
9CAF	0E08	0750		LD	с,8
9CB1	C5	0760	LINE1	PUSH	BC
9CB2	E5	0770		PUSH	HL
9CB3	CB3E	0780		SRL	(HL)
9CB5	5D	0790	BYTE1	LD	E,L
9CB6	54	0800		LD	D,H
9CB7	7D	0810		LD	A,L
9CB8	2D	0820		DEC	L
9CB9	B7	0830		OR	А
9CBA	200A	Ø84Ø		JR	NZ,0K1
9CBC	7C	0850		LD	А,Н
9CBD	25	0860		DEC	н
9CBE	E607	0870		AND	7
9000	2004	0880		JR	NZ,0K1
9002	7C	0890		LD	А,Н
9003	C608	0900		ADD	А,8
9005	67	0910		LD	Н,А
9006	1A	0920	0K1	LD	A,(DE)
9007	CB3E	0930		SRL	(HL)
9009	3804	0940		JR	C,SETBIT

9CCB	CB9F	0950		RES	3,A
9CCD	1802	0960		JR	BITOK
9CCF	CBDF	0970	SETBIT	SET	З,А
9CD1	CB5E	0980	BITOK	BIT	3,(HL)
9CD3	2802	0990		JR	Z,BITSET
9CD5	CBFF	1000		SET	7,A
9CD7	12	1010	BITSET	LD	(DE),A
9CD8	10DB	1020		DJNZ	BYTE1
9CDA	CB9E	1030		RES	3,(HL)
9CDC	E 1	1040		POP	HL
9CDD	7C	1050		LD	А,Н
9CDE	C608	1060		ADD	А,8
9CE0	67	1070		LD	H,A
9CE1	C1	1080		POP	BC
9CE2	0D	1090		DEC	ç
9CE3	2000	1100		JR	NZ,LINE1
9CE5	7C	1110		LD	А,Н
9CE6	D640	1120		SUB	64
9CE8	67	1130		LD	н,А
9CE9	7D	1140		LD	A,L
9CEA	C650	1150		ADD	A,80
9CEC	6F	1160		LD	L,A
9CED	300A	1170		JR	NC,DONE1
9CEF	24	1180		INC	н
9CF0	7C	1190		LD	А,Н
9CF1	E607	1200		AND	7
9CF3	2004	1210		JR	NZ,DONE1

9CF5	7C	1220		LD	А,Н
9CF6	D608	1230		SUB	8
9CF8	67	1240		LD	Н,А
90F9	D1	1250	DONE 1	POP	DE
9CFA	15	1260		DEC	D
9CFB	20B1	1270		JR	NZ,ROW1
9CFD	C9	1280		RET	
9CFE	05	1290	MODEØ	DEC	в
9CFF	D5	1300	ROWØ	PUSH	DE
9D00	0E08	1310		LD	с,8
9D02	C5	1320	LINEØ	PUSH	BC
9D03	E5	1330		PUSH	HL
9D04	C5	1340	BYTE0	PUSH	BC
9D05	5D	1350		LD	E,L
9D06	54	1360		LD	D,H
9D07	7D	1370		LD	A,L
9D08	2D	1380		DEC	L
9D09	B7	1390		OR	Α
9DØA	200A	1400		JR	NZ,0KØ
9D0C	7C	1410		LD	А,Н
9DØD	25	1420		DEC	н
9D0E	E607	1430		AND	7
9D10	2004	1440		JR	NZ,0K0
9D12	7C	1450		LD	А,Н
9D13	C608	1460		ADD	Α,8
9D15	67	1470		LD	H,A

9D16	7E	1480	0KØ	LD	A,(HL)
9D17	17	1490		RLA	
9D18	4F	1500		LD	C,A
9D19	1 A	1510		LD	A,(DE)
9D1A	1 F	1520		RRA	
9D1B	1F	1530		RRA	
9D1C	0604	1540		LD	В,4
9D1E	1F	1550	NXTROT	RRA	
9D1F	CB39	1560		SRL	С
9D21	CB39	1570		SRL	С
9D23	1F	1580		RRA	
9D24	10F8	1590		DJNZ	NXTROT
9D26	12	1600		LD	(DE),A
9D27	C1	1610		POP	BC
9D28	10DA	1620		DJNZ	ΒΥΤΕØ
9D2A	7E	1630		LD	A,(HL)
9D2B	1 F	1640		RRA	
9D2C	0604	1650		LD	в,4
9D2E	ØF	1660	LSTBYT	RRCA	
9D2F	CB3F	1670		SRL	А
9D31	10FB	1680		DJNZ	LSTBYT
9D33	77	1690		LD	(HL),A
9D34	E 1	1700		POP	HL
9D35	7C	1710		LD	А,Н
9D36	C608	1720		ADD	Α,8
9D38	67	1730		LD	н,А
9D39	C1	1740		POP	BC

9D3A	ØD	1750	DEC	С
9D3B	2005	1760	JR	NZ,LINEØ
9D3D	7C	1770	LD	А,Н
9D3E	D640	1780	SUB	64
9D40	67	1790	LD	H,A
9D41	7D	1800	LD	A,L
9D42	C650	1810	ADD	A,80
9D44	6F	1820	LD	L,A
9D45	300A	1830	JR	NC,DONE0
9D47	24	1840	INC	н
9D48	7C	1850	LD	А,Н
9D49	E607	1860	AND	7
9D4B	2004	1870	JR	NZ,DONEØ
9D4D	7C	1880	LD	А,Н
9D4E	D608	1890	SUB	8
9D50	67	1900	LD	Н,А
9D51	D1	1910 DONEØ	POP	DE
9D52	15	1920	DEC	D
9D53	20AA	1930	JR	NZ,ROWØ
9D55	C9	1940	RET	
		1950	END	

Eleven – RSX Chords

The Amstrad is equipped with a marvellous feature, which can be used to create new commands which are then accepted as BASIC keywords. This feature is known as RSX and this routine uses this facility to create 24 new commands. Each of these commands plays a particular chord through the Amstrad's sound channels.

A chord, as I'm sure you know, is a combination of several notes in harmony, which are played simultaneously. Typing CALL 40000 will give you access to 24 of the more popular chords listed below (note that there are no spaces between parts of the new commands and that '#' equals sharp and b. equals flat):

New Command	Chord	
CMAJOR	C Major	
CMINOR	C Minor	
CSMAJOR	C# Major or Db. Ma;	jor
CSMINOR	C# Minor or Db. Min	nor
DMAJOR	D Major	
DMINOR	D Minor	
DSMAJOR	D# Major or Eb. Ma;	jor
DSMINOR	D# Minor or Eb. Min	nor

EMAJOR	E Major
EMINOR	E Minor
FMAJOR	F Major
FMINOR	F Minor
FSMAJOR	F# Major or Gb. Majoı
FSMINOR	F# Minor or Gb. Mino:
GMAJOR	G Major
GMINOR	G Minor
GSMAJOR	G# Major or Ab. Majoı
GSMINOR	G# Minor or Ab. Mino:
AMAJOR	A Major
AMINOR	A Minor
ASMAJOR	A# Major or Bb. Major
ASMINOR	A# Minor or Bb. Minor
BMAJOR	B Major
BMINOR	B Minor

All of the chords are made up of three notes, and use the three Amstrad sound channels to produce the chord.
To get a chord, you simply type in one of the commands. It is a good idea to put the volume control at around three quarters of maximum as with all multi-channel sound on the Amstrad, the sound output is a bit too powerful for the built-in speaker to handle without creating a fair amount of distortion.

CHORDS 1 7 10 SYMBOL AFTER 256:MEMORY 39999:SYMBOL AFTER 240 20 FOR n=40000 TO 40617 30 READ a\$:POKE n.VAL("&"+a\$) 40 NEXT 50 DATA 01,4A,9C,21,94,9C,CD,D1,BC,C9 60 DATA 98,90,03,33,9D,03,37,9D,03,3B 70 DATA 9D,C3,3F,9D,C3,43,9D,C3,47,9D 80 DATA C3.4B.9D.C3.4F.9D.C3.53.9D.C3 90 DATA 57,9D,C3,5B,9D,C3,5F,9D,C3,63 100 DATA 9D,C3,67,9D,C3,6B,9D,C3,6F,9D 110 DATA C3.73,9D,C3,77,9D,C3,7B,9D,C3 120 DATA 7F,9D,C3,83,9D,C3,87,9D,C3,8B 130 DATA 9D.C3.8F.9D.00.00.00.00.43.4D 140 DATA 41,4A,4F,D2,43,4D,49,4E,4F,D2 150 DATA 43,53,4D,41,4A,4F,D2,43,53,4D 160 DATA 49.4E.4F.D2.44.4D.41.4A.4F.D2 170 DATA 44,4D,49,4E,4F,D2,44,53,4D,41 180 DATA 4A.4F.D2.44.53.4D.49.4E.4F.D2 190 DATA 45,4D,41,4A,4F,D2,45,4D,49,4E 200 DATA 4F, D2, 46, 4D, 41, 4A, 4F, D2, 46, 4D

210	DATA	49,4E,4F,D2,46,53,4D,41,4A,4F
220	DATA	D2,46,53,4D,49,4E,4F,D2,47,4D
230	DATA	41,4A,4F,D2,47,4D,49,4E,4F,D2
240	DATA	47,53,4D,41,4A,4F,D2,47,53,4D
250	DATA	49,4E,4F,D2,41,4D,41,4A,4F,D2
260	DATA	41,4D,49,4E,4F,D2,41,53,4D,41
270	DATA	4A,4F,D2,41,53,4D,49,4E,4F,D2
280	DATA	42,4D,41,4A,4F,D2,42,4D,49,4E
290	DATA	4F,D2,00,1E,00,18,5C,1E,06,18
300	DATA	58,1E,0C,18,54,1E,12,18,50,1E
310	DATA	18,18,4C,1E,1E,18,48,1E,24,18
320	DATA	44,1E,2A,18,40,1E,30,18,3C,1E
330	DATA	36,18,38,1E,3C,18,34,1E,42,18
340	DATA	30,1E,48,18,2C,1E,4E,18,28,1E
350	DATA	54,18,24,1E,5A,18,20,1E,60,18
360	DATA	1C,1E,66,18,18,1E,6C,18,14,1E
370	DATA	72,18,10,1E,78,18,0C,1E,7E,18
380	DATA	08,1E,84,18,04,1E,8A,18,00,21
390	DATA	1A,9E,16,00,19,5E,23,56,23,ED
400	DATA	53,02,9E,5E,23,56,23,ED,53,ØB
410	DATA	9E,5E,23,56,ED,53,14,9E,F5,B7
420	DATA	28,03,DD,7E,00,32,00,9E,32,09
430	DATA	9E,32,12,9E,B7,28,12,CD,C2,BC
440	DATA	D0,23,23,7E,B7,28,08,CB,7F,20
450	DATA	04,3E,00,18,02,3E,0F,32,05,9E
460	DATA	32,0E,9E,32,17,9E,F1,FE,02,CC

```
470 DATA A7.BC.21.FF.9D.CD.AA.BC.30,FB
480 DATA 21,08,9E,CD,AA,BC,30,FB,21,11
490 DATA 9E.CD.AA.BC.30,FB.C9.31,00,00
500 DATA 00.00.00.00.00.00.2A.00.00.00
510 DATA 00.00.00.00.00.1C.00.00.00.00
520 DATA 00,00,00,00,7E,02,DE,01,7B,01
530 DATA 7E,02,DE,01,92,01,5A,02,C3,01
540 DATA 66.01.5A.02.C3.01.7B.01.A4.02
550 DATA 38,02,AA,01,38,02,AA,01,66,01
560 DATA 7E.02.18.02.92.01.A4.02.18.02
570 DATA 92,01,5A,02,FA,01,7B,01,7E,02
580 DATA FA.01.7B.01.38.02.DE.01.66.01
590 DATA 5A,02,DE,01,66,01,A4,02,18,02
600 DATA C3.01,A4.02,38.02,C3.01,7E,02
610 DATA FA,01,AA,01,7E,02,18,02,AA,01
620 DATA 5A,02,DE,01,92,01,5A,02,FA,01
630 DATA 92,01,38,02,C3,01,7B,01,38,02
640 DATA DE.01.7B.01.18.02.AA.01.66.01
650 DATA 18,02,C3,01,66,01,A4,02,FA,01
660 DATA 92,01,A4,02,FA,01,AA,01
```

Envelopes

Let's move onto the more complex part of using these chords, using them with sound envelopes. If you 'play' the basic chord without additional parameters, you get a basic "organ-like" chord. The routine automatically sets the command to envelope \emptyset which makes each chord last for about two seconds.

It is possible to create your own defined envelopes which can be incorporated within the chord sound. To do this, set up a volume envelope as normal (if you are not familiar with the Amstrad's sound features then check the user manual). Next, assign the envelope number to the end of the chord command. For example, if your envelope was ENV 1,3,5,1,15,-1,10 with the first number corresponding to the envelope's number then you would add '1' to the end of the chord command, i.e. |CMAJOR, 1.

The initial volume for the chord is dependent upon the envelope being used. It is assessed by the program in the following way. If the step size in the first section of the envelope being used is positive (excluding \emptyset) then the initial volume is set to \emptyset . If it is negative then the initial volume is set to its maximum value, 15. This gives the envelope the greatest volume range to work with.

Using an envelope step size of \emptyset can make the chord play for as long as you wish. For example:

ENV 2,10,0,100 |CMAJOR,2

This plays a C Major chord for ten seconds while ENV $2,n,\emptyset,1\emptyset\emptyset$ would play the same chord for n seconds.

When the commands are executed the routine will wait until all three of the sound buffers are available before adding the notes that make up the chord to the sound queues. If you want the sound queues to be cleared and the chord to sound immediately when the command is executed then add an extra parameter **before** the envelope number, e.g.

In this example, only the D Major chord would sound as the first chord would be cut off as soon as it began. Any value can be used for this extra first parameter. The routine only scans for the presence of an extra parameter, and not for the value of that parameter.

Obviously, you may not be happy at the thought of losing your first chord immediately. You can get around this simply by setting up a delay loop between each chord. An example is given below:

> ENV 5,8,9,199 |EMAJOR,9,5 FOR T=1 TO 399:NEXT |CMAJOR 9,5 FOR T=1 TO 259:NEXT |DMAJOR,9,5

		0001	;		CHORDS
		0002	;		
9040		0010		ORG	40000
BCA7		<i>0</i> 020	SNDRES	DEFL	ØBCA7H
BCAA		0030	ADDSND	DEFL	ØBCAAH
BCC2		<i>0</i> 040	AMPADR	DEFL	ØBCC2H
BCD1		0050	LOGEXT	DEFL	ØBCD1H
9040	014A9C	0060		LD	BC,TABLE
9043	219490	0070		LD	HL,SPACE
9046	CDD1BC	0080		CALL	LOGEXT
9049	C9	0090		RET	
9C4A	989C	0100	TABLE	DEFW	NAMETB
9C4C	C3339D	0110		JP	CMAJ
9C4F	C3379D	0120		JP	CMIN
9052	C33B9D	0130		JP	CSMAJ

9055	C33F9D	0140		JP	CSMIN
9058	C3439D	0150		JP	DMAJ
9C5B	C3479D	0160		JP	DMIN
9C5E	C34B9D	0170		JP	DSMAJ
9061	C34F9D	0180		JP	DSMIN
9064	C3539D	0190		JP	EMAJ
9067	C3579D	0200		JP	EMIN
9C6A	C35B9D	0210		JP	FMAJ
9C6D	C35F9D	0220		JP	FMIN
9070	C3639D	0230		JP	FSMAJ
9073	C3679D	0240		JP	FSMIN
9076	C36B9D	0250		JP	GMAJ
9079	C36F9D	0260		JP	GMIN
9070	C3739D	0270		JP	GSMAJ
907F	C3779D	0280		JP	GSMIN
9082	C37B9D	0290		JP	AMAJ
9085	C37F9D	0300		JP	AMIN
9088	C3839D	0310		JP	ASMAJ
9C8B	C3879D	0320		JP	ASMIN
9C8E	C38B9D	0330		JP	BMAJ
9091	C38F9D	0340		JP	BMIN
9094	00	0350	SPACE	DEFB	0,0,0,0
	00 00 00				
9098		0360	NAMETB	DEFM	"CMAJO"
9C9D	D2	0370		DEFB	"R"+80H
9C9E		0380		DEFM	"CMINO"

9CA3	D2	0390	DEFB	"R"+80H
9CA4		0400	DEFM	"CSMAJO"
9CAA	D2	0410	DEFB	"R"+80H
9CAB		0420	DEFM	"CSMINO"
9CB1	D2	0430	DEFB	"R"+80H
9CB2		0440	DEFM	"DMAJO"
9CB7	D2	0450	DEFB	"R"+80H
9CB8		0460	DEFM	"DMINO"
9CBD	D2	0470	DEFB	"R"+80H
9CBE		0480	DEFM	"DSMAJO"
9004	D2	0490	DEFB	"R"+80H
9005		0500	DEFM	"DSMINO"
9CCB	D2	0510	DEFB	"R"+80H
9000		0520	DEFM	"EMAJO"
9CD1	D2	0530	DEFB	"R"+80H
9CD2		0540	DEFM	"EMINO"
9CD7	D2	0550	DEFB	"R"+80H
9CD8		0560	DEFM	"FMAJO"
9CDD	D2	0570	DEFB	"R"+80H
9CDE		0580	DEFM	"FMINO"
9CE3	D2	0590	DEFB	"R"+80H
9CE4		0600	DEFM	"FSMAJO"
9CEA	D2	0610	DEFB	"R"+80H
9CEB		0620	DEFM	"FSMINO"
9CF1	D2	0630	DEFB	"R"+80H
9CF2		0640	DEFM	"GMAJO"
9CF7	D2	0650	DEFB	"R"+80H

9CF8		0660		DEFM	"GMINO"
9CFD	D2	0670		DEFB	"R"+80H
9CFE		0680		DEFM	"GSMAJO"
9D04	D2	0690		DEFB	"R"+80H
9D05		0700		DEFM	"GSMINO"
9D0B	D2	0710		DEFB	"R"+80H
9D0C		0720		DEFM	"AMAJO"
9D11	D2	0730		DEFB	"R"+80H
9D12		0740		DEFM	"AMINO"
9D17	D2	0750		DEFB	"R"+80H
9D18		0760		DEFM	"ASMAJO"
9D1E	D2	0770		DEFB	"R"+80H
9D1F		0780		DEFM	"ASMINO"
9D25	D2	0790		DEFB	"R"+80H
9D26		0800		DEFM	"BMAJO"
9D2B	D2	0810		DEFB	"R"+80H
9D2C		0820		DEFM	"BMINO"
9D31	D2	0830		DEFB	"R"+80H
9D32	00	0840		DEFB	Ø
9D33	1E00	0850 C	CMAJ	LD	Ε,0
9D35	185C	0860		JR	CHORD
9D37	1E06	0870 C	MIN	LD	Ε,6
9D39	1858	0880		JR	CHORD
9D3B	1E0C	0890 C	SMAJ	LD	E,12
9D3D	1854	0900		JR	CHORD
9D3F	1E12	0910 C	SMIN	LD	E,18
9D41	1850	0920		JR	CHORD

9D43	1E18	0930	DMAJ	LD	E,24
9D45	184C	0940		JR	CHORD
9D47	1E1E	0950	DMIN	LD	Е,30
9D49	1848	0960		JR	CHORD
9D4B	1E24	0970	DSMAJ	LŊ	Е,36
9D4D	1844	0980		JR	CHORD
9D4F	1E2A	0990	DSMIN	LD	E,42
9D51	1840	1000		JR	CHORD
9D53	1E30	1010	EMAJ	LD	E,48
9D55	183C	1020		JR	CHORD
9D57	1E36	1030	EMIN	LD	E,54
9D59	1838	1040		JR	CHORD
9D5B	1E3C	1050	FMAJ	LD	E,60
9D5D	1834	1060		JR	CHORD
9D5F	1E42	1070	FMIN	LD	Е,66
9D61	1830	1080		JR	CHORD
9D63	1E48	1090	FSMAJ	LD	E,72
9D65	182C	1100		JR	CHORD
9D67	1E4E	1110	FSMIN	LD	E,78
9D69	1828	1120		JR	CHORD
9D6B	1E54	1130	GMAJ	LD	E,84
9D6D	1824	1140		JR	CHORD
9D6F	1E5A	1150	GMIN	LD	Е,90
9D71	1820	1160		JR	CHORD
9D73	1E60	1170	GSMAJ	LD	Е,96
9D75	181C	1180		JR	CHORD

9077	1566	1190	GOMIN	LD	E 100
9079	1010	1000	CONTR	10	CUCOD
7077	1818	1200		ЛК	CHURD
9D7B	1E6C	1210	AMAJ	LD	E,108
9D7D	1814	1220		JR	CHORD
9D7F	1E72	1230	AMIN	LD	E,114
9D81	1810	1240		JR	CHORD
9D83	1E78	1250	ASMAJ	LD	E,120
9D85	180C	1260		JR	CHORD
9D87	1E7E	1270	ASMIN	LD	E,126
9D89	1808	1280		JR	CHORD
9D8B	1E84	1290	BMAJ	LD	E,132
9D8D	1804	1300		JR	CHORD
9D8F	1E8A	1310	BMIN	LD	E,138
9D91	1800	1320		JR	CHORD
9D93	211A9E	1330	CHORD	LD	HL,DATA
9D96	1600	1340		LD	D,Ø
9D98	19	1350		ADD	HL,DE
9D99	5E	1360		LD	E,(HL)
9D9A	23	1370		INC	HL
9D9B	56	1380		LD	D,(HL)
9 D 9C	23	1390		INC	HL
9D9D	ED53029E	1400		LD	(TONE1),DE
9DA1	5E	1410		LD	E,(HL)
9DA2	23	1420		INC	HL
9DA3	56	1430		LD	D,(HL)
9DA4	23	1440		INC	HL

9DA5	ED530B9E	1450		LD	(TONE2),DE
9DA9	5E	1460		LD	E,(HL)
9DAA	23	1470		INC	HL
9DAB	56	1480		LD	D,(HL)
9DAC	ED53149E	1490		LD	(TONE3),DE
9DBØ	F5	1500		PUSH	AF
9DB1	B7	1510		OR	А
9DB2	2803	1520		JR	Z,SETENV
9DB4	DD7E00	1530		LD	A,(IX+0)
9DB7	32009E	1540	SETENV	LD	(ENV1),A
9DBA	32099E	1550		LD	(ENV2),A
9DBD	32129E	1560		LD	(ENV3),A
9DCØ	B7	1570		OR	А
9DC1	2812	1580		JR	Z,NEGTVE
9DC3	CDC2BC	1590	FNDENV	CALL	AMPADR
9DC6	DØ	1600		RET	NC
9DC7	23	1610		INC	HL
9DC8	23	1620		INC	HL
9DC9	7E	1630		LD	A,(HL)
9DCA	B7	1640		OR	А
9DCB	2808	1650		JR	Z,NEGTVE
9DCD	CB7F	1660		BIT	7,A
9DCF	2004	1670		JR	NZ,NEGTVE
9DD1	3E00	1680		LD	Α,0
9DD3	1802	1690		JR	SETAMP
9DD5	3E0F	1700	NEGTVE	LD	A,15
9DD7	32059E	1710	SETAMP	LD	(AMPL1),A

9DDA	320E9E	1720		LD	(AMPL2),A
9DDD	32179E	1730		LD	(AMPL3),A
9DEØ	F1	1740		POP	AF
9DE1	FE02	1750		СР	2
9DE3	CCA7BC	1760		CALL	Z,SNDRES
9DE6	21FF9D	1770		LD	HL, SND1
9DE9	CDAABC	1780	AFULL	CALL	ADDSND
9DEC	30FB	1790		JR	NC,AFULL
9DEE	21089E	1800		LD	HL, SND2
9DF 1	CDAABC	1810	BFULL	CALL	ADDSND
9DF4	30FB	1820		JR	NC,BFULL
9DF6	21119E	1830		LD	HL, SND3
9DF9	CDAABC	1840	CFULL	CALL	ADDSND
9DFC	30FB	1850		JR	NC,CFULL
9DFE	C9	1860		RET	
9DFF	31	1870	SND1	DEFB	49
9E00	00	1880	ENV1	DEFB	0,0
	00				
9E02	00	1890	TONE 1	DEFB	0,0,0
	00 00				
9E05	00	1900	AMPL1	DEFB	0,0,0
	00 00				
9E08	2A	1910	SND2	DEFB	42
9E09	00	1920	ENV2	DEFB	0,0
	<i>00</i>				
9EØB	00	1930	TONE2	DEFB	0,0,0
	00 00				

RSX Chords

9E0E	00 [°]		1940	AMPL2	DEFB	0,0,0
	00 00)				
9E11	1 C		1950	SND3	DEFB	28
9E12	00		1960	ENV3	DEFB	0,0
	00					
9E14	<i>00</i>		1970	TONE3	DEFB	0,0,0
	00 00)				
9E17	00		1980	AMPL3	DEFB	0,0,0
	00 00)				
9E1A	7E02		1990	DATA	DEFW	638,478,379
	DEØ1	7801				
9E20	7E02		2000		DEFW	638,478,402
	DEØ1	9201				
9E26	5A02		2010		DEFW	602,451,358
	C301	6601				
9E2C	5A02		2020		DEFW	602,451,379
	C301	7801				
9E32	A402		2030		DEFW	676,568,426
	3802	AA@1				
9E38	3802		2040		DEFW	568,426,358
	AA01	6601				
9E3E	7EØ2		2050		DEFW	638,536,402
	1802	9201				
9E44	A402		2060		DEFW	676,536,402
	1802	9201				
9E4A	5A02		2070		DEFW	602,506,379
	FA01	7801				

9E50	7E02	2080
	FA01	7801
9E56	3802	2090
	DEØ1	6601
9E5C	5A02	2100
	DEØ1	6601
9E62	A402	2110
	1802	C301
9E68	A402	2120
	3802	C301
9E6E	7EØ2	2130
	FA01	AAØ1
9E74	7E02	2140
	1802	AA@1
9E7A	5A02	2150
	DEØ1	9201
9E80	5A02	2160
	FAØ1	9201
9E86	3802	2170
	C301	7801
9E8C	3802	2180
	DEØ1	7801
9E92	1802	2190
	AA01	6601
9E98	1802	2200
	C301	6601

- DEFW 638,506,379
- DEFW 568,478,358
- DEFW 602,478,358
- DEFW 676,536,451
- DEFW 676,568,451
- DEFW 638,506,426
- DEFW 638,536,426
- DEFW 602,478,402
- DEFW 602,506,402
- DEFW 568,451,379
- DEFW 568,478,379
- DEFW 536,426,358
- DEFW 536,451,358

9E9E	A402	2210	DEFW	676,506,402
	FA01	9201		
9EA4	A402	2220	DEFW	676,506,426
	FAØ1	AAØ1		
		2230	END	

Chord Organ

We couldn't leave this routine without providing a chord organ program. It is not, however, just a matter of using lines like IF INKEY\$="A" THEN |AMAJOR, as there are other matters such as interrupts and envelopes which have to be controlled. Our program plays all the major chords, natural notes, sharps and flats.

Here is the 'keyboard':

2 3 5 6 7 Q W E R T Y U I

And these are the notes they represent:

C# D# F# G# A# C D E F G A B C 10 ENV 1,10,0,100 20 DIM k(12):m=-130 FOR n=0 TO 12:READ k(n):NEXT40 n=0

```
50 a=INKEY(k(n))
60 IF a=0 AND n=m GOTO 40
70 IF a=0 GOTO 120
80 n=n+1
90 IF n<13 GOTO 50
100 SOUND 135,0,0,0:m=-1
110 GOTO 40
120 ON n+1 GOSUB 140,150,160,170,180,190
,200,210,220,230,240,250,260
130 m=n:GOTO 40
140 CMAJOR.0.1:RETURN
150 [DMAJOR,0,1:RETURN
160 'EMAJOR.0.1:RETURN
170 [FMAJOR,0,1:RETURN
180 [GMAJOR,0,1:RETURN
190 ; AMAJOR.0.1:RETURN
200 BMAJOR.0.1:RETURN
210 CMAJOR.0.1:RETURN
220 CSMAJOR.0.1:RETURN
230 [DSMAJOR.0.1:RETURN
240 ;FSMAJOR,0,1:RETURN
250 GSMAJOR.0.1:RETURN
260 ; ASMAJOR, 0, 1: RETURN
270 DATA 67,59,58,50,51,43,42,35,65,57,4
9,48,41
```

Twelve – Screen Compactors

Compactor.1

The Amstrad's screen is stored in 16K of memory. This can be a lot of RAM to store what is often a simple screen display. Much of a screen is usually blank and has a value of \emptyset . A blank section of the screen which is $2\emptyset$ pixels is represented in the screen RAM by twenty zeros. This is a lot of memory wasted. Our compactor routine makes much better use of the screen memory.

This program works by first looking at a section of screen RAM. If it has a value other than \emptyset , then it stores that value as it is without alteration. If the value is \emptyset , it counts how many zeros follow and then stores one zero and a second value equalling the number of zeros.

1 ' COMPACT 10 SYMBOL AFTER 256:MEMORY 20000:SYMBOL AFTER 240 20 FOR n=43800 TO 43888 30 READ a\$:POKE n,VAL("&"+a\$) 40 NEXT 50 DATA 11,00,C0,21,17,AB,7E,12,13,2B 60 DATA CB,7A,C3,B7,20,F6,46,2B,05,28 70 DATA F1,12,13,10,FC,CB,7A,20,E9,C9 80 DATA FE,01,C0,11,00,C0,21,17,AB,1A 90 DATA 77,13,2B,CB,7A,28,18,B7,20,F5 100 DATA 06,01,1A,B7,20,08,13,CB,7A,28 110 DATA 07,04,20,F4,70,2B,18,E3,04,70 120 DATA 2B,EB,21,17,AB,B7,ED,52,EB,DD 130 DATA 6E,00,DD,66,01,73,23,72,C9

To compact a screen, use the following:

The number of bytes taken up by the screen is returned to the A% variable, so that you can check on the memory used and thus calculate the approximate memory saving. (This also applies to the second compactor routine following this one.)

To de-compact and display a previously compacted screen, type:

CALL 438ØØ

To save the screen, firstly compact it in the normal way and then type:

```
SAVE "NAME", B,43800-A%, A%+30
```

To load and display, simply LOAD "" and then CALL 43800.

When saving the screen, you have to save the de-compactor routine as well.

	0001 ;		COMPACT1
	0002 ;		
AB18	0010	ORG	43800
	0020 ;		ROUTINE1

AB18	110000	0030		LD	DE,0C000H
AB1B	2117AB	0040		LD	HL,43799
AB1E	7E	0050	UNCMPT	LD	A,(HL)
AB1F	12	0060		LD	(DE),A
AB20	13	0070		INC	DE
AB21	2B	0080		DEC	HL
AB22	CB7A	0090		BIT	7,D
AB24	C8	0100		RET	Z
AB25	B7	0110		OR	A
AB26	20F6	0120		JR	NZ,UNCMPT
AB28	46	0130		LD	B,(HL)
AB29	2B	0140		DEC	HL
AB2A	05	0150		DEC	В
AB2B	28F1	0160		JR	Z,UNCMPT
AB2D	12	0170	REPEAT	LD	(DE),A
AB2E	13	0180		INC	DE
AB2F	10FC	0190		DJNZ	REPEAT
AB31	CB7A	0200		BIT	7,D
AB33	20E9	0210		JR	NZ,UNCMPT
AB35	C9	0220		RET	
		0230	;		ROUTINE2
AB36	FEØ1	0240		СР	1
AB38	CØ	0250		RET	NZ
AB39	110000	0260		LD	DE,0C000H
AB3C	2117AB	0270		LD	HL,43799
AB3F	1 A	0280	COMPCT	LD	A,(DE)
AB40	77	0290		LD	(HL),A

AB41	13	0300		INC	DE
AB42	2B	0310		DEC	HL
AB43	CB7A	0320		BIT	7,D
AB45	2818	0330		JR	Z,TOTAL
AB47	B7	0340		OR	А
AB48	20F5	0350		JR	NZ,COMPCT
AB4A	0601	0360		LD	B,1
AB4C	1 A	0370	COUNT	LD	A,(DE)
AB4D	B7	0380		OR	А
AB4E	2008	0390		JR	NZ,GOTLEN
AB50	13	0400		INC:	DE
AB51	CB7A	0410		BIT	7,D
AB53	2807	0420		JR	Z,LAST
AB55	04	0430		INC	в
AB56	20F4	0440		JR	NZ,COUNT
AB58	70	0450	GOTLEN	LD	(HL),B
AB59	2B	Ø46Ø		DEC	HL
AB5A	18E3	0470		JR	COMPCT
AB5C	04	0480	LAST	INC	в
AB5D	70	0490		LD	(HL),B
AB5E	2B	0500		DEC	HL
AB5F	EB	0510	TOTAL	ΕX	DE,HL
AB60	2117AB	0520		LD	HL,43799
AB63	87	0530		OR	А
AB64	ED52	0540		SBC	HL,DE
AB66	EB	0550		ЕX	DE,HL

AB67	DD6E00	0560	LD	L,(IX+0)
AB6A	DD6601	0570	LD	H,(IX+1)
AB6D	73	0580	LD	(HL),E
AB6E	23	0590	INC	HL
AB6F	72	0600	LD	(HL),D
AB70	C9	0610	RET	
		0620	END	

Compactor.2

Unlike the previous routine, this one compacts **all** multiples of the same value and not just zeros. If, for example, a portion of the screen read as:

3, 160, 2, 2, 2, 2, 2, 2, 8, 8, 8, 8

This routine stores it as:

3, 1, 160, 1, 2, 6, 8, 4

The major drawback with this routine is that if a screen is too complicated then it doesn't compact at all and may take even more than the 16K used for storing the screen in normal circumstances. However, as it will work with most screens, you should always give it a try first.

1 ' COMPACT
10 SYMBOL AFTER 256:MEMORY 20000:SYMBOL
AFTER 240
20 FOR n=43800 TO 43875
30 READ a\$:POKE n,VAL("&"+a\$)

40 NEXT

50 DATA 11,00,C0,21,17,AB,7E,2B,46,2B 60 DATA 12,13,10,FC,CB,7A,20,F4,C9,FE 70 DATA 01,C0,11,00,C0,21,17,AB,1A,77 80 DATA 13,2B,06,01,CB,7A,28,12,4F,1A 90 DATA B9,20,08,13,CB,7A,28,07,04,20 100 DATA F4,70,2B,18,E5,04,70,2B,EB,21 110 DATA 17,AB,B7,ED,52,EB,DD,6E,00,DD 120 DATA 66,01,73,23,72,C9

To compact a screen:

A%=Ø:CALL 43819, @A%

To de-compact a screen:

CALL 438ØØ

To save the screen and the uncompactor (which is needed for displaying the screen in its normal way) type:

SAVE "NAME", B, 43800-A%, A%+19

B is not a variable in this case of course, but is the sign telling the computer that you are saving a binary file.

	0001 ;		COMPACT2
	0002 ;		
AB18	0010	ORG	43800
	0020 ;		ROUTINE1

AB18	1100C0	0030		LD	DE,0C000H
AB1B	2117AB	0040		LD	HL,43799
AB1E	7E	0050	UNCMPT	LD	A,(HL)
AB1F	2B	0060		DEC	HL
AB20	46	0070		LD	B,(HL)
AB21	2B	0080		DEC	HL
AB22	12	0090	REPEAT	LD	(DE),A
AB23	13	0100		INC	DE
AB24	10FC	0110		DJNZ	REPEAT
AB26	CB7A	0120		BIT	7,D
AB28	20F4	0130		JR	NZ, UNCMPT
AB2A	C9	0140		RET	
		0150	;		ROUTINE2
AB2B	FE01	0160		СР	1
AB2D	CØ	0170		RET	NZ
AB2E	110000	0180		LD	DE,0C000H
AB31	2117AB	0190		LD	HL,43799
AB34	1A	0200	COMPCT	LD	A,(DE)
AB35	77	0210		LD	(HL),A
AB36	13	0220		INC	DE
AB37	2B	0230		DEC	HL
AB38	0601	0240		LD	B,1
АВЗА	CB7A	0250		BIT	7,D
AB3C	2812	0260		JR	Z,TOTAL
AB3E	4F	0270		LD	с,А
AB3F	1 A	0280	COUNT	LD	A,(DE)
AB40	B9	0290		СР	С

AB41	2008	0300		JR	NZ,GOTLEN
AB43	13	0310		INC	DE
AB44	CB7A	0320		BIT	7,D
AB46	2807	0330		JR	Z,LAST
AB48	04	0340		INC	в
AB49	20F4	0350		JR	NZ,COUNT
AB4B	70	0360	GOTLEN	LD	(HL),B
AB4C	2B	0370		DEC	HL
AB4D	18E5	0380		JR	COMPCT
AB4F	04	0390	LAST	INC	в
AB50	70	0400	TOTAL	LD	(HL),B
AB51	2B	0410		DEC	HL
AB52	EB	0420		ΕX	DE,HL
AB53	2117AB	0430		LD	HL,43799
AB56	B7	0440		OR	А
AB57	ED52	0450		SBC	HL,DE
AB59	EB	0460		ЕX	DE,HL
AB5A	DD6E00	0470		LD	L,(IX+0)
AB5D	DD6601	0480		LD	H,(IX+1)
AB60	73	0490		LD	(HL),E
AB61	23	0500		INC	HL
AB62	72	0510		LD	(HL),D
AB63	C9	0520		RET	
		0530		END	

Thirteen – DOKE and DEEK

In machine code programming, you often have to POKE and PEEK 16-bit numbers. Whereas 8-bit numbers have a range of 256, 16-bit numbers have a range of 65536. 16-bit numbers use two bytes and to POKE a two byte number, you have to use a little formula:

POKE byte 1 + 256*byte 2

This also applies to PEEKing two bytes of a memory location. Some computers are fortunate enough to have Double-PEEK and Double-POKE commands, known as DEEK and DOKE. This routine provides you with two new BASIC commands, |DOKE and |DEEK.

```
1 ' DOKE/DEEK
10 SYMBOL AFTER 256:MEMORY 39999:SYMBOL
AFTER 240
20 FOR n=40000 TO 40071
30 READ a$:POKE n,VAL("&"+a$)
40 NEXT
50 DATA 01,4A,9C,21,52,9C,CD,D1,BC,C9
60 DATA 56,9C,C3,5F,9C,C3,72,9C,00,00
70 DATA 00,00,44,4F,4B,C5,44,45,45,CB
80 DATA 00,FE,02,C0,DD,6E,02,DD,66,03
90 DATA DD,7E,00,77,23,DD,7E,01,77,C9
```

```
100 DATA FE,02,C0,DD,6E,02,DD,66,03,DD
110 DATA 5E,00,DD,56,01,7E,23,12,13,7E
120 DATA 12,C9
```

You must first CALL 40000 to set up the new commands. To use |DOKE you need to follow it with a valid memory address and a 16-bit number. Both of these values should be in decimal. |DOKE, 30000,343 would split the value 343 up using the reverse of the formula above and place the two values into memory locations 30000 and 30001.

DEEK is slightly more complicated. You use it in the format:

 $N\% = \emptyset$: DEEK, A, @N%

A stands for the first address to be PEEKed () DEEK will look at A and A+1) while N% contains the value returned by the |DEEK. N% must be defined before a |DEEK. As you can see from the percent sign, it must be an integer variable. Remember this if you want to print out the actual value of DEEK. If the value returned is negative, you need to print 65536+N% to get the true value.

		0001	;	DOKE	DEEK
		0002	;		
9040		0010		ORG	40000
BCD1		0020	LOGEXT	DEFL	ØBCD1H
9040	Ø14A9C	0030		LD	BC,TABLE
9043	215290	0040		LD	HL,SPACE
9046	CDD1BC	0050		CALL	LOGEXT
9049	C9	0060		RET	
904A	569C	0070	TABLE	DEFW	NAMETB

9C4C	C35F9C	0080		JP	DOKE
9C4F	C3729C	0090		JP	DEEK
9052	00	0100	SPACE	DEFB	0,0,0,0
	00 00 00				
9056		0110	NAMETB	DEFM	"DOK"
9059	C5	0120		DEFB	"E"+80H
9C5A		0130		DEFM	"DEE"
9C5D	СВ	0140		DEFB	"K"+80H
9C5E	00	0150		DEFB	0
9C5F	FEØ2	0160	DOKE	СР	2
9061	C0	0170		RET	NZ
9062	DD6E02	0180		LD	L,(IX+2)
9065	DD6603	0190		LD	H,(IX+3)
9068	DD7E00	0200		LD	A,(IX+0)
9C6B	77	0210		LD	(HL),A
9060	23	0220		INC	HL
9C6D	DD7EØ1	0230		LD	A,(IX+1)
9070	77	0240		LD	(HL),A
9071	C9	0250		RET	
9072	FE02	0260	DEEK	СР	2
9074	CØ	0270		RET	NZ
9075	DD6E02	0280		LD	L,(IX+2)
9078	DD6603	0290		LD	H,(IX+3)
9C7B	DD5E00	0300		LD	E,(IX+0)
9C7E	DD5601	0310		LD	D,(IX+1)
9081	7E	0320		LD	A,(HL)

9082	23	0330	INC	HL
9083	12	0340	LD	(DE),A
9084	13	0350	INC	DE
9085	7E	0360	LD	A,(HL)
9086	12	0370	LD	(DE),A
9C87	C9	0380	RET	
		0390	END	

Fourteen – The Games Writing Package

Finally in this book we have a program which gives you a selection of machine code routines all in one easily-entered package. The routines create new BASIC keywords, using RSX calls.

|EXPLODE, |READCHAR, |BIGPRINT, |USCROLL, |DSCROLL, |LSCROLL, |RSCROLL, |COMMANDS

Let's deal with each in turn. EXPLODE needn't be followed by any parameters. It simply creates an explosion sound of great use in games programs. READCHAR and BIGPRINT are two commands corresponding to the titles of routines elsewhere in this book. BIGPRINT creates double-size characters while READCHAR calculates what ASCII character is in a specified screen position.

USCROLL, DSCROLL, LSCROLL and RSCROLL are the four box scrolls featured in this routines section with the first letter standing for the direction of the scroll. COMMANDS simply gives a list of all the above available commands.

To understand how each feature works and the parameters which must follow the command word, you'll need to turn to the appropriate individual routine.

The features in this package offer most of the machine code routines needed to write a good "hybrid" game, partly in BASIC and partly in machine code.

```
1 7
          GAMES WRITING PACKAGE
10 SYMBOL AFTER 256:MEMORY 39999:SYMBOL
AFTER 240
20 DIM x(12)
30 FOR c=0 TO 12:READ a$:x(c)=VAL("&"+a$
):NFXT
40 \ r=0:sum=0
50 FOR n=40000 TO 41207
60 READ a$:v=VAL("&"+a$)
70 sum=sum+v:POKE n.v
80 IF (n+1-40000) MOD 100<>0 GOTO 110
90 IF sum<>x(c) THEN PRINT"*ERROR IN DAT
A*";CHR$(7):PRINT"Check lines";210+100*c
:"to":300+100*c:END
100 \text{ sum} = 0:c = c + 1
110 NEXT
120 IF sum<>x(c) THEN PRINT"*ERROR IN DA
TA*":CHR$(7):PRINT"Check line 1410"
130 '
140 '
                CHECKSUMS
150 '
160 DATA 2879,2F09,2A7A,27A3,263B,2434,2
C42
170 DATA 2782,2593,2A89,278C,29D3,04DD
180 '
```

190 MACHINE CODE

200 '

210 DATA 01.4A,9C.21.64.9C.CD.D1.BC.C9 220 DATA 68,9C,C3,A4,9C,C3,C3,9C,C3,D7 230 DATA 9C.C3.77.9D.C3.09.9E.C3.A1.9E 240 DATA C3, B1, 9F, C3, C7, A0, 00, 00, 00, 00 250 DATA 45,58,50,4C,4F,44,C5,52,45,41 260 DATA 44,43,48,41,D2,42,49,47,50,52 270 DATA 49,4E,D4,55,53,43,52,4F,4C,CC 280 DATA 44.53.43.52.4F.4C.CC.4C.53.43 290 DATA 52,4F,4C,CC,52,53,43,52,4F,4C 300 DATA CC,43,4F,4D,4D,41,4E,44,D3,00 310 DATA CD.A7.BC.3E.01.21.B6.9C.CD.BC 320 DATA BC,21,BA,9C,CD,AA,BC,C9,01,0F 330 DATA FF, 19, 01, 01, 00, 00, 00, 0F, 0F, 00 340 DATA 00,DD,6E,02,DD,66,04,CD,75,BB 350 DATA CD,60,BB,DD,6E,00,DD,66,01,77 360 DATA C9,CD,93,BB,F5,DD,6E,04,DD,66 370 DATA 05,46,23,5E,23,56,DD,6E,06,DD 380 DATA 66,08,C5,D5,E5,1A,47,CD,06,B9 390 DATA 78,CD,34,9D,47,CD,09,B9,E1,5D 400 DATA 54,CD,75,BB,DD,7E,02,CD,90,BB 410 DATA 78,CD,5A,BB,3C,CD,5A,BB,6B,62 420 DATA 2C,CD,75,BB,DD,7E,00,CD,90,BB 430 DATA 78,3C,3C,CD,5A,BB,3C,CD,5A,BB 440 DATA 6B,62,24,24,D1,13,C1,10,BD,F1 450 DATA CD, 90, BB, C9, CD, A5, BB, EB, CD, AE

460	DATA	BB,F5,0E,02,06,04,C5,1A,0F,0F
470	DATA	0F,0F,06,04,1F,CB,1E,CB,2E,10
480	DATA	F9,7E,23,77,06,07,23,10,FD,1A
490	DATA	06,04,1F,CB,1E,CB,2E,10,F9,7E
500	DATA	23,77,06,07,28,10,FD,13,C1,10
510	DATA	D3,06,08,23,10,FD,0D,20,C9,F1
520	DATA	C9,DD,6E,02,DD,66,06,2D,25,DD
530	DATA	7E,04,94,4F,DD,7E,00,95,57,CD
540	DATA	1A,BC,3E,00,81,10,FD,5F,18,38
550	DATA	D5,43,5D,7C,C6,38,57,7D,C6,50
560	DATA	6F,30,0A,24,7C,E6,07,20,04,7C
570	DATA	D6,08,67,E5,7E,12,1C,20,0A,14
580	DATA	7A,E6,07,20,04,7A,D6,08,57,2C
590	DATA	20,0A,24,7C,E6,07,20,04,7C,D6
600	DATA	08,67,10,E2,E1,D1,D5,E5,4B,E5
610	DATA	06,07,5D,54,7C,C6,08,67,7E,12
620	DATA	10,F6,E1,2C,20,0A,24,7C,E6,07
630	DATA	20,04,7C,D6,08,67,0D,20,E2,E1
640	DATA	D1,15,20,A2,7C,C6,38,67,43,36
650	DATA	00,2C,20,0A,24,7C,E6,07,20,04
660	DATA	7C,D6,08,67,10,EF,C9,DD,6E,00
670	DATA	DD,66,06,2D,25,DD,7E,04,94,4F
680	DATA	7D,C6,02,DD,96,02,57,CD,1A,BC
690	DATA	7C,C6,38,67,3E,00,81,10,FD,5F
700	DATA	18,38,D5,43,5D,7C,D6,38,57,7D
710	DATA	D6,50,6F,30,0A,7C,25,E6,07,20

720	DATA	04,7C,C6,08,67,E5,7E,12,1C,20
730	DATA	ØA,14,7A,E6,07,20,04,7A,D6,08
740	DATA	57,2C,20,0A,24,7C,E6,07,20,04
750	DATA	7C,D6,08,67,10,E2,E1,D1,D5,E5
760	DATA	4B,E5,06,07,5D,54,7C,D6,08,67
770	DATA	7E,12,10,F6,E1,2C,20,0A,24,7C
780	DATA	E6,07,20,04,7C,D6,08,67,0D,20
790	DATA	E2,E1,D1,15,20,A2,7C,D6,38,67
800	DATA	43,36,00,2C,20,0A,24,7C,E6,07
810	DATA	20,04,7C,D6,08,67,10,EF,C9,DD
820	DATA	6E,02,DD,66,06,2D,25,DD,7E,04
830	DATA	94,4F,DD,7E,00,95,57,CD,1A,BC
840	DATA	3E,00,81,10,FD,47,CD,11,BC,DA
850	DATA	5B,9F,28,48,DD,6E,02,DD,66,04
860	DATA	2D,25,C5,CD,1A,BC,C1,0E,08,C5
870	DATA	E5, B7, CB, 16, F5, 7D, 2D, B7, 20, 0A
880	DATA	7C,25,E6,07,20,04,7C,C6,08,67
890	DATA	F1,10,EB,E1,7C,C6,08,67,C1,0D
900	DATA	20,DF,7C,D6,40,67,7D,C6,50,6F
910	DATA	30,0A,24,7C,E6,07,20,04,7C,D6
920	DATA	08,67,15,20,C6,C9,05,D5,0E,08
930	DATA	C5,E5,CB,26,5D,54,2C,20,0A,24
940	DATA	7C,E6,07,20,04,7C,D6,08,67,1A
950	DATA	CB,26,38,04,CB,A7,18,02,CB,E7
960	DATA	CB,66,28,02,CB,C7,12,10,DD,CB
970	DATA	A6,E1,7C,C6,08,67,C1,0D,20,CE

980 DATA 7C.D6.40.67.7D.C6.50.6F.30.0A 990 DATA 24,7C,E6,07,20,04,7C,D6,08,67 1000 DATA D1,15,20,B3,C9,05,D5,0E,08,C5 1010 DATA E5,C5,5D,54,2C,20,0A,24,7C,E6 1020 DATA 07,20,04,7C,D6,08,67,7E,1F,4F 1030 DATA 1A.17.17.06.04.17.CB.21.CB.21 1040 DATA 17,10,F8,12,C1,10,DC,7E,17,06 1050 DATA 04,07,CB,27,10,FB,77,E1,7C,C6 1060 DATA 08,67,C1,0D,20,C7,7C,D6,40,67 1070 DATA 7D,C6,50,6F,30,0A,24,7C,E6,07 1080 DATA 20,04,7C.D6,08,67,D1,15,20,AC 1090 DATA C9,DD,6E,02,DD,66,04,2D,25,7C 1100 DATA C6,02,DD,96,06,4F,DD,7E,00,95 1110 DATA 57,CD.1A.BC.2B.3E.00.81.23.10 1120 DATA FC,47,CD,11,BC,DA,6F,A0,28,46 1130 DATA DD,6E,02,DD,66,06,2D,25,C5,CD 1140 DATA 1A, BC, C1, 0E, 08, C5, E5, B7, CB, 1E 1150 DATA F5,2C,20,0A,24,7C,E6,07,20,04 1160 DATA 7C, D6, 08, 67, F1, 10, ED, E1, 7C, C6 1170 DATA 08,67,C1,0D,20,E1,7C,D6,40,67 1180 DATA 7D,C6,50,6F,30,0A,24,7C,E6,07 1190 DATA 20,04,7C,D6,08,67,15,20,C8,C9 1200 DATA 05,D5,0E,08,C5,E5,CB,3E,5D,54 1210 DATA 7D,2D,87.20.0A.7C.25.E6.07.20 1220 DATA 04,7C,C6,08,67,1A,CB,3E,38,04 1230 DATA CB,9F,18,02,CB,DF,CB,5E,28,02 1240 DATA CB, FF, 12, 10, DB, CB, 9E, E1, 7C, C6

```
1250 DATA 08,67,C1,0D,20,CC,7C,D6,40,67
1260 DATA 7D.C6.50.6F.30.0A.24.7C.E6.07
1270 DATA 20,04,7C,D6,08,67,D1,15,20,B1
1280 DATA C9,05,D5,0E,08,C5,E5,C5,5D,54
1290 DATA 7D,2D,87,20,0A,7C,25,E6,07,20
1300 DATA 04,7C,C6,08,67,7E,17,4F,1A,1F
1310 DATA 1F,06,04,1F,CB,39,CB,39,1F,10
1320 DATA F8,12,C1,10,DA,7E,1F,06,04,0F
1330 DATA CB, 3F, 10, FB, 77, E1, 7C, C6, 08, 67
1340 DATA C1,0D,20,C5,7C,D6,40,67,7D,C6
1350 DATA 50,6F,30,0A,24,7C,E6,07,20,04
1360 DATA 7C, D6, 08, 67, D1, 15, 20, AA, C9, 21
1370 DATA 68,9C,3E,0D,CD,5A,BB,3E,0A,CD
1380 DATA 5A, BB, 3E, 7C, CD, 5A, BB, 7E, 23, CB
1390 DATA 7F,20,05,CD,5A,BB,18,F5,CB,BF
1400 DATA CD, 5A, BB, 3E, 0D, CD, 5A, BB, 3E, 0A
1410 DATA CD.5A,BB,7E,B7,20,DD,C9
```

City Bomber

To demonstrate the package, we have the game City Bomber. It is, in our totally unbiased opinion, one of the best versions of the popular "bomb the buildings to make a landing runway" games that we've seen on the Amstrad. The program features a large eight-character UDG-designed plane, a big graphical bomb and good sound effects.

The program shows what can be achieved in a relatively short time using the games writing package.

```
10 SYMBOL AFTER 256:MEMORY 39999:SYMBOL
AFTER 236
20 LOAD "",40000
30 CALL 40000
40 ENT 1,230,1,3,230,1,3
50 ENT -2,6,2,2,4,-2,2
60 RANDOMIZE TIME
70 SYMBOL 240.120.124.62.63.31.31.63.62
80 SYMBOL 241.0.0.0.0.255,255,255,253
90 SYMBOL 242,0,0,0,0,224,24,206,255
100 SYMBOL 243,31,60,112.0,0.0,0.0
110 SYMBOL 244,251,127,15,31,31,62,60.0
120 SYMBOL 245,251,255,240,192,128,0.0.0
130 SYMBOL 246.254.0.0.0.0.0.0.0
140 SYMBOL 247,231,90,231,129,66,36,60.1
26
150 SYMBOL 248,126,255,255,255,255,126,1
26.60
160 SYMBOL 249,255,255,241,241,255,241,2
41,255
170 SYMBOL 250,255,255,143,143,255,143,1
43,255
180 s$=" ":' 4 spaces
190 a$=CHR$(240)+CHR$(241)+CHR$(241)+CHR
$(242)
200 b$=CHR$(243)+CHR$(244)+CHR$(245)+CHR
(246)
```
```
210 INK 0,14:INK 1,0:INK 2,6:INK 3,24
220 BORDER 14:MODE 1
230 m$="City Bomber": BIGPRINT, 10, 1, @m$,
3.2
240 FOR n=5 TO 37 STEP 2
250 \times = 15 - INT(RND \times 10)
260 PEN 2+n/2 MOD 2
270 FOR v=x TO 25
280 LOCATE n.y:PRINT CHR$(249);CHR$(250)
290 NEXT:NEXT
300 ch%=0:cr%=0:PEN 1:LOCATE 1,1
310 PRINT a$:PRINT b$
320 SOUND 129.500.2000.3.0.0.1
330 v=1:w=1:c=1
340 IF w<23 AND INKEY(47)=0 GOTO 470
350 IF V=30 AND w=23 GOTO 1120
360 FOR n=1 TO 3:NEXT
370 IF c=1 THEN ;READCHAR,v+4,w+1,@ch%:I
F ch%>248 GOTO 920
380 |RSCROLL,v,v+4,w,w+1
390 c=c+1:IF c=9 THEN c=1:v=v+1
400 IF v<37 GOTO 340
410 LOCATE 37,w:PRINT s$
420 LOCATE 37.w+1:PRINT s$
430 LOCATE 1,w+2:PRINT a$
440 LOCATE 1.w+3:PRINT b$
450 SOUND 129,500,2000,3,0,0,1
```

```
460 v=1:w=w+2:GOT0 340
470 d=c:x=v+2:v=w+2
480 {READCHAR,x,y,@ch%:IF ch%>32 THEN t=
V:GOTO 710
490 {READCHAR, x, y+1, @ch%: {READCHAR, x, y+2
.@cr%
500 SOUND 130,30,500,5,0,1
510 LOCATE x,y:PRINT CHR$(247):LOCATE x.
y+1:PRINT CHR$(248)
520 IF ch%+cr%>64 GOTO 700
530 IF c=1 THEN {READCHAR, v+4, w+1, @ch%:I
F ch%>248 GOTO 920
540 [RSCROLL.v.v+4.w.w+1
550 (DSCROLL,x,x,y,y+2:;DSCROLL,x,x,y,y+
2
560 c=c+1:IF c=9 THEN c=1:v=v+1
570 IF v<37 GOTO 640
580 LOCATE 37.w:PRINT s$
590 LOCATE 37, w+1:PRINT s$
600 LOCATE 1,w+2:PRINT a$
610 LOCATE 1,w+3:PRINT 5$
620 SOUND 129,500,2000,3,0,0,1
630 v=1:w=w+2
640 b=v
650 \text{ y=y-(d=c)-(d-1=(c+3) MOD 8)}
660 IF b<y THEN ;READCHAR,x,y+2,@ch%:IF
ch%>32 GOTO 700
```

```
670 IF v<24 GOTO 530
680 LOCATE x,24:PRINT" ":LOCATE x,25:PRI
NT" "
690 SOUND 130.0.0.0:GOTO 340
700 t=y+2:LOCATE x,y:PRINT" ":LOCATE x+(
x MOD 2=0).v+1:PRINT" ":' 2 spaces
710 IF t<11 OR t<23 AND RND<0.7 THEN b=t
+INT((24-t)/3) ELSE b=25
720 IF RND>0.6 OR b=25 AND RND>0.3 THEN
q=4 ELSE q=8
730 m=(1+b-t)*q:z=x+1
740 IF x MOD 2=0 THEN z=x:x=x-1
750 SOUND 130,100*q,800,15,0,2
760 FOR n=1 TO m
770 ;DSCROLL.x.z.t.b
780 IF c=1 THEN !READCHAR, v+4, w+1,@ch%:I
F ch%>248 GOTO 920
790 {RSCROLL, v, v+4, w, w+1
800 IF g=4 THEN ;DSCROLL,x,z,t,b
810 c=c+1:IF c=9 THEN c=1:v=v+1
820 IF v<37 GOTO 890
830 LOCATE 37.w:PRINT s$
840 LOCATE 37, w+1:PRINT s$
850 LOCATE 1,w+2:PRINT a$
860 LOCATE 1,w+3:PRINT b$
870 SOUND 129,500,2000,3,0,0,1
880 v=1:w=w+2
```

```
890 NEXT
900 SOUND 130,0,0,0
910 GOTO 340
920 SOUND 135,0,0,0
930 [EXPLODE:INK 2,6,24:INK 3,24,6:SPEED
 INK 2.2
940 IF w>23 GOTO 1050
950 t=w:b=t+2
960 {READCHAR, v, b, @ch%: IF ch%>32 GOTO 98
õ
970 b=b+1:IF b<26 GOTO 960
980 \ b=b-1:m=(b-t-1)*8
990 FOR n=1 TO m: DSCROLL, v, v+1, t, b:NEXT
1000 t=w:b=t+2
1010 {READCHAR, v+2, b, @ch%: IF ch%>32 GOTO
 1030
1020 b=b+1:IF b<26 GOTO 1010
1030 \text{ b=b-1:m=(b-t-1)*8}
1040 FOR n=1 TO m: DSCROLL, v+2, v+3, t, b:N
EXT
1050 PEN 1:FOR m=1 TO 2000:NEXT
1060 LOCATE 12,2:PRINT"Play again? (Y/N)
...
1070 BORDER 1
1080 r$=UPPER$(INKEY$)
1090 IF r$="Y" GOTO 210
```

```
1100 IF r$<>"N" GOTO 1080
1110 END
1120 FOR n=1 TO 88
1130 SOUND 135,100-n,2,7
1140 ;RSCROLL,30,40,19,24
1150 IF n MOD 3=0 THEN ;USCROLL,30,40,19
,24
1160 NEXT
1170 PEN 2
1180 m$="CONGRATULATIONS":;BIGPRINT,6,12
,@m$,2,3:GOTO 1050
```

The spacebar drops the bombs for you and you must clear the screen of buildings before you can land and take off again. When playing the game, note the smooth scroll across the screen of the aircraft and see if you can see where the new commands have been used.

Appendices

- A Memory Map
 B Z8Ø Op Codes
 C Hexadecimal To Decimal Conversion

Appendix A

Memory Map

There are several points concerning the memory map which need to be mentioned.

Firstly, it is complicated because in the 64K memory available, 64K RAM and 32K of ROM operate – hence the bank switching. Secondly, all the addresses given on the screen map are in hex. Third, notice the six-digit address at the top of the screen map. $\emptyset 10000$ is the upper boundary and is, in fact, one above the highest actual address, which is FFFF.

010000	DEFAULT SCREEN MEMORY	Ø1ØØØØ	UPPER ROM's		
CØØØ	STACK, FIRMWARE DATA	CØØØ	(bank switcheu)		
B1ØØ	AND JUMPBLOCK				
ACØØ	FOREGROUND DATA				
????	BACKGROUND DATA				
	MEMORY POOL				
????					
????	BACKGROUND DATA				
ØØ4Ø	FOREGROUND DATA	4000			
ØØØØ	FIRMWARE AREA	ØØØØ	LOWER ROM		

Appendix B

Z8Ø Op Codes

Here is a complete list of Z8Ø Op Codes and instructions.

SUMMARY OF Z80 INSTRUCTION FUNCTIONS

ADC	A, operand l	;Add operand and carry to accumulator
ADC	HL,regpair	;Add register-pair and carry to accumulator
ADD	A, operand l	;Add operand to accumulator
ADD	HL,regpair	;Add register-pair to accumulator
ADD	IX,regpair	;Add register-pair to IX
ADD	IX,IX	;Add IX to IX
ADD	IY,regpair	;Add register-pair to IY
ADD	IY,IY	;Add IY to IY
AND	operand l	;Logically AND operand to accumulator
BIT	b,operand	;Test bit b of operand
CALL	addr	;Call address
CALL	cc.addr	:Call address if cc is valid

CCF		;Complement carry flag
СР	operand l	;Compare operand to accumulator
CPD		;Compare (HL) to accumulator, decrement HL ;And BC
CPDR		;Compare (HL) to accumulator, decrement HL ;And BC, repeat until BC = 0 or match found
СРІ		;Compare (HL) to accumulator, increment HL ;And decrement BC
CPIR		;Compare (HL) to accumulator, increment HL ;And decrement BC, repeat until BC = 0 or ;Match found
CPL		;Complement (invert) accumulator
DAA		;Decimal adjust accumulator
DEC	operand	;Decrement operand
DEC	IX	;Decrement IX
DEC	IY	;Decrement IY
DEC	SP	;Decrement stack pointer
DI		;Disable interrupts
DJNZ	dis	;Decrement B, jump to displacement if $B = 0$
EI		;Enable interrupts

.

EX	(SP),HL	;Exchange (SP) and HL
EX	(SP),IX	;Exchange (SP) and IX
EX	(SP),IY	;Exchange (SP) and IY
EX	AF,AF'	;Exchange AF with alternate AF
EX	DE,HL	;Exchange DE and HL
EXX		;Exchange general-purpose registers with ;Alternates
HALT		;Suspect CPU operation
IM	mode	;Activate interrupt mode
IN	A,(data)	;Input into accumulator port data
IN	reg,(C)	;Input into register port C
INC	operand	;Increment operand
INC	regpair	;Increment register-pair
INC [.]	IX	;Increment IX
INC	IY	;Increment IY
IND		;Load (HL) with input from port C, ;Decrement HL and B
INDR		;Load (HL) with input from port C, ;Decrement HL and B, repeat until $B = 0$

INI		;Load (HL) with input from port C, ;Decrement B and increment HL
INIR		;Load (HL) with input from port C, ;Decrement B and increment HL, repeat until ;B = 0
JP	(HL)	;Jump to address of HL
JP	(IX)	;Jump to address of IX
JP	(IY)	;Jump to address of IY
JP	addr	;Jump address
JP	cc,address	;Load program counter with address ;If cc is valid
JR	dis	;Jump displacement
JR	c,dis	;Jump displacement if c valid
LD	A,I	;Load interrupt vector into accumulator
LD	A, operand l	;Load operand into accumulator
LD	A,R	;Load refresh into accumulator
LD	(BC),A	;Load accumulator into (BC)
LD	(DE),A	;Load accumulator into (DE)
LD	(HL),data	;Load data into (HL)
LD	HL,(addr)	;Load (address) into HL

LD	(HL),reg	;Load register into (HL)	
----	----------	--------------------------	--

- LD I,A ;Load accumulator into interrupt vector
- LD IX,addr ;Load address into IX
- LD IX,(addr) ;Load (address) into IX
- LD (IX + dis),data ;Load into IX + displacement data
- LD IY,addr ;Load address into IY
- LD IY,(addr) ;Add (address) into IY
- LD (IY + dis), data ;Load into IY + displacement data
- LD (addr),A ;Load accumulator into (address)
- LD (addr), regpair ;Load register-pair into (address)
- LD (addr),IX ;Load IX into (address)
- LD (addr),IY ;Load IY into (address)
- LD regpair,addr ;Load address into register-pair
- LD R,A ;Load accumulator into refresh
- LD reg, operand 1 ;Load operand into register
- LD SP,HL ;Load HL into stack pointer
- LD SP,IX ;Load IX into stack pointer
- LD SP,IY ;Load IY into stack pointer

LDD		;Load (HL) into (DE), decrement DE, BC and HL
LDDR		;Load (HL) into (DE), decrement DE, BC and HL, ;Repeat until BC = 0
LDI		;Load (HL) into (DE), increment DE and HL, ;Decrement BC
LDIR		;Load (HL) into (DE), increment DE and HL, ;Decrement BC
NEG		;Negate accumulator
NOP		;No operation
OR	operand l	;Logically OR operand to accumulator
OTDR		;Load (HL) into port C, decrement B and HL, ;Repeat until $B = 0$
OTIR		;Load (HL) into port C, increment HL, ;Decrement B, repeat until $B = 0$
OUT	(C),reg	;Load register into port C
OUT	(data),A	;Load accumulator into port data
OUTD		;Load (HL) into port C, decrement HL and B
OUTI		;Load (HL) into port C, increment HL and ;Decrement B
POP	regpair	;Recover register-pair from stack
POP	IX	;Recover IX from stack

POP	IY	;Recover IY from stack
PUSH	regpair	;Put register-pair on stack
P U S H IX	I	;Put IX on stack
P U S H IY	[;Put IY on stack
RES	b,operand	;Reset bit b of operand
RET		;Return
RET	сс	Return if condition valid
RETI		;Return from interrupt
RETN		;Return from non maskable interrupt
RL	operand	;Rotate operand left through carry
RLA		;Rotate accumulator left through carry
RLC	operand	;Rotate operand left circular
RLCA		;Rotate accumulator left circular
RLD		;Rotate digit left and right between (HL) ;And accumulator
RR	operand	;Rotate operand right through carry
RRA		;Rotate accumulator right through carry

RRC	operand	;Rotate operand right circular
RRCA		;Rotate accumulator right circular
RRD		;Rotate digit right and left between (HL) ;And accumulator
RST	vector	;Restart at location vector
SBC	A, operand	;Subtract operand from accumulator with carry
SBC	HL,repair	;Subtract register-pair from HL with carry
SCF		;Set carry flag
SET	b,operand	;Set bit b of operand
SLA	operand	;Shift operand left arithmetic
SRA	operand	;Shift operand right arithmetic
SRL	operand	;Shift operand right logical
SUB	operand	;Subtract operand from accumulator
XOR	operand	;Exclusive OR operand to accumulator

SYMBCLS —

reg A, B, C, D, E, H, L.

regpair HL, BC, DE, SP

mode	1, 2, 3
operand l	A, B, C, D, E, H, L, (HL), (IX + dis), (IY + dis), data
operand	A, B, C, D, E, H, L, (HL), (IX + dis), (IY + dis)
dis	displacement offset, a number between 0-255
data	number between 0-255
addr	address between 0-65535
сс	conditions Z, NZ, C, NC, P, M, PE, PO
с	conditions Z, NZ, C, NC
vector	addresses 0, 8, 16, 24, 32, 40, 48 and 56

ALPHABETICAL LIST OF Z80 INSTRUCTIONS

MNEMONIC		OBJECT CODE	Fl	LAO	GS	
ADC	A,(HL)	8E	С	Z	P/V	S
ADC	A,(IX+Ø)	DDSEØØ	С	Z	P/V	S
ADC	A,(IY+@)	FD8E00	С	Z	P/V	S
ADC	A,Ø	CEØØ	с	z	P/V	s

ADC	Α,Α	8F	С	Z	P/V	S
ADC	A, P	88	С	Z	P/V	S
ADC	А,С	89	С	Z	P/V	S
ADC	A,D	8A	С	Z	P/V	S
ADC	A,E	8 B	С	Z	P/V	S
ADC	А,Н	80	С	Z	P/V	S
ADC	HL, EC	ED4A	С	Z	P/V	S
ADC	HL,DE	ED5A	С	Z	P/V	S
ADC	HL, HL	ED6A	С	Z	P/V	S
ADC	HL, SP	ED7A	С	Z	P/V	S
ADD	A,(HL)	86	C	Z	P/V	S
ADD	A,(IX+Ø)	DD8600	С	Z	P/V	S
ADD	A,(IY+Ø)	FD86 00	с	Z	P/V	S
ADD	Α,Ø	C600	С	Z	P/V	S
ADD	Α,Α	87	С	Z	P/V	5
ADD	A, B	80	С	Z	P/V	S
ADD	A,C	81	С	Z	P/V	S
ADD	A,D	82	С	Ź	P/V	S

ADD	A,E	83	C Z P/V S
ADD	А,Н	84	C Z P/V S
ADD	HL,BC	Ø 9	С
ADD	HL,DE	19	С
ADD	HL,HL	29	С
ADD	HL,SP	39	С
ADD	IX,BC	DDØ9	С
ADD	IX,DE	DD 1 9	C
ADD	IX,IX	DD29	С
ADD	IX,SP	DD39	С
ADD	IY, BC	FDØ9	С
ADD	IY,DE	FD19	С
ADD	IY,IY	FD29	C
ADD	IY,SP	FD39	С
AND	(HL)	A6	Z P/V S
AND	(IX+Ø)	DDA600	Z P/V S
AND	(IY+0)	FDA600	Z P/V S
AND	Ø	E600	Z P/V S

AND	A	A7	Z P/V S
AND	E.	AØ	Z P/V S
AND	С	A1	Z P/V S
AND	D	A2	Z P/V S
AND	E	A3	Z P/V S
AND	н	A4	Z P/V S
AND	L.	A5	Z P/V S
BIT	Ø,(HL)	CB46	Z
BIT	Ø,(IX+Ø)	DDCB 00 46	Z.
BIT	0,(IY+0)	FDCE0046	Z
BIT	Ø, A	CB47	Z
BIT	Ø, B	CB40	Z
BIT	Ø, C	CB41	Z
BIT	Ø , D	CB42	Z
BIT	Ø,E	CE43	Z
BIT	0, H	CE44	Z
BIT	Ø,L	CB45	Z
BIT	1,(HL)	CB4E	Z

BIT	1,(IX+Ø)	DDCE004E	Z
BIT	1,(IY+Ø)	FDCB004E	Z
BIT	1,A	CE4F	Z
BIT	1,B	CB48	Z
BIT	1,C	CE49	Z
BIT	1,D	CB4A	Z
BIT	1,E	CB4B	Z
BIT	1,H	CB4C	Z
EIT	1 , L.	CB4D	Z
BIT	2,(HL)	CB56	Z
BIT	2,(IX+Ø)	DDCB0056	Z
BIT	2,(IY+Ø)	FDCB0056	Z
BIT	2,A	CB57	Z
PIT	2,B	CB5Ø	Z
BIT	2,C	CB5 1	Z
BIT	2,D	CB52	Z
BIT	2,E	CB53	Z
BIT	2,H	CB54	Z

EIT	2,L	CE55	Z
BIT	3,(HL)	CE5E	Z
BIT	3,(IX+Ø)	DDCE005E	Z
BIT	3,(IY+Ø)	FDCBØØ5E	Z
BIT	3,A	CB5F	Z
BIT	З,В	CB58	Z
BIT	З,С	CE59	Z
BIT	3,D	CB5A	Z
BIT	3,E	CB5B	Z
BIT	з,Н	CB5C	Z
BIT	3,L	CE5D	Z
BIT	4,(HL)	CB66	Z
BIT	4,(IX+Ø)	DDCB0066	Z
BIT	4,A	CB67	Z
BIT	4,E	CB6Ø	Z
BIT	4,C	CE61	Z
BIT	4,D	CB62	Z
BIT	4,E	CB63	Z

BIT	4,H	CB64	Z
BIT	4,L.	CE65	Z
BIT	5,(HL)	CB6E	Z
BIT	5,(IX+@)	DDCB006E	Z
BIT	5,(IY+Ø)	FDCB006E	Z
BIT	5,A	CB6F	Z
BIT	5,B	CB68	Z
BIT	5,C	CB69	Z
BIT	5,D	CE6A	Z
BIT	5,E	CB6B	Z
BIT	5,H	CB6C	Z
EIT	5,L	CBGD	Z
BIT	6,(HL)	CE76	Z
EIT	6,(IX+Ø)	DDCB0076	Z
BIT	6,(IY+Ø)	DDCBØ076	Z
BIT	6,A	CB77	z
BIT	6,B	CB7Ø	z
BIT	6,D	CB72	Z

BIT	6,E	CE73	Z
BIT	6,H	CB74	Z
BIT	6,L	CB75	Z
BIT	7,(HL)	CE7E	Z
BIT	7,(IX+Ø)	DDCB007E	Z
BIT	7,(17+0)	FDCB007E	Z
BIT	7,A	CE7F	Z
BIT	7,E	CB78	Z
BIT	7,C	CB79	Z
ΒIΤ	7,D	CB7A	Z
BIT	7,E	CE7E	Z
BIT	7,H	CB7C	Z
EIT	7,L	CE7D	Z
CALL	Ø	CDØØØØ	
CALL	С,Ю	DC ØØ ØØ	
CALL	Μ,Ø	FCØØØØ	
CALL	NC,Ø	D4000	
CALL	NZ,Ø	C40000	

CALL	Ρ,Ø	F40000				
CALL	PE,Ø	EC0000				
CALL	P0,0	E40000				
CALL	Z,Ø	CCØØØØ				
CCF		3F	С			
CP	Ø	FEØØ	С	Z	P/V	S
СР	(HL)	BE	С	Z	P/V	S
СР	(IX+Ø)	DDEEØØ	С	Z	P/V	S
СР	(IY+Ø)	FDBEØØ	С	Z	P/V	S
СР	A	BF	С	Z	P/V	S
СР	В	B 8	С	Z	P/V	S
СР	С	B 9	С	Z	P/V	S
СР	D	BA	с	Z	P/V	S
СР	E	BB	с	Z	P/V	s
CP	н	BC	с	Z	P/V	s
СР	L	BD	с	Z	P/V	S
CPD		EDA9		Z	P/V	S
CPDR		EDE9		Z	P/V	S

CPI		EDA1		Ζ	P/V	S
CPIR		EDB1		Z	P/V	S
CPL		2F				
DAA		27	С	Z	P/V	S
DEC	(HL)	35		Z	P/V	S
DEC	(IX+Ø)	DD3500		Z	P/V	S
DEC	(IY+Ø)	FD3500		Z	P/V	S
DEC	A	3D		Z	P/V	S
DEC	B	05		Z	P/V	S
DEC	BC	ØB				
DEC	С	ØD		Z	P/V	S
DEC	D	15		Z	P/V	S
DEC	DE	1 E				
DEC	E	1 D		Z	P/V	S
DEC	Н	25		z	P/V	S
DEC	HL	2 B				
DEC	IX	DD2B				
DEC	IY	FD2B				

DEC	L	2D		Z	P/V	S
DEC	SP	3B				
DI		F3				
DJNZ	Ø	1000				
EI		FB				
ΕX	(SP),HL	E3				
EX	(SP),IX	DDE3				
ΕX	(SP),IY	FDE3				
ΕX	AF, AF'	Ø8	с	Z	P/V	S
EX	DE,HL	EB				
EXX		D9				
HAL.T		76				
IM	Ø	ED46				
IM	1	ED56				
IM	2,	ED5E				
IN	A,(Ø)	DBØØ				
IN	A,(C)	ED78	7	Ζ	P/V	S
IN	B, (C)	ED4Ø	-	Z	P/V	S

IN	C,(C)	ED48	Z	P/V	S
IN	D,(C)	ED50	Z	P/V	S
IN	E,(C)	ED58	Z	P/V	S
IN	Н,(С)	ED6Ø	Z	P/V	S
IN	L,(C)	ED68	Z	P/V	S
INC	(HL)	34	Z	P/V	S
INC	(IX+Ø)	DD3400	Z	P/V	S
INC	(IY+Ø)	FD3400	Z	P/V	S
INC	A	30	Z	P/V	S
INC	P.	Ø4	Z	P/V	S
INC	BC	Ø 3			
INC	С	ØC	Z	P/V	S
INC	D	14	Z	P/V	S
INC	DE	13			
INC	E	1 C	Z	P/V	S
INC	Н	24	Z	P/V	S
INC	HL.	23	Z	P/V	5
INC	IX	DD23			

INC	IY	FD23			
INC	L	20	Z	P/V	S
INC	SP	33			
IND		EDAA	Z		
INDR		EDBA			
INI		EDA2	Z		
INIR		EDB2			
JP	Ø	C30000			
JF	(HL)	E9			
JP	(IX)	DDE9			
JP	(IY)	FDE9			
JF	С,Ø	DAØØØØ			
JP	М,Ø	FA0000			
JP	NC,Ø	D20000			
JP	NZ,Ø	C20000			
JP	Ρ,Ø	F20000			
JP	PE,Ø	EA0000			
JP	PO,0	E20000			

JP	Ζ,Ø	CAØØØØ
JR	Ø	1800
JR	с, Ø	3800
JR	NC,Ø	3000
JR	NZ,Ø	2000
JR	Ζ,Ø	2800
LD	(BC),A	Ø 2
LD	(DE),A	12
LD	(HL),Ø	3600
LD	(HL),A	77
LD	(HL),B	70
LD	(HL),C	71
LD	(HL),D	72
LD	(HL),E	73
LD	(HL),H	74
LD	(HL),L	75
LD	(IX+0),0	DD360000
LD	(IX+0),A	DD7700

LD	(IX+Ø),B	DD7000
LD	(IX+Ø),C	DD7100
LD	(IX+0),D	DD7200
LD	(IX+Ø),E	DD7300
LD	(IX+0),H [`]	DD7400
LD	(IX+0),L	DD7500
LD	(IY+0),0	FD360000
LD	(IY+0),A	FD7700
LD	(IY+0),B	FD7000
LD	(IY+Ø),C	FD7100
LD	(IY+0),D	FD7200
LD	(IY+Ø),E	FD7300
LD	(IY+0),H	FD7400
LD	(IY+Ø),L	FD7500
LD	(0000),A	320000
LD	(0000),EC	ED430000
LD	(0000),DE	ED530000
LD	(0000),HL	220000

LD	(0000),IX	DD220000
LD	(0000),IY	FD220000
LD	(0000),SP	ED730000
LD	A, (BC)	ØA
LD	A,(DE)	1A
LD	A,(HL)	7E
LD	A,(IX+Ø)	DD7E00
LD	A,(IY+Ø)	FD7EØØ
LD	A, (0000)	3AØØØØ
LD	A,Ø	3E00
LD	Α,Α	7F
LD	А,В	78
LD	A, C	79
LD	A,D	7 A
LD	Α,Η	7¢
LD	Α,Ι	ED57
LD	Α,L.	70
LD	A,R	ED5F

Z P/V S

Z P/V S

LD	B,(HL)	46
LD	B,(IX+∅)	DD4600
LD	E, (IY+Ø)	FD4600
LD	E , Ø	0600
LD	Е,А	47
LD	E, B	40
LD	Е,С	41
LD	E,D	42
LD	B,E	43
LD	в,н	44
LD	E., L.,	45
LD	BC,(0000)	ED4B0000
LD	BC,0000	010000
LD	C,(HL)	4E
LD	C,(IX+Ø)	DD4EØØ
LD	C,(IY+Ø)	FD4EØØ
LD	с,ø	ØEØØ
LD	С,А	4F

LD	С, В	48
LD	с,с	49
LD	C,D	4A
LD	C,E	4 B
LD	с,н	4C
LD	C,L	4D
LD	D, (HL.)	56
LD	D,(IX+Ø)	DD56 00
LD	D,(IY+Ø)	FD5600
LD	D,Ø	1600
LD	D,A	57
LD	D, P	50
LD	D,C	51
LD	D,D	52
LD	D,E	53
LD	D,H	54
LD	D,L	55
LD	DE,(0000)	ED580000

LD	DE,0000	110000
LD	E,(HL)	5E
LD	E,(IX+Ø)	DD5EØØ
LD	E,(IY+Ø)	FD5EØØ
LD	Ε,Ø	1EØØ
LD	E,A	5F
LD	Е, Б	58
LD	E,C	59
LD	E,D	5A
LD	E,E	5B
LD	E,H	5 C
LD	E,L	5D
LD	H,(HL)	66
LD	H,(IX+Ø)	DD66 00
LD	H,(IY+Ø)	FD6600
LD	Н,0	2600
LD	Н,А	67
LD	н,в	60
LD	H,C	61
----	-----------	----------
LD	н, D	62
LD	H,E	63
LD	н,н	64
LD	Η,L	65
LD	HL,(0000)	2A0000
LD	HL,0000	210000
LD	Ι,Α	ED47
LD	IX,(0000)	DD2A0000
LD	IX,0000	DD210000
LD	IY,(0000)	FD2A0000
LD	IY,0000	FD210000
LD	L,(HL)	6E
LD	L,(IX+Ø)	DD6EØØ
LD	L,(IY+Ø)	FD6EØØ
LD	L,Ø	2EØØ
LD	L,A	6F
LD	L, B	68

LD	L,C	69				
LD	L,D	6A				
LD	L,E	6B				
LD	L,H	6C				
LD	L,L	4D				
LD	R,A	ED4F				
LD	SP,(0000)	ED7B0000				
LD	SP,0000	310000				
LD	SP,HL	F9				
LD	SP,IX	DDF9				
LD	SP,IY	FDF9				
LDD		EDA8			P/V	
LDDR		EDB8				
LDI		EDAØ			PZV	
LDIR		EDPØ				
NEG		ED44	С	Z	P/V	S
NOP		00				
OR	(HI_)	B6		Z	P/V	g

OR	(IX+Ø)	DDB&ØØ	Z	P/V	S
OR	(17+0)	FDB6 00	Z	P7V	9
OR	Ø	F600	Z	P/V	S
OR	А	B7	Z	P/V	8
OR	В	БØ	Z	Ρ/V	9
OR	С	B-1	Z	P/V	S
OR	D	B2	Ζ	P/V	S
OR	E	ЕЗ	Z	P7V	S
OR	Н	B4	Ζ	₽/V	S
OR	L.	P 5	Z	P/V	S
OTDR		ED8B			
OTIR		EDB3			
ουτ	(C),A	ED79			
ουτ	(C),B	ED41			
OUT	(C),C	ED49			
OUT	(C),D	ED5 1			
OUT	(C),E	ED59			
OUT	(C),H	ED61			

OUT	(C),L	ED69				
OUT	(Ø),A	D300				
ουτρ		EDAB		Z		
OUTI		EDA3		Z		
POP	AF	F1	С	Z	P/V	S
POP	ВC	C1				
POP	DE	D1				
POP	HL	E1				
POP	IX	DDE 1				
POP	IY	FDE1				
PUSH	AF	F5				
PUSH	BC	C5				
PUSH	DE	D5				
PUSH	H1L	ES				
PUSH	IX	DDE5				
PUSH	IΥ	FDE5				
RES	0,(HL)	CB86				
RES	0,(IX+@)	DDCE0086				

RES	Ø,(IY+Ø)	FDCBØØ86
RES	0,A	CB87
RES	0, B	CB80
RES	Ø, C	CB81
RES	Ø, D	C B 82
RES	Ø,E	CB83
RES	Ø,H	CE 84
RES	Ø,L.	CB85
RES	1,(HL)	CESE
RES	1,(IX+Ø)	DDCB 00 8E
RES	1,(IY+Ø)	FDCBØØ8E
RES	1 ,A	CE8F
RES	1, B	CB88
RES	1,C	CE89
RES	1,D	CE8A
RES	1,E	CBSB
RES	1, H	CESC
RES	1 , L.	CESD

RES	2,(HL)	CB96
RES	2,(IX+@)	DDCB0096
RES	2,(IY+Ø)	FDCBØØ96
RES	2,A	CE97
RES	2,B	CB9Ø
RES	2,C	CE91
RES	2,D	CB92
RES	2,E	CB93
RES	2,H	CE94
RES	2,L	CB95
RES	3,(HL)	CB9E
RES	3,(IX+Ø)	DDCB009E
RES	3,(IY+Ø)	FDCB009E
RES	З,А	CB9F
RES	3,E	CB98
RES	3,C	CB99
RES	3,D	СВ9А
RES	3,E	CB9B

RES 3,H CB9C RES 3,L CE9D RES 4,(HL) CBA6 RES 4, (IX+Ø) DDCE00A6 RES 4, (IY+Ø) FDCB00A6 RES 4, A CBA7 RES CEAØ 4, P RES 4,C CBA1 RES 4,D CBA2 RES 4,E CBA3 4,H CEA4 RES RES 4, [... CEA5 RES 5,(HL) CBAE RES 5,(IX+Ø) DDCEØØAE RES FDCEØØAE 5,(IY+0) RES 5,A CBAF RES 5,E CBAS RES 5,C CBA9

RES	7,L	CBBD				
RET		C9				
RET	с	DB				
RET	М	F8				
RET	NC	DØ				
RET	NZ	CØ				
RET	P	FØ				
RET	PE	E8				
RET	PO	EØ				
RET	2	C8				
RETI		ED4D				
RETN		ED45				
RL	(HL)	CE16	С	Z	P/V	S
RL.	(IX+Ø)	DDCEØØ1ć	С	Z	P/V	S
RL	(IY+Ø)	FDCB0016	С	Z	P/V	9
RL.	A	CE17	С	Ζ	P/V	S
RL	В	CE 1 Ø	С	Z	P/V	S
RL	с	CB11	С	Z	P/V	S

RL.	D	CE12	С	Z	P/V	S
RL	E	CE13	С	Z	P/V	S
RL	н	CB14	С	Z	₽/V	S
RL.	L.	CB15	С	Z	P/V	S
RLA		17	С			
RES	5,D	CBAA				
RES	5,E	CBAB				
RES	5,H	CBAC				
RES	5,L.	CBAD				
RES	6,(HL)	CBB6				
RES	6,(IX+Ø)	DDCBØØEG				
RES	6,(IY+Ø)	FDCBØØB6				
RES	6,A	CBB7				
RES	6, E	CBBØ				
RES	6,C	CFB1				
RES	6,D	CBB2				
RES	6,E	CBBB				
RES	6,H	CEB4				

RES	6,L	CBB5					
RES	7,(HL)	CBFE					
RES	7,(IX+Ø)	DDCB00EE					
RES	7,(IY+Ø)	FDCEØØEE					
RES	7,A	CEEF					
RES	7,E	CEBS					
RES	7,C	CEE9					
RES	7,D	CBBA					
RES	7 , E	CEEE					
RES	7,H	CBBC					
RLC	(HL.)	CBØ6	С	Z	P/V	S	
RL C	(IX+Ø)	DDCB0016	С	Z	P/V	S	
RLC	(IY+Ø)	FDCB0016	С	Z	P/V	S	
RLC	А	CBØ7	С	Z	P/V	S	
RLC	B	CEØØ	С	Z	P/V	S	
RLC	С	CEØ1	С	Z	P/V	S	
RLC	D	CBØ2	С	Z	P/V	S	
RL C	E	CBØ3	С	Z	P/V	S	

RLC	Н	CBØ4	С	Ζ	P/V	S
RLC	L	CBØ5	С	Z	P/V	S
RLCA		07	ç			
RLD		ED6F		Z	P/V	S
RR	(HL)	CBIE	С	Z	P/V	S
RR	(IX+Ø)	DDCBØØ1E	С	Z	P/V	S
RR	(17+0)	FDCB001E	С	Ζ	P/V	S
FR	A	CB1F	С	Ζ	₽7V	S
RF	F	CB18	С	Z	P/V	S
RR	С	CE19	С	Z	₽/V	S
RR	D	CE-20	С	Z	P/V	S
RR	E	CB21	С	Z	P/V	S
RR	Н	CB22	С	Z	P/V	S
RR	L.	CB23	С	Z	P/V	S
RRA		1F	С			
RFC	(HL)	CBØE	С	Z	P/V	S
RRC	(1×+Ø)	DDCE000E	С	Z	P/V	S
RRC	(IY+Ø)	FDCBØØØE	С	Z	P/V	S

RRC	Α	CBØF	С	Z	P/V	S
RRC	B	СВØЗ	С	Z	P/V	S
RRC	С	СВØЭ	С	Z	P/V	S
RRC	D	CBØA	С	Z	PZV	S
RRC	E	CBØF	С	Z	P/V	S
RRC	н	CEØC	С	Z	P/V	S
RRC	L	CBØD	С	Z	P/V	S
RRCA		ØF	С			
RRD				Z	P/V	S
RST	00H	C7				
RST	0 8H	CF				
RST	1ØH	D7				
RST	18H	DF				
RST	20H	E7				
RST	23H	EF				
RST	30H	F7				
RST	38H	FF				
SBC	Α,Ø	DEØØ	С	Z	P/V	S

SEC	A,(HL)	φE	С	Z	P/V	S
SBC	A,(IX+0)	DD9EØØ	С	Z	P/V	S
SBC	A,(IY+@)	FD9EØØ	С	Z	P/V	S
SEC	Α, Α	9F	С	Z	۶/۷	S
SBC	A,B	98	С	Z	P7V	S
SBC	Α, Ο	99	С	7.	P/V	S
SBC	A,D	9A	С	Ζ	P/V	S
SEC	A,E	9B	С	Ζ	P7V	S
SEC	А,Н	90	С	Z	P/V	S
SBC	A, L.	9D	C	Z	P/V	S
SBC	HL,BC	ED42	C	Z	P/V	S
SBC	HL, DE	ED52	С	Z	P/V	S
SBC	HL, HL	ED62	С	Z	P/V	S
SBC	HL,SP	ED72	С	Z	P/V	S
SCF		37	C			
SET	0,(HL)	CBC6				
SET	Ø,(IX+Ø)	DDCB00C6				
SET	Ø, (JY+Ø)	FDCB00C6				

SET	Ø ,A	CEC7
SET	Ø, B	CECØ
SET	Ø, C	CEC1
SET	Ø, D	CB02
SET	Ø,E	CBC3
SET	Ø,H	CBC4
SET	Ø,L	CBC5
SET	1,(HL)	CBCE
SET	1,(IX+Ø)	DDCE00CE
SET	1,(IY+@)	FDCBØØCE
SET	1 ,A	CBCF
SET	1 , E	CBC8
SET	1 , C	CBC9
SET	1,D	CBCA
SET	1,E	CBCE
SET	1,H	CBCC
SET	1 , L	CECD
SET	2,(HL)	CEDA

SET	2,(IX+Ø)	DDCB00D6
SET	2,(17+0)	FDCB00D6
SET	2,A	CBD7
SET	2,B	CBDØ
SET	2,C	CBD 1
SET	2,D	CBD2
SET	2,E	CBD3
SET	2,H	CBD4
SET	2,L	CED5
SET	3,(HL)	CBDE
SET	3,(IX+Ø)	DDCBØØDE
SET	3,(IY+0)	FDCBØØDE
SET	З,А	CBDF
SET	3, B	CBD8
SET	З,С	CBD9
SET	3,D	CEDA
SET	3,E	CEDB
SET	З,Н	CBDC

SET	3,1	CEDD
SET	4,(HL)	CBE6
SET	4,(IX+0)	DDCB00E6
SET	4,(IY+Ø)	FDCB00E6
SET	4,A	CBE7
SET	4, P	CBEØ
SET	4,C	CBE 1
SET	4,D	CBE2
SET	4,E	CBE3
SET	4 , H	CBE4
SET	4,L	CBE5
SET	5,(HL)	CBEE
SET	5,(IX+@)	DDCEØØEE
SET	5,(IY+Ø)	FDCBØØEE
SET	5,A	CBEF
SET	5,E	CBE8
SET	5,C	CBE9
SET	5,D	CBEA

SET 5,H CBEE CBEC 5,L SET CBF6 SET 6, (HL) SET 6.(IX+0) DDCB00F6 SET 6, (IY+0) FDCBØØF6 SET 6,A CBF7 SET 6, B CBFØ SET 6,C CBF1 SET 6,D CBF2 SET 6,E CBF3 SET 6,H CEF4 CBF5 SET 6, L. 7,(HL) CEFE SET SET 7,(1X+0)**DDCBØØFE** 7,(IY+Ø) FDCBØØFE SET SET 7,A CEFF 7,B CBF8 SET SET 7,C CBF9

SET	7,D	CBFA				
SET	7,E	CBFE				
SET	7,H	CBFC				
SET	7,L	CBFD				
SLA	(HL)	CB26	С	Z	P/V	S
SLA	(IX+Ø)	DDCB0026	С	Z	P/V	S
SLA	(IY+Ø)	FDCB0026	С	Z	P/V	S
SLA	А	CB27	C	Z	P/V	9
SLA	P.	CB 20	С	Z	P/V	S
SLA	С	CB ⊇1	С	Z	P/V	S
SLA	D	CB22	С	Ζ	P/V	S
SLA	E	CB23	С	Z	P/V	9
SLA	н	CB24	с	Z	P/V	S
SLA	L.	CB25	С	Z	P/V	S
SRA	(HL.)	CB2E	С	Z	P/V	S
SRA	(IX+Ø)	DDCBØØ2E	С	Z	P/V	S
SRA	(IY +0)	FDCB002E	С	Z	P/V	S
SRA	A	CB2F	С	Z	P/V	s

SRA	P.	CB28	С	Z	P/V	S
SRA	С	CB29	С	Z	P/V	S
SRA	D	CB2A	С	Z	P/V	S
ŞRA	E	CE 2B	С	Z	P/V	9
SRA	Н	CE2C	С	Z	P/V	S
SRA	L	CE20	С	Z	P/V	S
SRL	(HL)	CBBE	С	Z	P/V	S
SRL.	(IX+Ø)	DDCE003E	С	Z	P/V	S
SRL.	(IY+ 0)	FDCB003E	С	Z	P/V	S
SRL	А	CB3F	C	Z	P/V	S
SRL	P	CB38	с	Z	P/V	S
SRL	C	CB39	С	Z	P/V	S
SRL.	D	CBBA	С	Z	P/V	S
SRL	E	CB3B	С	Z	P/V	S
SRL.	н	CB3C	С	Z	P/V	S
SRL	L	CB3D	С	Z	P/V	S
SUB	Ø	D600	с	z	P/V	S

SUB	(HL)	96	С	Z	P/V	S
SUE	(IX+Ø)	DD9600	С	Z	P/V	S
SUB	(IY+Ø)	FD9600	С	Z	P/V	S
SUB	Α	97	С	Z	P/V	S
SUB	E	90	С	Z	P/V	S
SUB	с	91	С	Z	P/V	S
SUB	D	92	с	Z	P/V	S
SUE	E	93	с	Z	P/V	S
SUE	н	94	С	Z	P/V	S
SUB	L	95	с	Z	P/V	S
XOR	(HL)	AE		Z	P/V	S
XOR	(IX+Ø)	DDAEØØ		Z	P/V	S
XOR	(IY+Ø)	FDAEØØ		Z	P/V	S
XOR	Ø	EEØØ		Z	P/V	S
XOR	A	AF		Z	P/V	S
XOR	E	AS		Z	P/V	S
XOR	С	A9		Z	P/V	S
XOR	D	AA		z	P/V	S

XOR	E	AB	Z	P/V	S
XOR	Н	AC	Z	P/V	S
XOR	L.	AD	Z	P/V	S

Appendix C

Hexadecimal to Decimal Conversion

Hex	0	1	2	3	4	5	6	7	8	9	Α	B	С	D	Ε	F
0	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
1	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31
2	32	33	34	35	36	37	38	39	40	41	42	43	44	45	46	47
3	48	49	50	51	52	53	54	55	56	57	58	59	60	61	62	63
4	64	65	66	67	68	69	70	71	72	73	74	75	76	77	78	79
5	80	81	82	83	84	85	86	87	88	89	90	91	92	93	94	95
6	96	97	98	99	100	101	102	103	104	105	106	107	108	109	110	111
7	112	113	114	115	116	117	118	119	120	121	122	123	124	125	126	127
8	128	129	130	131	132	133	134	135	136	137	138	139	140	141	142	143
9	144	145	146	147	148	149	150	151	152	153	154	155	156	157	158	159
A	160	161	162	163	164	165	166	167	168	169	170	171	172	173	174	175
B	176	177	178	179	180	181	182	183	184	185	186	187	188	189	190	191
C	192	193	194	195	196	197	198	199	200	201	202	203	204	205	206	207
D	208	209	210	211	212	213	214	215	216	217	218	219	220	221	222	223
E	224	225	226	227	228	229	230	231	232	233	234	235	236	237	238	239
F	240	241	242	243	244	245	246	247	248	249	250	251	252	253	254	255

Now you can really get to grips with machine code on your Amstrad. Written by Clive Gifford and Scott Vincent, two expert programmers who've proved in previous books they can write clear, easy-to-understand English, this book is a certain path to machine code mastery.

The book also includes an invaluable collection of machine code routines (including pixel by pixel scrolling in four directions) which you can incorporate into your BASIC programs, even if you haven't a clue how they work! These 'pre-packaged' machine code routines are designed to make creating arcade and animated games as simple as possible. A complete arcade game is also included in the book to show them in action.

Chapter headings include:

 Your first machine code program

 Passing parameters to and from machine code
 Simple arithmetic
 Stacking and jumping
 Logical operators, and manipulating bits
 The Amstrad's screen and ROM routines
 The Games Writing Package
 Z80 Op Codes

Designed to work with the 464, 664 and 6128 models, this book will get you into the driver's seat in a matter of hours. You can now command the wizardry of machine code magic for yourself.











jı



